

VORLESUNG **NETZWERKSICHERHEIT**

SOMMERSEMESTER 2020 MO. 10-12 UHR





INTERNET ROUTING (ANOMALIEERKENNUNG)



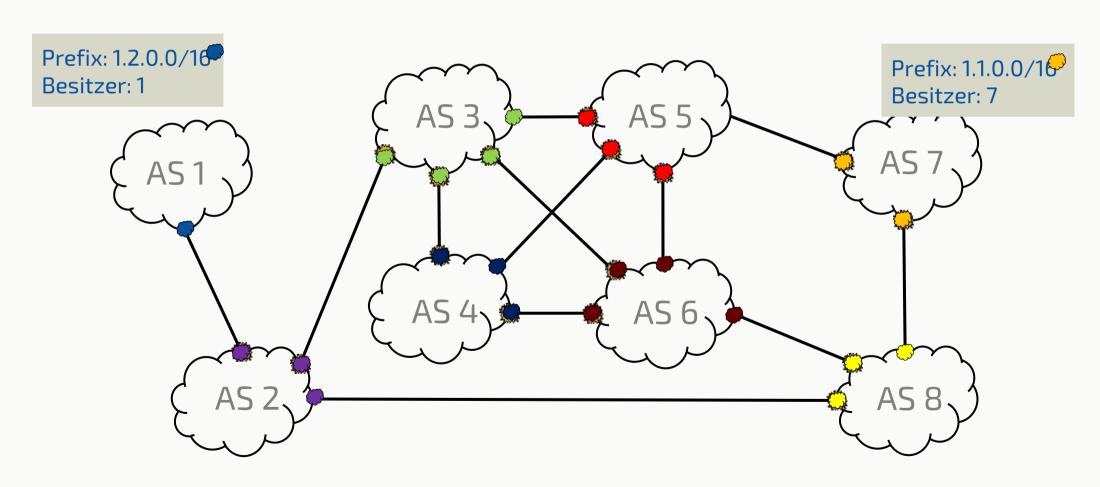
INHALTE

- Border Gateway Protocol
- Routinganomalien, Angriffsvektoren & Angreifermodel
- Topological Disorder
 - Angriffszenario "Quantuminsert"
- Prefix-Hijacking
- Analyse von Routinganomalien





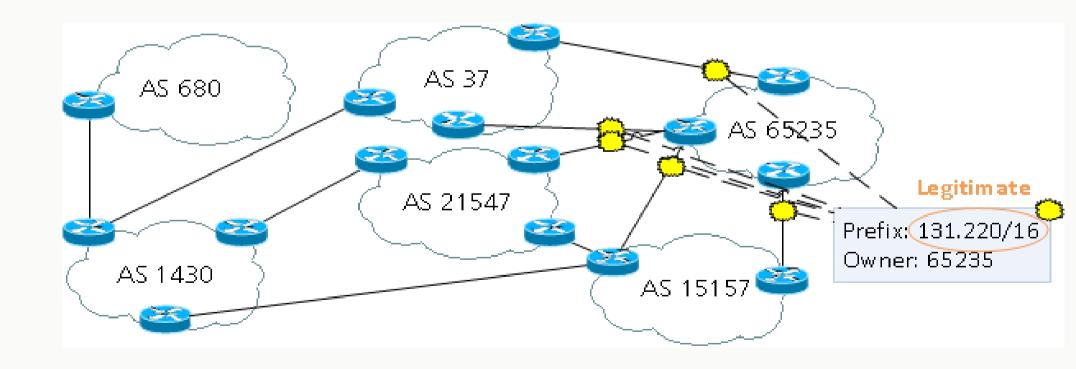
INTERNETROUTING





Annahme: AS 680 ist Angreifer (und möchte Datenverkehr abgreifen)

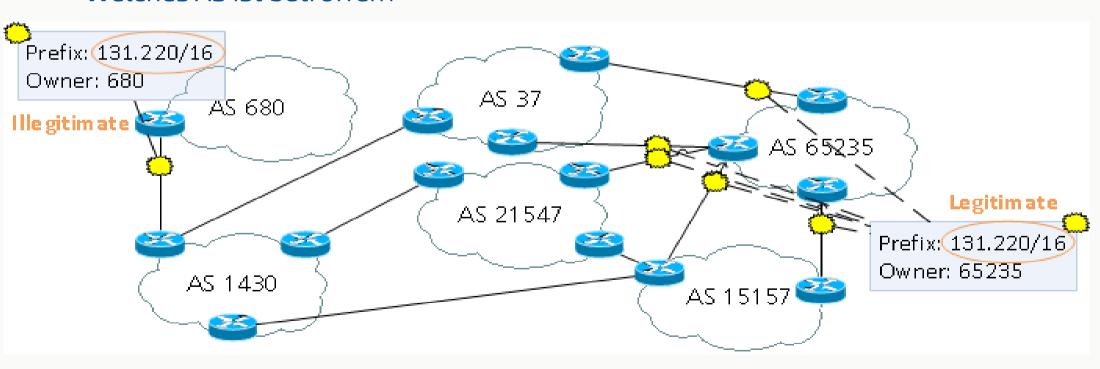
Ziel: Datenverkehr, der eigentlich für AS 65235 gedacht ist





Annahme: AS 680 ist Angreifer (und möchte Datenverkehr abgreifen)

