From Source Code to Executable

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Pre-process / Compile / Link

- Creating an executable includes multiple steps
- The "compiler" (gcc) is a wrapper for <u>several</u> commands that are executed in succession
- The "compiler flags" similarly fall into categories and are handed down to the respective tools
- The "wrapper" selects the compiler language from source file name, but links "its" runtime
- We will look into a C example first, since this is the language the OS is (mostly) written in

A simple C Example

```
    Consider the minimal C program 'hello.c':

 #include <stdio.h>
 int main(int argc, char **argv)
        printf("hello world\n");
        return 0;
• i.e.: what happens, if we do:
 > gcc -o hello hello.c
 (try: gcc -v -o hello hello.c)
```

Step 1: Pre-processing

- Pre-processing is mandatory in C (and C++)
- Pre-processing will handle '#' directives
 - File inclusion with support for nested inclusion
 - Conditional compilation and Macro expansion
- In this case: /usr/include/stdio.h
 - and all files are included by it are inserted and the contained macros expanded
- Use -E flag to stop after pre-processing:
 - > cc -E -o hello.pp.c hello.c

Step 2: Compilation

- Compiler converts a high-level language into the specific instruction set of the target CPU
- Individual steps:
 - Parse text (lexical + syntactical analysis)
 - Do language specific transformations
 - Translate to internal representation units (IRs)
 - Optimization (reorder, merge, eliminate)
 - Replace IRs with pieces of assembler language
- Try:> gcc -S hello.c (produces hello.s)

Compilation cont'd

```
.file "hello.c"
                                gcc replaced printf with puts
        .section
                  .rodata
. LC0:
       .string "hello, world!"
                                 try: gcc -fno-builtin -S hello.c
       .text
.globl main
                               #include <stdio.h>
              main, @function
        .type
main:
                              int main(int argc,
               %ebp
       pushl
                                         char **argv)
               %esp, %ebp
       movl
       andl $-16, %esp
       subl $16, %esp
                                printf("hello world\n");
       movl
               $.LCO, (%esp)
                                return 0;
       call puts
       movl
               $0, %eax
       leave
       ret
       .size
               main, .-main
              "GCC: (GNU) 4.5.1 20100924 (Red Hat 4.5.1-4)"
       .ident
                       .note.GNU-stack, "", @progbits
       .section
```

Step 3: Assembler / Step 4: Linker

- Assembler (as) translates assembly to binary
 - Creates so-called object files (in ELF format)

```
Try: > gcc -c hello.c
Try: > nm hello.o
00000000 T main
U puts
```

- Linker (Id) puts binary together with startup code and required libraries
- Final step, result is executable.
 Try: > gcc -o hello hello.o

Adding Libraries

```
    Example 2: exp.c

#include <math.h>
#include <stdio.h>
int main(int argc, char **argv)
\{ double a=2.0;
     printf("exp(2.0)=%f\n", exp(a));
     return 0;
• > gcc -o exp exp.c
    Fails with "undefined reference to 'exp'". Add: -Im
• > gcc -03 -o exp exp.c
    Works due to inlining at high optimization level.
```

Symbols in Object Files & Visibility

- Compiled object files have multiple sections and a symbol table describing their entries:
 - "Text": this is executable code
 - "Data": pre-allocated variables storage
 - "Constants": read-only data
 - "Undefined": symbols that are used but not defined
 - "Debug": debugger information (e.g. line numbers)
- Entries in the object files can be inspected with either the "nm" tool or the "readelf" command

Example File: visbility.c

```
static const int val1 = -5;
const int val2 = 10;
static int val3 = -20;
int val4 = -15;
extern int errno;
static int add abs(const int v1, const int v2) {
   return abs(v1)+abs(v2);
                                      nm visibility.o:
                                      00000000 t add_abs
int main(int argc, char **argv) {
    int val5 = 20;
                                                  U errno
    printf("%d / %d / %d\n",
                                      00000024 T main
           add abs(val1, val2),
           add_abs(val3,val4),
                                                  U printf
           add abs(val1,val5));
                                       00000000 r val1
    return 0;
                                       00000004 R val2
                                       00000000 d val3
                                       00000004 D val4
```

What Happens During Linking?

- Historically, the linker combines a "startup object" (crt1.o) with all compiled or listed object files, the C library (libc) and a "finish object" (crtn.o) into an executable (a.out)
- With current compilers it is more complicated
- The linker then "builds" the executable by matching undefined references with available entries in the symbol tables of the objects
- crt1.o has an undefined reference to "main" thus C programs start at the main() function

Static Libraries

- Static libraries built with the "ar" command are collections of objects with a global symbol table
- When linking to a static library, object code is <u>copied</u> into the resulting executable and all direct addresses recomputed (e.g. for "jumps")
- Symbols are resolved "from left to right", so circular dependencies require to list libraries multiple times or use a special linker flag
- When linking only the <u>name</u> of the symbol is checked, not whether its argument list matches

Shared Libraries

- Shared libraries are more like executables that are missing the main() function
- When linking to a shared library, a marker is added to load the library by its "generic" name (soname) and the list of undefined symbols
- When resolving a symbol (function) from shared library all addresses have to be recomputed (relocated) on the fly.
- The shared linker program is executed first and then loads the executable and its dependencies

Differences When Linking

- Static libraries are fully resolved "left to right"; circular dependencies are only resolved between explicit objects or inside a library -> need to specify libraries multiple times or use: -WI,--start-group (...) -WI,--end-group
- Shared library symbols are <u>not</u> fully resolved at link time, only checked for symbols required by the object files. <u>Full check</u> only at runtime.
- Shared libraries may depend on other shared libraries whose symbols will be globally visible

Dynamic Linker Properties

- Linux defaults to dynamic libraries:
 - > ldd hello
 linux-gate.so.1 => (0x0049d000)
 libc.so.6 => /lib/libc.so.6
 (0x005a0000)
 /lib/ld-linux.so.2 (0x0057b000)
- /etc/ld.so.conf, LD_LIBRARY_PATH define where to search for shared libraries
- gcc -Wl,-rpath,/some/dir will encode
 /some/dir into the binary for searching

What is Different in Fortran?

- Basic compilation principles are the same
 => preprocess, compile, assemble, link
- In Fortran, symbols are <u>case insensitive</u>
 => most compilers translate them to lower case
- In Fortran symbol names may be modified to make them different from C symbols (e.g. append one or more underscores)
- Fortran entry point is not "main" (no arguments)
 PROGRAM => MAIN___ (in gfortran)
- C-like main() provided as startup (to store args)

Pre-processing in C and Fortran

- Pre-processing is <u>mandatory</u> in C/C++
- Pre-processing is optional in Fortran
- Fortran pre-processing enabled implicitly via file name: name.F, name.F90, name.FOR
- Legacy Fortran packages often use /lib/cpp: /lib/cpp -C -P -traditional -o name.f name.F
 - -C : keep comments (may be legal Fortran code)
 - -P: no '#line' markers (not legal Fortran syntax)
 - -traditional : don't collapse whitespace (incompatible with fixed format sources)

Common Compiler Flags

- Optimization: -00, -01, -02, -03, -04, ...
 - Compiler will try to rearrange generated code so it executes faster
 - Aggressive compiler optimization may not always execute faster or may miscompile code
 - High optimization level (> 2) may alter semantics
- Preprocessor flags: -I/some/dir -DSOM_SYS
- Linker flags: -L/some/other/dir -lm
 - -> search for libm.so/libm.a also in /some/dir