

# Binary Tree Representation



# What are the rooted tree representations?

- \* Single array representation
  - \* Complete Binary Tree
- \* Linked List representation



# Representation of Rooted Trees

- Representing rooted trees by **linked data structures**
  - Binary trees
  - Rooted trees, in which nodes having **arbitrary number of children**



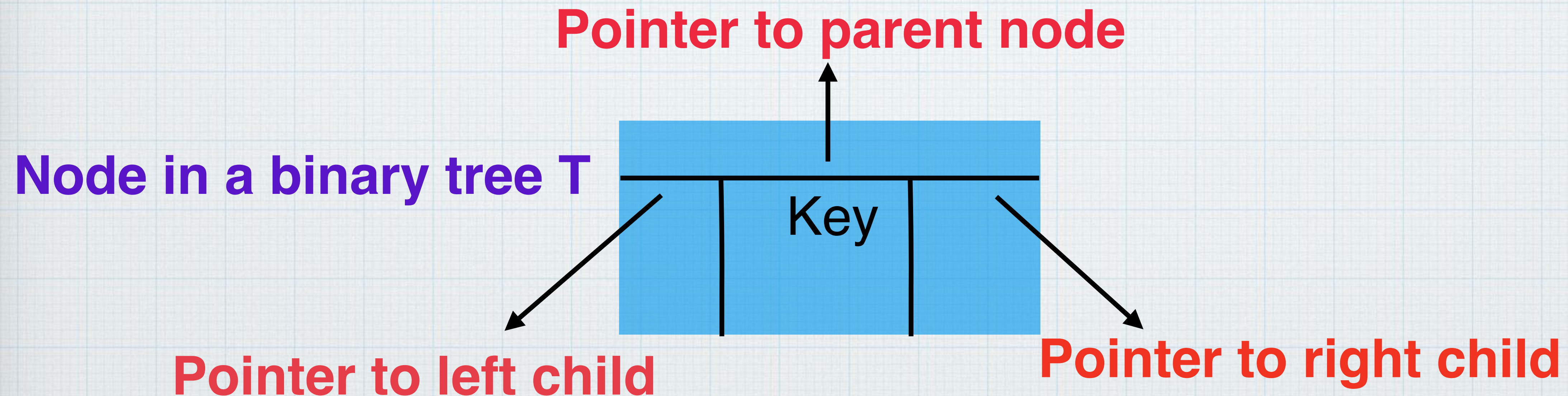
# Node in a Rooted Tree

- Represent **each node** of a tree **by an object**
  - Each **node** contains a **key attribute**
  - **Pointers** to other nodes (vary according to the type of trees)