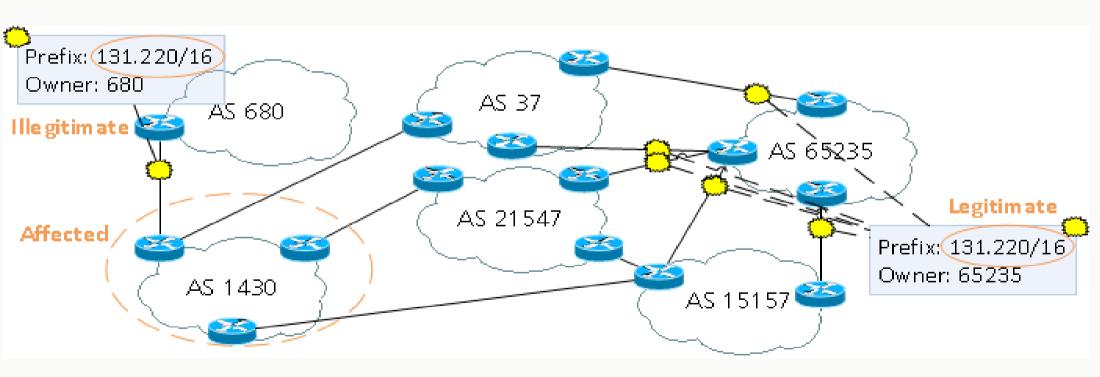
- Was kann passieren?
- Welches AS ist betroffen?





Annahme: AS 680 ist Angreifer (und möchte Datenverkehr abgreifen)

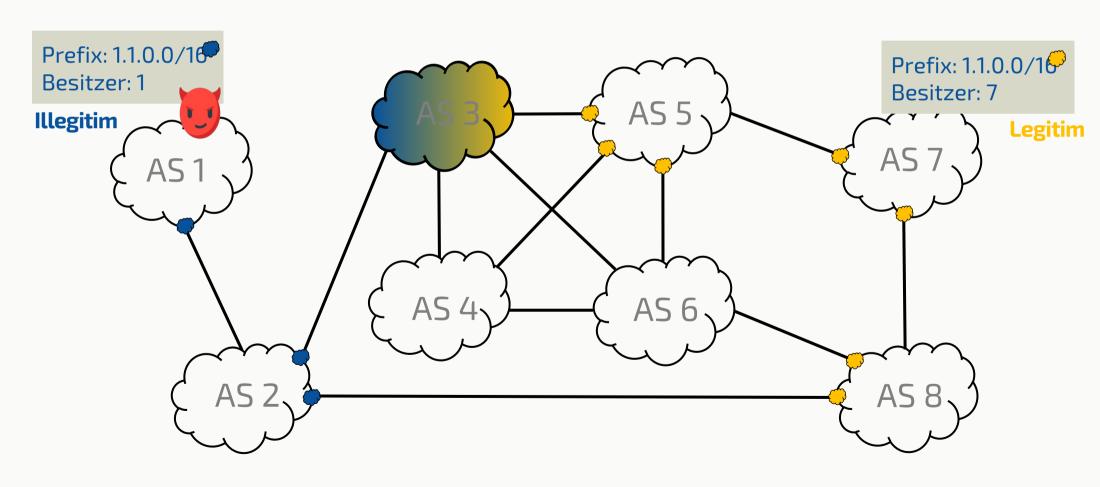
Ist dies ein realistisches Szenario (= eine echte Bedrohung)?







PREFIX-HIJACKING





- Annahme: AS 680 ist Angreifer (und möchte Datenverkehr abgreifen)
 - Ist dies ein realistisches Szenario (= eine echte Bedrohung)?





Connect with us



Reviews

News

Video

How To

Games

Download

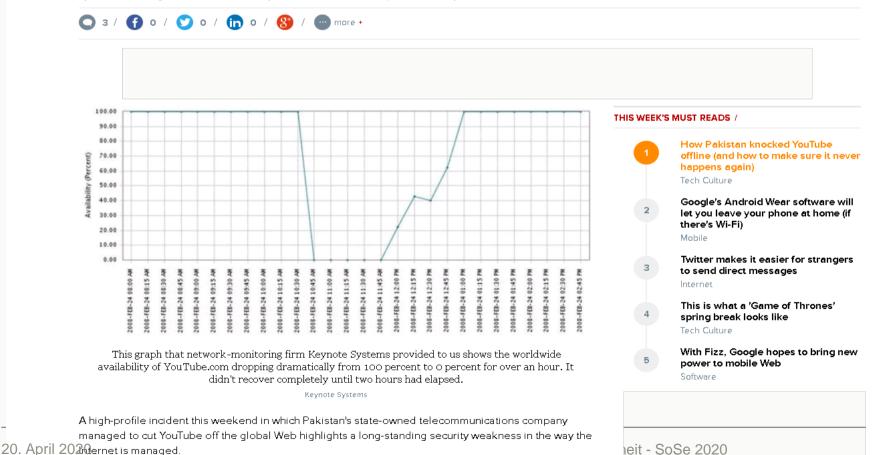


CNET > Tech Culture > How Pakistan knocked YouTube offline (and how to make sure it never happens again)

How Pakistan knocked YouTube offline (and how to make sure it never happens again)

YouTube becoming unreachable isn't the first time that Internet addresses were hijacked. But if it spurs interest in better security, it may be the last.

by Declan McCullagh y @declanm / February 25, 2008 2:30 PM PST / Updated: February 25, 2008 4:28 PM PST



45

Video

Games

Download

Log In / Join





RISK ASSESSMENT / SECURITY & HACKTIVISM

How China swallowed 15% of 'Net traffic for 18 minutes

In April 2010, 15 percent of all Internet traffic was suddenly diverted ...

by Nate Anderson - Nov 17, 2010 8:45pm CET

f Share

▼ Tweet 54

In a 300+ page report (PDF) today, the US-China Economic and Security Review Commission provided the US Congress with a detailed overview of what's been happening in China—including a curious incident in which 15 percent of the world's Internet traffic suddenly passed through Chinese servers on the way to its destination.

