

# Pseudocode of Linear Search

**LINEAR SEARCH(A, key)**    // Pseudocode of Linear Search

1. found = 0
2. for i = 1 to A.length
3.     if A[ i ] = key
4.       found = 1
5.       return i
6. if found = 0
7.    return 0

# Binary Search – Algorithm

```
BinarySearch (A, m, n, k)
    if (m>n) return -1;    //Base Case
    mid=(m+n)/2
    if A[mid]=k return mid;    //Base Case
    else if k < A[mid]
        BinarySearch(A, m, mid-1, k)
    else if k > A[mid]
        BinarySearch(A, mid+1, n, k)
```

# Insertion Sort

INSERTION SORT(A)

1. for  $j=2$  to  $A.length$
2.        $key = A[j]$ ;
3.       //Insert  $A[j]$  into the sorted sequence  $A[1...j-1]$
4.        $i = j-1$
5.       while  $i > 0$  and  $A[i] > key$
6.                $A[i+1] = A[i]$
7.                $i = i-1$
8.        $A[i+1] = key$

# Merge Sort - Recursive Algorithm

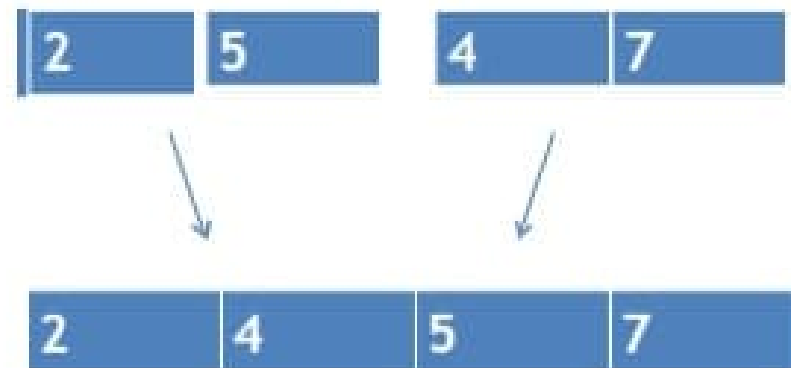
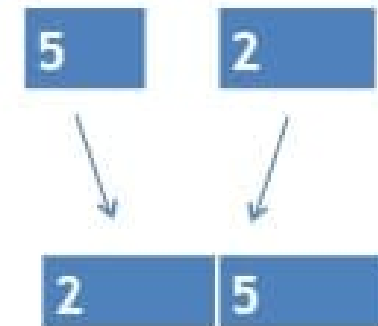
MERGE-SORT( $A, p, r$ )

```
1  if  $p < r$ 
2       $q = \lfloor (p + r) / 2 \rfloor$ 
3      MERGE-SORT( $A, p, q$ )
4      MERGE-SORT( $A, q + 1, r$ )
5      MERGE( $A, p, q, r$ )
```

# Merge function

MERGE( $A, p, q, r$ )

```
1   $n_1 = q - p + 1$ 
2   $n_2 = r - q$ 
3  let  $L[1..n_1 + 1]$  and  $R[1..n_2 + 1]$  be new arrays
4  for  $i = 1$  to  $n_1$ 
5       $L[i] = A[p + i - 1]$ 
6  for  $j = 1$  to  $n_2$ 
7       $R[j] = A[q + j]$ 
8   $L[n_1 + 1] = \infty$       Sentinel - a special value
9   $R[n_2 + 1] = \infty$ 
10  $i = 1$ 
11  $j = 1$ 
12 for  $k = p$  to  $r$ 
13     if  $L[i] \leq R[j]$ 
14          $A[k] = L[i]$ 
15          $i = i + 1$ 
16     else  $A[k] = R[j]$ 
17          $j = j + 1$ 
```



# Pseudocode of Quick Sort

QUICKSORT(A, p, r)

```
1  if  $p < r$ 
2      then  $q \leftarrow \text{PARTITION}(A, p, r)$ 
3           QUICKSORT(A, p,  $q-1$ )
4           QUICKSORT(A,  $q+1, r$ )
```

## Pseudo code : *PARTITION*

**PARTITION** (A, p, r)

1  $x = A[r]$

2  $i = p - 1$

3 **for**  $j = p$  **to**  $r - 1$

4     **do if**  $A[j] \leq x$

5         **then**  $i = i + 1$

6         Exchange  $A[i]$  with  $A[j]$

7     Exchange  $A[i + 1]$  with  $A[r]$

8 **return**  $i + 1$

# MAX-HEAPIFY

MAX-HEAPIFY( $A, i$ )

- 1  $l = \text{LEFT}(i)$
- 2  $r = \text{RIGHT}(i)$
- 3 **if**  $l \leq A.\text{heap-size}$  and  $A[l] > A[i]$
- 4      $\text{largest} = l$
- 5 **else**  $\text{largest} = i$
- 6 **if**  $r \leq A.\text{heap-size}$  and  $A[r] > A[\text{largest}]$
- 7      $\text{largest} = r$
- 8 **if**  $\text{largest} \neq i$
- 9     exchange  $A[i]$  with  $A[\text{largest}]$
- 10    MAX-HEAPIFY( $A, \text{largest}$ )