# Binary Trees

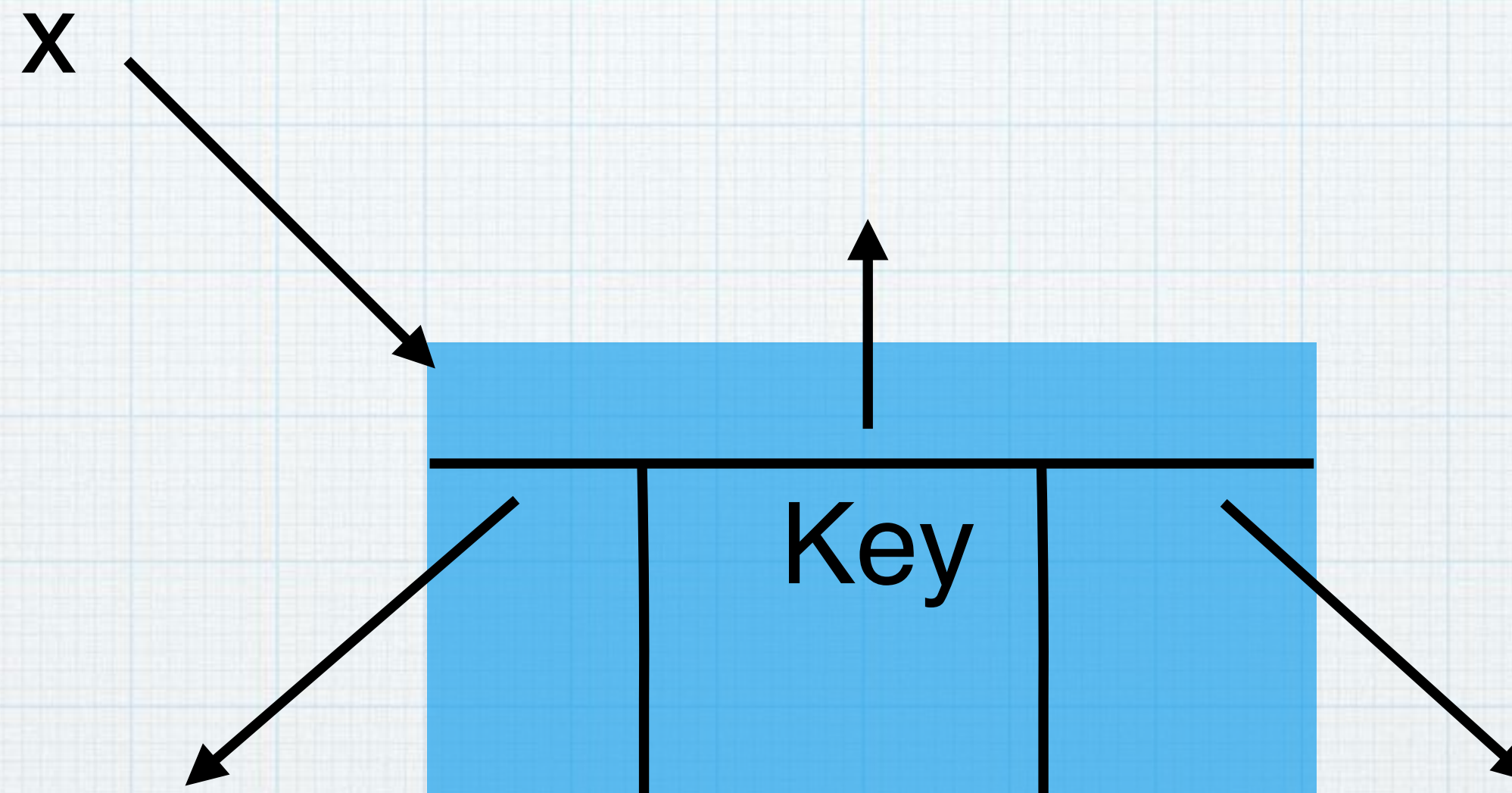
- Represent **each node** of a tree by an object
  - Each **node** contains a **key attribute**
  - **Pointers** to **Parent node**, **left child** and **right child**





# Special nodes in the Binary tree

- Node x
- $x.p = \text{NIL} \longrightarrow ?$
- $x.\text{left} = \text{NIL}$
- $x.\text{right} = \text{NIL}$
- If both  $x.\text{left}$  and  $x.\text{right}$  are NIL, then ?



