	Last Name	First Name & Initial
	Student No.	_
	NO AIDS ALLOWED. Answer ALL questions on test paper. Use backs of sheets for scratch work.	
		Total Marks: 40
[5]	1. Give a specific formula $A$ such that	$\exists x A \not\models A(c/x)$
	where $c$ is a constant symbol not occ	urring in A. Prove your answer, based on the definition of $\models$ .

[10] 2. Recall that **TA** (True Arithmetic) is the set of all sentences A in the language  $\mathcal{L}_A = [0, s, +, \cdot; =]$  of arithmetic such that A is true in the standard model  $\underline{\mathbb{N}}$ . Suppose that A(x) is a formula of  $\mathcal{L}_A$  whose only free variable is x, such that  $A(s^n0)$  is in **TA** for arbitrarily large  $n \in \mathbb{N}$ . Show that the infinite set of sentences

$$\mathbf{TA} \cup \{A(c), c \neq 0, c \neq s0, c \neq ss0, \cdots\}$$

is satisfiable, where c is a new constant.

[15] 3. Give an **LK**- $\Phi$  proof showing that  $A \models B$ , where

$$A =_{syn} \forall x \exists y Pxy$$
  
$$B =_{syn} \exists x \exists y \exists z (Pxy \land Pyz)$$

You do not need to indicate weakenings and exchanges.

You may use abbreviations for formulas (including A, B) in your proof.

**Hint:** You will need two instances of A in your proof. The first part of your proof will not have quantifiers.

## [10] 4. Give an LK proof of the sequent

$$s0 \neq ss0 \rightarrow \forall x(x \neq s0 \lor x \neq ss0)$$

(Here  $t \neq u$  stands for  $\neg t = u$ .)

You may leave out weakenings and exchanges.

Start by giving the specific instances of the LK equality axioms that you need in your proof.

Here are the LK equality axioms:

EL1: 
$$\rightarrow t = t$$

EL2: 
$$t = u \rightarrow u = t$$

EL3: 
$$t = u, u = v \rightarrow t = v$$

EL4: 
$$t_1 = u_1, ..., t_n = u_n \rightarrow ft_1...t_n = fu_1...u_n$$
, for each  $f$  in  $\mathcal{L}$ 

EL5: 
$$t_1 = u_1, ..., t_n = u_n, Pt_1...t_n \rightarrow Pu_1...u_n$$
, for each  $P$  in  $\mathcal{L}$