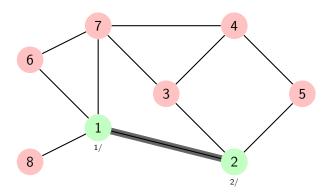
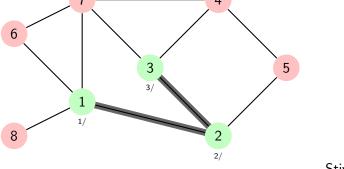


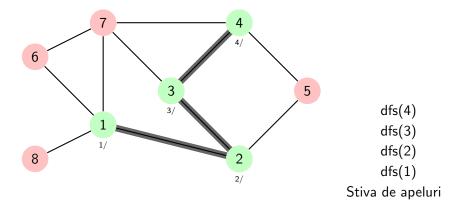
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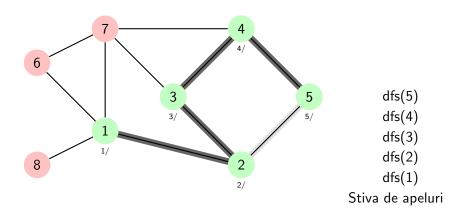


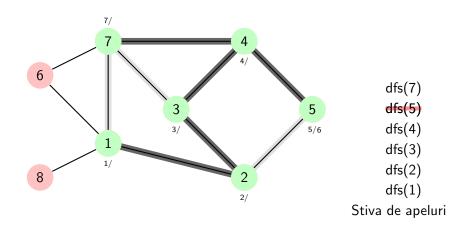
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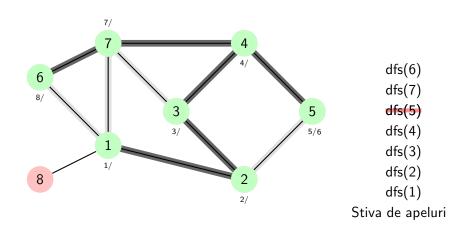


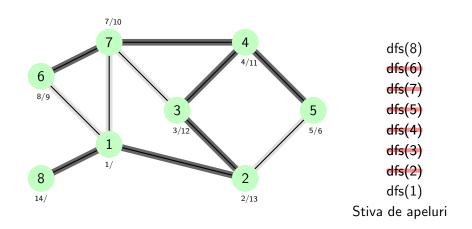
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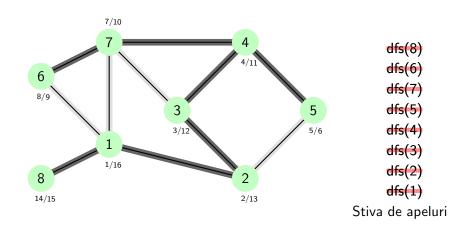


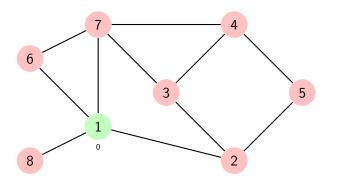




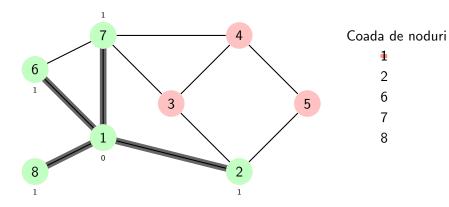


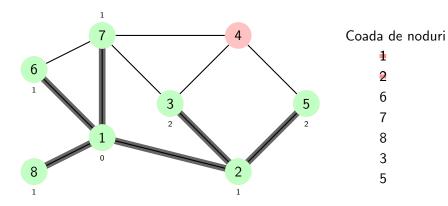


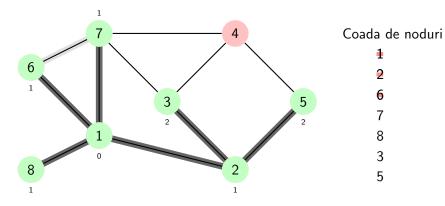


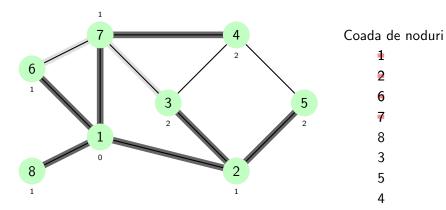


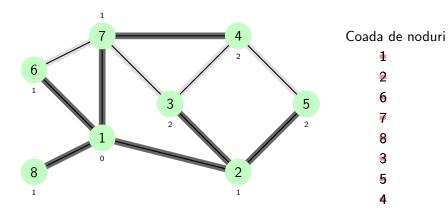
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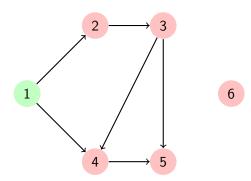




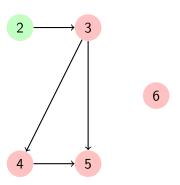








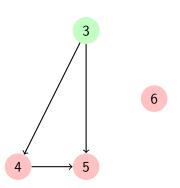
Sortare topologică



Sortare topologică

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Sortare topologică

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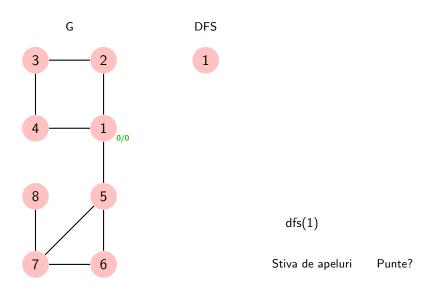
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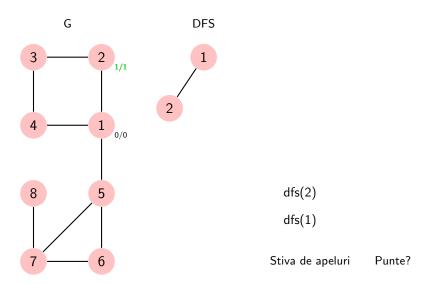
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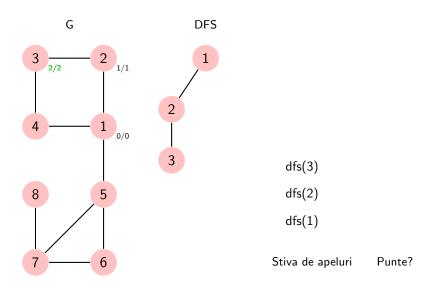


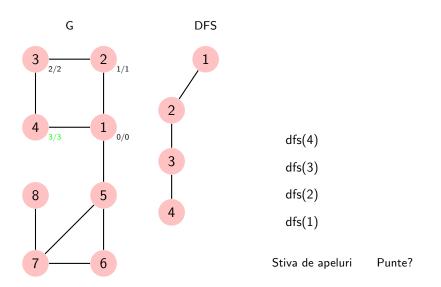
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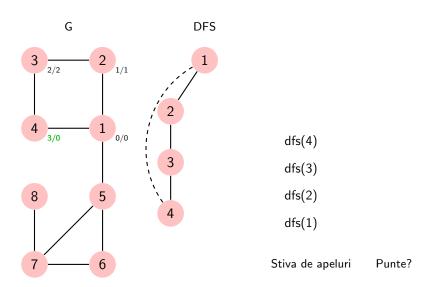
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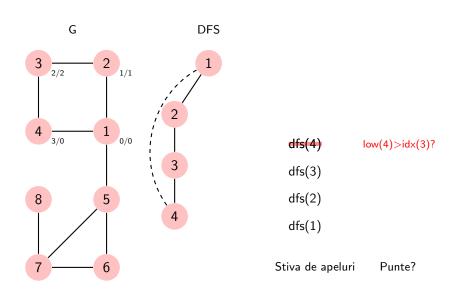