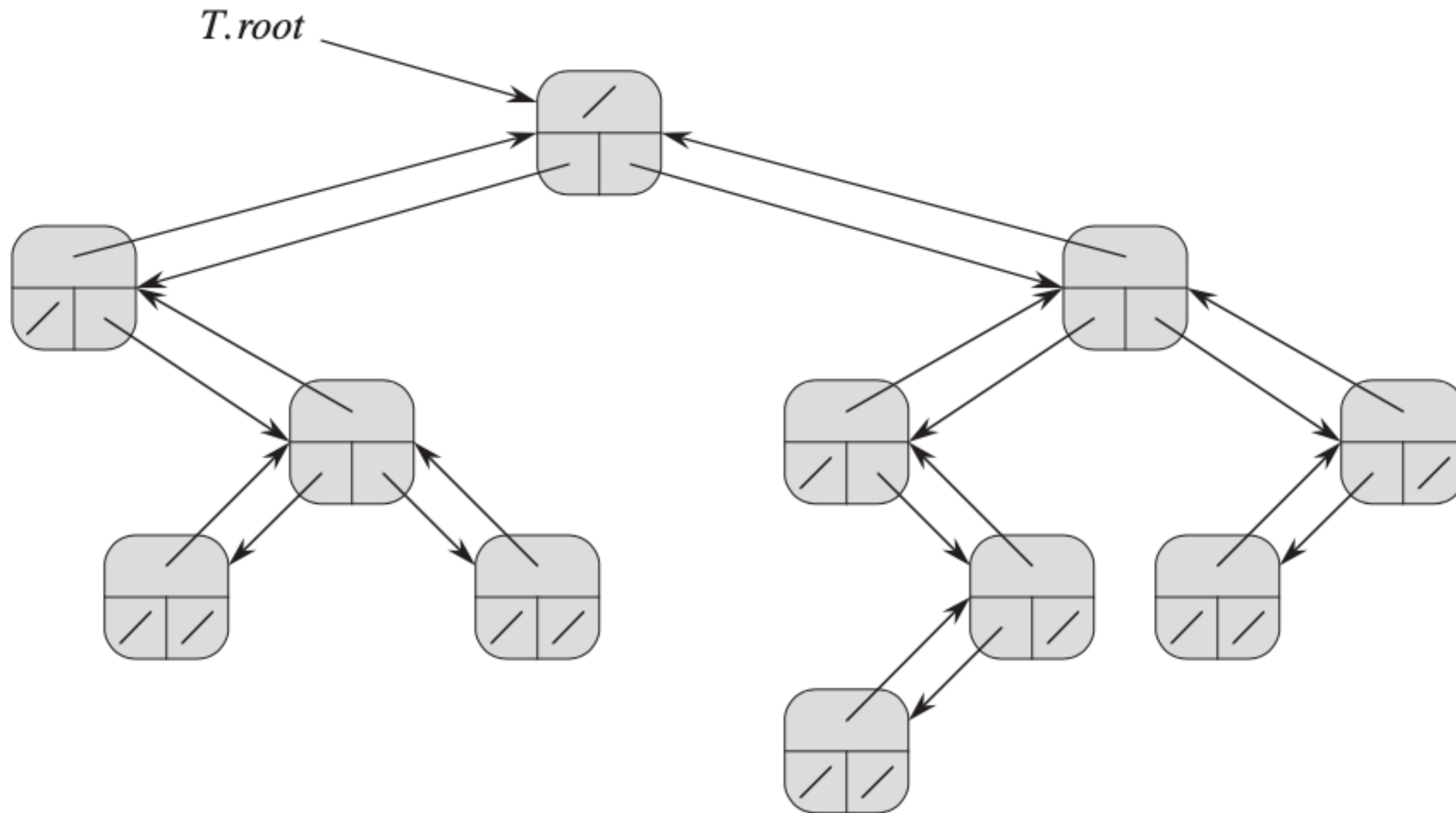


# Root of the Binary Tree

- The **root of the entire tree T** is pointed to by the attribute ***T.root***.
- If ***T.root* = NIL**, then the **tree is empty**.



