# CompSci 150 Homework 2

## Theo Koss

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- 1. (a) Similar: They both temporarily store data. They both can be accessed by the CPU, and they are both "volatile", meaning they lose their data when the computer's power is turned off.
  - (b) Difference: They differ in capacity, and access speed. Registers fast but limited in size, main memory larger capacity but slower. They also differ in their role, registers serve as storage for data being changed by the CPU, when a program is excecuted. Main memory cells serve as the main storage for programs.
- 2. (a) 0x2304:

0010 0011 0000 0100

(b) Opcode of 0xB2A5:

B = 1011

(c) Operand of 0xB2A5:

 $2A5 = 0010 \ 1010 \ 0101$ 

3. Start:  $0x98 = 152_{10}$  end:  $0xA2 = 162_{10}$ . There are 11 memory cells in this block. The addresses are counting

up in hex from 0x98 to 0xA2:

0x98 0x99 0x9A 0x9B 0x9C 0x9D 0x9E 0xA0 0xA1 0xA2

- 4. Value of the program counter after 0xB0CD is the operand, CD.
- 6. (a) First, load each value from memory into its own register.
  - (b) Add together the registers containing x and y (only 2 at a time).
  - (c) Add the result to the register containing z.
  - (d) Store the result in memory.

For (2x) + y, load each value into its own register, multiply the register containing x by 2, then add the result to the register containing y. Then store the result in memory.

#### 7. Translate to English:

- (a) 0x7123: OR the bit patterns in registers 2 and 3, placing the result in register 1.
- (b) 0x40E1: MOVE the bit pattern in register E to register 1.
- (c) OxA304: ROTATE the bit pattern in register 3 to the right, 4 times.
- (d) OxB100: JUMP to the instruction at address 00 if register 1 has the same bit pattern as register 0. Otherwise, continue.

- (e) 0x2BCD: LOAD register B with bit pattern CD.
- 9. Translate from English to Vole:
  - (a) 0x2677
  - (b) 0x1677
  - (c) 0xBA24
  - (d) 0xA403
  - (e) 0x81E2
- 11. Classify whether the instruction changes the contents of the memory cell at location 0x3C, retrieves the contents of 0x3C, or is independent of 0x3C.
  - (a) 0x353C: Changes.
  - (b) 0x253C: Retrieves.
  - (c) 0x153C: Retrieves.
  - (d) 0x3C3C: Changes.
  - (e) 0x403C: Independent.
- 12. (a) Translate first instruction: 0x2655 = LOAD register 6 with the bit pattern 55.
  - (b) If PC started containing 0x00, what bit patterns is in register 0x6 when the machine halts?: Bit pattern: 55.
- 13. (a) First instruction: 0x1221 = LOAD register 2 with the bit pattern found in 21.
  - (b) First instruction starting with 0x01: 0x2134 = LOAD register 1 with bit pattern 34.
- 15. (a) When the machine halts, the bit pattern found in register A will be in address 0x00.
  - Register A contains: ADD bit patterns of register B and C. B contains: bit pattern 03. C also contains bit pattern 03. So A is the sum of 0x03 and 0x03.
  - (b) Program counter has bit pattern 0x0A.

- 18. Value in register 0x0:  $0010 + 0010 = 0100 = 4_{10}$ .
- 23. Write a prograam in Vole to perform:
  - (a) Move the value at memory location 0xD8 to memory location 0xB3.

Address	Contents
0x00	0x11
0x01	0xD8
0x02	0x31
0x03	0xB3
0x04	0xC0
0x05	0x00

(b) Interchange the values stored at memory locations 0xD8 and 0xB3.

Address	Content
0x00	0x11
0x01	0xD8
0x02	0x12
0x03	0xB3
0x04	0x31
0x05	0xB3
0x06	0x32
0x07	0xD8
0x08	0xC0
0x09	0x00

(c) If the value stored in 0x44 is 0x00, then place the value 0x01 in memory location 0x46, otherwise put the value 0xFF in memory

### location 0x46.

Address	Conten	$\operatorname{ts}$
0x00	0x11	
0x01	0x44	(Load reg. 1 with value of mem. loc. 0x44)
0x02	0x20	
0x03	0x00	(Load reg. 2 with bit pattern 0x00)
0x04	0x23	
0x05	0x01	(Load reg. 3 with bit pattern 0x01)
0x06	0xB1	
0x07	0x0E	
(Check if reg. 1 is equa	al to reg.	0, if true, jump to 0x0E. If false, continue.)
0x08	0x24	
0x09	0xFF	(Load reg. 4 with bit pattern 0xFF)
0x0A	0x34	
0x0B	0x44	(Store reg. 4 to mem. loc. $0x44$ )
0x0C	0xC0	
0x0D	0x00	(Halt)
0x0E	0x33	
0x0F	0x44	(Store reg. 3 to mem. loc. 0x44)
0x10	0xC0	
0x11	0x00	(Halt)

34. (a)

 $\begin{array}{c}
111001 \\
\text{AND } 101001 \\
=101001
\end{array}$ 

c.

$$001110$$
AND 010101
=000100

i.

$$\begin{array}{r}
111001 \\
XOR \ 101001 \\
= 010000
\end{array}$$

(P.S. A link between math and CS: "AND" is multiplication in the boolean ring  $\mathbb{Z}/2\mathbb{Z}$ , and "XOR" is addition (mod 2))

$$0 \cdot 1 = 0$$
  $1 \cdot 1 = 1$   $0 \cdot 0 = 0$   
 $0 + 1 = 1$   $1 + 1 = 0$   $0 + 0 = 0$ 

and:

$$1 \cdot (1+0) = (1 \cdot 1) + (1 \cdot 0) = 1 + 0 = 1$$
 (Multiplication distributes over +)

Corresponds to the truth tables for AND and XOR.

- 35. Identify mask and logical operation needed:
  - (a) Put 1s in the upper 4 bits of an 8-bit pattern without disturbing the other bits:

$$Mask = 11110000$$

$$Operation = OR$$

(b) Complement the most significant bit of an 8-bit pattern without changing the other bits:

$$Mask = 100000000$$

$$Operation = XOR$$

d. Put a 0 in the least significant bit of an 8-bit pattern without disturbing the other bits:

 $\mathrm{Mask}\ = 111111110$ 

Operation = AND

(Essentially, multiply every bit by 1, except least significant bit which gets multiplied by 0)