

Comp Architecture - Assignment 1

• Soln 1:

Power & Energy (30 Points)

→ Processor gives at: (CPV Bound assumption)

freq: 2 GHz

Voltage: 1.1 V

Total Power: 110 W

$$\begin{array}{c} 20 \text{ W} \\ \text{Leakage Power} \end{array} \quad \begin{array}{c} 90 \text{ W} \\ \text{Dynamic power} \end{array}$$

Time taken: 50 s

Assumption: Voltage & freq follows linear relation
 $V_{orange} (0.9V - 1.2V)$

(i) Smallest time it takes to execute the program

→ Since, V & freq have a linear relation,
if $V \uparrow$ then freq \uparrow proportionally.For min. time, freq should be max (freq \propto 1
under safety levels) exec timeif V increases $\rightarrow 1.1V$ to $1.2V$

$$f_{\text{req new}} = 2 \times \frac{1.2}{1.1} = 2.182 \text{ GHz}$$

$$\therefore \text{exec time} = 50 \times \frac{2}{2.182} = 45.83 \text{ secs}$$

(By inverse proportionality)

i) Highest Power to execute the program

Since Dynamic Power $\propto V^2 \times f$

To \uparrow D.P., we would maximize V , hence f

As per previous part:

$$V : 1.1V \rightarrow 1.2V$$

$$f : 2 \text{ GHz} \rightarrow 2.182 \text{ GHz} \quad (\text{proportional to } V)$$

$$\frac{\text{D.P}_{\text{old}}}{\text{D.P}_{\text{max}}} = \frac{V_0^2 \times f_0}{V_{\text{max}}^2 \times f_{\text{max}}}$$

$$\frac{90}{\text{D.P}_{\text{max}}} = \frac{1.1 \times 1.1 \times 2}{1.2 \times 1.2 \times 2.182}$$

$$\text{D.P}_{\text{max}} = 116.84 \text{ W}$$

Also,

Leakage Power \propto Supply Voltage

$$\begin{aligned} \text{Leakage Power}_{\text{max}} &= 20 \times \frac{1.2}{1.1} \\ &= 21.818 \text{ W} \end{aligned}$$

$$\text{Total max Power} = 116.84 + 21.818$$

$$= \boxed{138.66 \text{ W}}$$

iii Lowest Energy to execute the Program

→ we use DVFS to minimize energy

$$\therefore \text{minimum } V : 1.1V \rightarrow 0.9V$$

Hence,

$$\text{minimum } f : 2 \times \frac{0.9}{1.1} = 1.636 \text{ MHz}$$

$$\frac{D.P_{\text{total}}}{D.P_{\text{min}}} = \frac{1.1 \times 1.1 \times 0.2}{0.9 \times 0.9 \times 1.636} = \frac{90}{D.P_{\text{min}}}$$

$$D.P_{\text{min}} = 49.28 \text{ W}$$

Also,

$$L.P' \propto \text{Supply } V$$

$$L.P_{\text{min}} = 20 \times \frac{0.9}{1.1} = 16.364 \text{ W}$$

$$\text{Total } P_{\text{min}} = 49.28 + 16.364 \\ = 65.64 \text{ W}$$

Since,

$$\text{exec time} \propto \frac{1}{f}$$

$$\text{exec time}_{\text{new}} = 50 \times \frac{2}{1.636} = 61.125 \text{ secs}$$

$$\begin{aligned} \text{Min energy} &= T.P_{\text{min}} \times \text{exec time}_{\text{new}} \\ &= 65.64 \times 61.125 \\ &= \boxed{4012.25 \text{ J}} \end{aligned}$$

• Soln 2: Power & Energy (30 Pts)

→ Processor

$$f = 3 \text{ GHz}, V = 1V, T = 100s, D.P = 80W, L.P = 20W$$

Above values are baseline values.

