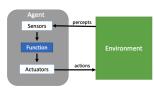
CS2109S Midterms AY24/25 SEM 1

1 Introduction to Al

1.1 Intelligent Agents



· PFAS

- 1. Performance Measure: Best for whom, what are we optimizing, what information is available, any unintended effects, what are the costs
- 2 Environment: Refer to Environment section
- 3. Actuators: Allow intelligent agent to take actions or affect its environment 4. Sensors: Allow intelligent agent to perceive information about its
- Agent Function: Maps from percept histories to actions, refer to Agent Function section

1.2 Task Environment

environment

1 Fully Observable VS Partially Observable

- . Fully Observable: Agent has complete & accurate info about env state at all times (Eq. Chess) Partially Observable: Agent has access to incomplete, uncertain or noisy
- info about env state (Eq. Self-driving cars)

2 Deterministic VS Stochastic VS Strategic

- Deterministic: Next env state is completely determined by current state & agent action | Outcome is fully predictable (Eq. Sudoku)
- . Stochastic: Next env state is not completely determined by current state & agent action | Outcome is uncertain (Eq. Self-driving car) Strategic: Env is deterministic, but outcomes depend on other agents'
- actions, requiring agent to consider strategies & behaviors of others (Eg. Chess)

3. Episodic VS Sequential

- · Episodic: Agent actions are divided into discrete periods, each episode is independent of one another, agent makes decisions based on current enisode (En. Classification task)
- Sequential: Agent actions are inter-dependent, each action affects future states & decisions, agent considers sequence of actions over time (Eg. Chess)

4. Static VS Dynamic

- . Static: Env state does not change while agent is deliberating Dynamic: Env state changes over time even when agent is deliberating
- Semi-dynamic: Env state does not change, but agent's performance

5. Discrete VS Continuous

- Discrete: Finite # of distinct, clearly defined percepts & actions
- Continuous: Infinite # of percepts & actions
- 6. Single Agent VS Multi Agent
- Single Agent: Agent operating by itself in an env
- Multi Agent: Multiple agents in an env

1.3 Agent Structures

Note: Agent is completely specified by Agent Function mapping percept sequences to actions

- 1. Simple Reflex Agent: Operates based on a set of predefined rules or conditions → Reacts to current state of env with a corresponding action — Does not have memory of past states or actions & does not consider future consequences
- 2 Model-based Reflex Agent: Extends simple reflex agent by maintaining internal model of world -> Allows agent to keep track of current env state & handle situations where envistate is partially observable or changes over
- 3 Goal-based Agent: Operates with specific goals in mind → Selects actions based on ability to achieve these goals -> Considers future & plans its actions to achieve desired end state | Uses goal representation & perform search and planning
- 4. Utility-based Agent: Extends goal based agent by considering not just whether goals are achieved, but how well they are achieved -> Assigns utility value to different states & chooses actions that maximize overall utility 5. Learning Agent: Improves performance over time by learning from its
- experiences | Can be reflex, model, goal & utility based · Exploitation: Maximize expected utility according to current knowledge
- about world
- · Exploration: Trying to learn more about the world

2 Solving Problems by Searching

2.1 Designing an Agent

- · Assumptions: Goal-based agent | Env is fully observable, deterministic, static discrete
- Problem-solving Agent: Agent that plans ahead (considers a seg. of actions that form a path to a goal state), undertakes SEARCH process
- 1. Goal Formulation: (What do we want?)

- 2. Problem Formulation: (How the world works?) → States (state space), Initial State(initial state of agent), Goal State/Test (goal state of agent), Actions (things that agent can do in a given state). Transition Model (specifies outcome of an action to a given state & how it leads to new states). Action Cost Function (cost of performing an action)
- 3. Search: (How to achieve it?) → Path (seq. of actions). Solution (path to
- Execute
- Representation Invariant: A condition that must be true over all valid concrete representations of a class

2.2 Search Algorithms (Introduction)

- · Search Algorithm: Takes in search problem (input), returns solution/failure (output) | Defined by Order of Expansion (FRONTIER)
- **Evaluation Criteria:**
- 1. Time Complexity: # of nodes generated/expanded
- Space Complexity: Max # of nodes in memory 3. Completeness: Does it return solution if it exists?
- 4. Optimality: Does it always find least cost solution?

