

CS301

Embedded System and Microcomputer Principle

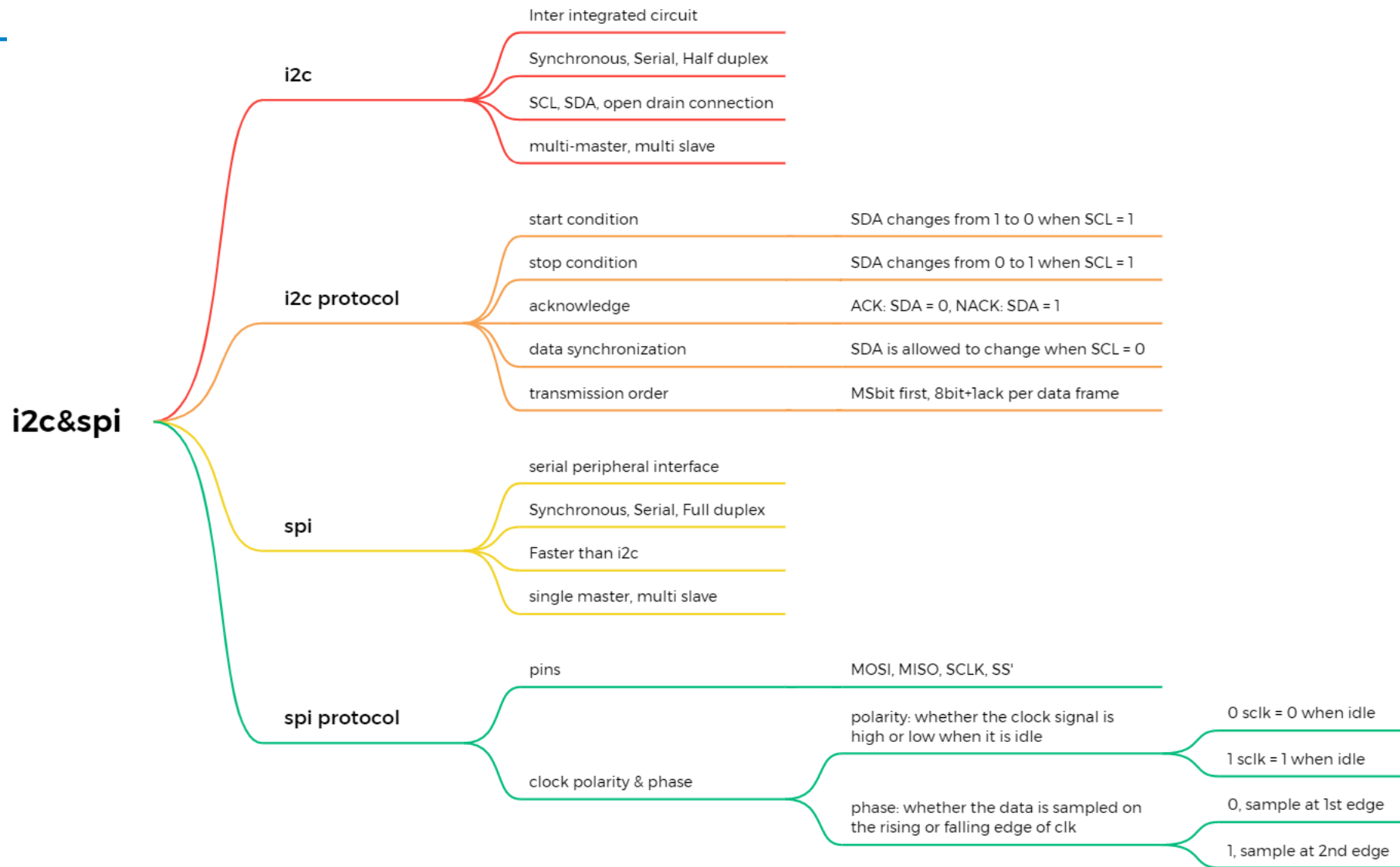
Lecture 11: SD Card & File System

2024 Fall

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Recap

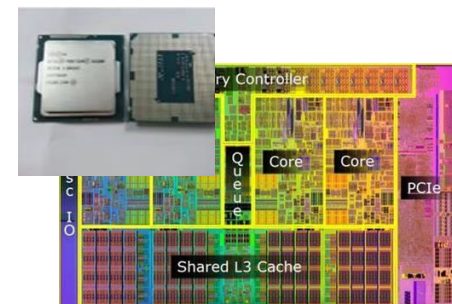
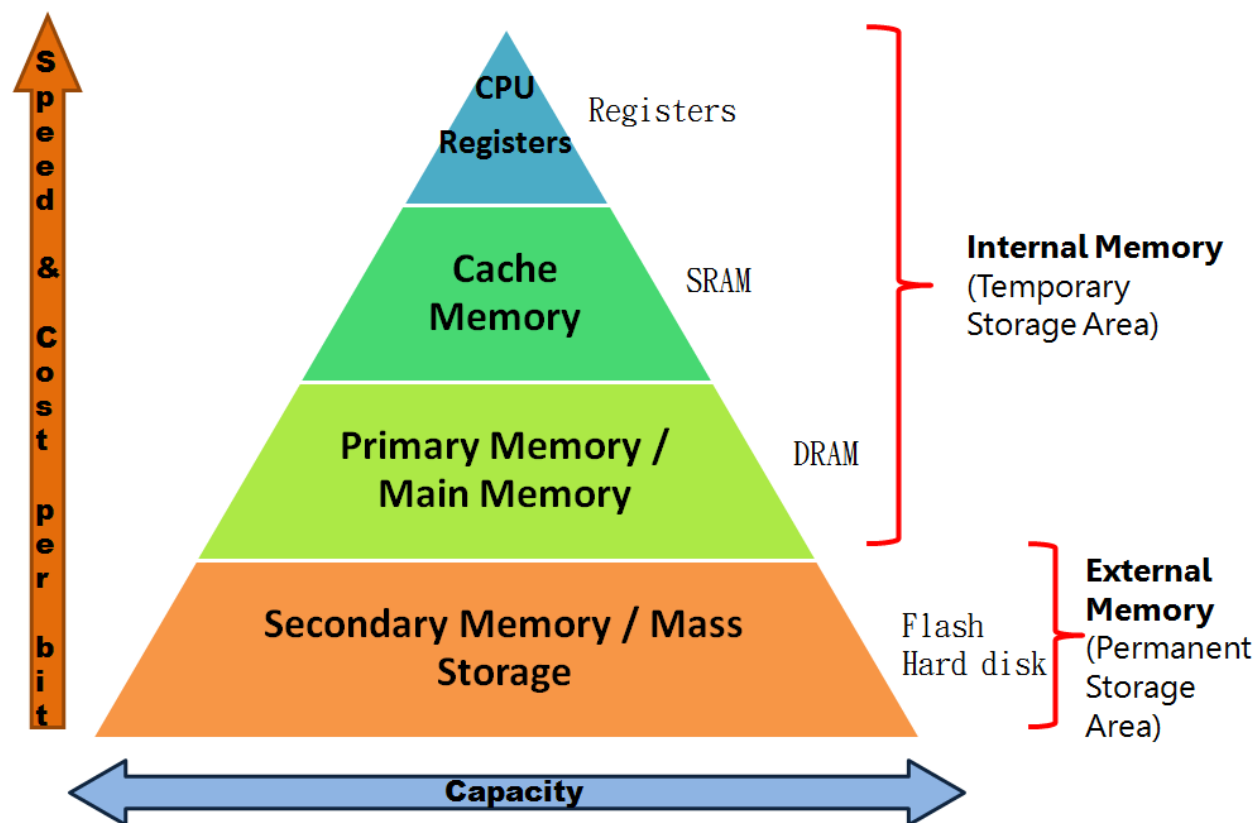


Outline

- **Massive Storage**
- SD Card
- File System

Memory Layout

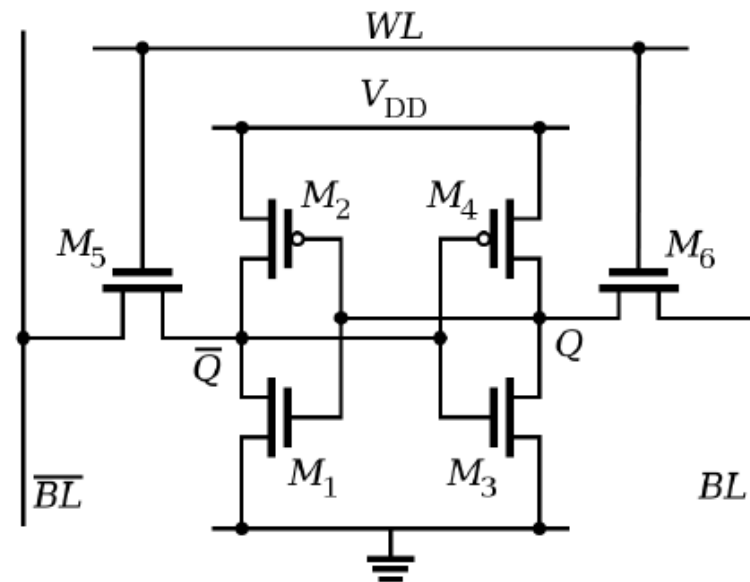
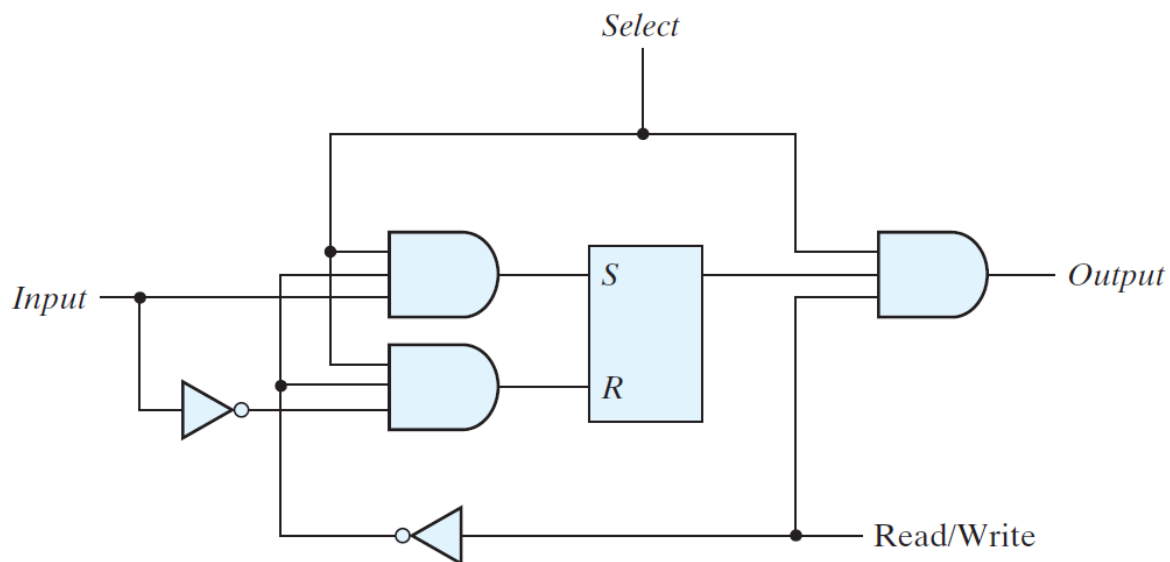
- Memory
 - Internal Memory
 - SRAM
 - DRAM
 - External
 - Hard disk
 - Flash