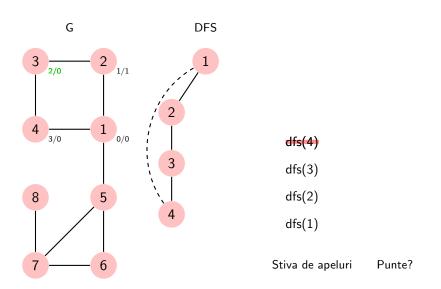


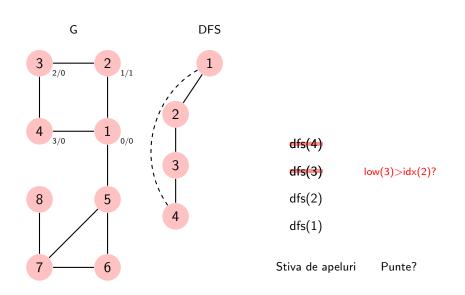


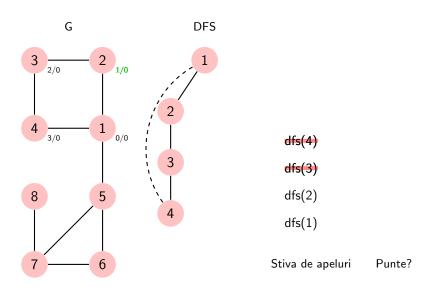


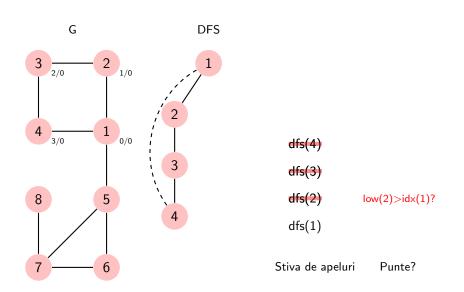


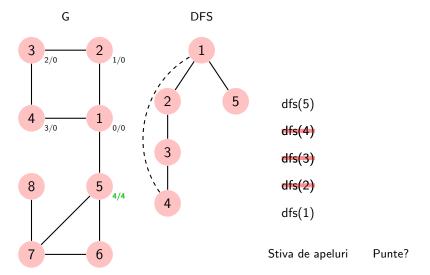


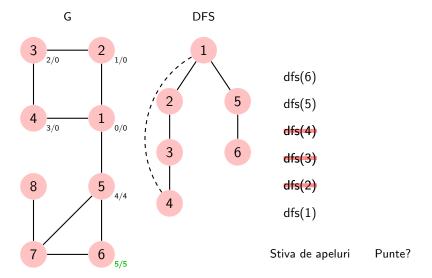


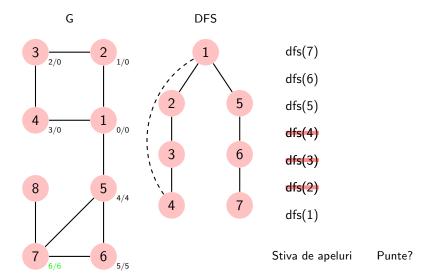


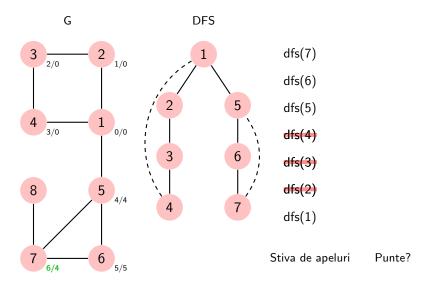


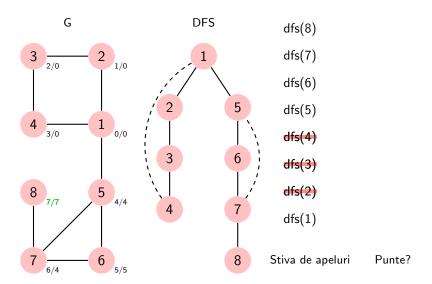


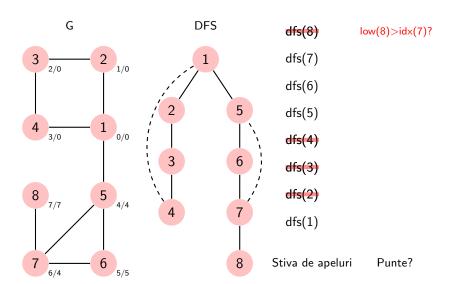


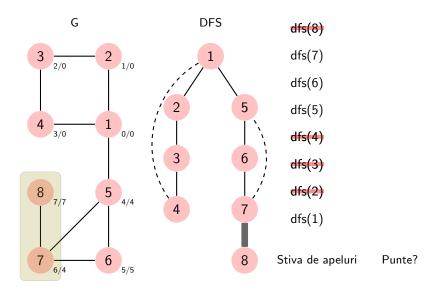


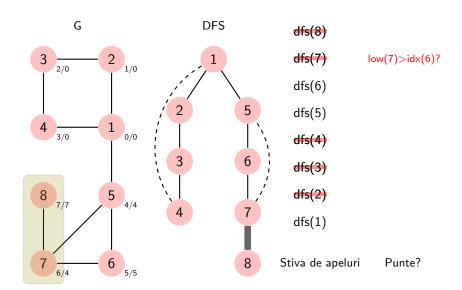


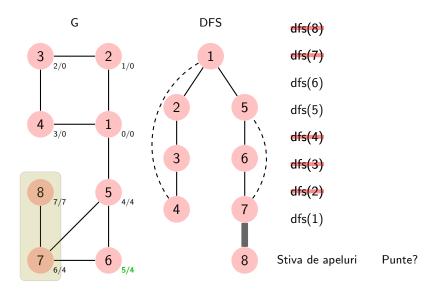


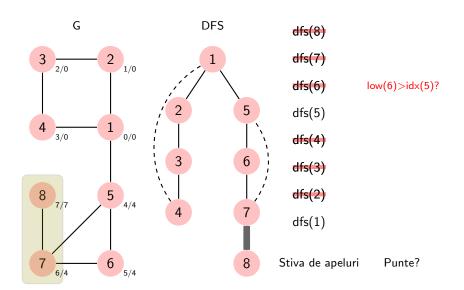


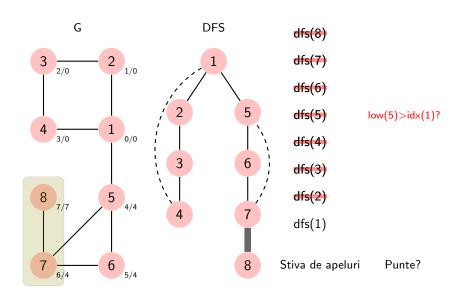


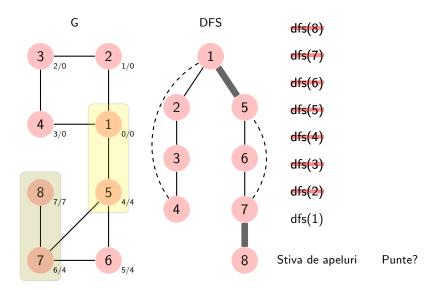


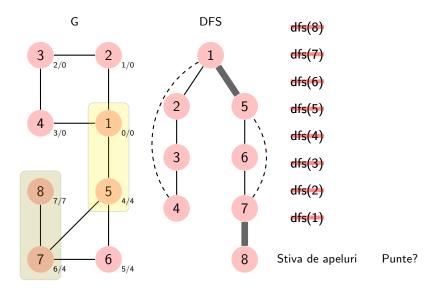


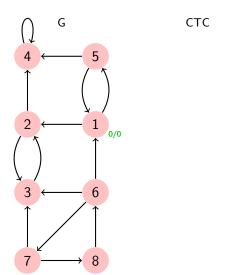




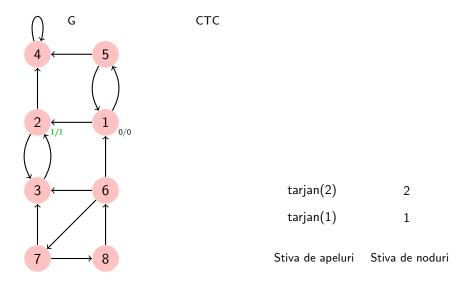


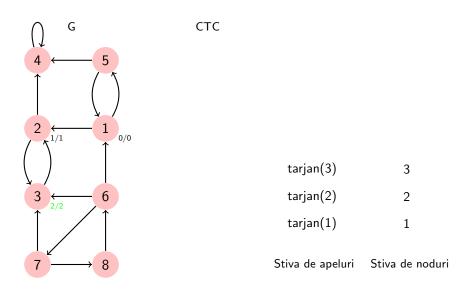


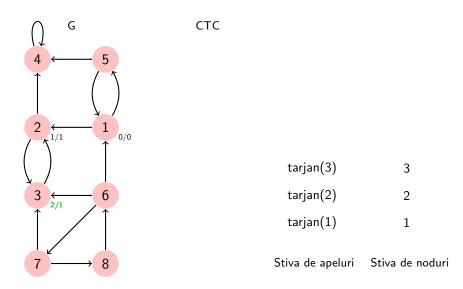


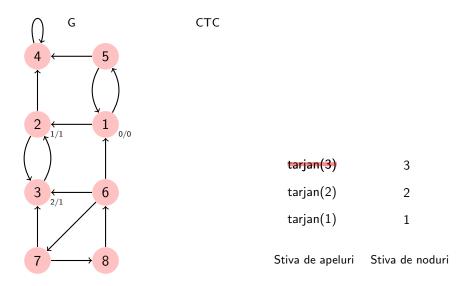


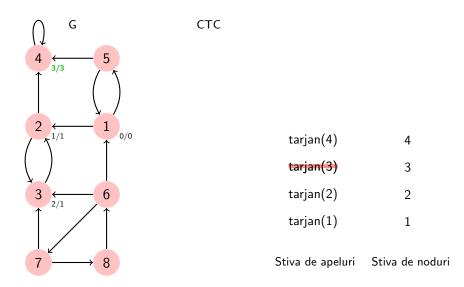
tarjan(1) 1
Stiva de apeluri Stiva de noduri