# BUILD-MAX-HEAP

## Pseudocode for BUILD-MAX-HEAP

BUILD-MAX-HEAP( $A$ )

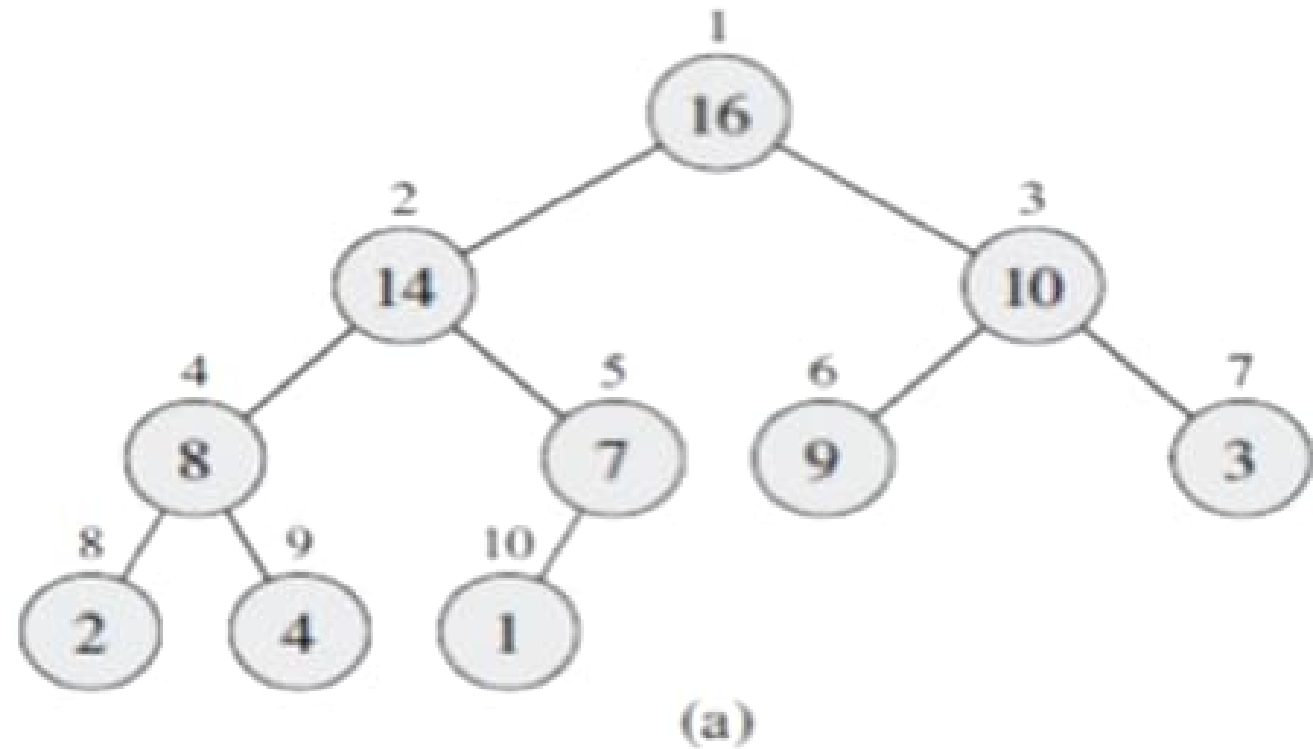
- 1  $A.heap-size = A.length$
- 2 **for**  $i = \lfloor A.length/2 \rfloor$  **downto** 1
- 3     MAX-HEAPIFY( $A, i$ )

# Heap Sort: ALGORITHM

HEAPSORT( $A$ )

- 1 BUILD-MAX-HEAP( $A$ )
- 2 **for**  $i = A.length$  **downto** 2
- 3     exchange  $A[1]$  with  $A[i]$
- 4      $A.heap-size = A.heap-size - 1$
- 5     MAX-HEAPIFY( $A, 1$ )

# MAXIMUM(A) ?



HEAP-MAXIMUM(A)  
return A[1]

# Extract-Max operations

HEAP-EXTRACT-MAX(A)

1 if  $A.heap-size < 1$

2 error “heap underflow”

3  $max = A[1]$

4  $A[1] = A[A.heap-size]$

5  $A.heap-size = A.heap-size - 1$

6 MAX-HEAPIFY(A, 1)

7 return  $max$

## HEAP-INCREASE-KEY

HEAP-INCREASE-KEY ( $A, i, key$ )

- 1   if  $key < A[i]$
- 2       error “new key is smaller than current key”
- 3    $A[i] = key$
- 4   while  $i > 1$  and  $A[\text{PARENT}(i)] < A[i]$
- 5       exchange  $A[i]$  with  $A[\text{PARENT}(i)]$
- 6        $i = \text{PARENT}(i)$

# MAX-HEAP-INSERT

MAX-HEAP-INSERT( $A, key$ )

- 1  $A.heap-size = A.heap-size + 1$
- 2  $A[A.heap-size] = -\infty$
- 3 HEAP-INCREASE-KEY( $A, A.heap-size, key$ )