Tree Search:

while frontier is not empty: state = frontier.pop() if state is goal: return solution

for action in actions(state) next state = transition(state, action) frontier.add(next state)

create frontie insert initial state to fronti while frontier is not empty if state is goal: return solution visited.add(state)

next state = transition(state, action)

frontier.add(next state)

Granh Search

- Checking of Goal State: New state is checked for goal state before new states are PUSHED to frontier → Expand less states, may skip states with less cost
- State is checked for goal state after state is POPPED from frontier → Expand more states, will not skip states with less cost

2.3 Search Algorithms (Uninformed Search)

- · Key Idea: Search Algo is given no clue about how close a state is to the goal | Can be Tree or Graph Search
- **BFS:** Queue Frontier | Time Complexity: $O(b^d) = 1 + b + b^2 + b^2$ $\ldots + b^d$, where b is branching factor, d is depth of optimal solution Space Complexity: $O(b^d)$ when expanded until last child in worst case Completeness: Complete if b is finite | Optimality: Optimal if step cost is same everywhere
- UCS: Priority Queue (path cost) Frontier, where path cost == cost from root to a state | Time Complexity: $O(b^{C^*/\epsilon})$, where C^* is cost of optimal solution, ϵ is minimum edge cost $\to C^*/\epsilon$ is est. depth of optimal solution in worst case | Completeness: Complete if $\epsilon > 0$ and C^* is finite (if $\epsilon=0$, zero cost cycle may occur) | Optimality: Optimal if $\epsilon>0$ Note: BFS is special case of UCS where step cost == 1 for every edge
- **DFS:** Stack Frontier | Time Complexity: $O(b^m)$ where b is branching factor, m is max depth | Space Complexity; O(bm) as only 1 path is expanded at one time | Completeness: Not complete (when depth is infinite or can go back or forth) | Optimality: Not optimal (there can be paths with less cost not explored yet)
- DLS (Depth Limited Search): Limit the search depth to l where l <= m, backtrack once depth limit is reached | Time Complexity: $O(b^l)$ | Space Complexity: O(bl) | Completeness: Not complete when soln lies deeper 1 | Ontimality: Not optimal when soln lies deeper than 1 Note: We dk the denth of solution, which is a downside
- IDS (Iterative Deenening Search): Do DLS with may depth of 0 → return soln if found, otherwise increase depth | Time Complexity
- $O(b^d)$, $Overhead = (n_{IDS} n_{DLS})/n_{DLS} \mid Space$ Complexity: O(bd) | Completeness: Complete | Optimality: Optimal if step cost is same everywhere Note: IDS is not always faster than DES → Consider state space s.t. each state have only single successor & goal node is at depth $n \to {\it IDS}$ will run
- in $O(n^2)$. DFS will run in O(n)Backward Search: Search from goal
- Bidirectional Search: Combine forward search & backward search, stop when 2 searches meet | Time Complexity: $2 * O(b^{d/2}) < O(b^{d})$

2.4 Search Algorithms (Informed Search)

- Key Idea: Search Algo has a clue on how close a state is to the goal Best First Search: Priority Queue (f(n)) Frontier, where f(n)estimates the goodness of a state (Node with lowest f(n) is selected first to be expanded) $\mid f(n)$ can be purely heuristic (estimated cost from n to goal) or a combi of path cost & heuristic
- Greedy Best First Search: Priority Queue (f(n) = h(n)) Frontier, where h(n) is heuristic function that est, cost from n to goal (Expands node that seems closest to goal according to h(n) without considering path cost so far) | Time Complexity: $O(b^m)$ | Space Complexity: $O(b^m)$ | Completeness: Not complete since GBFS might keep expanding nodes based on h(n) without ever finding goal | Optimality: Not optimal since GBFS selects nodes based on h(n) without considering
- A^* Search: Priority Queue (f(n) = g(n) + h(n)) where g(n) is cost so far to reach n | Time Complexity: $O(b^m)$ | Space Complexity: $O(b^m)$ | Completeness: Complete | Optimality: Optimal