SRAM

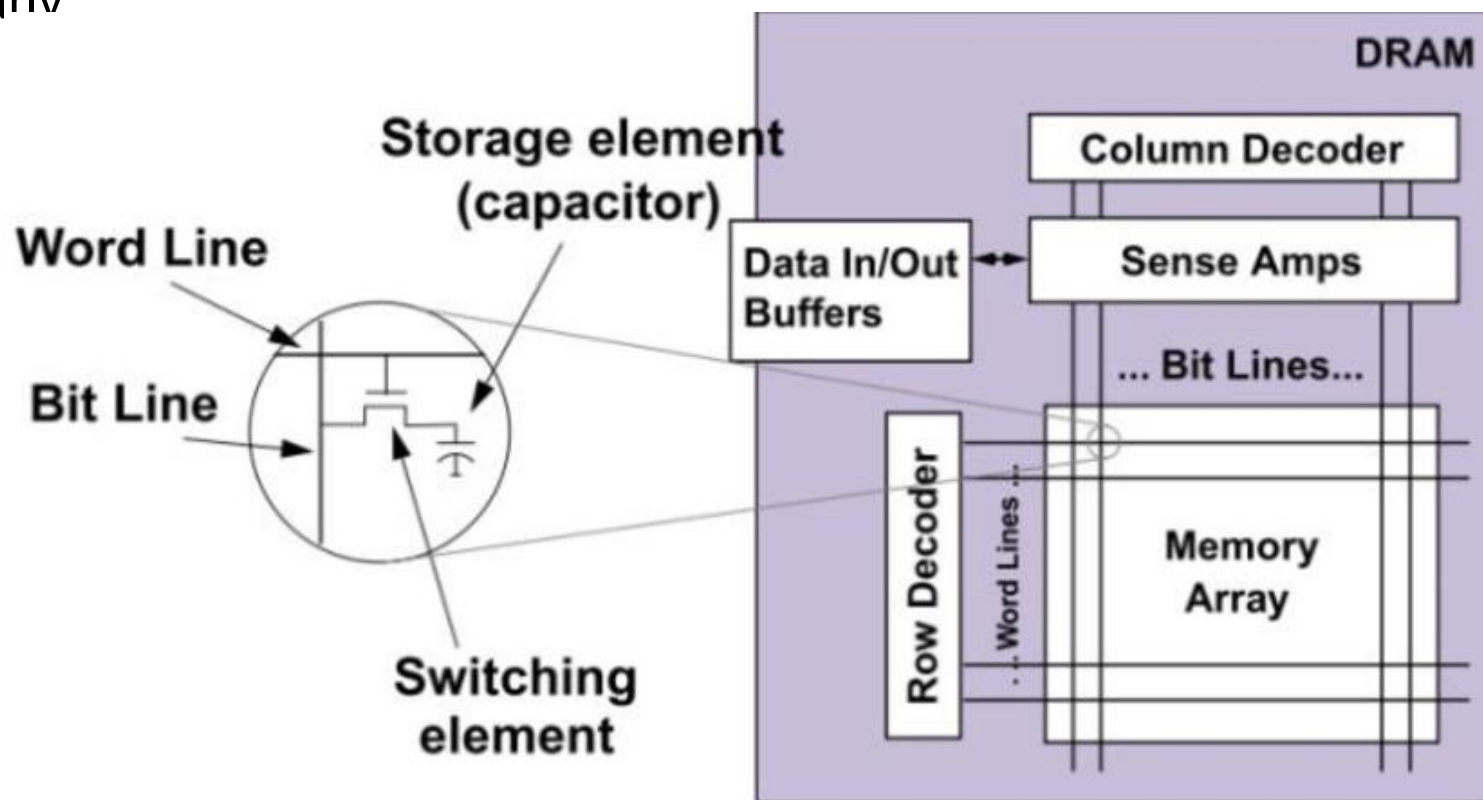
- SRAM (Static Random Access Memory)
 - Low density, high power, expensive, fast
 - Content will last until lose power
 - Often used for caches

一个SRAM需要六个晶体管，因此低密度，高功率，昂贵



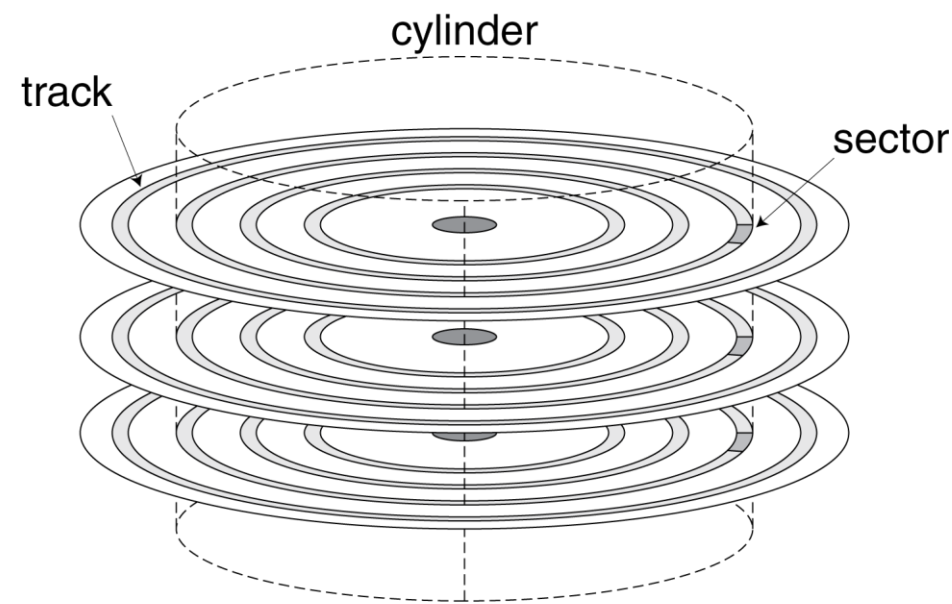
DRAM

- DRAM (Dynamic Random Access Memory)
 - High density, low power, cheap, slow
 - Need to be refreshed regularly
 - Used for main memory



Disk Storage

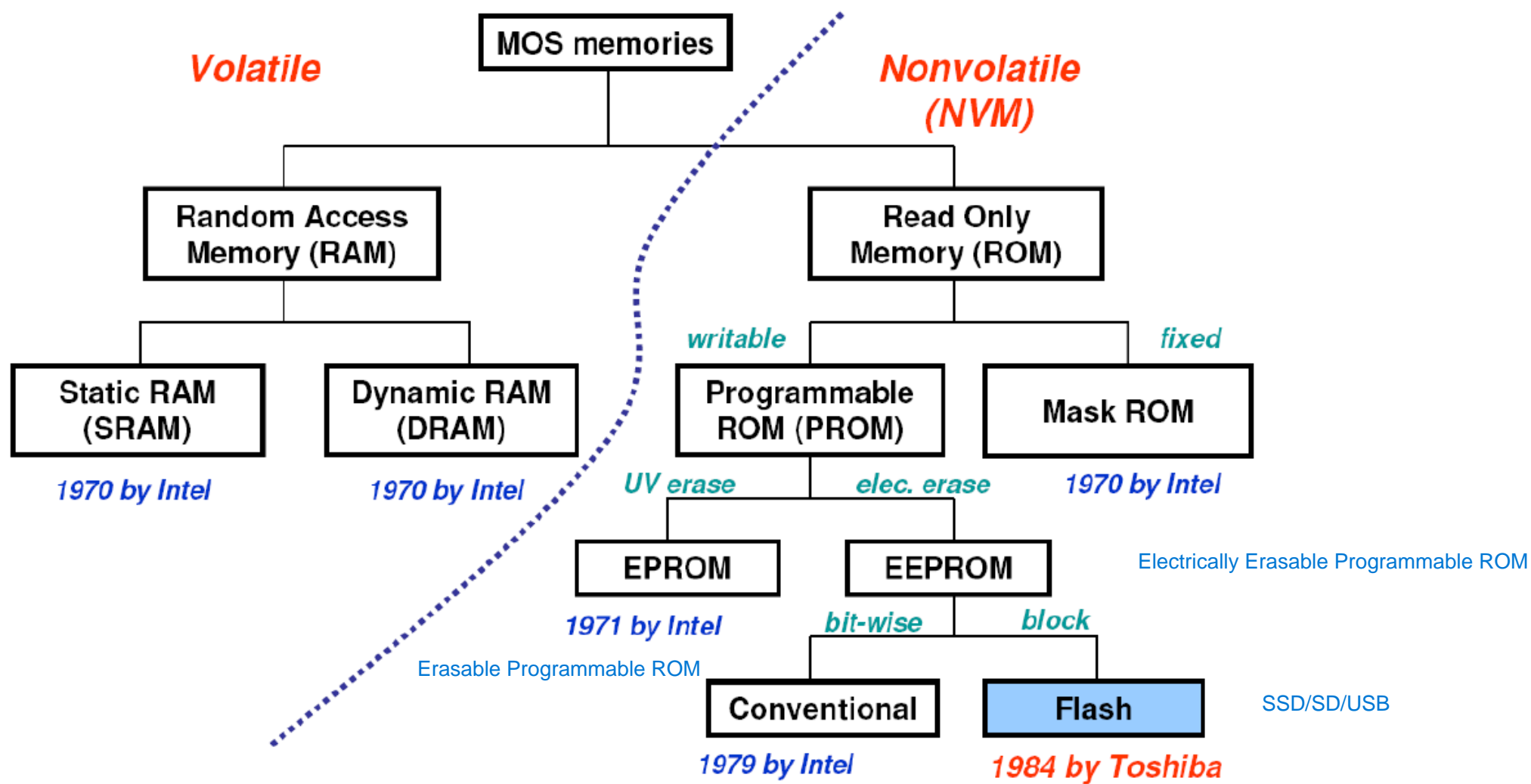
- Each sector records
 - Sector ID
 - Data
- Access to a sector involves
 - Queuing delay if other accesses are pending
 - Seek: move the heads
 - Rotational latency
 - Data transfer
 - Controller overhead



