

集合 $\{x_1, \dots, x_n\} \subset \mathbb{R}^n$
1. 凸集: $\forall x_1, x_2 \in C, \theta \in [0, 1], \theta x_1 + (1-\theta)x_2 \in C$
仿射集: $\forall x_1, x_2 \in C, \theta \in \mathbb{R}, \theta x_1 + (1-\theta)x_2 \in C$
锥: $\forall x \in C, \theta \geq 0, \theta x \in C$

凸函数: $f: C \rightarrow \mathbb{R}$
定义: $\forall x, y \in C, \theta \in [0, 1], f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$
仿射函数: $f(x) = a^T x + b$
凹函数: $f: C \rightarrow \mathbb{R}$
定义: $\forall x, y \in C, \theta \in [0, 1], f(\theta x + (1-\theta)y) \geq \theta f(x) + (1-\theta)f(y)$

范数: $\|x\|_1, \|x\|_2, \|x\|_\infty$
内积: $\langle x, y \rangle$
正交: $x \perp y$
投影: $P_C(x)$
距离: $\text{dist}(x, C) = \inf_{y \in C} \|x - y\|$

子集: $C \subseteq D$
交集: $C \cap D$
并集: $C \cup D$
差集: $C \setminus D$
对称差: $C \oplus D$

线性规划: $\min c^T x$
s.t. $Ax = b, x \geq 0$
对偶规划: $\max b^T y$
s.t. $A^T y \leq c$
强对偶性: $\min c^T x = \max b^T y$

拉格朗日乘子法: $L(x, \lambda) = f(x) + \lambda(g(x) - c)$
KKT条件: $\nabla L(x, \lambda) = 0, g(x) = c, \lambda \geq 0$
凸优化: f 凸, C 凸, λ^* 最优

二次规划: $\min \frac{1}{2}x^T Qx + c^T x$
s.t. $Ax = b, x \geq 0$
半正定: $Q \succeq 0$

对偶性: $f^* = \max_{\lambda} g(\lambda)$
强对偶性: $f^* = \min f(x)$
弱对偶性: $f^* \leq \min f(x)$
强对偶性成立条件: Slater's condition

子模函数: $f: 2^S \rightarrow \mathbb{R}$
定义: $f(S) + f(T) \geq f(S \cup T) + f(S \cap T)$
拟阵: (E, \mathcal{I})

图论: $G = (V, E)$
最短路径: $\min \sum_{e \in P} w_e$
最大流: $\max \sum_{e \in E} f_e$

网络流: $\max \sum_{e \in E} f_e$
s.t. $\sum_{e \in \delta^-(v)} f_e = \sum_{e \in \delta^+(v)} f_e$
最小费用最大流: $\min \sum_{e \in E} c_e f_e$

匹配: $M \subseteq E$
完美匹配: $|M| = |V|/2$
匈牙利算法: $\max |M|$

背包问题: $\max \sum_{i \in S} v_i$
s.t. $\sum_{i \in S} w_i \leq W$
动态规划: $dp[i][j]$

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状态转移方程: $dp[i][j] = \max_k \{dp[i-1][j-k] + v_k\}$
最长公共子序列: $\max_{i,j} \{dp[i-1][j] + 1, dp[i][j-1] + 1, dp[i-1][j-1]\}$

图论: $G = (V, E)$
最小生成树: $\min \sum_{e \in T} w_e$
Kruskal's algorithm

图论: $G = (V, E)$
最短路径: $\min \sum_{e \in P} w_e$
Dijkstra's algorithm

图论: $G = (V, E)$
最大流: $\max \sum_{e \in E} f_e$
Ford-Fulkerson algorithm

图论: $G = (V, E)$
匹配: $\max |M|$
Hungarian algorithm

图论: $G = (V, E)$
最小生成树: $\min \sum_{e \in T} w_e$
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Dijkstra's algorithm

1. 1706. minimize $f(x)$ s.t. $f_1(x) \leq 0$ $h_1(x) = 0$

2. 第1问 ① (线性规划) LP (线性规划) $\min c^T x$ s.t. $Gx \leq h$ $Ax = b$

② (二次规划) QP $\min \frac{1}{2} x^T P x + q^T x + r$ s.t. $Gx \leq h$ $Ax = b$

③ (半正定规划) SDP $\min c^T x$ s.t. $x_i \in \mathbb{R}$ $x_i^T F_i + \dots \leq B$

④ (凸优化) SOCP $\min c^T x$ s.t. $\|A_i x + b_i\|_2 \leq c_i$

⑤ (几何规划) GP $\min \prod c_i x_i^{a_i}$ s.t. $\prod x_i^{b_i} = 1$

3. 几何规划 minimize $f(x)$ s.t. $f_1(x) \leq 0$ $h_1(x) = 0$

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1. Lagrange 对偶: $L(x, \lambda, \nu) = f(x) + \sum \lambda_i f_i(x) + \sum \nu_i h_i(x)$

2. 对偶问题: $\max_{\lambda, \nu} g(\lambda, \nu) = \inf_x L(x, \lambda, \nu)$

3. 强对偶: $\max_{\lambda, \nu} g(\lambda, \nu) = \min_x f(x)$

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