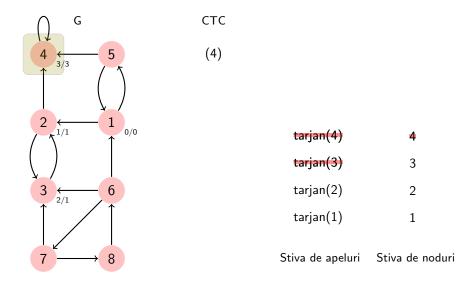


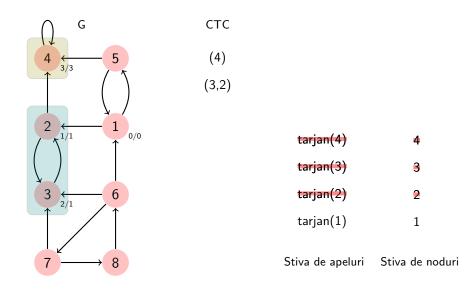


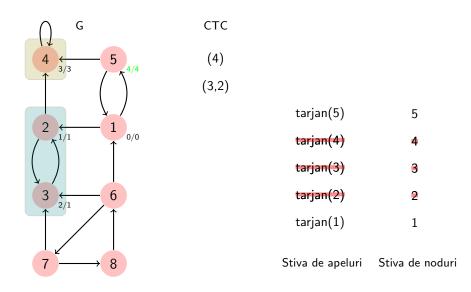


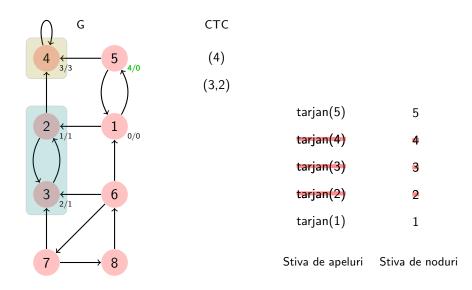


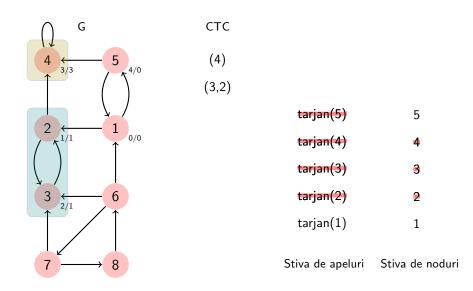


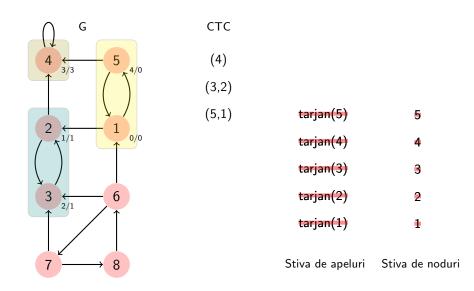


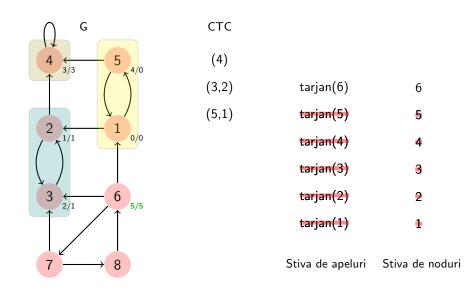


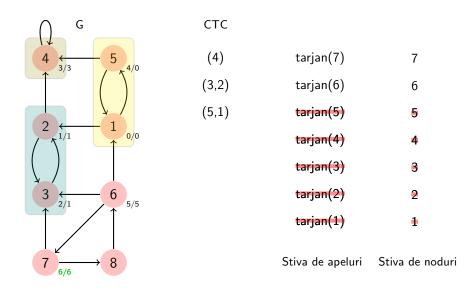


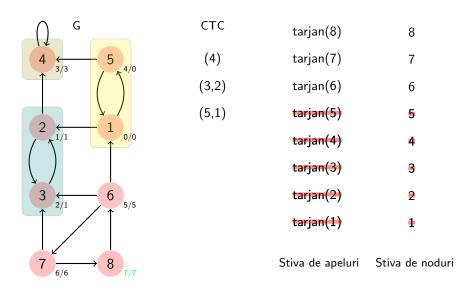


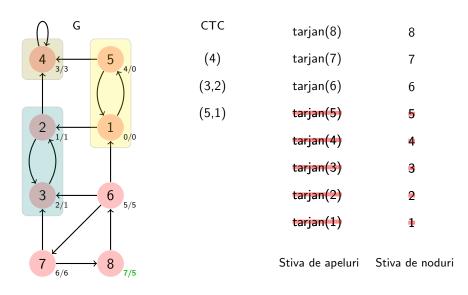


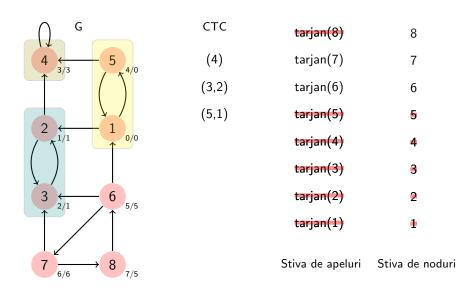


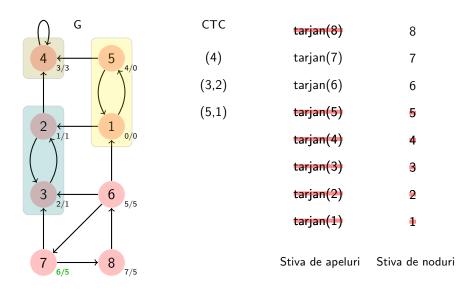


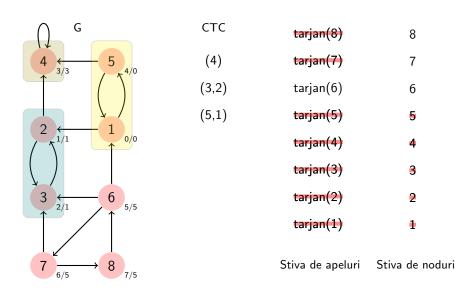


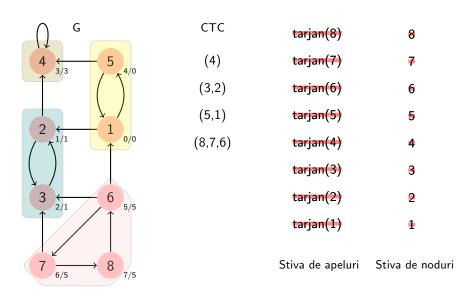


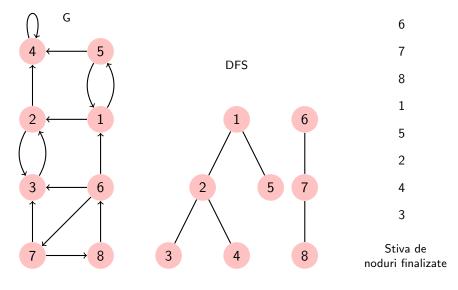


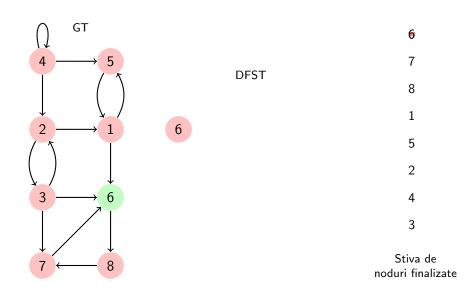


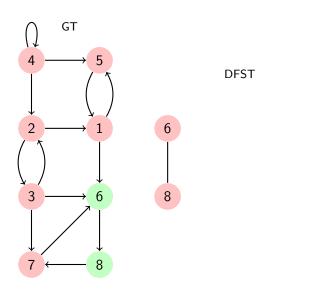






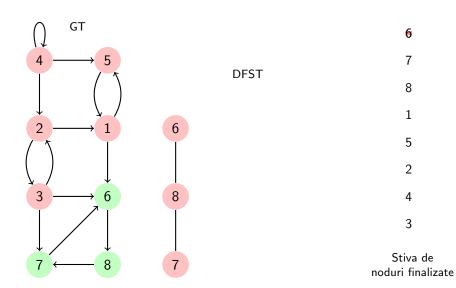