Here's how the Commission describes the incident, which took place earlier this year;

For about 18 minutes on April 8, 2010, China Telecom advertised erroneous network traffic routes that instructed US and other foreign Internet traffic to travel through Chinese servers. Other servers around the world quickly adopted these paths, routing all traffic to about 15 percent of the Internet's destinations through servers located in China. This incident affected traffic to and from US government (".gov") and military (".mil") sites, including those for the Senate, the army, the navy, the marine corps, the air force, the office of secretary of Defense, the National Aeronautics and Space Administration, the Department of Commerce, the National Oceanic and Atmospheric Administration, and many others. Certain commercial websites were also affected, such as those for Dell, Yahool, Microsoft, and IBM.

The culprit here was "IP hijacking," a well-known routing problem in a worldwide system based largely on trust. Routers rely on the Border Gateway Protocol (BGP) to puzzle out the best route between two IP addresses; when one party advertises incorrect routing information, routers across the globe can be convinced to send traffic on geographically absurd paths.

This happened famously in 2008, when Pakistan blocked YouTube. The block was meant only for internal use, and it relied on new routing information that would send YouTube requests not to the company's servers but into a "black hole."

As we described the situation at the time, "this routing information escaped from Pakistan Telecom to its ISP PCCW in Hong Kong, which propagated the route to the rest of the world. So any packets for YouTube would end up in Pakistan Telecom's black hole instead." The mistake broke YouTube access from across much of the Internet.

The China situation appears to have a similar cause. The mistaken routing information came from IDC China Telecommunications, and it was then picked up by the huge China Telecom. As other routers around the world accepted the new information, they began funneling huge amounts of US traffic through Chinese servers, for 18 minutes.

As with many things involving cyberattacks and Internet security, it's hard to know if anything bad

A high-prome madent this weekend in which akistans state-owned telecommunications company managed to cut YouTube off the global Web highlights a long-standing security weakness in the way the OPP and to managed

'ouTube offline (and er happens again)

et addresses were hijacked. But if it spurs interest in

08 4:28 PM PST THIS WEEK'S MUST READS / How Pakistan knocked YouTube offline (and how to make sure it never happens again) Tech Culture Google's Android Wear software will let you leave your phone at home (if there's Wi-Fi) Mobile Twitter makes it easier for strangers to send direct messages Internet FEB-24 2008-FEB-24 This is what a 'Game of Thrones' spring break looks like Tech Culture With Fizz, Google hopes to bring new s the worldwide power to mobile Web or over an hour. It Software

neit - SoSe 2020





ars te

BLOG ABOUT US PRODUCTS AND SERVICES

NEWS AND PRESS

CLIENT PORTAL

Connect with us ownload.

MAIN MENU

RISK ASS

How Chi 18 minu

In April 2010, 15 j

by Nate Anderson - No

In a 300+ page rep provided the US Co curious incident in servers on the way

Here's how the Cor

For about 18 m routes that instr Other servers a percent of the Ir traffic to and fro Senate, the arm the National Aei Oceanic and At also affected, si

The culprit here wa on trust. Routers re IP addresses; wher be convinced to se

This happened farr internal use, and it company's servers

As we described th its ISP PCCW in Ho YouTube would end from across much (

The China situation China Telecommun around the world a through Chinese se

As with many thing

Anigrepion managed to

20. April 2020 Gernet is n

How the Internet in Australia went down under

Posted by Andree Toonk - February 27, 2012 - BCP instability - 2 Comments

This Wednesday for about 30 minutes many Australians found themselves without Internet access All these users were relying either directly of indirectly on the Telstra network, which at that point was isolated from the Internet. This story quickly hit the local headlines, in this blog we'll look at the technical details of this event and what the cause of this outage likely was.

Telstra is one of Australia's major Internet providers. It normally originates approximately 500 lpv4 prefixes and 3 lpv6 prefixes. Telstra also provides Transit for many ISP's and enterprises such as for example AS38285 'Dodo' an Australian ISP and AS10235 'National Australia Bank'. So how could such a large provider go down, surely it has lots of redundant hardware and multiple connections in and out of the country?

As it turns out Wednesday's outage was caused by a routing error many network engineers have first hand experience with, a simple routing leak. A routing leak can happen when small ISP X buys transit from ISP A and also from ISP B. ISP X receives a full bgp routing table from A and because of incorrect filtering relays these messages to ISPB. As a result ISPB now learns all Internet routes via ISP X to ISP B and ISP X (the customers) now became an upstream provider for ISP B.

The above is likely what happened last Wednesday between Telstra en Dodo (AS38285). Dodo a Telstra customer, re-announced all Internet routes to Telstra, which because it prefers customer routes now thinks the best way to the Internet is through Dodo. This post on the Ausnog mailings list shows how Telstra was using Dodo (a customer) as transit to reach a network in India.

This is not a new zero day attack scenario or anything like it. Instead it's probably the number one mistake when configuring BGP routing. I remember when I was just learning about BGP my mentor always used to tell me.. Filter, Filter, Filter, filter!! Which is exactly what didn't happen here. Because it is so easy to accidentally leak routes in BGP you have to explicitly define filters that prevent this. In this case Dodo should have had filters to make sure they would only announce their prefixes and Telstra should have had these filters as well to prevent hijacks but more importantly to protect its own infrastructure. In this case these filters did not seem to be in place, which allowed this leak to happen.

However, this alone should not have brought down all of Telstra's International connections. So what happened? It's likely that Telstra now tagged all routes learned from Dodo (all 400,000 of them) as customer routes and faithfully announced this to all of its peers and upstream providers.

