

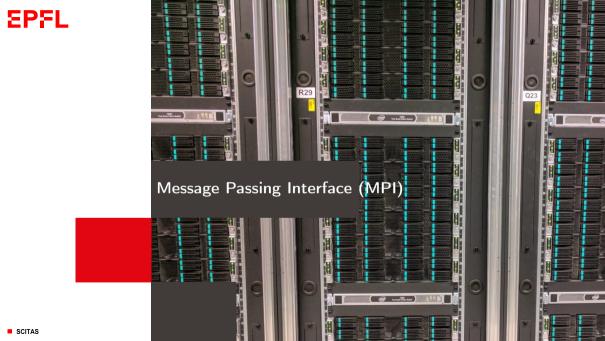
MATH-454 Parallel and High Performance Computing

Lecture 4: MPI basics

Pablo Antolin

Slides of N. Richart, E. Lanti, V. Keller's lecture notes

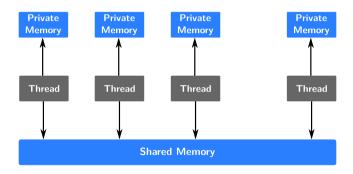
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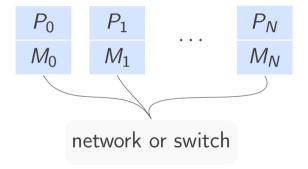




- Introduce distributed memory programming paradigm.
- Point-to-point communications.
- Collective communications.





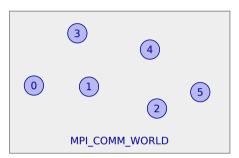


- MPI is a *Message-Passing Interface* specification.
- There are many implementations (Intel MPI, OpenMPI, MPICH, MVAPICH, etc.)
- Library interface, not a programming language.
- It is standarized:
  - ▶ Defined by the MPI forum
  - ► Current version is MPI 4.1
- As such, it is portable, flexible, and efficient.
- Interface to C and Fortran in standard. C interface usable with C++. Also MPI bindings for other languages (Python, Julia, MATLAB, etc.).

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```
mpi/hello mpi.cc
1 #include <iostream>
  # include <mpi.h>
3
4 int main(int argc, char *argv[]) {
    MPI_Init(&argc, &argv);
    int size, rank;
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
9
10
    std::cout << "I am process " << rank << " out of " << size <<
11
     \rightarrow std::endl:
12
    MPI Finalize():
    return 0;
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```





- MPI code is bounded by a MPI\_Init and a MPI\_Finalize.
- MPI starts N processes numbered 0, 1, ..., N-1. The number of every process is usually denoted as rank.
- They are grouped in a *communicator* of *size N*.
- After init, MPI provides a default communicator called MPI\_COMM\_WORLD.



- Point-to-Point (One-to-One).
- Collectives (One-to-All, All-to-One, All-to-All).
- One-sided/Shared memory (One-to ...).
- Blocking and Non-Blocking of all types.



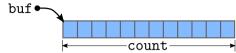
#### Syntax

- **buf**: pointer to the data to send/receive.
- **count**: number of elements to send/receive.
- **datatype**: datatype of the data to send/receive.
- **dest**, **source**: the rank of the destination/source of the communication.
- **tag**: a message tag to differentiate the communications.
- **comm**: communicator in which the communication happens.
- **status**: object containing information on the communication.

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# Send/Receive Details on the buffer



- Buffer is a pointer to the first data (buf), a size (count), and a datatype.
- Datatypes (extract):
  - ► MPI\_INT
  - ► MPI\_UNSIGNED
  - ► MPI\_FLOAT
  - ► MPI\_DOUBLE
- For std::vector<double> vect:
  - buf = vect.data()
  - count = vect.size()
  - datatype = MPI\_DOUBLE

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## EPFL

### $\mathsf{Send}/\mathsf{Receive}$

Useful constants and status

#### Constants:

- MPI\_STATUS\_IGNORE: to state that the status is ignored.
- MPI\_PROC\_NULL: placeholder for the source or destination.
- MPI\_ANY\_SOURCE: is a wildcard for the source of a receive.
- MPI\_ANY\_TAG: is a wildcard for the tag of a receive.