Considering memory bound $Y\%$

$Y\% \text{ of } 100s \text{ (baseline time)}$

$$\text{memory access time} = \frac{Y}{100} \times 100 = Y \text{ sec}$$

$$CPU \text{ access time} = 100 - Y$$

Since freq \downarrow to 1.5 GHz (Half)

CPU time will double = $2(100 - Y)$

memory time unaffected = Y

$$\begin{aligned} \text{Baseline Energy} &= (80 + 20) \times 100 \\ &= 10000 J \end{aligned}$$

Now if frequency is halved (Applying DFS)

$$D.P : 80W \rightarrow 40W$$

L.P will be constant

$$\text{New T.P} = 40 + 20 = 60W$$

Eqn for minimum $Y\%$ memory bound such that energy is less than baseline

$$\text{New Power} \times \text{New time} < 10000 J$$

$$60 \times (2(100 - Y) + Y) < 10000 J$$

$$60 \times (200 - Y) < 10000$$

$$200 - Y < \frac{500}{3}$$

$$Y > 33.33$$

The System should be $> 33.33\%$ memory bound to reduce energy than baseline

$$Y\% = \sim 34\%$$

• 3) Sum of Execution time (20 pts)

	P	Q	R	S
→ Sys A (sec)	100	120	80	400
Sys B (sec)	120	80	100	300
Sys C (sec)	90	110	90	500

a) Sum of Execution time

$$\begin{aligned} \text{Total exec time for B} &= 120 + 80 + 100 + 300 \\ &= 600 \text{ s} \end{aligned}$$

$$\begin{aligned} \text{Total exec time for C} &= 90 + 110 + 90 + 500 \\ &= 790 \text{ s} \end{aligned}$$

$$\begin{aligned} \text{Speedup (C to B)} &= \frac{\text{exec time of B}}{\text{exec time of C}} = \frac{600}{790} \\ &= 0.76 \end{aligned}$$

∴ System C has 0.76 times the performance than B

b) Sum of weighted exec time

Ref machine = System A

For P, Normalized factor = $\frac{1}{100}$

For Q, Normalized factor = $\frac{1}{120}$

For R, " = $\frac{1}{80}$

For S, " = $\frac{1}{400}$

$$SWET(B) = \frac{1}{100} \times 120 + \frac{1}{120} \times 80 + \frac{1}{80} \times 100 + \frac{1}{400} \times 300 \\ = 3.87$$

$$SWET(C) = \frac{1}{100} \times 90 + \frac{1}{120} \times 110 + \frac{1}{80} \times 90 + \frac{1}{400} \times 500 \\ = 4.192$$

$$\text{Speedup} = \frac{3.87}{4.192} = 0.923$$

System C has speedup 0.923 times that of B

(less performance)

c) GM of execution times

$$GM = \sqrt{T_1 \times T_2 \times \dots \times T_n}$$

$$\begin{aligned} GM_{(B)} &= \sqrt[4]{120 \times 80 \times 100 \times 300} \\ &= 130.27 \end{aligned}$$

$$\begin{aligned} GM_{(C)} &= \sqrt[4]{90 \times 110 \times 90 \times 500} \\ &= 145.28 \end{aligned}$$

$$\text{Speedup} = \frac{130.27}{145.28} = 0.897$$

System C has 0.897 times performance than B
(speed down)

Q4 Performance eqn. (20 pts)

	P	Q	R	S	T
→ IPC _{old}	0.6	0.8	0.4	1.1	1.4
IPC _{new}	0.7	0.6	0.3	1.1	1.5

New laptop has 10% clock speed

Some binaries are run for same no. of cycles on each machine

Considering AM of IPCs

$$\begin{aligned} \text{AM IPCs old} &= 0.6 + 0.8 + 0.4 + 1.1 + 1.4 \\ &= \frac{4.3}{5} \end{aligned}$$

$$\begin{aligned} \text{AM IPCs new} &= 0.7 + 0.6 + 0.3 + 1.1 + 1.5 \\ &= \frac{4.2}{5} \end{aligned}$$

Clock speed of new = 1.1 if Clock > new old = 1

$$(\text{PV time})_{\text{old}} = 1 \times \frac{4.3}{5} = \frac{4.3}{5}$$

$$(\text{PV time})_{\text{new}} = 1.1 \times \frac{4.2}{5} = \frac{4.62}{5}$$

$$\text{Perf improvement} = \frac{4.62 - 4.3}{4.3} \times 100$$

= 7.44% increase in performance

$$\frac{4.62}{4.3} = 1.074 \text{ Speed up.}$$