

# Parallel and High Performance Computing

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#### Solution Series 4

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## **MPI**

## Different parallelizations of the $\pi$ reduction

If you pull from the git exercises repository, the solutions can be found in the lecture\_04/pi/solutions folder.

## 1 First steps with MPI

#### Exercise 1.1: MPI: Hello, World!

- For the initialization you can either uncomment the parameters argc and argv, and pass them to MPI\_Init, or, according to the documentation in C, you can also simply do MPI\_Init(NULL, NULL).
- To compile, in Makefile CXX and CC have to be changed from g++ to mpicxx (or to mpiicpc for using intel MPI compiler). To know what mpicxx does you can test mpicxx -show (or mpiicpc -show).

# 2 Using a Ring to communicate

#### Exercise 2.1: $\pi$ MPI Ring (point to point synchronous)

- Split the integral calculation among the processes by partitioning the indices of the for loop according to the rank of every process.
- As explained in the course, MPI\_Ssend is blocked as long as the corresponding receive is not posted. In this case if you write MPI\_Ssend followed by MPI\_Recv, all the processes will stay stuck on MPI\_Ssend (deadlock). That is why you should ensure that at least one process is doing a MPI\_Recv before the MPI\_Ssend.
  - In the solution, all the even process ranks are doing send and then receive, while all the odd ranks do receive followed by send.

If you use MPI\_Send instead of MPI\_Send in this example, it will not deadlock since MPI\_Send acts as MPI\_Bsend on small buffers, and as MPI\_Ssend for large ones<sup>1</sup>. Small and large are implementation dependent.

### Exercise 2.2: $\pi$ MPI Ring (point to point synchronous sendrecv)

With MPI\_Sendrecv there is no need to do a conditional choice based on the rank.

#### Exercise 2.3: $\pi$ MPI Ring (point to point asynchronous)

With MPI\_Isend you can overlap the computation of the sum. But you have to be careful to wait for the end of the communication before changing the send buffer.

It is not because you already received data from the previous process that it means the send to the next process is already finished.

### 3 MPI collectives

### Exercise 3.1: $\pi$ MPI (collective gather)

In MPI\_Gather you have to be careful with the arguments: The send buffer is of size 1, while the receive buffer on the root process, here the process 0, is a buffer of size (at least) psize, even though the argument specifying the receive size is 1. It is the size per process and not the size of the actual buffer.

#### Exercise 3.2: $\pi$ MPI (collective reduce)

The exercise asks to use MPI\_Reduce. This would get the sum stored on the process 0. If you want the solution on every process, you can use MPI\_Allreduce, that is equivalent to a reduce followed by a broadcast.

In the solution we also use MPI\_IN\_PLACE as the send buffer, this makes the receive buffer act as a send and receive buffer. It can also by used with MPI\_Reduce.

<sup>&</sup>lt;sup>1</sup>Check the documentation in https://rookiehpc.github.io/mpi/docs/mpi\_send/index.html.