- If h(n) is admissible $\to A^*$ using Tree search is optimal
- If h(n) is consistent $\to A^*$ using Graph search is optimal • Note: UCS is special case of A^* search where h(n) = 0

2.5 Heuristics

- Estimate cost from node n to goal
- Admissible Heuristics: For every node $n, h(n) \leq h^*(n)$, where $h^*(n)$ is true cost to reach goal state from n (Never over-estimate)
- Consistent Heuristics: For every node n, every successor n'generated by action $a, h(n) \leq c(n, a, n') + h(n')$ and $h(G) = 0 (Proof h(n) - h(n') \le c(n, a, n'))$ Note: If h(n) is consistent, $f(n') > f(n) \rightarrow f(n)$ is non-decreasing along any path → Nodes are expanded in order of increasing f cost
- **Dominance:** If $h_2(n) \geq h_1(n)$ for all $n \to h_2$ dominates $h_1 \mid$ If h_2 is admissible $o h_2$ is better for search Creating Admissible Heuristics:
- Problem with fewer restrictions on actions is called a relaxed problem. Cost of an optimal soln to a relaxed problem is an admissible h for original problem

3 Local Search & Adversarial Search

3.1 Local Search

- Assumptions: Agent is a Goal/Utility-based agent, Env has a very large state space
- Informed & Uninformed Search VS Local Search
- 1. IUS: Low to moderate state space | Optimal or no soln | Search path is usually the soln
- 2. LS: Very large state space | Good enuf soln is preferable rather than no soln | State is the soln (don't care about search path)
- . Local Search Overview
- Basic Idea: Start somewhere in state space, move towards a better spot Problem Formulation: States(state space) Initial State(initial state of agent). Goal test (ontional, coz we actually dk the goal state, rely on eval function instead). Successor Function (possible states from a state). Evaluation Function (Output value/goodness of a state)
- Hill Climbing Algorithm

current = initial state

neighbor = a highest-valued successor of current

if value(neighbor) <= value(current):

return current

- current = neighbo
- · Known as Greedy Local Search (pick best amongst neighbors, repeat)
- Best Soln: State space where eval. function has a max value (global max) Disadvantages: Cannot reach global max if it enters local max, plateau Sensitive to choice of initial state, poor initial state may result in poor final state (Can overcome with random restarts, walks)

Simulated Annealing current = initial state

T = a large positive value

while T > 0:

next = a randomly selected successor of current

if value(next) > value(current): current = next

else with probability P(current, next , T): current = next

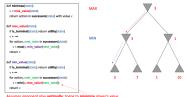
decrease T return current

- P(current, next, T) =
- (value(next)-value(current))/T
- . More exploration of bad states is allowed when T is high, more exploitation is done when T is low \rightarrow basically choosing worse
- successor may lead to a better max - Theorem: If $\stackrel{\cdot}{T}$ decreases slowly enough, SA will find global optimum with high probability

3.2 Adversarial Search

 Assumptions: Agent is Utility-based | Env is a game (game cannot be) single player partially observable stochastic but must be fully observable deterministic discrete terminal states exist 2 players zero-sum turn taking

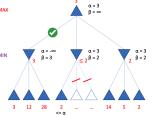
Minimax Algorithm:



- . Intuition: MAX wins when utility is high, MIN wins when utility is low | Assign utility values to all terminal states & start tracing from terminal states

 Fventually, all states will have utility values, starting player can choose a state that will may/min his utility
- Analysis: Completeness: Complete if tree is finite. | Time Complexity: $\overline{O(b^m)}$ | Space Complexity: O(bm) depth first exploration | Optimality: Optimal against optimal opponent

Alpha-beta Pruning



- Definitions: α is best explored option to the root for MAX player (Highest value for MAX) | β is best explored option along path to the root for MIN player (Lowest value for MIN)
- Procedure: 1. Assign $\alpha = -\infty$, $\beta = \infty$ for root 2. Propagate values down to the terminal node 3. Update α value at MAX node, β value at MIN node 4. Propagate values up 5. Prune branches of nodes

Minimax with Cutoff

- Instead of calling is_terminal, call is_cutof f which returns TRUE if (1): State is terminal or (2): Cut-off is reached
- Instead of using utility, call eval which is an eval. function that returns (1): Utility for terminal states or (2): Heuristic value for

