

# 2019-1. Discrete Math Midterm Solutions.

P1. T, F, F, F, F. (1 point each)

P2. (1) 9 (2 pts) (2) 12 (3 pts) (3) 18 (5 pts)

P3. Let  $x \in X$ . Since the matrix of relation has a nonzero entry in the row  $x$ , we have at least one  $y \in X$  such that  $(x, y) \in R$ . (2 pts)

Since  $R$  is symmetric and  $(x, y) \in R$ , we have  $(y, x) \in R$ . (2 pts)

Since  $R$  is transitive and  $(x, y), (y, x)$ , we have  $(x, x) \in R$ . (2 pts)

Therefore  $R$  is reflexive. (2 pts)

Thus  $R$  is an equivalence relation. (2 pts)

P4. Let  $d$  be the number of digits.

$$\text{Then } m = a_d b^d + a_{d-1} b^{d-1} + \dots + a_0 b^0$$

$$\text{for } 0 \leq a_i \leq b-1, (0 \leq i \leq d-1) \text{ and } 1 \leq a_d \leq b-1 \quad (5 \text{ pts})$$

$$\text{Thus } b^d \leq m \leq (b-1) \cdot (b^d + b^{d-1} + \dots + 1) = b^{d+1} - 1 < b^{d+1} \quad (5 \text{ pts})$$

By taking  $\log_b$ , we obtain

$$d \leq \log_b m < d+1 \quad (5 \text{ pts})$$

$$\text{Thus } \lfloor \log_b m \rfloor = d.$$

P5.  $2019 = 7 \cdot 283 + 38$

$$283 = 7 \cdot 38 + 17$$

$$38 = 2 \cdot 17 + 4$$

$$17 = 4 \cdot 4 + 1 \quad (5 \text{ pts})$$

So,

$$1 = 17 - 4 \cdot 4 = 17 - 4(38 - 2 \cdot 17)$$

$$= -4 \cdot 38 + 9 \cdot 17 = -4 \cdot 38 + 9(283 - 7 \cdot 38)$$

$$= 9 \cdot 283 - 67 \cdot 38 = 9 \cdot 283 - 67(2019 - 7 \cdot 283)$$

$$= -67 \cdot 2019 + 478 \cdot 283. \quad (5 \text{ pts})$$

Thus the answer is 283. (5 pts)

P6.  $\sum_{i=1}^n i^2 \lg i \leq n \cdot n^2 \lg n \Rightarrow O(n^3 \lg n). \quad (5 \text{ pts})$

$$\sum_{i=1}^n i^2 \lg i \geq \left\lceil \frac{n}{2} \right\rceil^2 \lg \left\lceil \frac{n}{2} \right\rceil + \dots + n^2 \lg n \quad (3 \text{ pts})$$

$$\geq \left( \frac{n^2}{4} \lg \frac{n}{2} \right) \cdot \frac{n}{2} \quad (2 \text{ pts})$$

$$= \frac{n^3}{8} \lg \frac{n}{2} \geq \frac{1}{8} \cdot n^3 \lg n. \quad (3 \text{ pts})$$

So  $\sum_{i=1}^n i^2 \lg i = \Omega(n^3 \lg n). \quad (2 \text{ pts})$

Therefore  $\sum_{i=1}^n i^2 \lg i = \Theta(n^3 \lg n).$

P7 We prove by induction on the number  $n$  of letters in  $\alpha$ .

If  $n=0$ , then  $\alpha$  is the null string, so  $\#a's = \#b's = 0. \quad (5 \text{ pts})$

Suppose that the statement is true for all strings in  $L$  with less than  $n$  letters and that  $\alpha \in L$  has  $n$  letters.

By the construction of  $L$ , we have either

①  $\alpha = a\beta b$  or  $b\beta a$  for some  $\beta \in L$

②  $\alpha = \beta\gamma$  for some  $\beta, \gamma \in L$  with  $\beta, \gamma \neq \alpha$ .

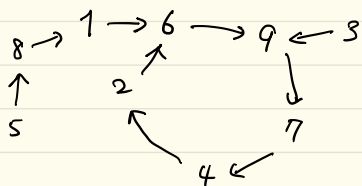
In the first case ①,  $\beta$  has the same numbers of  $a$ 's and  $b$ 's, then so does  $\alpha. \quad (5 \text{ pts})$

In case ②,  $\beta, \gamma$  have smaller number of letters.

Thus by induction hypothesis,  $\beta, \gamma$  have the same number of  $a$ 's and  $b$ 's, and so does  $\alpha. \quad (5 \text{ pts})$

Therefore by induction the statement is true for all  $n$ .

P8 Let's represent  $f$  by the following diagram with an arrow from  $i$  to  $f(i)$ :



Then for every  $i \in \{2, 6, 9, 7, 4\}$  we have  $f^5(i) = i.$

Thus  $f^2(1) = 9. \quad (2 \text{ pts})$

$f^{20}(3) = f^{19}(9) = f^4(9) = 6 \quad (2 \text{ pts})$

$f^{201}(5) = f^{198}(6) = f^3(6) = 4 \quad (5 \text{ pts})$

$f^{2019}(8) = f^{2017}(6) = f^2(6) = 7. \quad (5 \text{ pts})$