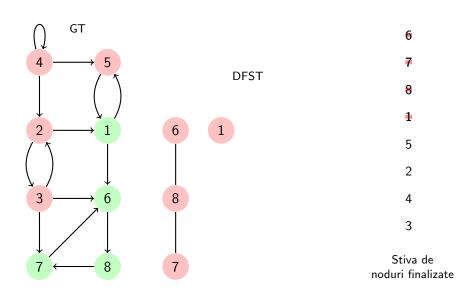


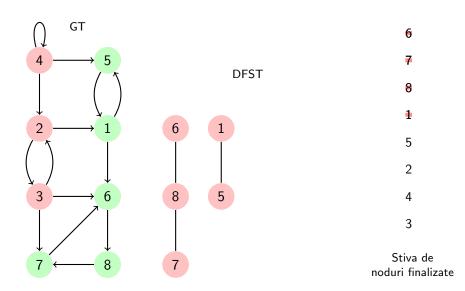


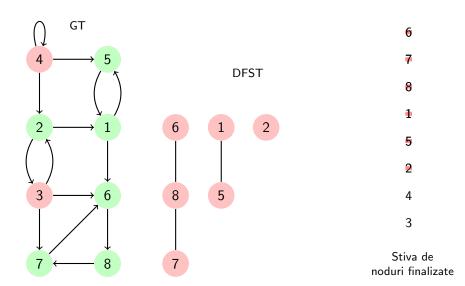
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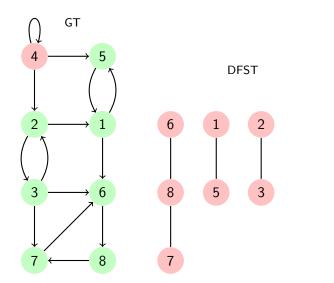
Stiva de noduri finalizate





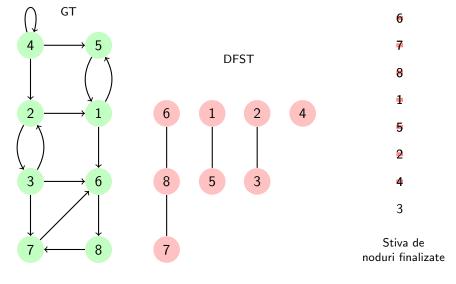




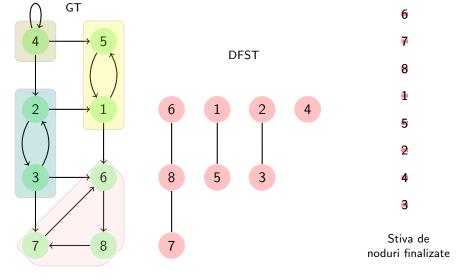


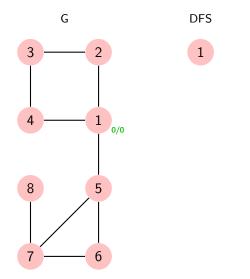
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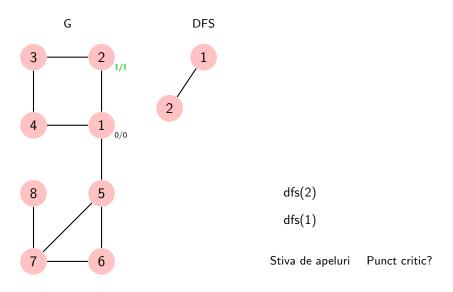


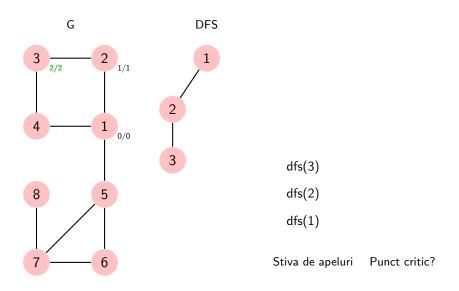
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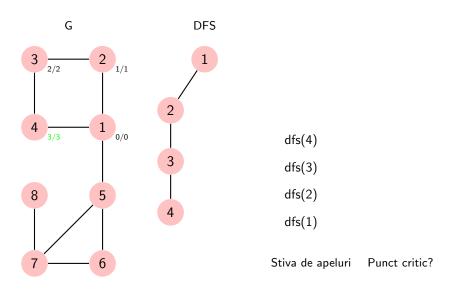


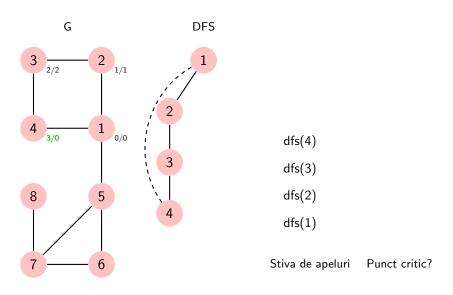


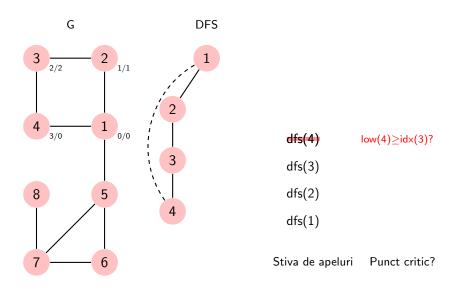
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Stiva de apeluri Punct critic?