#### MPI\_Status:

Structure containing tag and source

■ Size of the received buffer can be asked using the status:

```
mpi/send recv.cc
16 MPI_Init(NULL, NULL);
17 MPI_Comm_rank(MPI_COMM_WORLD, &rank);
18 MPI_Comm_size(MPI_COMM_WORLD, &size);
19
20 assert(size == 2 && "Works only with 2 procs");
21
22 int tag = 0:
23
24 if (rank == 0) {
    fill_buffer(buf);
25
    MPI_Ssend(buf.data(), buf.size(), MPI_INT, 1, tag, MPI_COMM_WORLD);
27 } else if (rank == 1) {
    MPI_Recv(buf.data(), buf.size(), MPI_INT, 0, tag, MPI_COMM_WORLD,
28
     → MPI STATUS IGNORE):
```

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# EPFL Send/Receive Example

Rank 0		Rank 1
send(1,0)	$\longrightarrow$	recv(0,0)
finalize()		finalize()

- MPI\_Ssend: (S for Synchronous) function returns when other end posted matching recv and the buffer can be safely reused.
- MPI\_Bsend: (B for Buffer) function returns immediately, send buffer can be reused immediately.
- MPI\_Rsend: (R for Ready) can be used only when a receive is already listening (waiting).
- MPI\_Send: acts like MPI\_Bsend on small arrays, and like MPI\_Ssend on bigger ones.

- MPI\_Ssend: (S for Synchronous) function returns when other end posted matching recv and the buffer can be safely reused.
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- MPI\_Send: acts like MPI\_Bsend on small arrays, and like MPI\_Ssend on bigger ones.

Be careful with deadlocks and race conditions!!

```
mpi/send recv deadlock.cc
25 assert(size == 2 && "Works only with 2 procs");
  int tag = 0;
27
_{28} if (rank == 0)
29 {
    fill_buffer(buf);
30
    MPI_Ssend(buf.data(), buf.size(), MPI_INT, 1, tag, MPI_COMM_WORLD);
31
    MPI_Recv(buf.data(), buf.size(), MPI_INT, 1, tag, MPI_COMM_WORLD,

→ MPI_STATUS_IGNORE):
33 }
_{34} else if (rank == 1)
35 {
    MPI_Ssend(buf.data(), buf.size(), MPI_INT, 0, tag, MPI_COMM_WORLD);
36
    MPI_Recv(buf.data(), buf.size(), MPI_INT, 0, tag, MPI_COMM_WORLD,
       MPI STATUS IGNORE):
```

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 $\begin{array}{ccc} \underline{\mathsf{Rank}\;0} & \underline{\mathsf{Rank}\;1} \\ \mathsf{send}(1,0) & \mathsf{send}(0,0) \\ \mathsf{barrier!} & \mathsf{barrier!} \end{array}$ 

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```
mpi/send recv deadlock 2.cc
16 MPI_Init(NULL, NULL);
17 MPI_Comm_rank(MPI_COMM_WORLD, &rank);
18 MPI_Comm_size(MPI_COMM_WORLD, &size);
19
20 // Let's now assume that we may have 2 or more processes.
21 int tag = 0;
22
_{23} if (rank == 0) {
   fill_buffer(buf):
24
    MPI_Ssend(buf.data(), buf.size(), MPI_INT, 1, tag, MPI_COMM_WORLD);
26 } else {
    MPI_Recv(buf.data(), buf.size(), MPI_INT, 0, tag, MPI_COMM_WORLD,

→ MPI_STATUS_IGNORE);
28 }
29 MPI Finalize():
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```



#### Send/Receive Deadlock example - 2

With 2 processes:

Rank 0 Rank 1  $send(1,0) \longrightarrow recv(0,0)$ finalize() finalize()

P. Antolin 16 / 30 SCITAS



#### With 3 processes:

Rank 0		Rank 1	Rank 2
send(1,0)	$\longrightarrow$	recv(0,0)	recv(0,0)
finalize()		finalize()	barrier!

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- Combines a send and a receive, to help mitigate deadlocks.
- Has a in-place variant MPI\_Sendrecv\_replace.



- I for *immediate*.
- request in addition to parameters from blocking version. It is an object attached to the communication.
- Receive does not have a status.
- The communication starts but is not completed.
- S. B. and R variants are also defined.



## Non-blocking send/receive Completion

#### Syntax

```
int MPI_Wait(MPI_Request *request, MPI_Status *status);
int MPI_Test(MPI_Request *request, int *flag, MPI_Status *status);
```

- For non-blocking communications, completion should be "manually" checked.
- MPI\_Test or MPI\_Wait.
- Send completed means the buffer can be reused.
- Receive completed means the buffer can be read.
- status is set at completion.
- flag is true if completed, false otherwise.