Volatile & Non-Volatile Memory

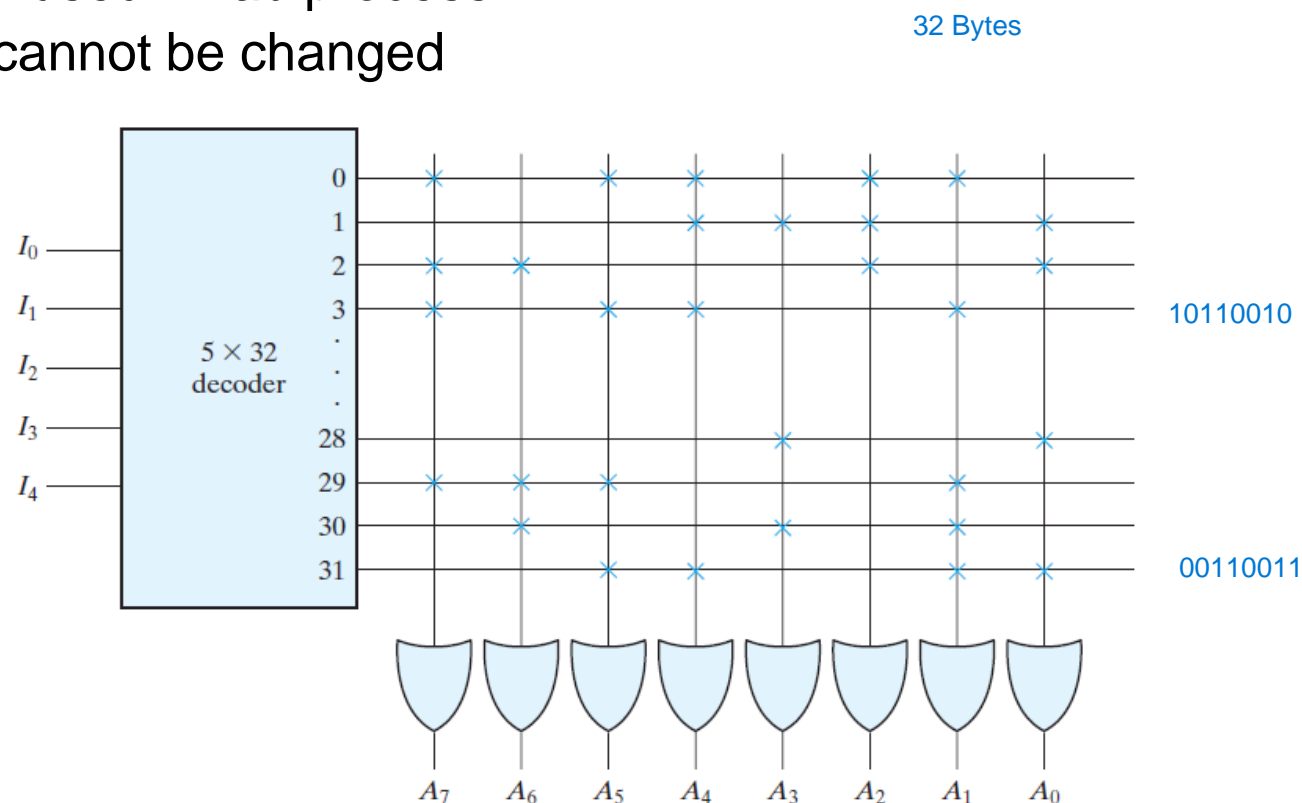
- Volatile memory: temporary storage that loses its content when power is turned off
 - SRAM
 - DRAM
- Non-volatile memory: retains stored information even when power is disconnected
 - Disk
 - ROM (Read-Only Memory):
 - Contains permanent, pre-programmed data.
 - Retains information across power cycles.
 - EPROM/Flash (Erasable Programmable ROM)
 - Rewritable, non-volatile storage used for firmware, configuration, and data storage.

Volatile & Non-Volatile Memory



Mask ROM

- The “simplest” memory technology
- Presence/absence of diode at each cell denote value
- Pattern of diodes(fuse) defined by mask used in fab process
- Contents are fixed when chip is made; cannot be changed
- Good for applications where
 - Upgrading contents not an issue
 - e.g. boot ROM
- Example:
 - What are the contents:
 - When Input = 00011?
 - When Input = 11111?

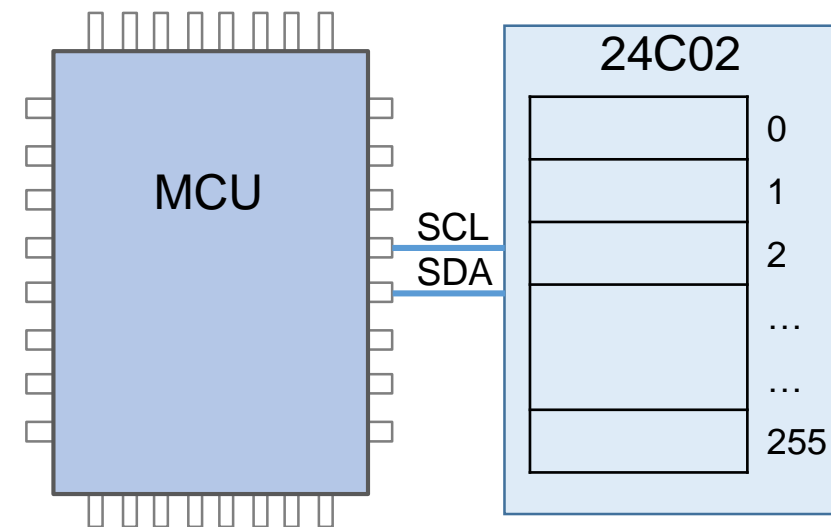


EPROM

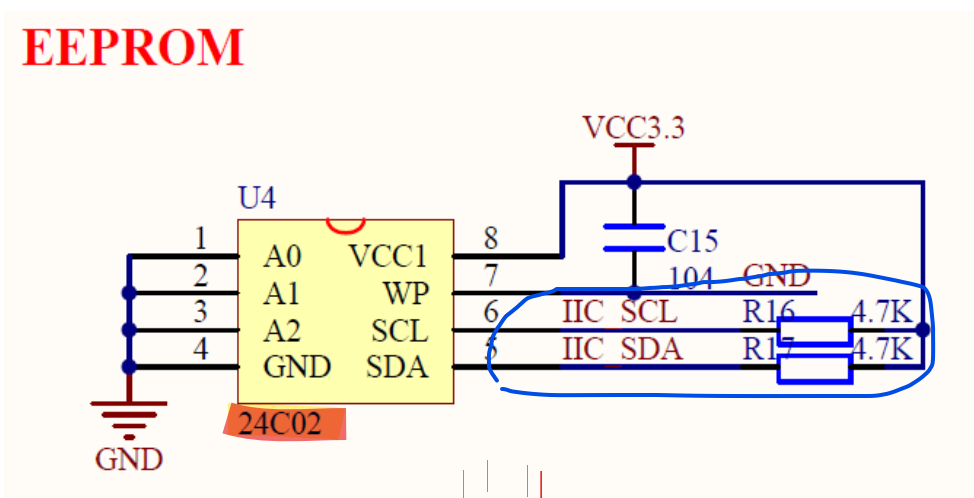
- Erasable & Programmable ROM
 - Erasing means changing from 0 \rightarrow 1
 - Uses UV light (not electrically!)
 - Writing means changing from 1 \rightarrow 0
- Erase unit is the whole device
- Retains data for 10-20 years
- Not used much these days
- Costly because
 - Use of quartz window (UV transparent)
 - Use of ceramic package

EEPROM

- **Electrically** Erasable & Programmable ROM
- Typical EEPROM: 24C02
 - 256 bytes
 - Organized and accessed in bytes
 - Random access
 - No need of erase before rewriting
- STM32 AT24C02 EEPROM (256 x 8 (2K))



EEPROM



IIC通信

baiyh@sustech.edu.cn

EEPROM Storage Example

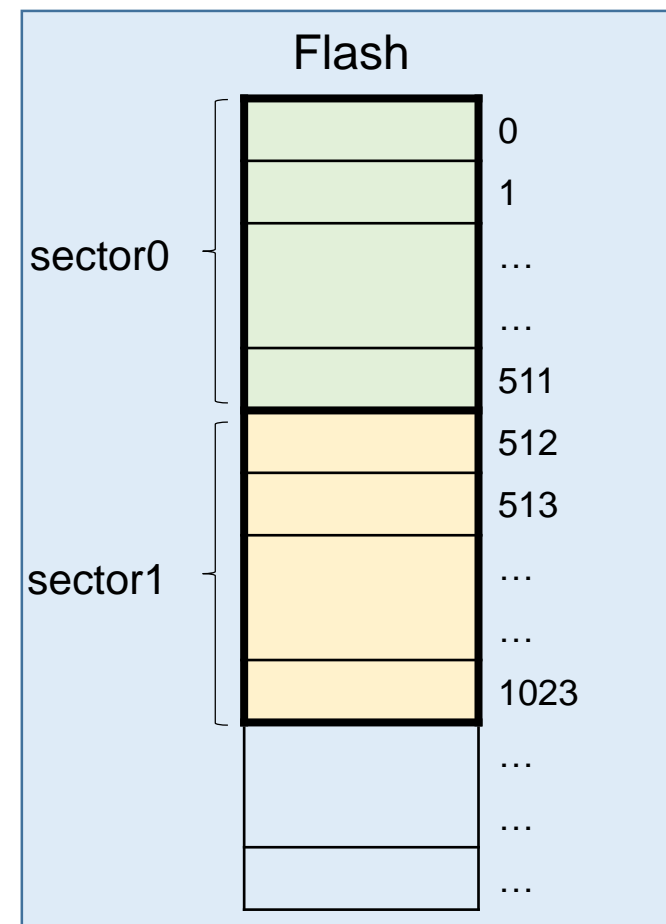
Address	Data
0	Volume
1	Brightness
2	Channel
...	...
255	...