As keeping large filters up to date can be tedious we often see large providers use a mechanism known as max prefix limits. Instead of explicitly defining which prefixes to allow the number of prefixes expected plus some extra is set as the maximum number of prefixes allowed. This is useful to prevent a sudden spike in announcements, often caused by leaks. In case the limit is reached the PCD exection is brought down to provent the look from enreading.

ie (and ain)

t spurs interest in

EADS /

Pakistan knocked YouTube e (and how to make sure it never ens again)

Julture

le's Android Wear software will u leave your phone at home (if 's Wi-Fi)

er makes it easier for strangers nd direct messages

s what a 'Game of Thrones' g break looks like

Julture

Fizz, Google hopes to bring new r to mobile Web

)20

47



BGPMON

ars te

MAIN MENU

RISK ASS

How Chi

18 minu

In April 2010, 15

by Nate Anderson - No

In a 300+ page rep

provided the US Co

curious incident in servers on the way

Here's how the Cor

For about 18 m routes that instr

Other servers a

percent of the li

traffic to and fro

Senate, the arm the National Aei

Oceanic and At

also affected, si

The culprit here wa on trust. Routers re IP addresses; wher

be convinced to se

This happened farr internal use, and it

company's servers

As we described th

its ISP PCCW in Ho

YouTube would end

from across much (

The China situation

China Telecommun

around the world a

through Chinese se

As with many thing

PRODUCTS

SUBSCRIBE

АБОНАМЕНТ

Home

Archive

Search

Sponsors

About us

Contact

Donation

Posted by Andree Toonk - February 27, 2012 - BCP ins

All these users were relying either directly o was isolated from the Internet. This story qu technical details of this event and what the

Telstra is one of Australia's major Internet p prefixes and 3 lpv6 prefixes. Telstra also pr for example AS38285 'Dodo' an Australian such a large provider go down, surely it has and out of the country?

As it turns out Wednesday's outage was ca transit from ISP A and also from ISP B. ISP X r ISP X to ISP B and ISP X (the customers) now

routes now thinks the best way to the Intern list shows how Telstra was using Dodo (a ci

However, this alone should not have broug what happened? It's likely that Telstra now t them) as customer routes and faithfully ann

As keeping large filters up to date can be te known as max prefix limits. Instead of explic prefixes expected plus some extra is set as to prevent a sudden spike in announcemen

How the Internet in Australia:

This Wednesday for about 30 minutes man

first hand experience with, a simple routing incorrect filtering relays these messages to

The above is likely what happened last Wed Telstra customer, re-announced all Internet

This is not a new zero day attack scenario mistake when configuring BGP routing. I rem always used to tell me.. Filter, Filter, Filter, fil Because it is so easy to accidentally leak ro prevent this. In this case Dodo should have prefixes and Telstra should have had these protect its own infrastructure. In this case the this leak to happen.

BCD excelor is brought down to provent the

CTPOTO CEKPETHO

STROGO SEKRETNO

Hijacked? UK's Nuclear Weapons Data Re-Routes and Travels via Ukraine



CURRENT ISSUE

Krassimir Ivandjiiski









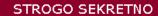
Sensitive internet data from British company Royal Mail and the UK Atomic Weapons Establishment (AWE) has passed through Russia and Ukraine via insecure connections, according to internet performance and analysis company Dyn.

An article published in technewstoday.com, suggests "web traffic originating from Texas, intended for certain addresses in the UK has been taking an unconventional route to its destination, through Ukraine and Russia".

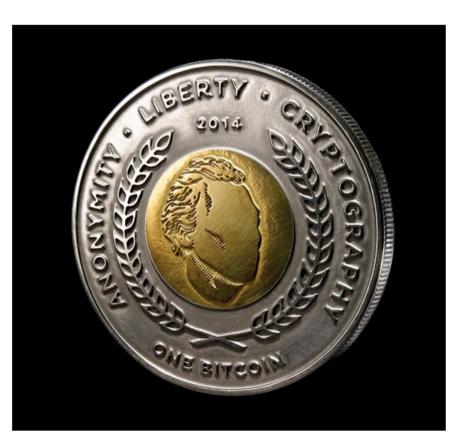
According to research carried out by Dyn, Ukrainian telecom provider Vega "began announcing 14 British Telecom (BT) routes, resulting in the redirection of Internet traffic through Ukraine for a handful of BT customers". This includes the UK's Atomic Weapons Establishment.

AWE is 'responsible for the design, manufacture and support of warheads for the United Kingdom's nuclear deterrent'.

Anigrepion managed to 20. April 2020 ernet is n



HACKER REDIRECTS TRAFFIC FROM 19 INTERNET PROVIDERS TO STEAL BITCOINS



Adam Voorhes S Gail Anderson + Joe Newton



AMONG ALL THE scams and thievery in the bitcoin economy, one recent hack sets a new bar for brazenness: Stealing an entire chunk of raw internet traffic from more than a dozen internet service

Nuclear Weapons Data Re-Routes and



om British company Royal Mail and the UK Atomic Weapons s passed through Russia and Ukraine via insecure connections, formance and analysis company Dyn.

nnewstoday.com, suggests "web traffic originating from Texas, isses in the UK has been taking an unconventional route to its ine and Russia".

ied out by Dyn, Ukrainian telecom provider Vega "began announcing utes, resulting in the redirection of Internet traffic through Ukraine ers". This includes the UK's Atomic Weapons Establishment.

e design, manufacture and support of warheads for the United



















CURRENT ISSUE

HACKER FROM 19

MAIN MENU

MY STORIES: 25

FORUMS

ARS CONSORTIUM

RISK ASSESSMENT / SECURITY & HACKTIVISM

WTF, Russia's domestic Internet traffic mysteriously passes through Chinese routers

Unexplained diversion underscores insecurity of Net's global routing system.

