CS 336 – EXTRA CREDIT HOMEWORK

Problem 1 Suppose that we have the following four tuples in a relation S with three attributes ABC: (1,2,3), (4,2,3), (5,3,3), (5,3,4). Which of the following functional (\rightarrow) and multivalued $(\rightarrow\rightarrow)$ dependencies can you infer does *not* hold over relation S?

- 1. $A \rightarrow B$
- $A \longrightarrow B$
- 3. $BC \rightarrow A$
- 4. $BC \rightarrow \rightarrow A$
- 5. $B \rightarrow C$
- 6. $B \rightarrow \rightarrow C$

Problem 2 Consider a relation R with five attributes ABCDE.

- 1. For each of the following instances of R, state whether it violates (a) the FD $BC \rightarrow D$ and (b) the MVD $BC \rightarrow D$:
 - (a) { } (i.e., empty relation)
 - (b) $\{(a,2,3,4,5), (2,a,3,5,5)\}$
 - (c) $\{(a,2,3,4,5), (2,a,3,5,5), (a,2,3,4,6)\}$
 - (d) $\{(a,2,3,4,5), (2,a,3,4,5), (a,2,3,6,5)\}$
 - (e) $\{(a,2,3,4,5), (2,a,3,7,5), (a,2,3,4,6)\}$
 - (f) $\{(a,2,3,4,5), (2,a,3,4,5), (a,2,3,6,5), (a,2,3,6,6)\}$
 - (g) $\{(a,2,3,4,5), (a,2,3,6,5), (a,2,3,6,6), (a,2,3,4,6)\}$
- 2. If each instance for R listed above is legal, what can you say about the FD $A \rightarrow B$?

Problem 3 Consider the following actions taken by transaction T1 on database objects X and Y:

- 1. Give an example of another transaction T2 that, if run concurrently to transaction T without some form of concurrency control, could interfere with T1.
- 2. Explain how the use of *Repeatable Read* Isolation Level would prevent interference between the two transactions.

Problem 4 Consider a database with objects X and Y and assume that there are two transactions T1 and T2. Transaction T1 reads objects X and Y and then writes object X. Transaction T2 reads objects X and Y and then writes objects X and Y.

- 1. Give an example schedule with actions of transactions T1 and T2 on objects X and Y that results in a write-read conflict.
- 2. Give an example schedule with actions of transactions T1 and T2 on objects X and Y that results in a read-write conflict.
- 3. Give an example schedule with actions of transactions T1 and T2 on objects X and Y that results in a write-write conflict.
- 4. For each of the three schedules show that the use of *Repeatable Read* Isolation Level disallows the schedule.