

American University in Cairo
Computer Science and Engineering Department

Date: March 22, 2015

CSCE 3104- Concepts of Programming Languages-Midterm Exam

Name: _____

ID: _____

Question	Full Mark	Grade
1	25	
2	10	
3	10	
4	10	
5	10	
Total	65	

Question 1: Select one correct answer from the following:

1. Syntax of programming languages is the study of:
 - a. Vocabulary.
 - b. Grammar.
 - c. Meaning.
 - d. (a) and (b).
 - e. All of the above.
2. Semantics of programming languages is the study of:
 - a. Meaning.
 - b. Meaning and Grammar.
 - c. Vocabulary.
 - d. All of the above.
 - e. None of the above
3. Regular expressions can be used to specify:
 - a. Grammatical structure.
 - b. Vocabulary.
 - c. Natural language statements.
 - d. Context sensitive grammar.
 - e. All the above.
4. Regular expressions are composed mainly of :
 - a. Concatenation.
 - b. Selection.
 - c. Repetition.
 - d. Conditions.
 - e. (a) and (b).
 - f. (a) (b) and (c)
5. A grammar is said to be ambiguous if:
 - a. Two production rules have the same meaning.
 - b. There are two ways of writing a given production rule.
 - c. A string in the language described by the grammar has two parse trees.
 - d. There are two non-terminals on the left hand side of the production.
 - e. None of the above.
6. Using natural language grammars in programming language can:
 - a. Eliminate ambiguity.
 - b. Create ambiguity.
 - c. Limit expressiveness.
 - d. Limit extensibility.
 - e. None of the above.
7. A programming language is:
 - a. A notational system.
 - b. Describes computation in machine readable format.
 - c. Describes computation in human readable format.
 - d. Part of its specification indicates whether it is compiled or interpreted.
 - e. All of the above.
 - f. (a) (b) and (c).
8. A GOTO statement can be catastrophic because:
 - a. It can create multiple exit points in blocks.
 - b. It can create multiple entry points in blocks.
 - c. Increases maintainability.
 - d. All the above.
 - e. (a) and (b).

9. Logic programming languages are considered:
- High level.
 - Low level.
 - Close to assembly.
 - None of the above.
10. Interpreted programming languages are typically more efficient than compiled languages:
- True.
 - False.
 - Doesn't make a difference.
11. Interpreted programming languages typically exhibit more of:
- Static features.
 - Dynamic features.
 - Polymorphic features.
 - Low level features.
12. Describe, in English, the language defined by the following grammar (non-terminals are in italics):
- $$\begin{aligned}
 S &\rightarrow m A B C d \\
 A &\rightarrow a A \mid a \\
 B &\rightarrow b B \mid b \\
 C &\rightarrow c C \mid c
 \end{aligned}$$
- One or more m's followed by zero or more a's, then zero or more b's then zero or more c's.
 - One m followed by zero or more a's, then zero or more b's then zero or more c's.
 - One m followed by one or more a's, then one or more b's then one or more c's followed by a d.
 - One or more m's followed by zero or more a's, then zero or more b's then zero or more c's.
 - None of the above.
13. Which string proves that the following grammar is ambiguous:
- $$\begin{aligned}
 S &\rightarrow A \\
 A &\rightarrow A - A \mid id \\
 id &\rightarrow letter\ id \mid letter \\
 letter &\rightarrow a \mid b
 \end{aligned}$$
- ab
 - aaaabbb
 - a+b
 - a-b-a
 - a*a*b
 - None of the above.
14. The following code is written in LISP using the functional paradigm. What does the following LISP code achieve (given $N \geq 1$)?
- ```

(define xyz (N)
 (if (= N 1) 1
 (* N (xyz (- N 1)))))

```
- Factorial (N)
  - Zero
  - $N - (N-1) - (N-2) - \dots - 0$
  - $N + (N-1) + (N-2) + \dots + 0$
  - Average (N)
  - None of the above.

15. What is the associativity of addition and multiplication in the following grammar:

$\text{expr} \rightarrow \text{term} + \text{expr} \mid \text{term}$   
 $\text{term} \rightarrow \text{factor} * \text{term} \mid \text{factor}$   
 $\text{factor} \rightarrow (\text{expr}) \mid \text{number}$   
 $\text{number} \rightarrow \text{number digit} \mid \text{digit}$   
 $\text{digit} \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

- a. Right, left.
  - b. Right, right.
  - c. Left, right.
  - d. Left, left.
16. Why is the following grammar considered to be context sensitive, and not context free:

$S \rightarrow abc \mid aSBc$   
 $cB \rightarrow Bc$   
 $bB \rightarrow bb$

- a. Upper case S in the grammar.
  - b. Upper case S and B in the grammar.
  - c. No single non-terminal on the left side.
  - d. No single terminal on the right side.
  - e. None of the above.
17. A regular expression for a language composed of **one or more** of the following: even number of 0's followed by an odd number of 1's.
- a.  $((00)^* 1 (11)^*)^*$
  - b.  $((00)^+ 1 (11)^*)^+$
  - c.  $(00)^+ 1 (11)^+$
  - d.  $(00) 1 (11)^*$
  - e. None of the above
18. A typical example of a delimiter used in free form languages is:
- a. #
  - b. //
  - c. spaces
  - d. semicolons
  - e. All of the above.
19. To eliminate ambiguity of precedence, the following can be used:
- a. Prefix notation.
  - b. Postfix notation.
  - c. Full parenthesis of expressions.
  - d. In order notation.
  - e. All of the above.
  - f. (a), (b), and (c).
20. Dynamic type binding is done when
- a. Variables are allocated by a programming language primitive
  - b. Variables are specified through an assignment statement
  - c. Variables are declared in functions or procedures
  - d. Variables are bound to storage during runtime time
21. The lifetime of a variable is the time during which
- a. The variable is bound to a specific type
  - b. The variable is bound to a specific name
  - c. The variable is bound to a specific value
  - d. None of the above

22. An implicit declaration is a default mechanism for:
- Specifying types of variables
  - Allocating a variable implicitly
  - Specifying a default value to a variable
  - None of the above
23. A binding is dynamic if
- It occurs before run time and remains unchanged throughout program execution
  - It occurs before run time and changes throughout program execution
  - It occurs during runtime and remains unchanged throughout program execution
  - It occurs during runtime and changes through program execution
24. A named constant is
- a variable that can change its value only during compilation
  - a variable that is bound to a value only when it is bound to storage
  - a variable that cannot change its location
  - all of the above
25. Dynamic scoping means:
- The variables are allocated statically
  - The variables are visible according to the location of calling a function
  - The variables are visible according to the lexical structure of the program
  - The variables are allocated dynamically

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**Question 2:**

- a. Write the formal definition in BNF for the variable(s) declaration that is defined informally as a list of variables preceded by the reserved word **var**, and followed by a colon (:), a type, and semicolon (;). The type is one of the reserved words **char**, or **bool**. The list of variables is defined as a set of variables separated by comma (.). The variable is an identifier. For example this a valid string driven from the above grammar: **var a,b,c: char;** [4 marks]

- b. Convert the grammar written in (a) to EBNF notation [3 marks]

- c. Draw the syntax diagram for the above grammar [3 marks]



**Question 4:**

Given the following program (in some hypothetical programming language):

```

Program main;
 var x, z : integer;

 procedure sub1;
 var a, b, z: integer;
 ----->

 begin {sub 1}
 writeln('The value of x is : ', x);
 ...
 end {sub 1}

 procedure sub 2;
 var a, y, z: integer;

 begin {sub 2}
 sub 1; ...
 end {sub 2}

 procedure sub3;
 var a, x, y: integer;
 ----->

 begin {sub 3}
 x := 5;
 sub 2;
 ...
 end {sub 3}

 begin {main}
 x:= 30;
 sub 3; ...
 end {main}

```

- Show the content of the symbol table using static scoping rules at the indicated locations  
[4 marks]
- Assume the above program was compiled and executed using static scoping rules. What value of x is printed in procedure sub1?  
[3 marks]
- What if dynamic scoping rules are used in the program above, what value of x is printed in procedure sub1  
[3 marks]



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**Question 5:**

Given the following skeletal program where BIGSUB is the main program

|                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   | Memory Stack at Position 1 | Memory Stack at Position 1 |
|-----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|----------------------------|----------------------------|
| <pre> → <b>Main</b> BIGSUB;    <b>var</b> a,b: integer;   → <b>Procedure</b> C;     <b>var</b> x,y :real;     → <b>Procedure</b> D;       <b>var</b> a,x,w: char;       /* start of D */       ..... ← 1     → <b>End</b>; {D}     /* start of C */.....     <b>call</b> D;    → <b>End</b>; {C}   → <b>Procedure</b> A(f: boolean);     <b>var</b> x: boolean;     → <b>Procedure</b> B;       <b>var</b> b: real;.....       /* start of B */       A(false)       ..... ← 2     → <b>End</b>; {B}     /*start of A */.....     <b>If</b> f <b>then</b> B <b>else</b> C    → <b>End</b>; {A}   /* start of BIGSUB */ .....   <b>call</b> A(true); → <b>End</b>; {BIGSUB} </pre> |                            |                            |

Show the stack with all activation record instances, when execution reaches position 1 and position 2 in the above skeletal program [10 marks]