by **Dan Goodin** - Nov 9, 2014 7:00pm CET Share Tweet 74

Traceroutes from Moscow, RU to Yaroslavl, RU



Enlarge Dvn Research

Domestic Internet traffic traveling inside the borders of Russia has repeatedly been rerouted outside of the country under an unexplained series of events that degrades performance and could compromise the security of Russian communications

The finding, reported Thursday in a blog post published by Internet monitoring service Dyn, underscores the fragility of the border gateway protocol (BGP), which forms the underpinning of the Internet's global routing system. In this case, domestic Russian traffic was repeatedly routed to routers operated by China Telecom, a firm with close ties to that country's government. When huge amounts of traffic are diverted to far-away regions before ultimately reaching their final destination, it increases the chances hackers with the ability to monitor the connections have monitored or even altered some of the communications. A similar concern emerged last year, when Dyn found big chunks of traffic



REPEATED ATTACKS HIJACK HUGE CHUNKS OF INTERNET TRAFFIC, RESEARCHERS WARN

Man-in-the-middle attacks divert data on scale never before seen in ata Re-Routes and









ail and the UK Atomic Weapons Ukraine via insecure connections, any Dyn.

reb traffic originating from Texas, g an unconventional route to its

com provider Vega "began announcing in of Internet traffic through Ukraine omic Weapons Establishment.

pport of warheads for the United



AMON econo

braze

Adam Voorhes & Gail

traffic

belonging to US banks, government agencies, and network

the wild.

20. April 2020















THESE OPINIONS ARE THAT OF THE AUTHOR ALONE, NOT OF THEIR EMPLOYER.



Sign in





Kevin Beaumont Follow

InfoSec, from the trenches of reality. Email kevin.beaumont@gmail.com | Twitter: @gossithedog

Apr 24 · 3 min read

Hijack of Amazon's internet domain service used to reroute web traffic for two hours unnoticed

Between 11am until 1pm UTC today, DNS traffic—the phone book of the internet, routing you to your favourite websites—was hijacked by an unknown actor.



The attackers used BGP—a key protocol used for routing internet traffic around the world —to reroute traffic to Amazon's Route 53 service, the largest commercial cloud provider who count major websites such as Twitter.com as customers.

They re-routed DNS traffic using a man in the middle attack using a server at Equinix in Chicago.

From there, they served traffic for over two hours.











tions,

as,

nouncing kraine

:ed

Mögliche Ursachen für Prefix-Hijacking?

Angriff?

Fehlkonfiguration?

Software-Bug?



Mögliche Ursachen für Prefix-Hijacking?

- Angriff?
 - Ja!
- Fehlkonfiguration?
 - Das kann tatsächlich passieren (etwa verursacht durch übermüdete Administratoren an einem Freitag Nachmittag) und scheint ein sehr häufiges Phänomen/Problem zu sein
- Software failure?
 - Möglich, aber wie Topology Disorder sehr unwahrscheinlich



- Folgen von Prefix-Hijacking:
 - Nachahmung / Immitation von Kommunikationspartnern / Online-Diensten
 - Traffic Blackholing (= eine Art von Denial-of-Service)
 - Man-in-the-Middle (schwierig, aber theoretisch möglich)



INHALTE

- Border Gateway Protocol
- Routinganomalien, Angriffsvektoren & Angreifermodel
- Topological-Disorder
 - Angriffszenario "Quantuminsert"
- Prefix-Hijacking
- Analyse von Routinganomalien



Wie können Topology-Disorder und Prefix-Hijacking erkannt werden?

- Möglichkeiten zur Erkennung von Topology-Disorder:
 - Basierend auf einer vollständigen Karte des Internets (mit allen Peering-Verbindungen zwischen den AS)
 - Überprüfe jeden Pfad aller Announcements
 - Problem: eine solche Karte existiert nicht!
 - Basierend auf einer Routing-Historie (und der AS-Verbindungen)
 - Eine neu erscheinende AS-Verbindung kann bereits bekannt sein (und lediglich für einen gewissen Zeitraum nicht genutzt)
 - Was passiert beim ersten Auftauchen einer AS-Verbindung
 - Problem: Diese Methode kann die Zuverlässigkeit von AS-Verbindungen zeigen, eignet sich aber nicht zur Anomalieerkennung



Wie können Topology-Disorder und Prefix-Hijacking erkannt werden:

- Möglichkeiten zur Erkennung von Topology-Disorder
 - Basierend auf Routing-Historie, verbunden mit zusätzlichen Quellen über AS-Verbindungen (z.B. IXP oder LG) Daten.
 - Nutzung der Wahrscheinlichkeit für die Erstellung einer Verbindung von AS auf Basis bereits bekannter Verbindungen der AS
 - Gibt zusätzliche Hinweise, aber nur für eine überschaubare Anzahl an AS-Verbindungen



Zusätzliche Informationen von IXPs (Internet eXchange Points)



About Products & Services Customers & Partners News & Events Contact

CUSTOMERS & PARTNERS

Customers
What our
customers say

DE-CIX Frankfurt

DE-CIX Hamburg

DE-CIX Munich

DE-CIX New York

UAE-IX Dubai

Partners

Customers DE-CIX Frankfurt

(Peering policy information from PeeringDB)

Download JSON file (opens a new window)

