# Cayley graphs of given degree and diameter on linear groups

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#### Overview

Degree/diameter

Graph lifting and Cayley graphs

3 Computer search of graphs

#### Motivation

- In it's simplest form, networks can be modeled by graphs with nodes as vertices and links between them as edges.
- In design of graphs we can take many restrictions into acount such degree, grith, diameter.
- Two important problems concerning degree and diameter and degree and grith of graph

## The degree/diameter problem

#### Degree/diameter problem

Find graph with biggest possible number of vertices with given degree and diameter.

#### Moore bound

There is theoretical upper bound for largest order of graph with d-degree and k-diameter.

$$n_{d,k} \le M_{d,k} = 1 + d + d(d-1) + \dots + d(d-1)^{k-1}$$

$$= 1 + d(1 + (d-1) + \dots + (d-1)^{k-1})$$

$$= \begin{cases} 1 + d\frac{(d-1)^k - 1}{d-2}, & \text{if } d > 2\\ 2k + 1, & \text{if } d = 2 \end{cases}$$

$$(1)$$

#### Moore bound

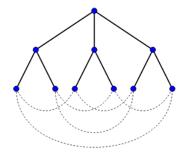


Figure: Peterssen graph is Moore graph with d=3 and k=2

## Moore graphs

Graphs with order equal Moore bound are called Moore graphs and are reached only in few cases.

- If d=2 for any  $k\geq 1$
- If k = 1 for any  $d \ge 2$
- For k = 2 for  $d \in \{3,7\}$ , and possibly 57

For other cases we try to construct graphs with order as close to Moore bound as possible.

## Moore graphs

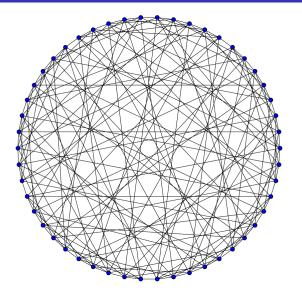


Figure: Hoffman-singleton graph is Moore graph with d=7 and k=2

## Graph lifting

Let G be an undirected graph. We will assign direction to every edge of graph and making them *arcs*. Arc with reversed direction of e is denoted by  $e^{-1}$ .

#### Definition (Graph lifting)

Let G be a graph as above and let  $\Gamma$  be a finite group. The mapping

$$\alpha: D(G) \to \Gamma$$

will be called a *voltage assignment* if  $\alpha(e^{-1}) = (\alpha(e))^{-1}$ , for any arc  $e \in D(G)$ .

## Graph lifting example

Obrazok zdvihu na petersenov graf.

## Cayley graphs

Let  $\Gamma$  be a group and let  $S \subset \Gamma$  be a symmetric unit-free generating set for  $\Gamma$ ; that is, we require that  $S = S^{-1}$  and  $1 \notin S$ .

#### Definition (Cayley graphs)

The Cayley graph  $C(\Gamma, S)$  is the graph with vertex set  $\Gamma$  in which vertices a, b are adjacent if  $a^{-1}b \in S$ .

## General linear and Special linear groups

#### Definition (General linear group)

Let q be a power of a prime and let GF(q) be the Galois field of order q. The general linear group GL(m,q) consists of all non-singular  $m \times m$  matrices over GF(q) under multiplication of matrices. Special linear group is subgroup of GL(m,q) consisting of matrices with determinant equal to 1.

#### Theorem (Order of GL(m, q))

$$|GL(m,q)| = (q^m - 1)(q^m - q) \cdots (q^m - q^{n-1})$$

## Theorem (Order of SL(m, q))

$$|SL(m,q)| = |GL(m,q)|/(q-1)$$

## Computer search of Cayley graphs

In order to find graphs with given d and k we generted Cayley graphs from random symmetric unit-free sets S of SL(2,5). With assumptions that:

- Generating sets yield all elements of |SL(2,5)| = 120
- Cayley graphs are regular with with d = |S|

we know by Moore bound that size of S must be at least 12 and all we have to check is diameter of graph.

## Computer search of Cayley graphs

With one involution in SL(2,5) and with unit-free property of symmetry we have 59 pairs of elements with its inverses. To check all graphs for d=12 we have to generate and check diameter on  $\binom{59}{6}=45057474$  graphs. In our computation we generated 150000 graphs for all listed degrees.  $\frac{d}{6}=\frac{13}{6}=\frac{14}{6}=\frac{15}{6}=\frac{17}{6}=\frac{18}{6}=\frac{19}{6}=\frac{21}{6}=\frac{19}$ 

## Algorithm for generation of cayley graph

Algoritmus na generovanie grafu

## Algorithm for diameter

Algoritmus na kontrolu priemeru grafu

## Output of program

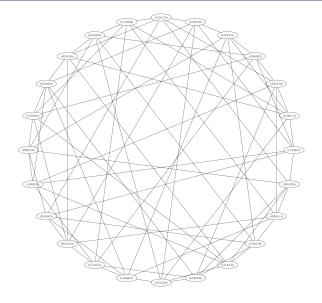


Figure: One of the generated graphs from SL(2,3)

#### Further research

- Finding smallest graph with given girth and degree is konwn as Degree/girth problem. With small changes to program we can systematically or randomly look for Cayley graphs and check its girth.
- Computations to find smallest graphs with k=2 in bigger Galois fields .

## The End