电视机嵌入式设备为例，保存亮度，音量和频道

每个单元可以独立去写

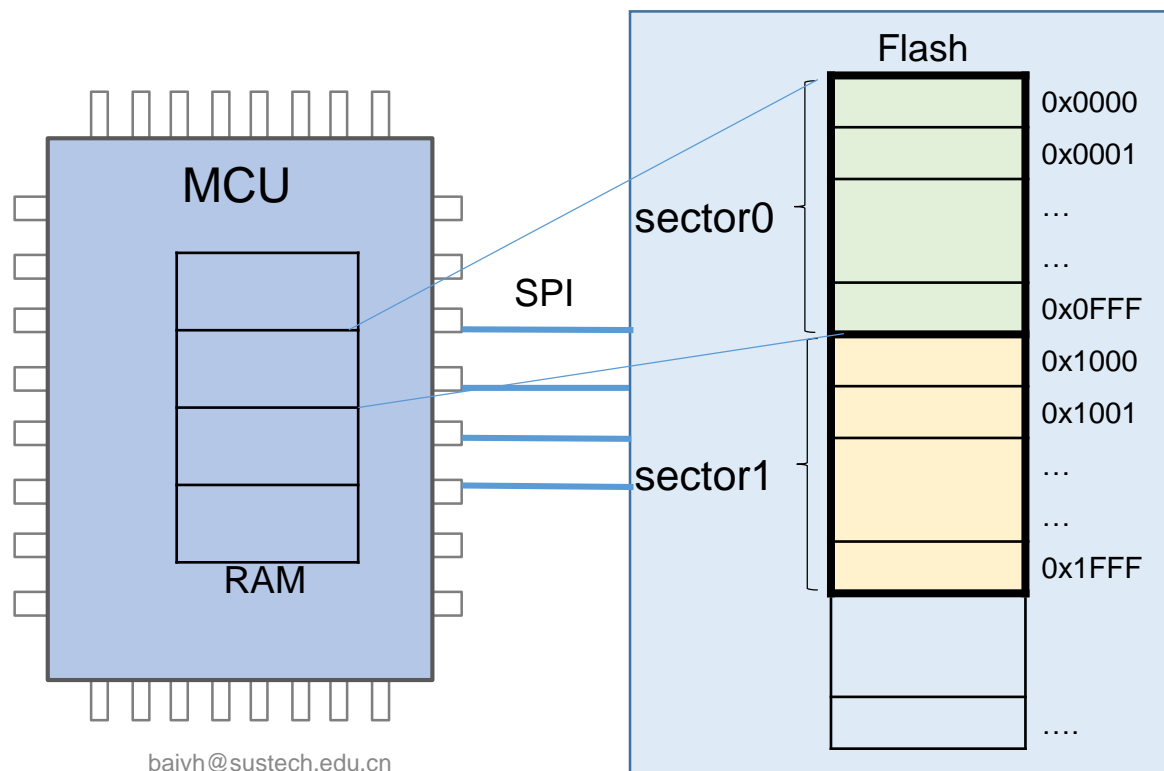
Flash Memory

- Electrically erasable (like EEPROM, unlike EPROM)
- Used in many reprogrammable systems these days
- **Sectors**(扇区):
 - refers to a fixed-size, contiguous block of storage
 - typically contains a specific number of bytes, e.g. 512
 - represents the smallest addressable unit for reading and writing data.
 - Smallest erasable unit



Flash Memory

- Read: are like standard RAM
- Write: must explicitly erase entire sector before writing
 - Erase sets entire sector contents to '1'
 - RAM can help transferring data without interference of neighbor contents
 - Fetch
 - Modify portion
 - Write
- Applications
 - Flash
 - Secure Digital Card
 - USB Flash Drive
 - Solid State Drive
 - etc



STM32 Flash Programming

- Erase before writing



Outline

- Massive Storage
- **SD Card**
- File System

SD Card

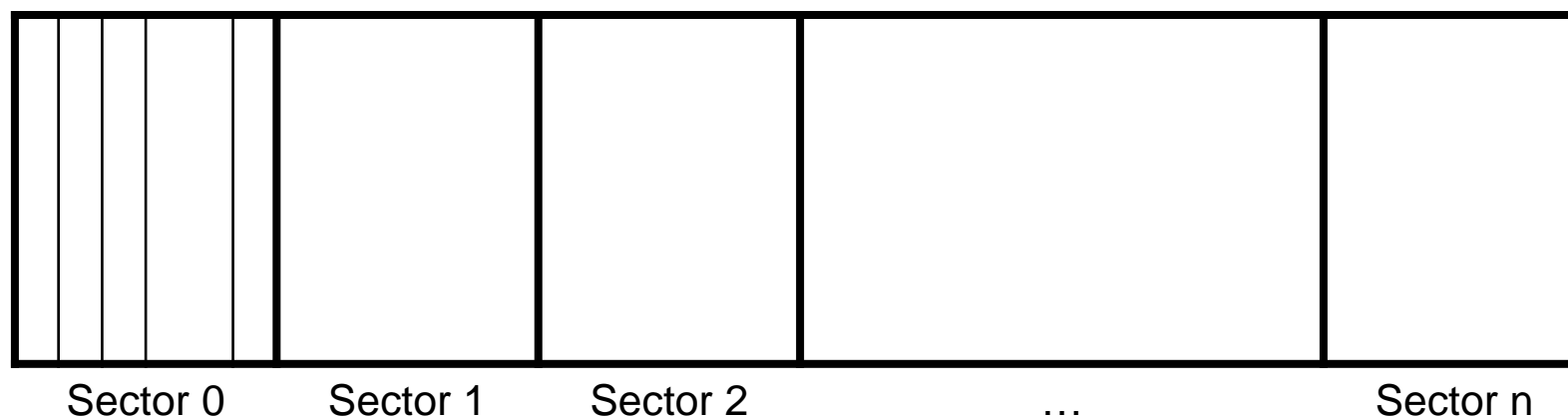
- SD Card: Secure Digital Card, **Small, portable, non-volatile** memory card for data storage.
- Underlying Technology:
 - SD cards primarily **use Flash memory** for data storage
 - Data is written and erased at the sector level.



SD Card Operation

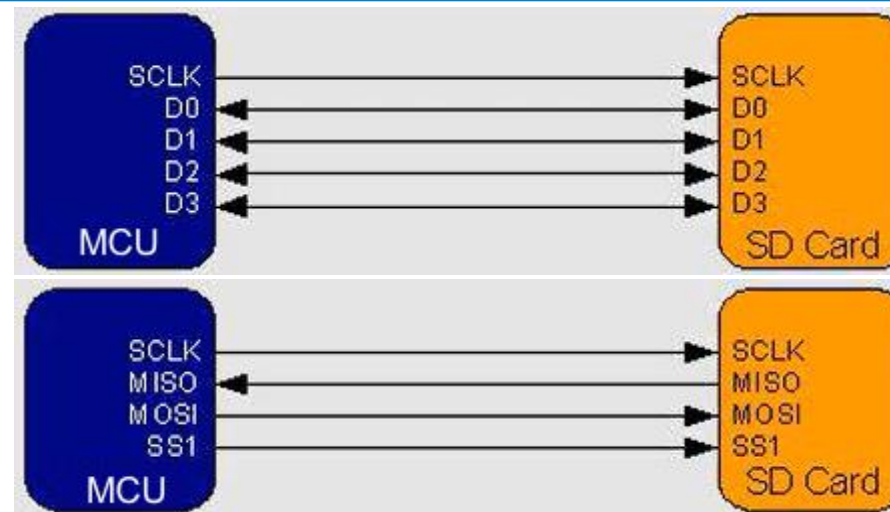
- Read a sector
- Write a sector
- Get SD Card status
- Initialize SD Card

} SD Card Driver



SD Card Interface

- SD Mode
 - Clock
 - 4 Data Lines
- SPI Mode
 - Clock
 - Card Select
 - 2 Data Lines

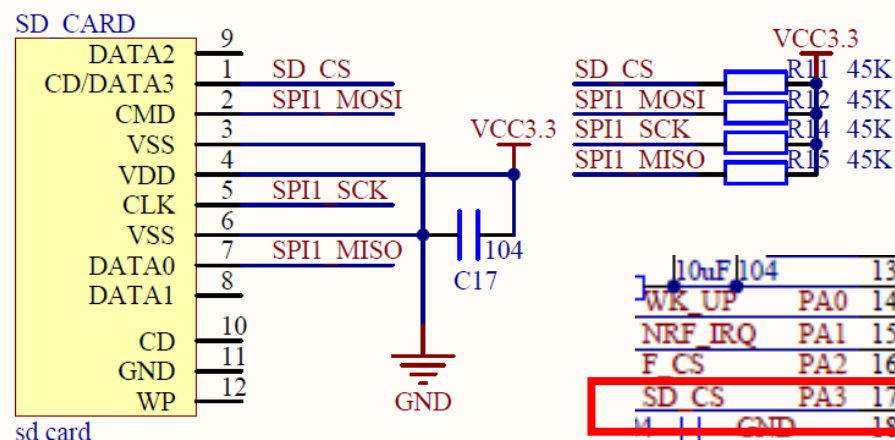


Pin	SD	SPI
1	DAT3	CS
2	CMD	DI(MOSI)
3	Vss	Vss
4	Vcc	Vcc
5	CLK	CLK
6	Vss2	Vss2
7	DAT0	DO(MISO)
8	DAT1	Reserved
9	DAT2	Reserved