CONNECTING

AS	Name	Macro	Peering-Policy
<u>4651</u>	CAT TELECOM PUBLIC COMPANY LIMITED	-	<u>Open</u>
<u>6682</u>	GNC-ALFA CJSC	-	Unknown
9038	Umniah Mobile Company	<u>AS-9038</u>	Unknown
<u>31055</u>	Consultix GmbH	AS-CONSULTIX	Open
<u>34086</u>	T-Systems International GmbH	<u>AS34086</u>	Unknown
<u>37100</u>	SEACOM Limited	-	Selective
<u>41059</u>	T-CIX Novatel	AS-T-CIX	Unknown
<u>42861</u>	JSC Universal Card Technologies	AS-PRIME-LINE	Open
<u>57858</u>	Inter Connects Inc	AS-FIBERGRID	Unknown
<u>62955</u>	eBay, Inc	-	Unknown
196922	Hofmeir Media GmbH	AS-HOFMEIR	Unknown

Customer I	_ogin
Username:	
Password:	
	Login



https://www.de-cix.net/customers-partners/customers/de-cix-frankfurt/



Zusätzliche Informationen von LGs (Looking Glass)



Looking Glass

Welcome to Hurricane Electric's Network Looking Glass. The information provided by and the support of this service are on a best effort basis. These are some of our routers at core locations within our network. We also operate a public route server accessible via telnet at route-server.he.net.

Show options

core1.ams1.he.net> show ip bgp summary								
Local AS Number			6939					
Number of Neighbors Configured			842, 776 up					
Number of Routes Installed			2942857 (253085702 bytes)					
Number of Routes Advertised			65124013 (5061468 entries) (242950464 bytes)					
Number of Attribute Entries			526348 (47371320 bytes)					
Neighbor Address	\$	ASN ♦	State ♦	Time ♦	Rt:Accepted 💠	Rt:Filtered ♦	Rt:Sent ♦	Rt:ToSend ♦
72.14.212.34		15169	ESTAB	173d 8h25m	342	0	23778	0
80.249.208.1		1200	ESTAB	147d 2h15m	3	1	66404	0
80.249.208.26		26496	ESTAB	12d10h 0m	493	0	66404	0
80.249.208.27		29075	ESTAB	173d 8h25m	226	0	66404	0
80.249.208.29		8304	ESTAB	4d 1h42m	27	0	66404	0
80.249.208.30		8529	ESTAB	50d 9h14m	4308	33	66404	0
80.249.208.31		16637	ESTAB	40d 9h19m	505	0	66404	0
80.249.208.32		12871	ESTAB	25d 9h57m	14	0	66404	0
80.249.208.33		559	IDLE	1h 9m25s	0	0	0	66404
80.249.208.34		1103	ESTAB	73d16h19m	189	5	66404	0
80.249.208.35		12859	ESTAB	47d18h16m	70	0	66404	0
80.249.208.37		2686	ESTAB	32d 7h33m	356	0	66404	0
80.249.208.38		4589	ESTAB	33d 9h57m	251	0	66404	0
80.249.208.39		112	ESTAB	3d22h 1m	1	0	66404	0
80.249.208.42		9145	ESTAB	41d 9h16m	219	0	66404	0
80.249.208.43		2611	ESTAB	50d 8h48m	55	0	66404	0
00 040 000 45		47045	CCT AD	70447540	0	0	CC 40 4	0

https://lg.he.net/





Wie lassen sich Topology-Disorder und Prefix-Hijacking erkennen?

- Prefix Hijacking, a.k.a. Multiple Origin ASes (MOAS) Konflikt:
 - Definition:

```
Let pre x p being associated with two paths asp_1 = (p_1; p_2; \dots; p_n) and asp_2 = (q_1; q_2; \dots; q_m) then, a MOAS con cit occurs if p_n \in q_m.
```

Drei Klassen von MOAS-Konflikten:

```
OrigTransAS: p_n = q_j (j < m)
SplitView: p_i = q_j (i < n; j < m)
DistinctPaths: p_i \in q_j (8i \ 2 \ [1::n]; j \ 2 \ [1::m]
```





Wie lassen sich Topology-Disorder und Prefix-Hijacking erkennen?

- Prefix Hijacking, a.k.a. Multiple Origin ASes (MOAS) Konflikt:
 - Im Bezug auf Prefixe könnte eine vierte Klasse "Sub-MOAS" definiert werden:

Let prefix p and prefix q be related in a form that q < p (i.e. q is a real subnet of p) and a MOAS conflict occurs for all IP addresses in $q \cap p$.



Wie lassen sich Topology-Disorder und Prefix-Hijacking erkennen?

- Prefix Hijacking, a.k.a. Multiple Origin ASes (MOAS) Konflikt:
 - Basierend auf annoncierten Prefixen
 - Teste für alle Announcements desselben Prefixes, ob es ein MOAS ist
 - Wo gibt es annoncierte Prefixe?



Wie lassen sich Topology-Disorder und Prefix-Hijacking erkennen?

- Prefix Hijacking, a.k.a. Multiple Origin ASes (MOAS) Konflikt:
 - Basierend auf annoncierten Prefixen
 - Teste für alle Announcements desselben Prefixes, ob es ein MOAS ist
 - Wo gibt es annoncierte Prefixe?
 - RIPE RIS, RouteViews, PCH, ... sammeln und archivieren BGP-Announcements



Wie lassen sich Topology-Disorder und Prefix-Hijacking erkennen?