4 Introduction to ML & Decision Trees

4.1 Introduction to ML

Definitions: Computer program is said to learn from experience ${\cal E}$ w.r.t. some class of tasks T & performance measure P, if its performance at tasks in T , as measured by P , improves with experience E

Types of Feedback:

- . Supervised Learning: Involves training a model on a labeled dataset, where input data is paired with correct output → Model learns to map inputs to outputs based on this labeled data, allowing it to make predictions on new data
- · Regression: Predict continuous input
- Classification: Predict discrete input
- 2. Unsupervised Learning: Deals with dataset that do not have labeled outputs → Goal is to identify patterns & structures within data
- 3. Reinforcement Learning: Agent learns to make decisions by interacting with an environment → Agent receives feedback in the form of rewards or penalties based on its actions → Learns optimal behaviors over time



4.2 Performance Measure

For a set of N examples $\{(x_1, y_1), ..., (x_N, y_N)\}$ we can compute the

$$MSE = \frac{1}{N} \sum_{i=1}^{N} (\hat{y}_i - y_i)^2$$

Where $\hat{y}_i = h(x_i)$ and $y_i = f(x_i)$.

For a set of N examples $\{(x_1, y_1), ..., (x_N, y_N)\}$ we can compute the

$$MAE = \frac{1}{N} \sum_{i=1}^{N} |\hat{y}_i - y_i|$$

Where $\hat{y}_i = h(x_i)$ and $y_i = f(x_i)$

Classification:

Classification is correct when the prediction $\hat{y} = y$ (true label).

For a set of N examples $\{(x_1, y_1), \dots, (x_N, y_N)\}$ we can compute the

 $Accuracy = \frac{1}{N} \sum_{i=y_i}^{N} \mathbf{1}_{\hat{y}_i = y_i}$

Where $\hat{y}_i = h(x_i)$ and $y_i = f(x_i)$.



many selected items are relevant. maximise if FP is costly) • Recall: $\frac{TP}{TP+FN}$ (How many relevant items are selected,

maximise if FN is dangerous)

• F1 Score: $\frac{2}{1/precision+1/recall}$

4.3 Decision Trees

- Traits of Decision Trees:
- · Decision Trees can express any function of input attributes
- · Consistent Decision Tree for any training set, but probably will not generalize to new examples
- # of distinct decision trees with n boolean attributes = $2^{2^{n}}$ Decision Tree Learning Algorithm

def DTL(examples, attributes, default): if examples is empty: return default

if examples have the same classification:

return classification

if attributes is empty: return mode(examples)

hest = choose attribute(attributes examples)

tree = a new decision tree with root best for each value 12: of best:

 $examples_i = \{rows in examples with best = v_i\}$

subtree = DTL(examples:, attributes - best, mode(examples)) add a branch to tree with label w and subtree subtree

· mode: Category with the highest number

- choose attribute: Chooses attribute with the highest information gain
- Choosing an attribute: · Ideally select an attribute that splits examples into "all positive" or "all
- negative" Entropy (Measure of randomness):

 $\overline{I(P(v_1),\ldots,P(v_n))} = -\sum_{i=1}^n P(v_i)log_2P(v_i),$ where for data set containing p positive k n negative examples,

where for data set containing p positive & n negation $I(\frac{p}{p+n}, \frac{n}{p+n}) = \frac{p}{p+n} \log_2 \frac{p}{p+n} - \frac{n}{p+n} \log_2 \frac{n}{p+n}$ Note: I(1,0) = I(0,1) = 0, I(0,5), 0,5) = 1

 Information Gain (Entropy of curr. node - Total Entropy of children nodes): $\overline{IG(A)} = I(\frac{p}{p+n}, \frac{n}{p+n}) - remainder(A)$

 $\begin{aligned} & \underset{\text{remainder(A)}}{\text{remainder(A)}} &= \sum_{i=1}^{v} \frac{p_i + n_i}{p + n} \, I(\frac{p_i}{p_i + n_i} \, , \, \frac{n_i}{p_i + n_i}), \text{ where} \\ & \text{examples are split into } v \text{ subsets by attribute } A \end{aligned}$ Dealing with continuous valued attributes: Define a discrete valued input

attribute to partition values into discrete set of intervals Dealing with missing values: Assign most common value of attribute,