# Representation of a Binary Tree





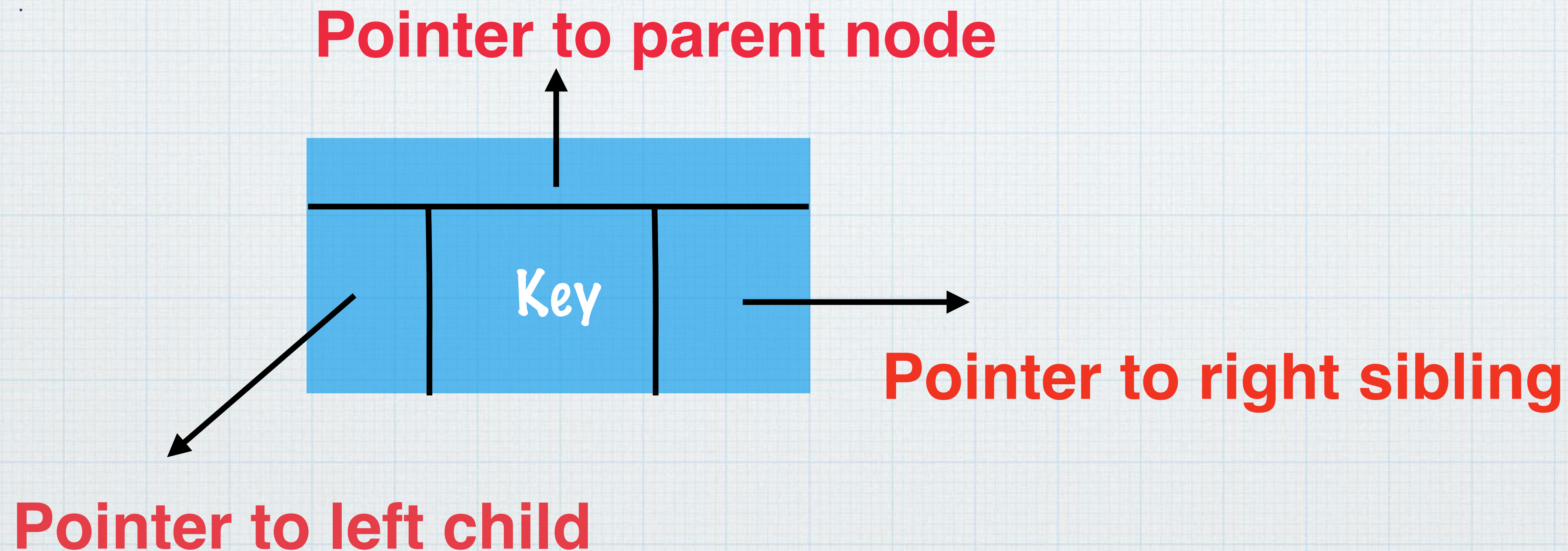
# Rooted trees with unbounded branching

- Extend the scheme of **representation of a binary tree** to **any class** of trees
- Trees in which the **number of children of each node** is **at most constant k**
- Replace the *left* and *right* attributes by  $child_1, child_2, \dots, child_k$
- Whether this scheme works, if the number of children of a node is unbounded ?
- Space requirement ?
- Even if the number of children k is bounded by a large constant but most nodes have a small number of children



# Left-child, right-sibling representation

- Scheme to represent trees with arbitrary numbers of children
  - **The left-child, right-sibling representation**
- Each node contains a parent pointer  $p$ , and  **$T.root$**  points to the root of tree  $T$ .