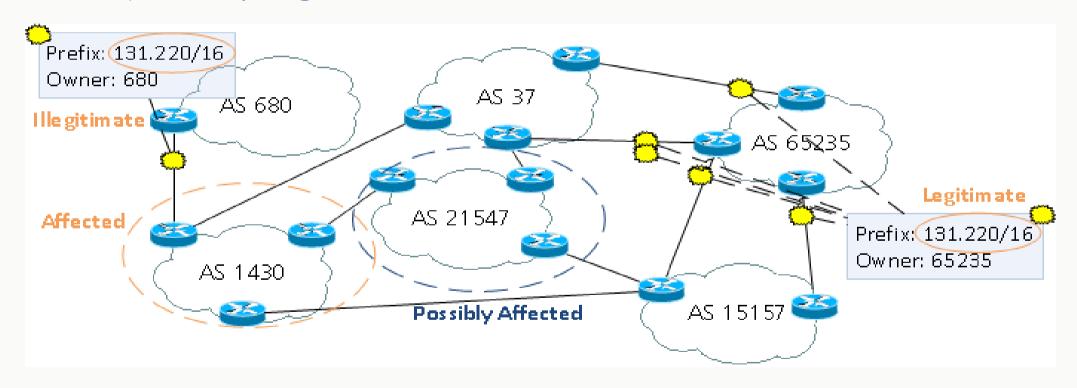
- Wo gibt es annoncierte Prefixe?
 - RIPE RIS
 - Standorte: Amsterdam, London, Paris, Geneva, Vienna, Otemachi (Japan),
 Stockholm, San Jose, Zurich, Milan, New York, Frankfurt, Moscow, Palo Alto,
 Soa Paulo
 - http://www.ripe.net/data-tools/stats/ris/
 - RouteViews
 - Oregon, Ashburn, Palo Alto, Nairobi (Kenya), London, Portland, Tokyo,
 Sydney, Sao Paulo, Atlanta, Fort Collins
 - http://www.routeviews.org/
 - PCH
 - Mehr than 50 Standorte.
 - https://www.pch.net/resources/data.php



Wie bereits gezeigt, teilt Prefix-Hijacking das Internet in verschiedene Partitionen, eine für jedes Ursprungs-AS

Wieviel Einfluss hat ein Angreifer?









Wie bereits gezeigt, teilt Prefix-Hijacking das Internet in verschiedene Partitionen, eine für jedes Ursprungs-AS

Wieviel Einfluss hat ein Angreifer?

V : the set of all ASs (vertices)

 $A = V \setminus \{t\}$: the set of all possible attackers

for a given true origin $t \in V$.

 $N = V \setminus \{t, a\}$: the remainder of V for chosen origin

 $t \in V$ and attacker $a \in A$

 $P_{(t,a,n)}$: the set of connected ASs providing

shortest paths to the origins $t \in V$ or

 $a \in A \text{ to } n \in N$





Wie bereits gezeigt, teilt Prefix-Hijacking das Internet in verschiedene Partitionen, eine für jedes Ursprungs-AS

Wieviel Einfluss hat ein Angreifer?

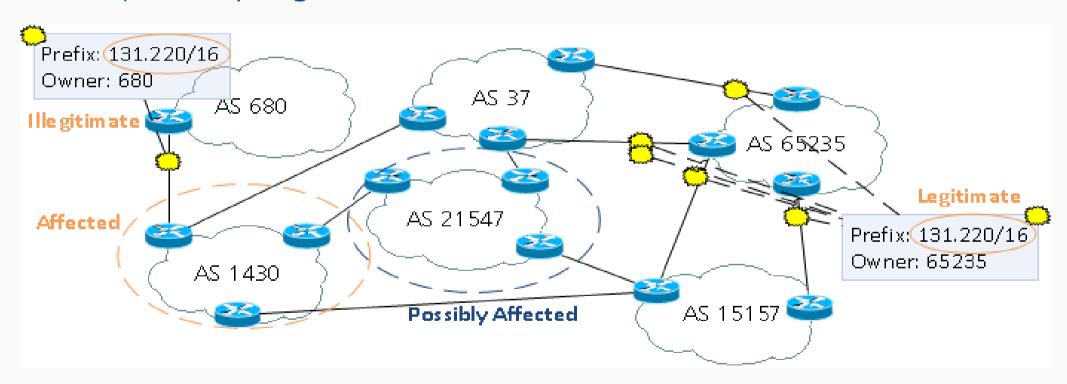
$$\beta(t, a, n) = \begin{cases} 0 & , \text{ if } Case \ 1 \\ \frac{1}{|P_{(t,a,n)}|} & , \text{ if } Case \ 2 \\ \sum_{p \in P_{(t,a,n)}} \frac{1}{|P_{(t,a,n)}|} \times \beta(t, a, p) & , \text{ if } Case \ 3 \end{cases}$$

Case 1: t is not, but a is directly connected to n.

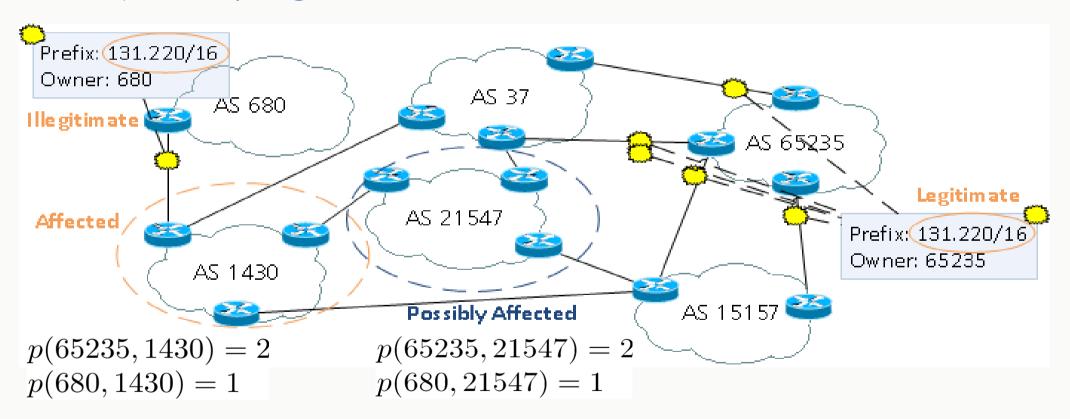
Case 2: t is and a might be directly connected to n.