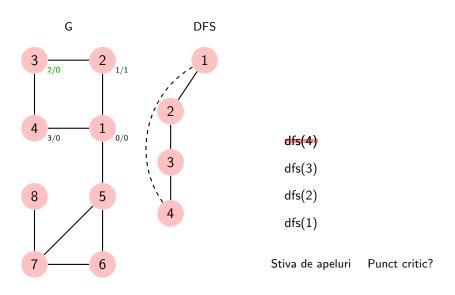


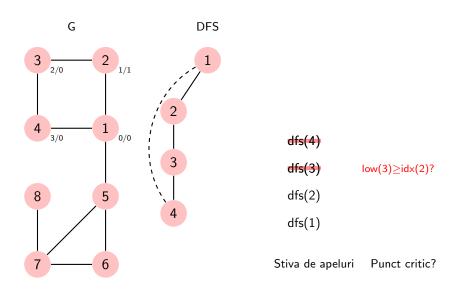


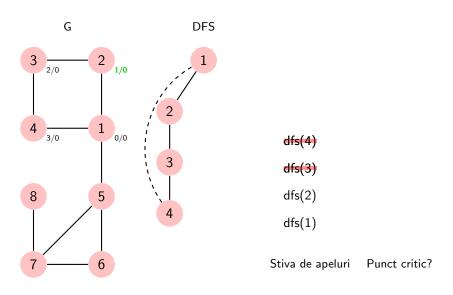


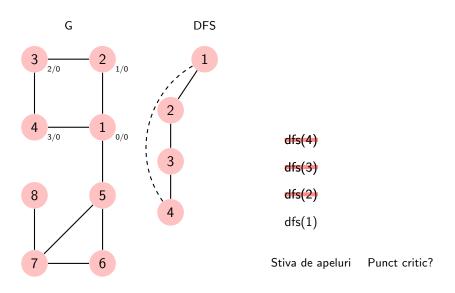


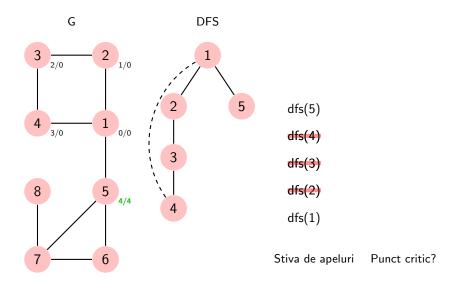


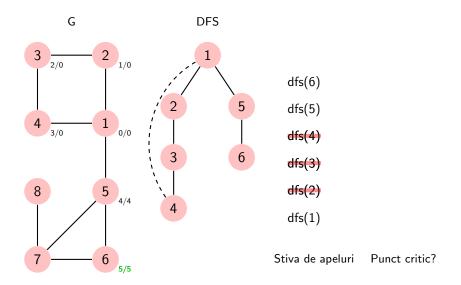


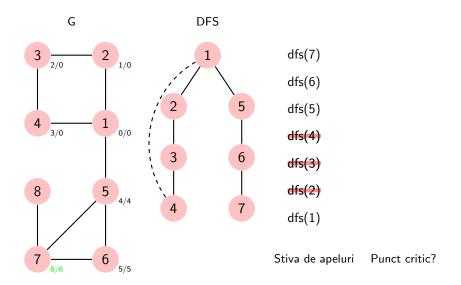


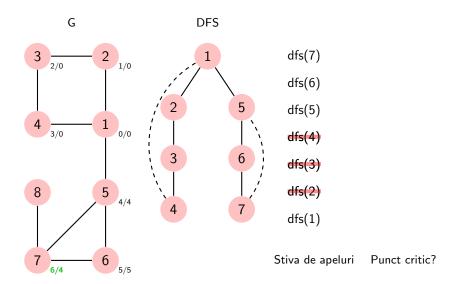


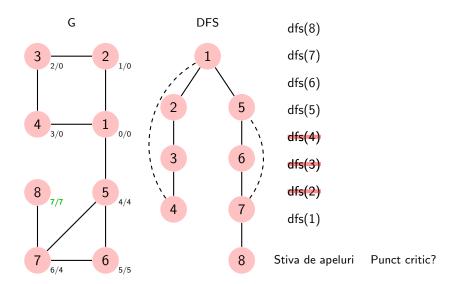


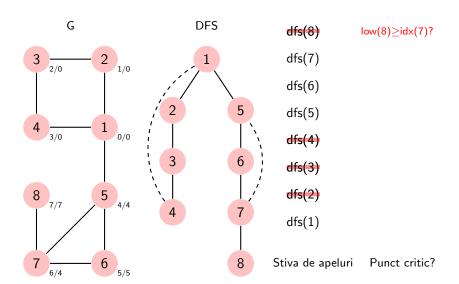


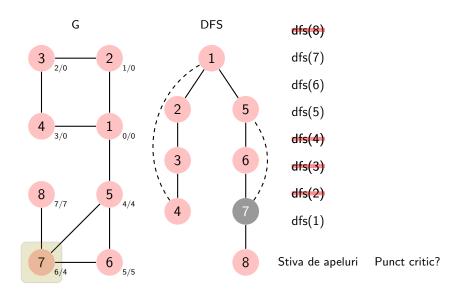


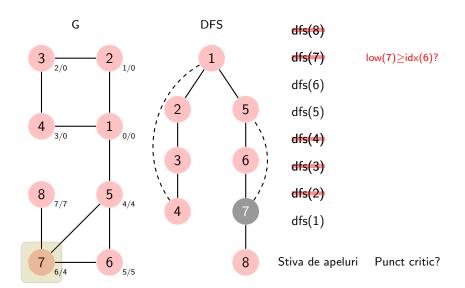


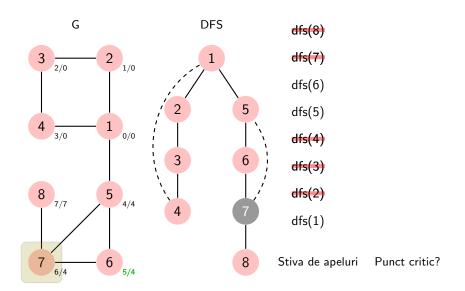


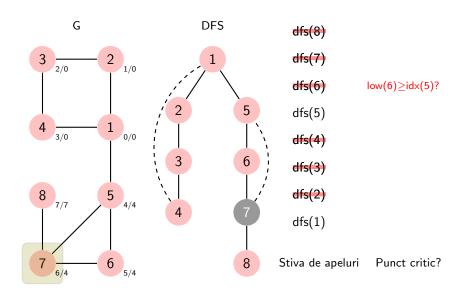


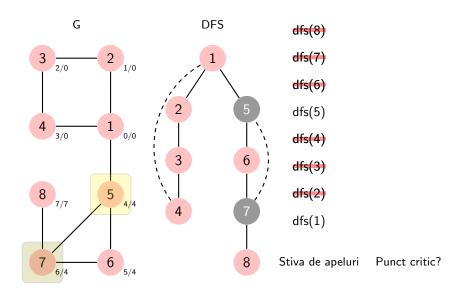


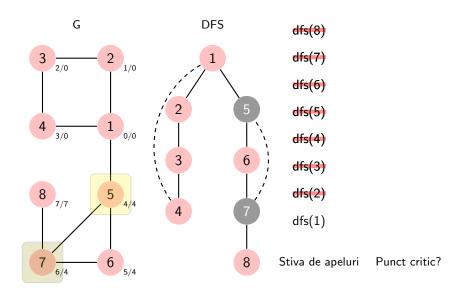


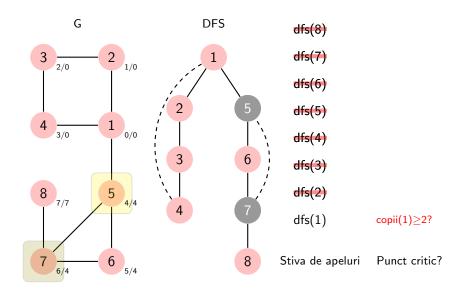


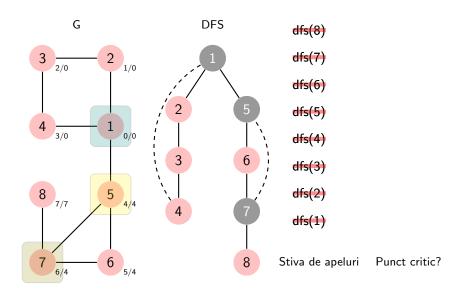


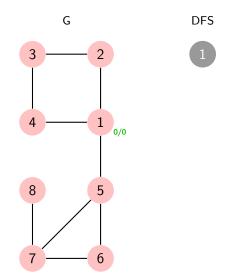




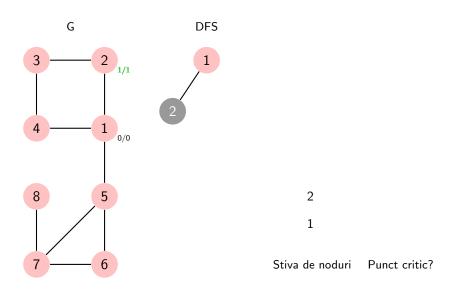


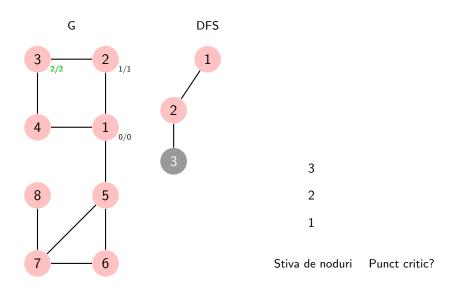


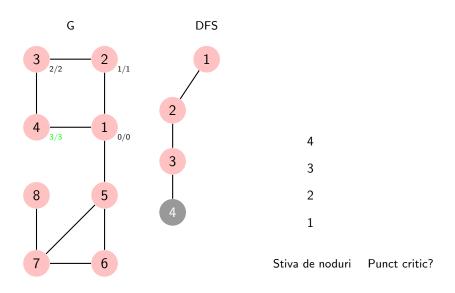


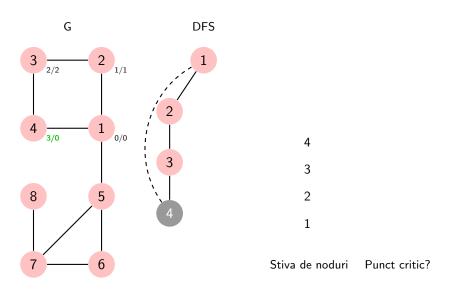


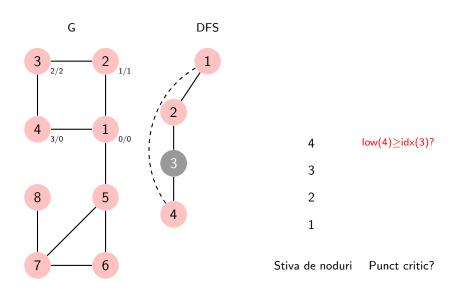
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Stiva de noduri Punct critic?