STM32 SD Card

- Alientek Mini board **only supports SPI Mode**, CS pin is PA3

SD_CARD



13	VSSA	PC1
14	VDDA	PC1
15	PA0/WKUP/ADC123_IN0/TIM2_CH1_ETR/TIM5_CH1/TIM8_ETR	PC1
16	PA1/ADC123_IN1/TIM2_CH2/TIM5_CH2	PA15/JTDI
17	PA2/U2_TX/ADC123_IN2/TIM2_CH3/TIM5_CH3	F
18	PA3/U2_RX/ADC123_IN3/TIM2_CH4/TIM5_CH4	
19	VSS	
20	VDD	
21	PA4/SPI1_NSS/ADC12_IN4/DAC1_OUT	PA12/CAN_TX
22	PA5/SPI1_SCK/ADC12_IN5/DAC2_OUT	PA11/CAN_RX
23	PA6/SPI1_MISO/ADC12_IN6/TIM3_CH1/TIM8_BKIN	PA10
24	PA7/SPI1_MOSI/ADC12_IN7/TIM3_CH2/TIM8_CH1N	PA9
25	PC4/ADC12_IN14	PA
26	PC5/ADC12_IN15	PC9/TI
27	PB0/ADC12_IN8/TIM3_CH3/TIM8_CH2N	PC8/TI
28	PB1/ADC12_IN9/TIM3_CH4/TIM8_CH3N	PC7/I2S3_MCK/TI
29	PB2/BOOT1	PC6/I2S2_MCK/TI
30	PB10/I2C2_SCL/U3_TX	PB15/SPI2_MOSI/I2S
31	PB11/I2C2_SDA/U3_RX	PB14/SPI2
32	VSS	PB13/SPI2_SCK/I2S
	VDD	PB12/SPI2_NSS/I2S2_WS/I2C2_S

STM32F103RCT6

SD Commands

C.1 SD Mode Command List

Table C- 1 and Table C- 2 show the commands that are supported by SD memory and SDIO devices in both SPI and SD modes. If a command is not identified as either mandatory or optional, then it is not supported by that device.

Supported Commands	Abbreviation	SDMEM System	SDIO System	Comments
CMD0	GO_IDLE_STATE	Mandatory	Mandatory	Used to change from SD to SPI mode
CMD2	ALL_SEND_CID	Mandatory		CID not supported by SDIO
CMD3	SEND_RELATIVE_ADDR	Mandatory	Mandatory	
CMD4	SET_DSR	Optional		DSR not supported by SDIO
CMD5	IO_SEND_OP_COND		Mandatory	
CMD6	SWITCH_FUNC	Mandatory ¹		Added in Part 1 v1.10
CMD7	SELECT/DESELECT_CARD	Mandatory	Mandatory	
CMD8	SEND_IF_COND	Optional	Optional	SDHC or SDXC

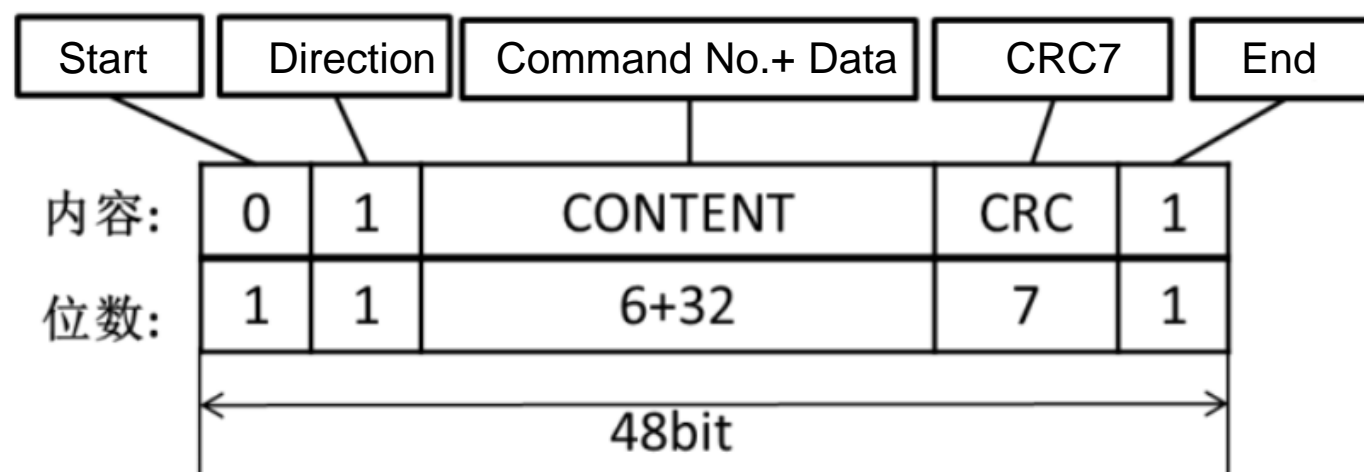
CMD9	SEND_CSD
CMD10	SEND_CID
CMD11	VOLTA
CMD12	STOP
CMD13	SEND_STATUS
CMD15	GO_IN
CMD16	SET_BLOCKLEN

C.2 SPI Mode Command List

Supported Commands	Abbreviation	SDMEM System	SDIO System	Comments
CMD0	GO_IDLE_STATE	Mandatory	Mandatory	Used to change from SD to SPI mode
CMD1	SEND_OP_COND	Mandatory		
CMD5	IO_SEND_OP_COND		Mandatory	
CMD6	SWITCH_FUNC	Mandatory ¹		Added in Part 1 v1.10
CMD9	SEND_CSD	Mandatory		CSD not supported by SDIO
CMD10	SEND_CID	Mandatory		CID not supported by SDIO
CMD12	STOP_TRANSMISSION	Mandatory		
CMD13	SEND_STATUS	Mandatory		Card Status includes only SDMEM information.
CMD16	SET_BLOCKLEN	Mandatory		
CMD17	READ_SINGLE_BLOCK	Mandatory		
CMD18	READ_MULTIPLE_BLOCK	Mandatory		
CMD24	WRITE_BLOCK	Mandatory		
CMD25	WRITE_MULTIPLE_BLOCK	Mandatory		

SD Commands

- Each Command has 48-bits



Outline

- Massive Storage
- SD Card
- **File System**

File System

- File System: a program used to manage and organize file data on a disk, facilitating operations such as searching, modifying, editing, etc.
- Typical File Systems

- **FAT** Win | 12,16,32是cluster address width
 - FAT12, FAT16, FAT32, ExFAT
- **NTFS** Win and Linux
- **ext2, ext3, ext4** Linux
- **CDFS** CD
 - CD-R, CD-RW Read/Write

SD card file systems

	Storage	File System
SD	128MB~2GB	FAT16
SDHC	4GB~32GB	FAT32
SDXC	64GB~2TB	exFAT

extended FAT

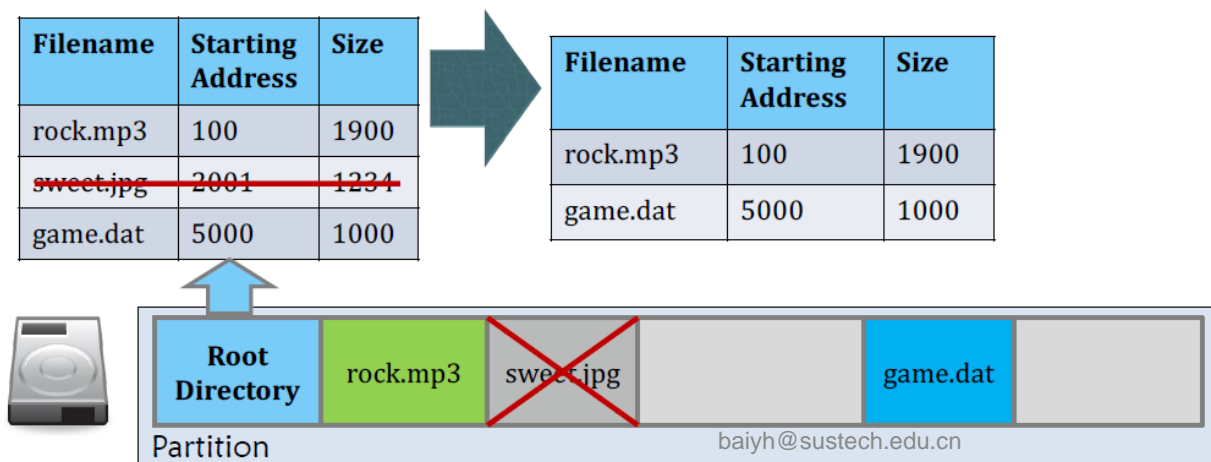
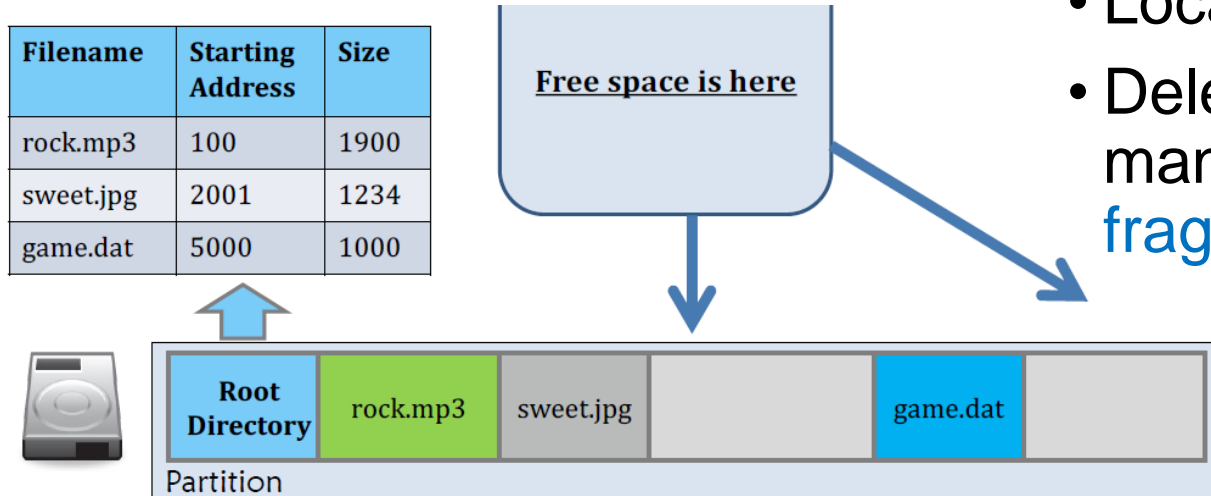