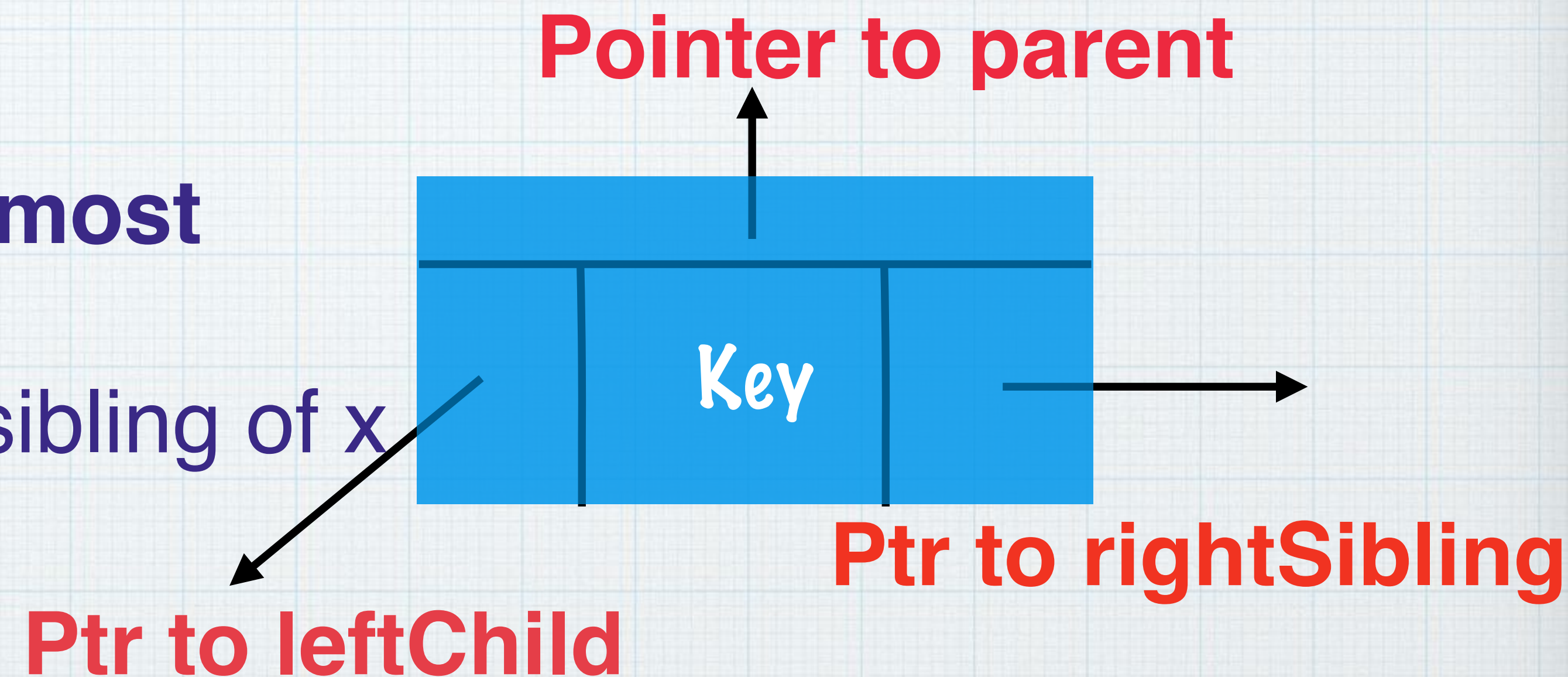




# Left Child, Right Sibling Representation

- Instead of having a pointer to each of its children, each **node x** has only **two pointers**:

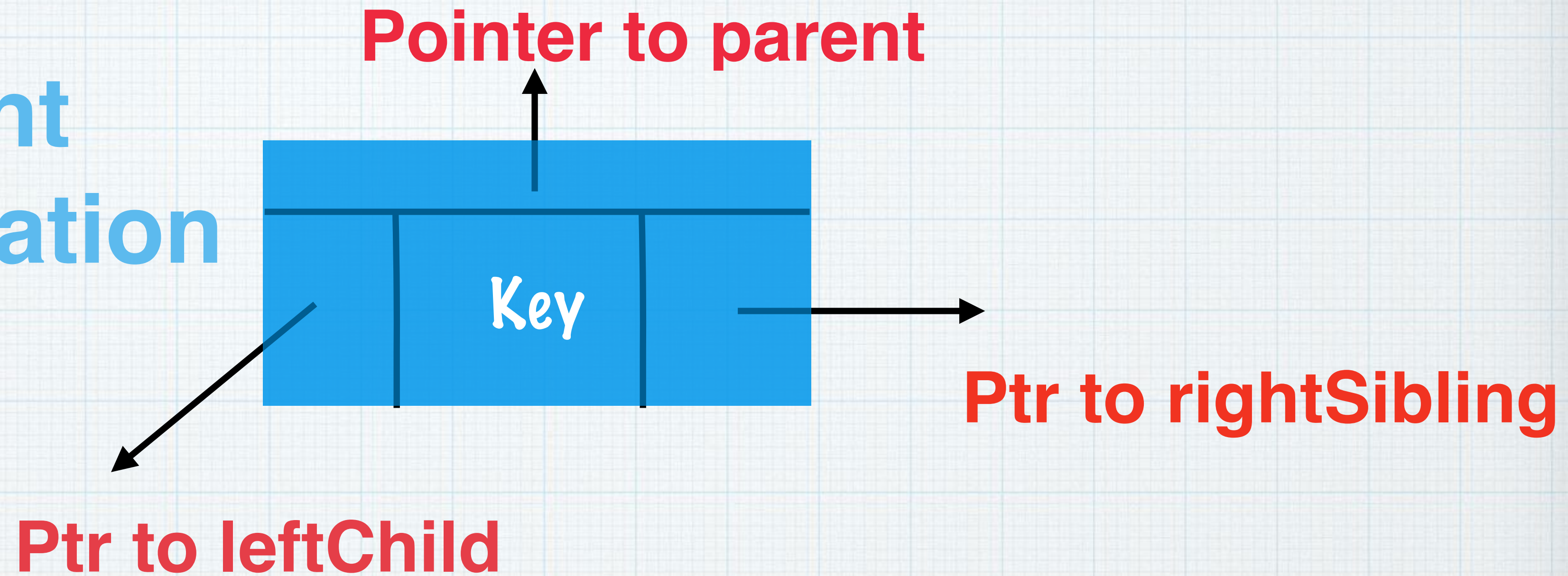
- 1. ***x.leftChild*** points to the **left most child of node x**
- 2. ***x.rightSibling*** points to the sibling of x immediately to its right.



- If **node x** has no children, ***x.leftChild* = NIL**
- If node x is the rightmost child of its parent, then ***x.rightSibling* = NIL**.