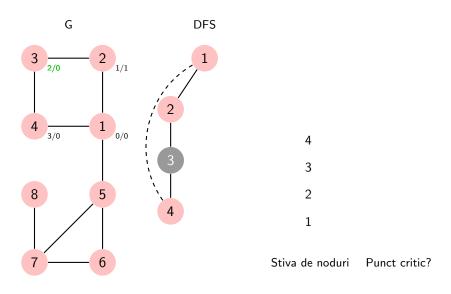


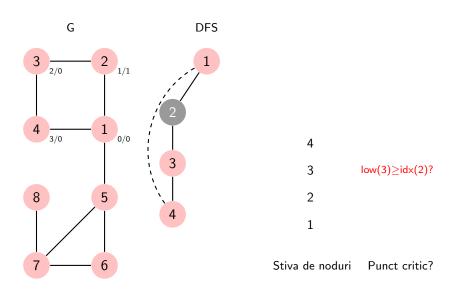


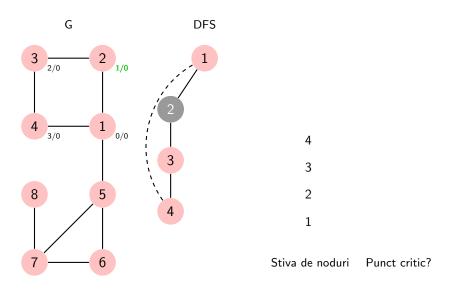


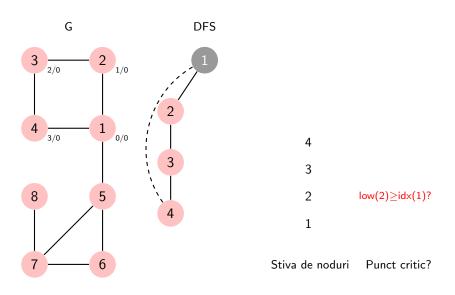


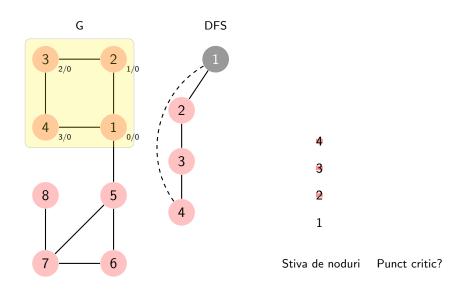


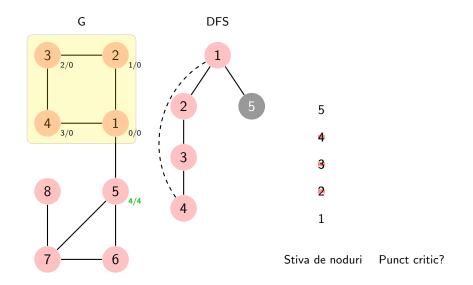


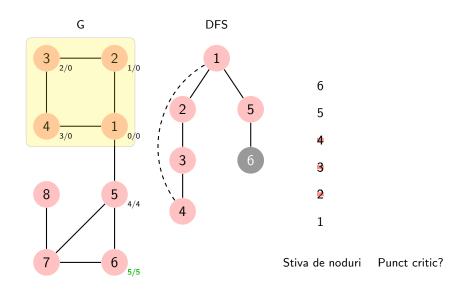


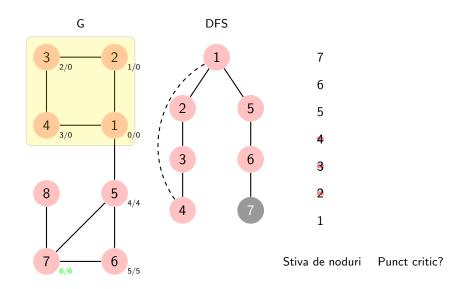


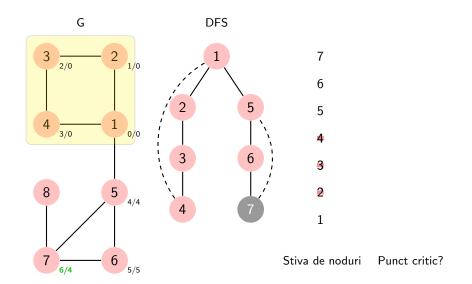


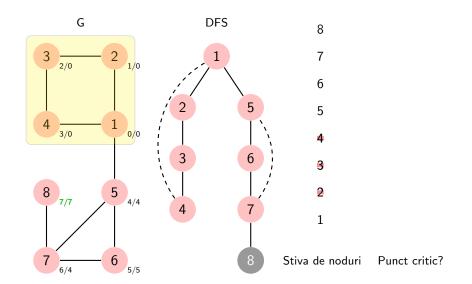


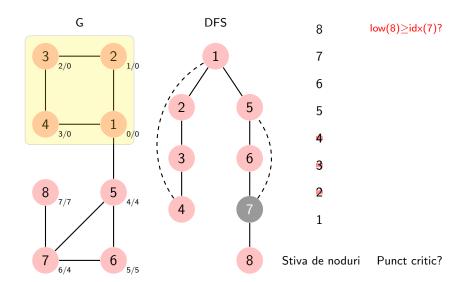


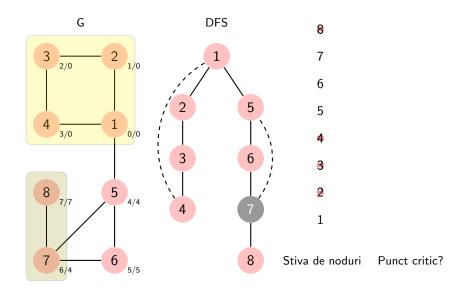


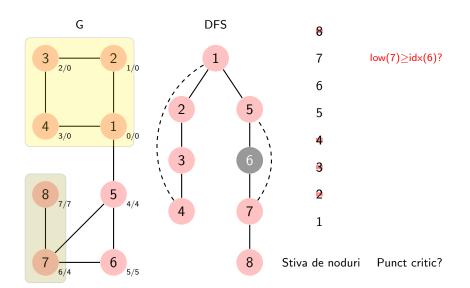


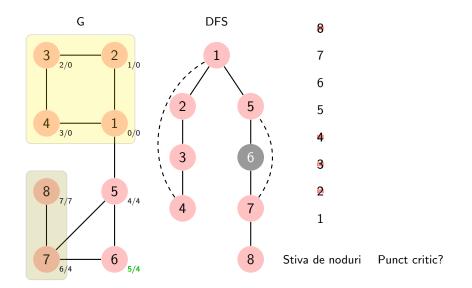


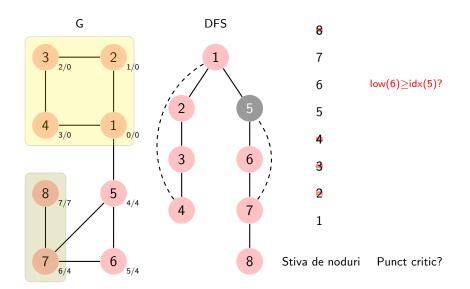


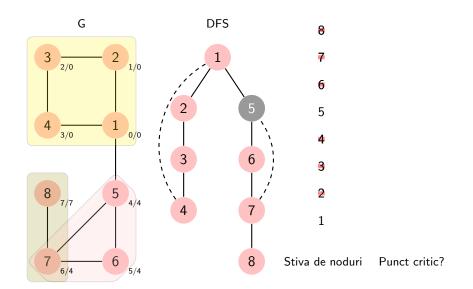


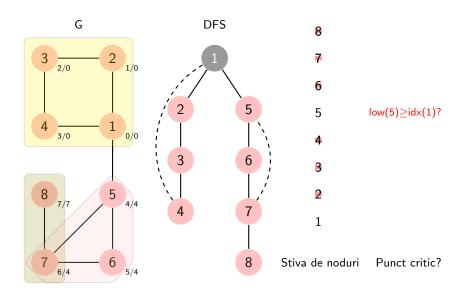


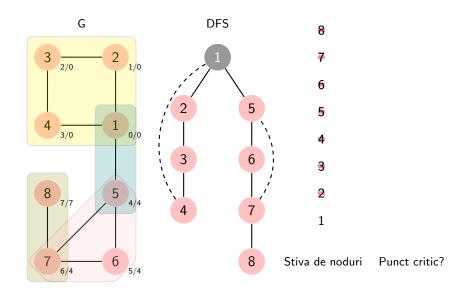


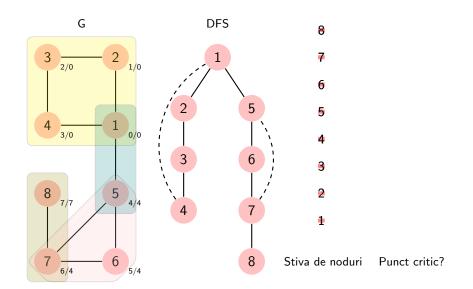










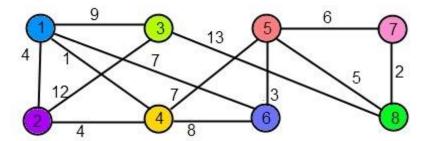


<u>Arbori minimi de acoperire – Kruskal</u>

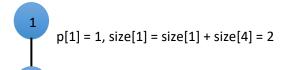
```
Kruskal(V, E)
    E = sort(E) după cost
    AMA = {}
    for (e : E)
        if AMA U {e} n-are cicluri
        AMA = AMA U {e}
```

Pentru a determina că AMA nu are cicluri, folosim o structură de tip union-find:

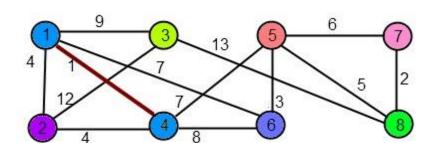
- Inițial fiecare nod face parte dintr-un arbore distinct cu rădăcina (root) în nodul respectiv
- Când adăugăm o muchie (u,v) la AMA
 - o ea trebuie să unească arbori distincți, adică: findroot(u) ≠ findroot(v)
 - o cei 2 arbori trebuie uniți (union), adică de acum încolo, pentru orice nod x din arbore cu u și orice nod y din arbore cu v: findroot(x) = findroot(y)
- Pentru a optimiza operațiile de union/find