- assign probability to each value and sample, drop attribute, drop rows Overfitting: Decision Tree is perfect on training data, but worse on test data Occam Razor: Prefer short/simple hypothesis (long/complex hypothesis
- that fits data may be coincidence) Pruning: Prevents nodes from being split even when it fails to cleanly separate examples (Min samples leaf: Merge until leaf node is above min. samples number | Max depth: Merge until leaf nodes are at depth less than

- Tutorials Pointers DFS: DFS utilizes O(bm) memory as DFS must store all nodes along the current path from root to the deepest node explored, along with branching
- factor b at each level UCS vs Dijkstra: 2 algos are the same (both traverse search space in the manner, using a PQ to keep track of which nodes/states to visit next), but can argue that Diikstra finds the shortest path to every node from a single source, while UCS only finds the shortest path to goal states

6 Midterms PYP Pointers

- Tree Search VS Graph Search: Assuming the problem consists of discrete states (eq. pitcher problem), search space using tree search may be infinite (Eg. we keep filling up and emptying pitcher) \rightarrow but using graph search will make search space finite (graph search will not revisit visited states)
- Tree Search:
- If optimal solution is needed: use BES (for problems where each action.) has same cost), use UCS (for problems where each action has different cost), use IDS (if space is more important concern than time, since BES
- is space inefficient) If optimal solution is not needed, but we need to preserve space; use DLS
- · Termination & Completeness on Tree Search Algorithms: . BES is complete, but may not terminate
- . DES is incomplete and may not terminate
- . UCS is complete, but may not terminate IDS is complete, but may not terminate (unless depth limiting condition)
- is applied) DI S is complete and terminates regardless of whether a solution exists due to its depth limiting property

Heuristics:

- Max of 2 admissible heuristics $(max(h_0,h_1))$ for a problem is also an admissible heuristic $\rightarrow max(h_0, h_1)$ is also dominant over h_0 & h 1 since it takes max of both heuristics
- Given h_0 is admissible for problem p_0 & h_1 is admissible for problem $p_1
 ightarrow h_0$ & h_1 are admissible for the combined problem $p_0 + p_1$ $ightarrow max(h_0,h_1)$ is also admissible and dominant for p_0+p_1
- . If heuristic is consistent, it must be admissible too! · If heuristic is non-admissible, it must be non-consistent too!
- Common Proving Techniques:
- Consistency: Decrease in heuristic value must not exceed action cost $(h(n) - h(n') \le c(n, a, n')) \to \text{Focus on maximum}$ decrease in heuristic value and showing it is less than actual action

- Relaxed Problem: True cost/admissible heuristic of a relaxed problem is an admissible heuristic to original problem, True cost of relaxed problem is consistent heuristic to original problem
- Dominance: Not useful to describe dominance between non-admissible heuristics since non-admissible heuristics lead to sub-optimal search \rightarrow for consistency sake, just treat it such that dominance concept can apply to both admissible & non-admissible heuristics

6.1 Log Values (Base 2)

6.1 Log Values (Ba: 1. log21 = 0 2. log22 = 1 3. log33 = 1.58496 4. log24 = 2 5. log25 = 2.32192 6. log26 = 2.58496 7. log27 = 2.80735 8. log28 = 3 9. log29 = 3.16993 10. log210 = 3.32193 6.2 Entropy Table

	0	1	2	3	4	5	6	7	8	9	10
0	0.0										
1	0.0	1.0									
2	0.0	0.918	1.0								
3	0.0	0.811	0.971	1.0							
4	0.0	0.722	0.918	0.985	1.0						
5	0.0	0.65	0.863	0.954	0.991	1.0					
6	0.0	0.592	0.811	0.918	0.971	0.994	1.0				
7	0.0	0.544	0.764	0.881	0.946	0.98	0.996	1.0			
8	0.0	0.503	0.722	0.845	0.918	0.961	0.985	0.997	1.0		
9	0.0	0.469	0.684	0.811	0.89	0.94	0.971	0.989	0.998	1.0	
10	0.0	0.439	0.65	0.779	0.863	0.918	0.954	0.977	0.991	0.998	1.0