Case 3: both, t and a are not directly connected to n.

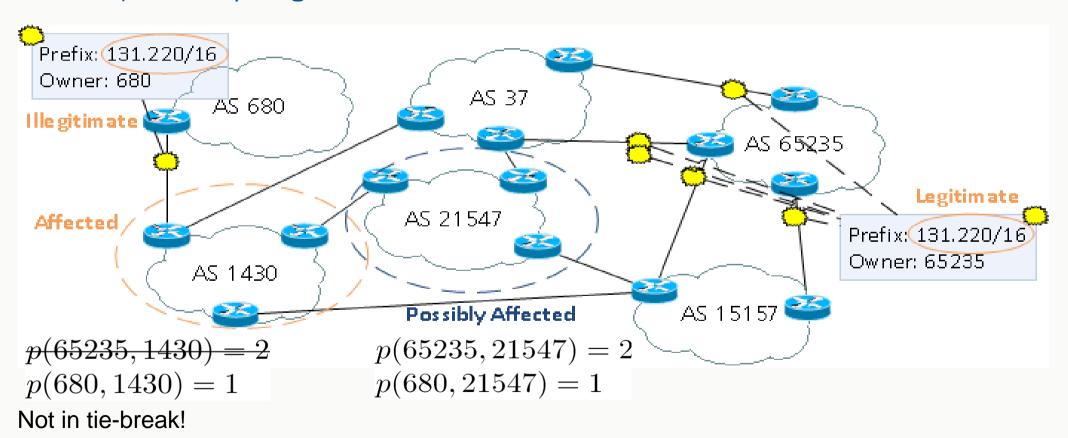




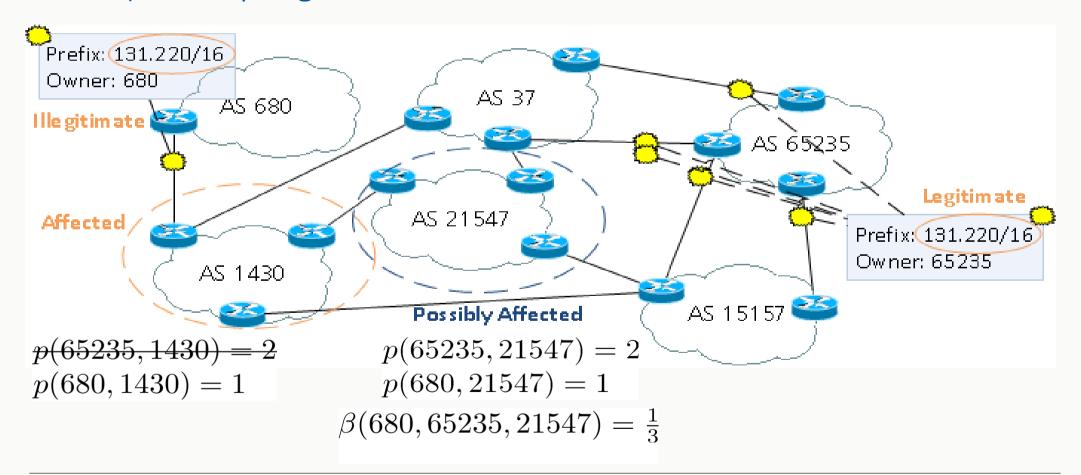














Wie bereits gezeigt, teilt Prefix-Hijacking das Internet in verschiedene Partitionen, eine für jedes Ursprungs-AS

Wieviel Einfluss hat ein Angreifer?

$$\beta(680, 65235, 21547) = \frac{1}{3}$$

- Das ist lediglich der Einfluss von AS 680 bezogen auf AS 21547 für Prefixe, die auch von AS 62543 annonciert werden
- Wie lässt sich der Einfluss (Impact) generalisieren?



Wie bereits gezeigt, teilt Prefix-Hijacking das Internet in verschiedene Partitionen, eine für jedes Ursprungs-AS

Wieviel Einfluss hat ein Angreifer?

$$\beta(680, 65235, 21547) = \frac{1}{3}$$

- Das ist lediglich der Einfluss von AS 680 bezogen auf AS 21547 für Prefixe, die auch von AS 62543 annonciert werden
- Wie lässt sich der Einfluss (Impact) generalisieren?

$$I(a) = \sum_{t \in V \setminus \{a\}} \sum_{v \in N} \frac{\beta(a, t, v)}{|V \setminus \{a\}||N|}$$

- Äußere Summe über |V| 1 (True Origins ohne den Angreifer).
- Innere Summer über |V| 2 ASes (ohne den Anfreifer und ohne True Origin)



Wie bereits gezeigt, teilt Prefix-Hijacking das Internet in verschiedene Partitionen, eine für jedes Ursprungs-AS

Wie widerstandsfähig ist ein AS (zur Verteidigung des eigenen Prefixes)?



Wie bereits gezeigt, teilt Prefix-Hijacking das Internet in verschiedene Partitionen, eine für jedes Ursprungs-AS

- Wie widerstandsfähig ist ein AS (zur Verteidigung des eigenen Prefixes)?
- Widerstandsfähigkeit (Resilience)

$$R(t) = \sum_{a \in A} \sum_{v \in N} \frac{\beta(t, a, v)}{|A||N|}$$

• Offensichtlich: Die Resilience eines AS ist gleich dem Impact, den das AS als Angreifer hätte, wenn die Rollen vertauscht wären.



ENDE

Vielen Dank für die Aufmerksamkeit!

Fragen?

Nächste Vorlesung:

Montag, 22. Juni 2020

Nächste Übung:

- Dienstag, 16. Juni 2020 16 Uhr
- Abgabe des Übungszettels 7 bis morgen 16 Uhr