File System Structure

- I/O transfers between memory and disk are performed in units of blocks (cluster)
 - one block is one or more sectors
 - one sector is usually 512 bytes
 - block: logically, sector: physically
- Two design problems in FS
 - interface to user programs (APIs: fopen, fwrite...)
 - interface to physical storage (drivers)

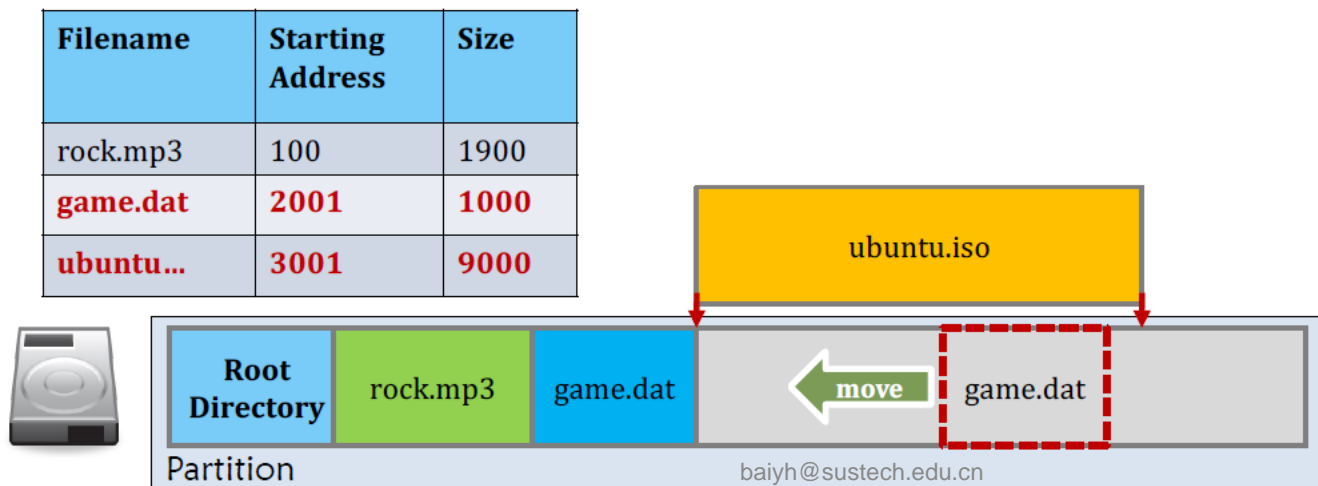
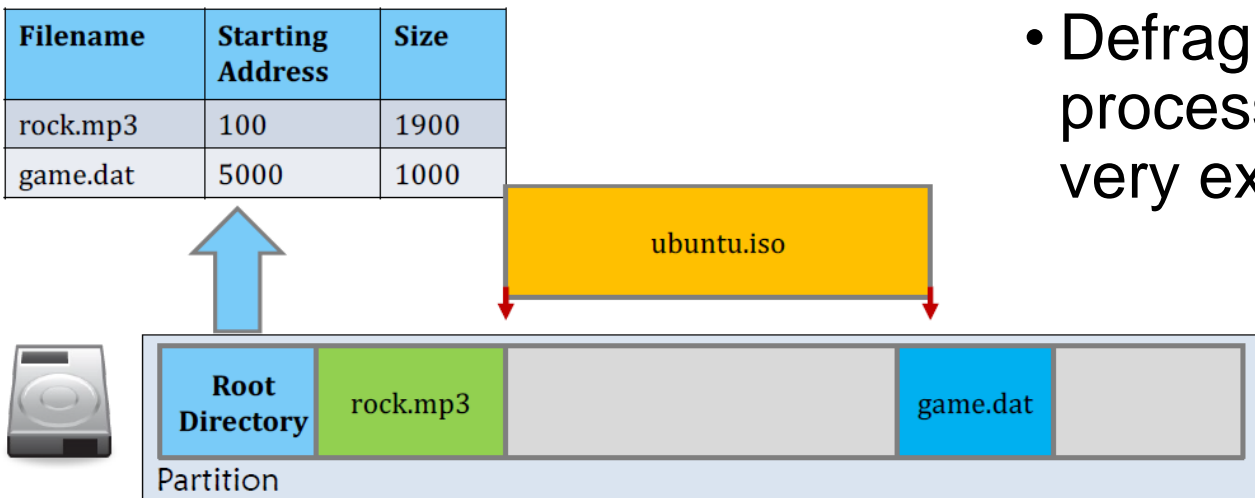
FAT Motivation

- Locate files is easy
- Deleting file will cause many **external fragmentation**