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33 MPI Finalize():

```
mpi/isend recv.cc
22 MPI_Request request;
23 if (rank == 0) {
    fill buffer(buf):
24
   MPI_Isend(buf.data(), buf.size(), MPI_INT, 1, 0, MPI_COMM_WORLD,
    26 } else if (rank == 1) {
    MPI_Recv(buf.data(), buf.size(), MPI_INT, 0, 0, MPI_COMM_WORLD,

→ MPI_STATUS_IGNORE):
28 }
29
  // While MPI communications are performed here I can do computations as
  → long as buf is not modified in rank 0, or read/modified in rank 1.
31
32 MPI_Wait(&request, MPI_STATUS_IGNORE);
```

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## Non-blocking send/receive Multiple completions

- MPI\_Waitall, MPI\_Testall waits or tests completion of all the pending requests.
- MPI\_Waitany, MPI\_Testany waits or tests completion of one out of many.
- MPI\_Waitsome, MPI\_Testsome waits or tests completion of all the specified requests.
- For arrays of statuses can use MPI\_STATUSES\_IGNORE.
- MPI\_Request\_get\_status equivalent to MPI\_Test but does not free completed requests.

- Check incoming messages without receiving.
- Immediate variant returns **true** if matching message exists.
- Can be used in combination with a successive MPI\_Get\_count for deducing the size of an incoming message before actually recieving it.
   Thus, we can allocate a buffer for holding the message.

```
int MPI_Barrier(MPI_Comm comm);
```

- Collective communications must be called by all processes in the communicator.
- Barrier is hard synchronization.
- Avoid as much as possible.



## Collective communications Broadcast

#### Syntax

■ The sender process sends data to every other process.





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## Collective communications Scatter

#### Syntax

- The sender process sends a piece of the data to all processes.
- The sendbuf, sendcount, and sendtype are only relevant on the sender.



 $P_3$ 



## Collective communications Scatter

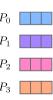
- The sender process sends a piece of the data to all processes.
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## Collective communications Gather

- Every process sends its data to the receiver process.
- The recvbuf, recvcount, and recvtype are only relevant on the receiver.
- recvcount is the size per process, not the total size.





## Collective communications Gather

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- Every process sends its data to the receiver process.
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## Collective communications Gather to all

#### Syntax

Every process sends its data to all the other processes.

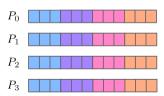




## Collective communications Gather to all

# Syntax int MPI\_Allgather(const void \*sendbuf, int sendcount, MPI\_Datatype → sendtype, void \*recvbuf, int recvcount, MPI\_Datatype recvtype, → MPI\_Comm comm);

Every process sends its data to all the other processes.

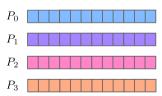




## Collective communications

All to all gather/scatter

• Every process sends a piece of its data to all the other processes.

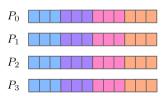




## Collective communications

All to all gather/scatter

• Every process sends a piece of its data to all the other processes.



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## Collective communications Reduction

#### Syntax

- Data from all processes are reduced on the receiver process.
- Common operations being MPI\_SUM, MPI\_MAX, MPI\_MIN, MPI\_PROD.
- A MPI\_Allreduce variant exists where all processes get the results.
- MPI\_IN\_PLACE can be passed to sendbuf of receiver, for reduce, and of all processes, for allreduce. P<sub>0</sub>
  - $P_1$
  - $P_2$
  - $P_3$

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  P0 P0

 $P_1$  +  $P_2$  +  $P_3$ 

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#### Collective communications Reduction

#### Syntax

```
int MPI_Reduce(const void *sendbuf, void *recvbuf, int count,
     MPI_Datatype datatype, MPI_Op op, int receiver, MPI_Comm comm);
```

- Data from all processes are reduced on the receiver process.
- Common operations being MPI\_SUM, MPI\_MAX, MPI\_MIN, MPI\_PROD.
- A MPI\_Allreduce variant exists where all processes get the results.
- MPI IN PLACE can be passed to sendbuf of receiver, for reduce, and of all processes, for allreduce.

 $P_1$ 

 $P_2$ 

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## Collective communications Partial reductions

- Performs the prefix reduction on data.
- MPI\_Scan: on process i contains the reduction of values from processes [0, i]
- MPI\_Exscan: on process i contains the reduction of values from processes [0, i[.
- MPI\_IN\_PLACE can be passed to sendbuf.