 - vom folosi un <u>vector de părinți</u> pentru fiecare nod (doar root-ul unui arbore va fi propriul său părinte), restul nodurilor trebuie să circule din părinte în părinte spre root
 - <u>la find</u>: nodul care circulă din părinte în părinte spre root va actualiza, pe parcursul acestei operații, părintele tuturor nodurilor din cale la root, astfel încât data viitoare să indice direct la acesta: p[nod] = findroot(p[nod])
 - <u>la union</u>: pentru a micșora lungimea căilor spre root, vom adăuga mereu arborele mai mic ca subarbore al arborelui mai mare (deci e nevoie de un <u>vector size</u>): <u>p[root_mic]</u> = <u>root_mare</u>



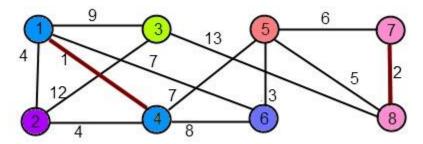
 $E = \{(1,4), (7,8), (5,6), (1,2), (2,4), (5,8), (5,7), (1,6), (4,5), (4,6), (1,3), (2,3), (3,8)\}$



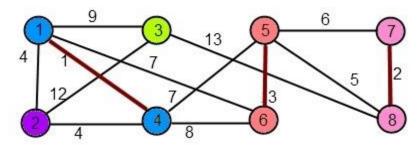
p[4] = 1



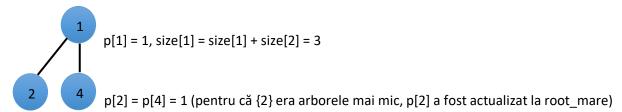
 $E = \{(1,4), (7,8), (5,6), (1,2), (2,4), (5,8), (5,7), (1,6), (4,5), (4,6), (1,3), (2,3), (3,8)\}$

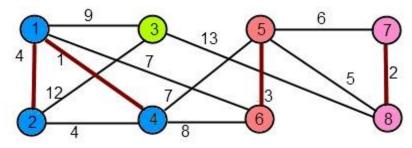


 $E = \{(1,4), (7,8), (5,6), (1,2), (2,4), (5,8), (5,7), (1,6), (4,5), (4,6), (1,3), (2,3), (3,8)\}$



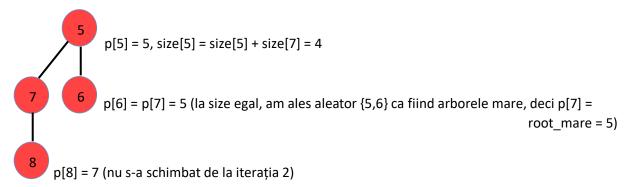
 $E = \{(1,4), (7,8), (5,6), (1,2), (2,4), (5,8), (5,7), (1,6), (4,5), (4,6), (1,3), (2,3), (3,8)\}$

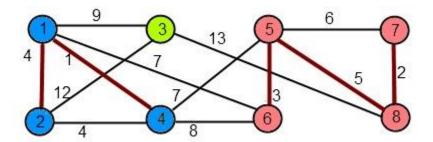




 $E = \{(1,4), (7,8), (5,6), (1,2), (2,4), (5,8), (5,7), (1,6), (4,5), (4,6), (1,3), (2,3), (3,8)\}$

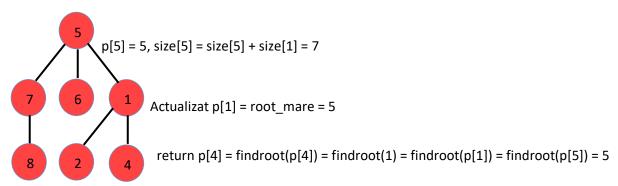
findroot(2) = findroot(4) = 1 (deci ignorăm muchia (2,4))

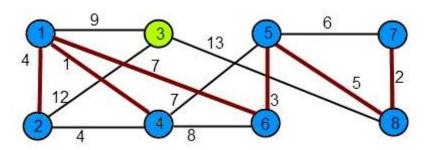




 $E = \{(1,4), (7,8), (5,6), (1,2), (2,4), (5,8), (5,7), (1,6), (4,5), (4,6), (1,3), (2,3), (3,8)\}$

findroot(5) = findroot(7) = 5 (deci ignorăm muchia (5,7))