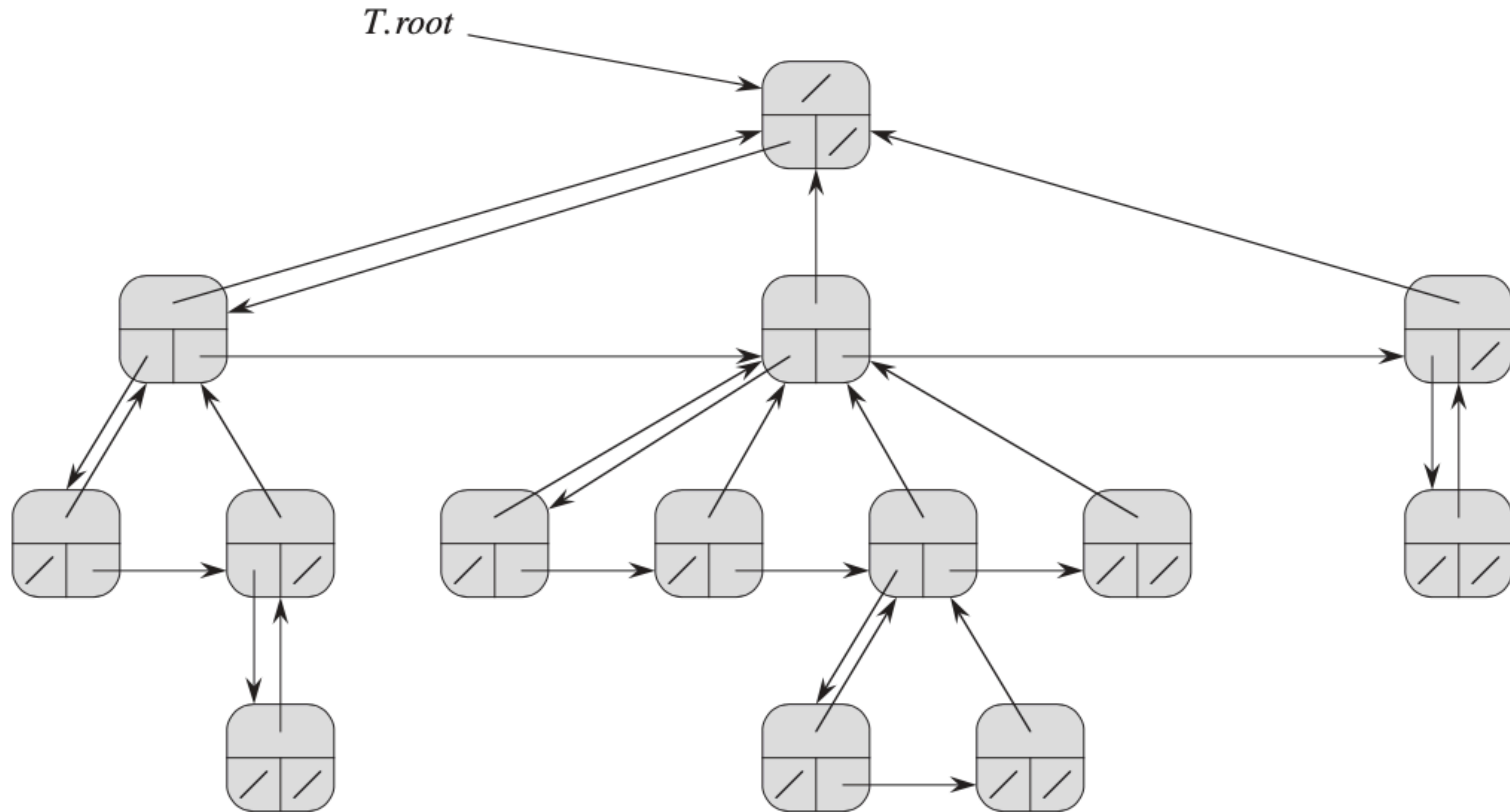
# Left Child, Right Sibling Representation



- Node  $x$
- $x.p = \text{NIL}$  and  $x.\text{rightSibling} = \text{NIL} \longrightarrow ?$
- If both  $x.\text{leftChild}$  and  $x.\text{rightSibling}$  are  $\text{NIL} \longrightarrow ?$



# Left-child and Right-Sibling Representation





# Implementation Details

```
struct node
```

```
{
```

```
    elemType Key;
```

```
    struct node *p;
```

```
    struct node *left;
```

```
    struct node *right;
```

```
};
```

```
struct binaryTree
```

```
{
```

```
    struct node *root;
```

```
};
```



# Reference

\* CLRS Book