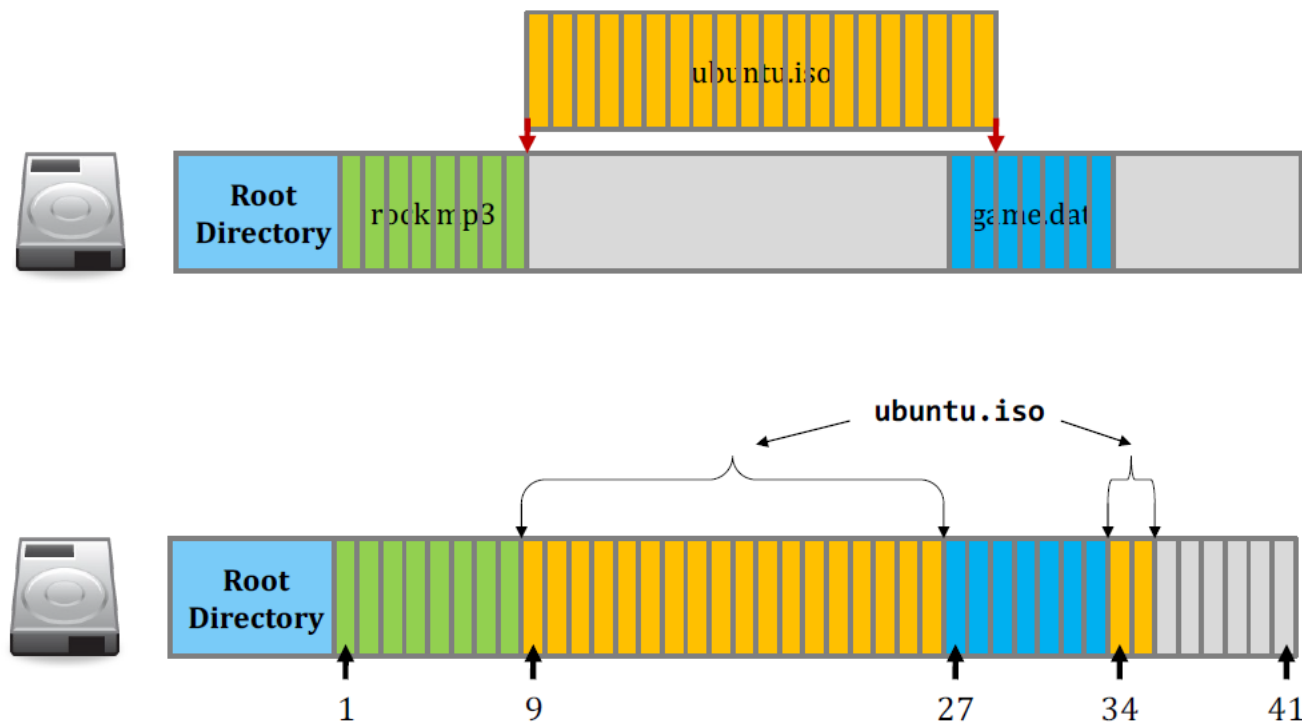
FAT Motivation

- Defragmentation process may help, but very expensive



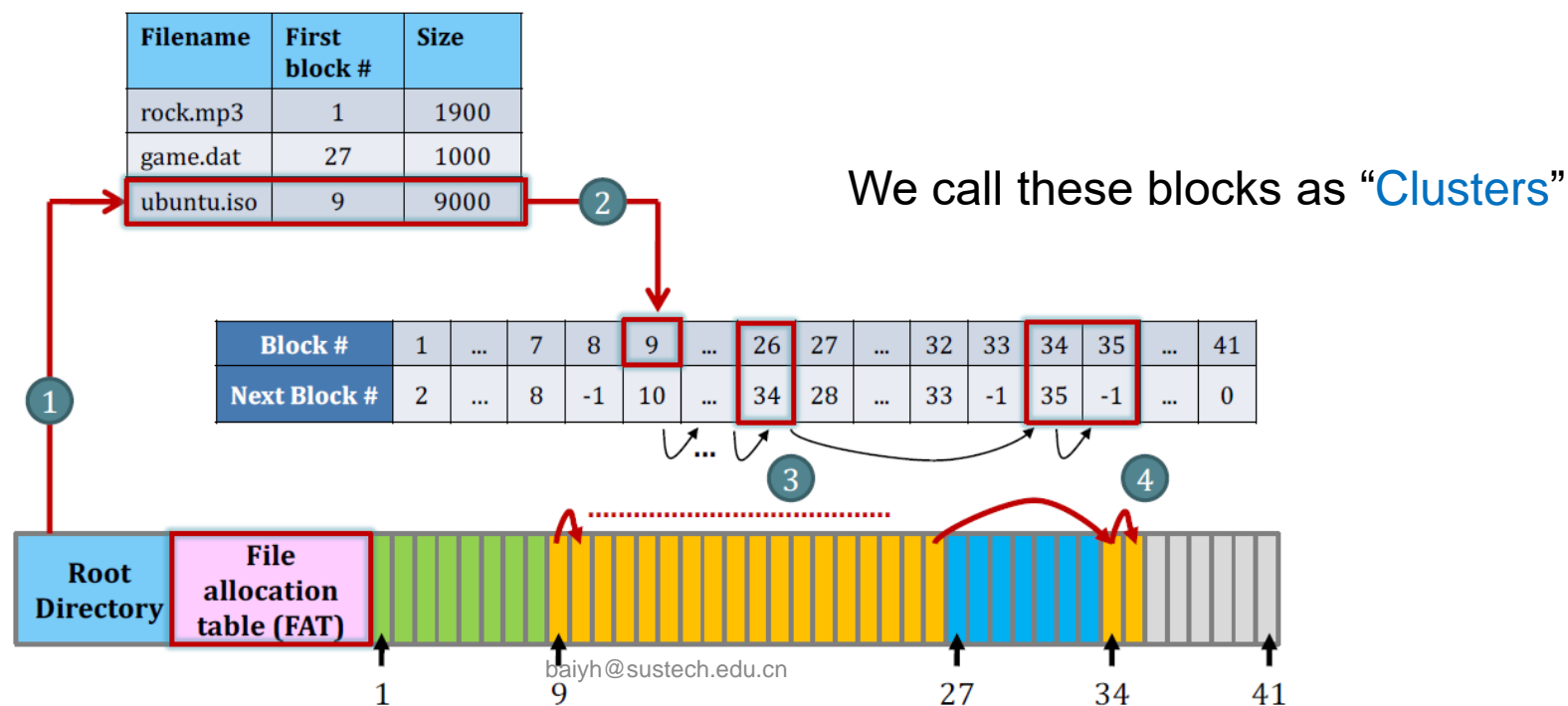
FAT Motivation

- Chop the storage device and data into equal sized blocks.
- Fill the empty space in a block-by-block manner



File Allocation Table (FAT)

- Task: read “ubuntu.iso” sequentially
 - Step 1. Read the root directory and retrieve the first block number.
 - Step 2. Read the FAT to determine the location of next block.
 - Step 3. After reading the 2nd block, the process continues. Note that the blocks may not be contiguously allocated.
 - Step 4. The process stops until the FAT says the next block # is -1.



FAT

- **Cluster(簇)**

- 1 cluster = 2^n sectors
- E.g. FAT12: 12-bit cluster address
- Can point up to $2^{12} = 4096$ clusters

Example:

Cluster size: 32KB

Cluster address width = 28bits

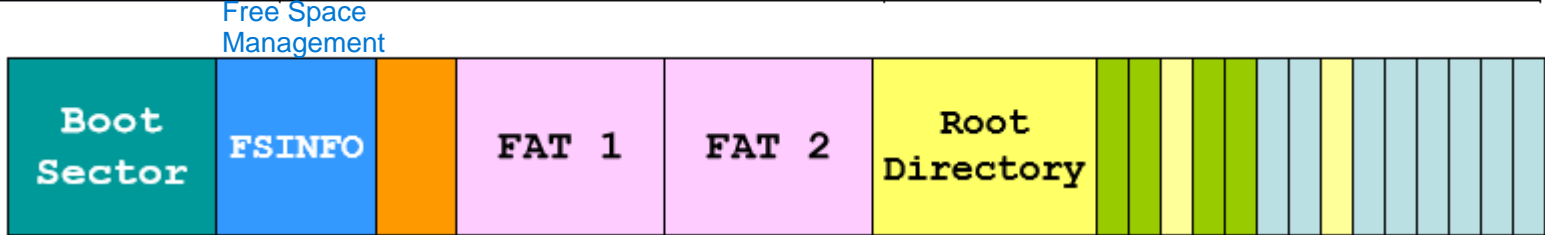
File system size = $(32 * 2^{10}) * 2^{28}$
 $= 2^5 * 2^{10} * 2^{28} = 2^{43}$ (8TB = $2^{(40+3)}$)

$$\begin{aligned}\text{FS Size} &= \text{number of Clusters} \times \text{Cluster size} \\ &= 2^{\text{cluster address width}} \times \text{Cluster size}\end{aligned}$$

	FAT12	FAT16	FAT32
Cluster address width	12 bits	16 bits	28 bits 4 bits reserved
Number of Clusters	2^{12} (4K)	2^{16} (64K)	2^{28} (256M)

FAT Layout

	Propose	Size
Reserved sectors	Boot sector	FS-specific parameters
	FSINFO	Free-space management
	More reserved sectors	Optional
	FAT (2 pieces)	1 copy as backup
	Root directory	Start of the directory tree.
		Variable, can be changed during formatting
		Variable, depends on disk size and cluster size.
		At least one cluster, depend on the number of directory entries.



第0扇区，主引导区
存储磁盘主要属性，
逻辑盘数量
分区状态
格式化信息
分区表信息

Reserved Sectors

copy of FAT1

File Read

Task: read "C:\windows\gamedata.dat" sequentially.

FAT1

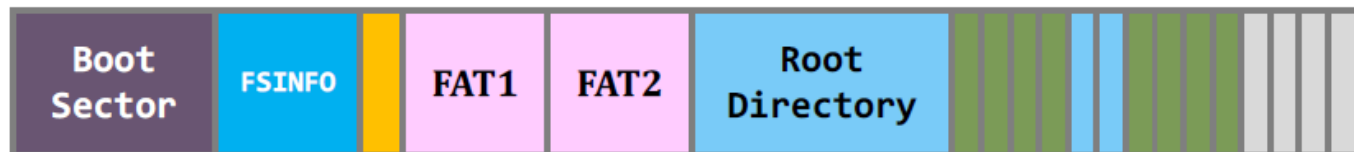
0	...
1	...
...	...
32	33
33	EOF
34	0
35	0

Damaged	= 0xffffffff7
EOF	>= 0xffffffff8
Unallocated	= 0x0

Filename	Attributes	Cluster #
gamedata.dat	32

Step 1. Read the content from Cluster #32.
Note. The **file size** may also help determining if the last cluster is reached.

Step 2. Look for the next cluster and it is Cluster #33.



File Read

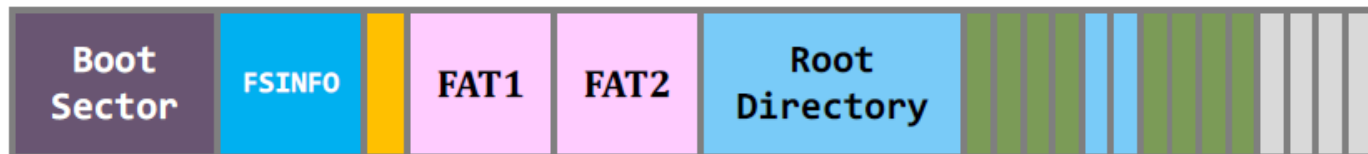
Task: read "C:\windows\gamedata.dat" sequentially.

0	...
1	...
...	...
32	33
33	EOF
34	0
35	0

Filename	Attributes	Cluster #
gamedata.dat	32

Step 3. Since the FAT has marked "EOF", we have reached the last cluster of that file.

Note. The file size help determining **how many bytes to read** from the last cluster.



File Write

Task: append data to "C:\windows\gamedata.dat".

0	...
1	...
...	...
32	33
33	EOF
34	0
35	0

Filename	Attributes	Cluster #
gamedata.dat	32

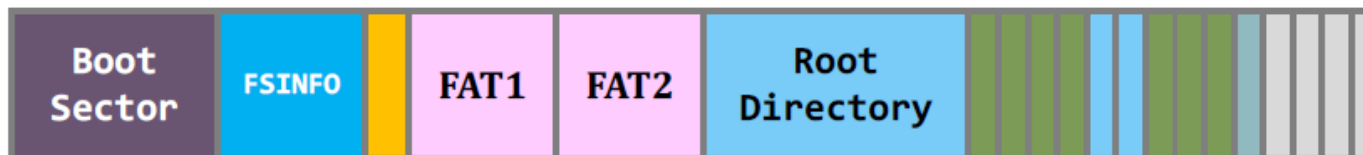
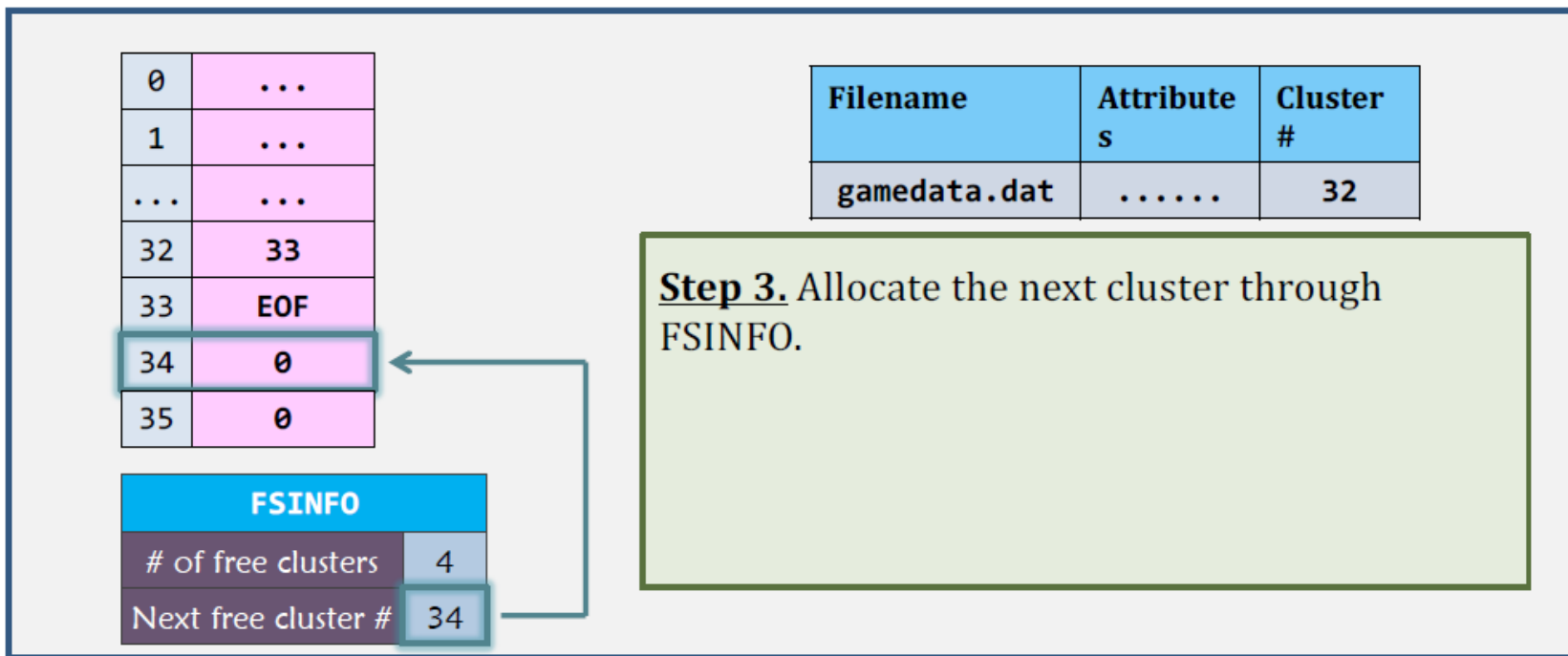
Step 1. Locate the last cluster.

Step 2. Start writing to the non-full cluster.



File Write

Task: append data to "C:\windows\gamedata.dat".



File Write

Task: append data to "C:\windows\gamedata.dat".

0	...
1	...
...	...
32	33
33	34
34	EOF
35	0

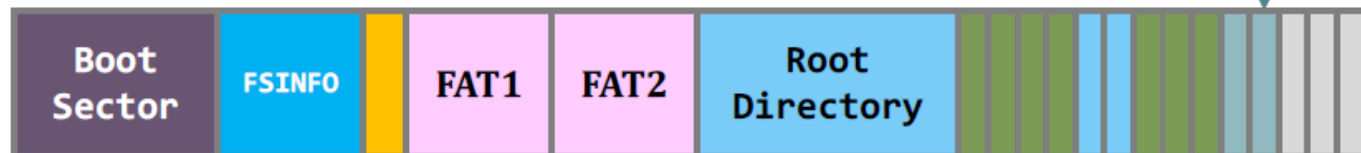
FSINFO	
# of free clusters	3
Next free cluster #	35

Filename	Attribute s	Cluster #
gamedata.dat	32

Step 3. Allocate the next cluster through FSINFO.

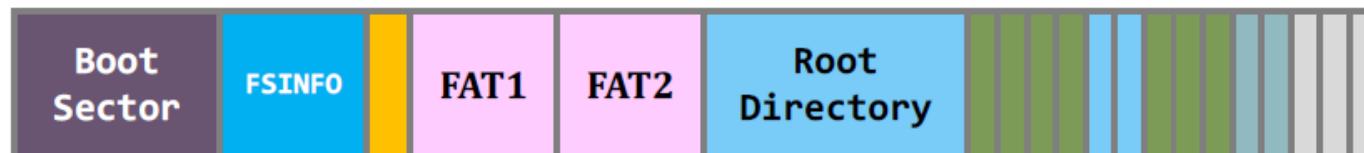
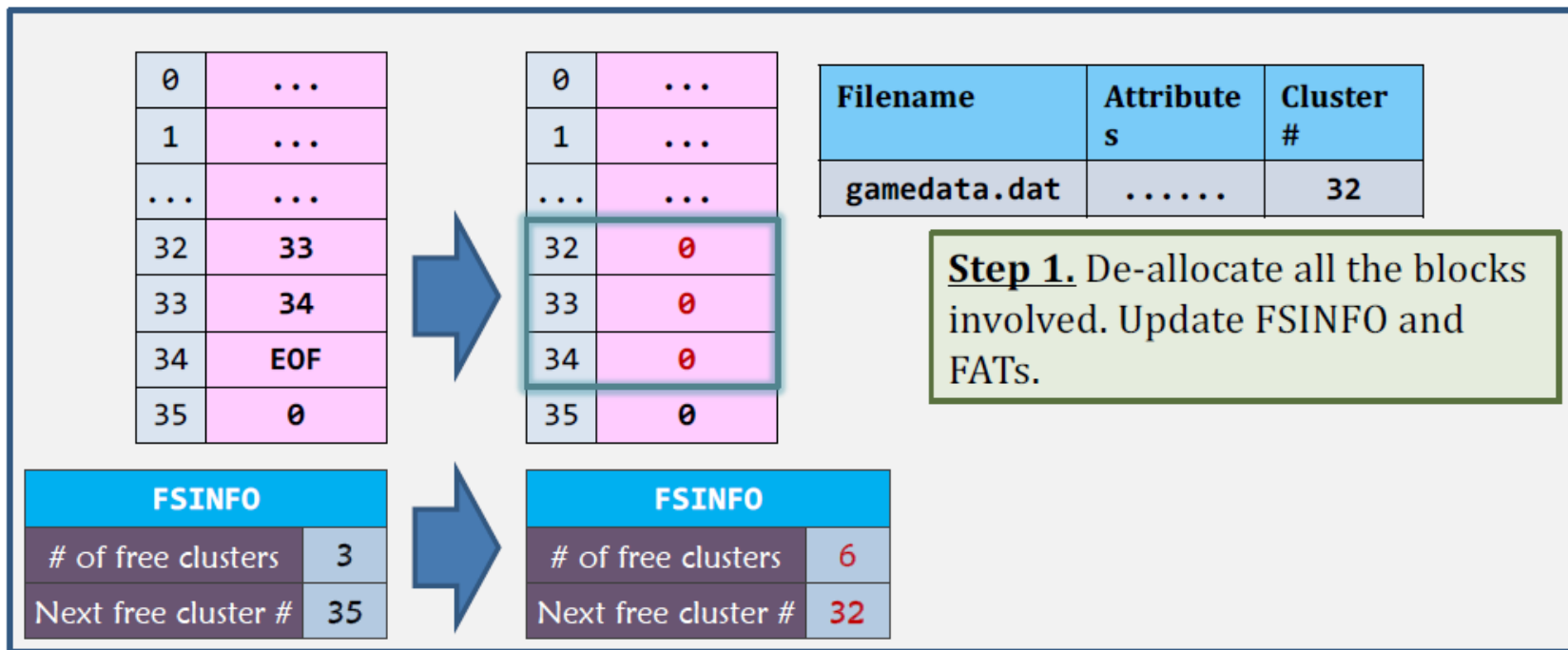
Step 4. Update the FATs and FSINFO.

Step 5. When write finishes, update the file size.



File Delete

Task: delete "C:\windows\gamedata.dat".



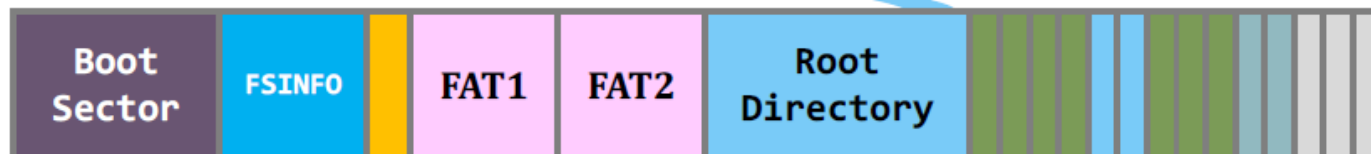
File Delete

Task: delete "C:\windows\gamedata.dat".

Directory "windows"		
Filename	Attributes	Cluster #
.	?
..	?
_amedata.dat	32
notepad.exe	456

Step 2. Change the first byte of the directory entry to _ (0xE5)

That's the end of deletion!



File Recovery

- “Deleted data” persists until the de-allocated clusters are reused.
- If you really care about the deleted file, then...
 - **PULL THE POWER PLUG AT ONCE!**
 - Pulling the power plug stops the target clusters from being over-written.

File size is within one block (cluster)	Because the first cluster address in the direct is still readable, the recovery is having a very high successful rate.
File size spans more than 1 Block	Because of the next-available search, clusters of a file are likely to be contiguous allocated. This provides a hint in looking for deleted blocks.

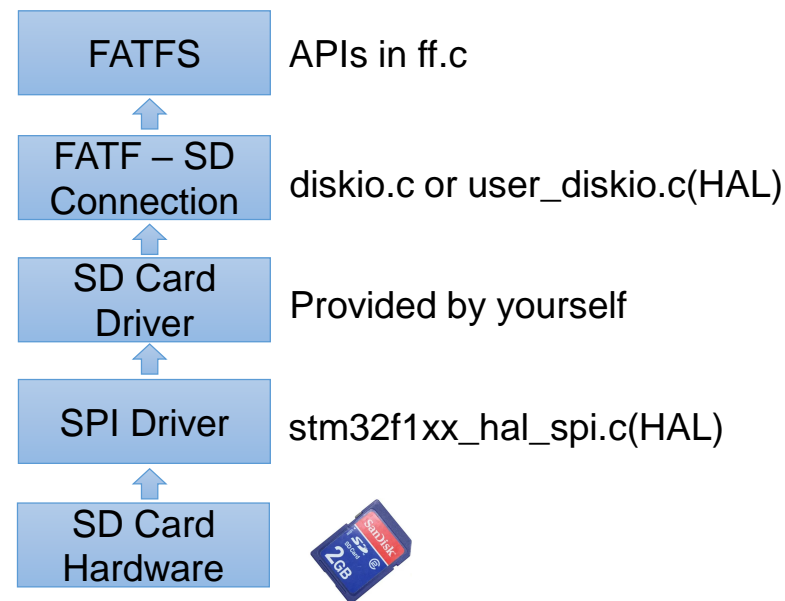
File System Implementation

- FATFS

- Generic FAT/exFAT filesystem module for small embedded systems
- Platform Independent. Easy to port.

- API

- f_mount - Register/Unregister the work area of the volume
- f_open - Open/Create a file
- f_close - Close an open file
- f_read - Read data from the file
- f_write - Write data to the file
- ...
- ...



File System Implementation

- The storage device control module is storage dependent (e.g. SD Card Driver), it needs to be **provided by implementer**
 - disk_status - Get SD card status
 - disk_initialize - Initialize SD card
 - disk_read - Read a sector
 - disk_write - Write a sector