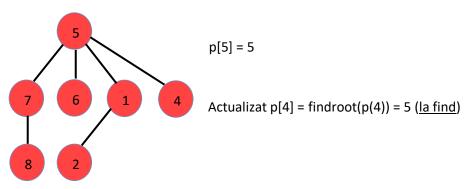




 $E = \{(1,4), (7,8), (5,6), (1,2), (2,4), (5,8), (5,7), (1,6), (4,5), (4,6), (1,3), (2,3), (3,8)\}$

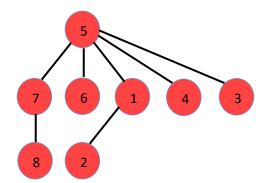
findroot(4) = findroot(5) = 5 (deci ignorăm muchia (5,7))

!!! În plus, in timpul findroot(4): p(4) = findroot(p(4)) = 5



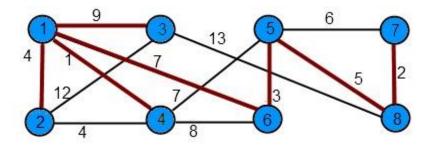
findroot(4) = findroot(6) = 5 (deci ignorăm muchia (4,6))

!!! Am profitat de calea scurtată de la 4 la 5



p[5] = 5, size[5] = size[5] + size[3]

Actualizat p[3] = root_mare = 5



 $E = \{(1,4), (7,8), (5,6), (1,2), (2,4), (5,8), (5,7), (1,6), (4,5), (4,6), (1,3), (2,3), (3,8)\}$

Putem să continuăm să respingem muchii sau, dacă am reținut numărul de noduri/muchii deja adăugate, să ne oprim aici pentru că știm că avem deja un AAM.

costAMA = suma costurilor muchiilor verzi = 1 + 2 + 3 + 4 + 5 + 7 + 9 = 31

<u>Arbori minimi de acoperire – Prim</u>

Prim(V, E)

alege un nod sursă root

pq = coadă de priorități inițializată cu vecinii lui root (distanțele minime au prioritate) while (pq nevidă)

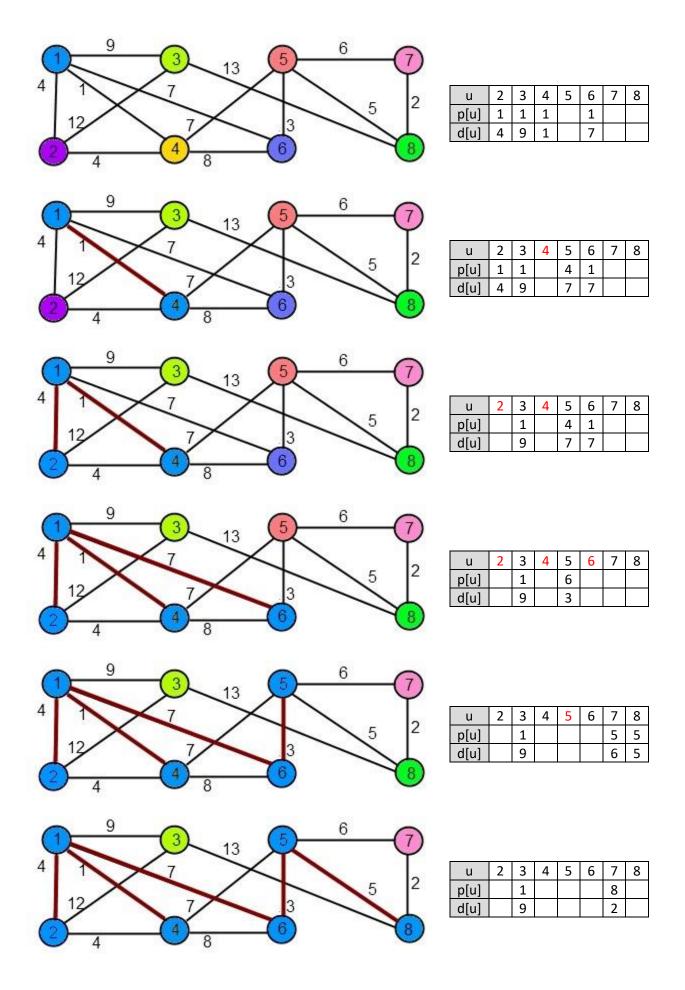
u = top(pq) // cel mai apropiat nod de AMA if $u \notin AMA$

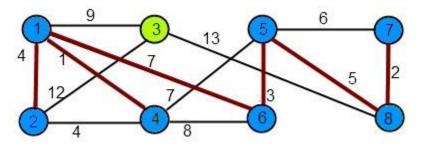
AMA = AMA U { (p(u),u) } // p(u) = cel mai apropiat vecin al lui u în AMA for (v : vecini(u))

actualizează în pq distanța lor față de AMA

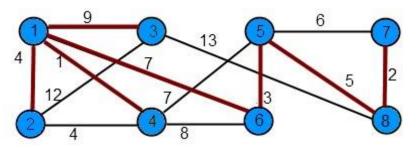
Pentru a construi AMA, elementele din pq vor conține toată informația utilă:

- nodul u de adăugat în AMA
- părintele p[u] (cel mai apropiat vecin al lui u care este deja în AMA), necesar dacă
 - o vrem să știm ce muchie adăugăm în AMA
- distanța d[u] = w(p[u],u) (costul muchiei cu care vrem să legăm u la AMA), necesar pentru
 - o a ordona elementele în pq
 - o a calcula rezultatul final costAMA

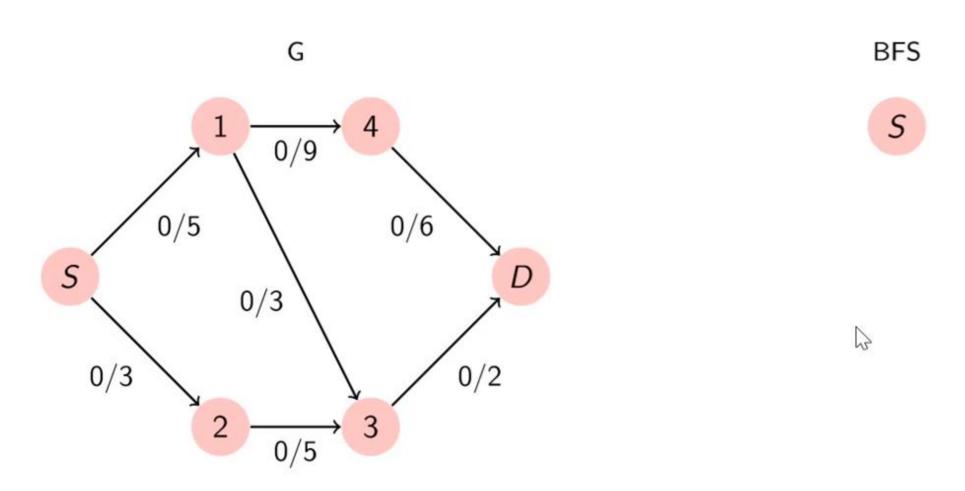


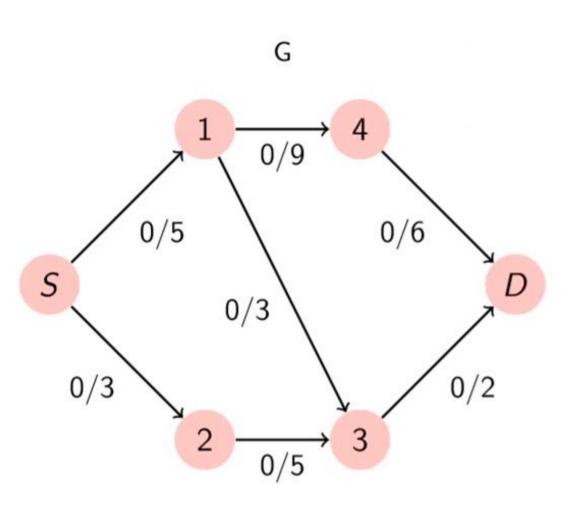


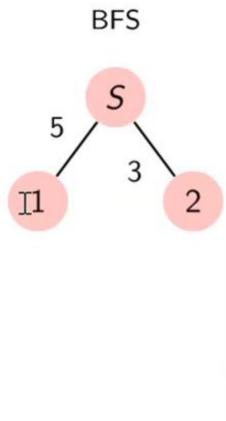
u	2	3	4	5	6	7	8
p[u]		1					
d[u]		9					

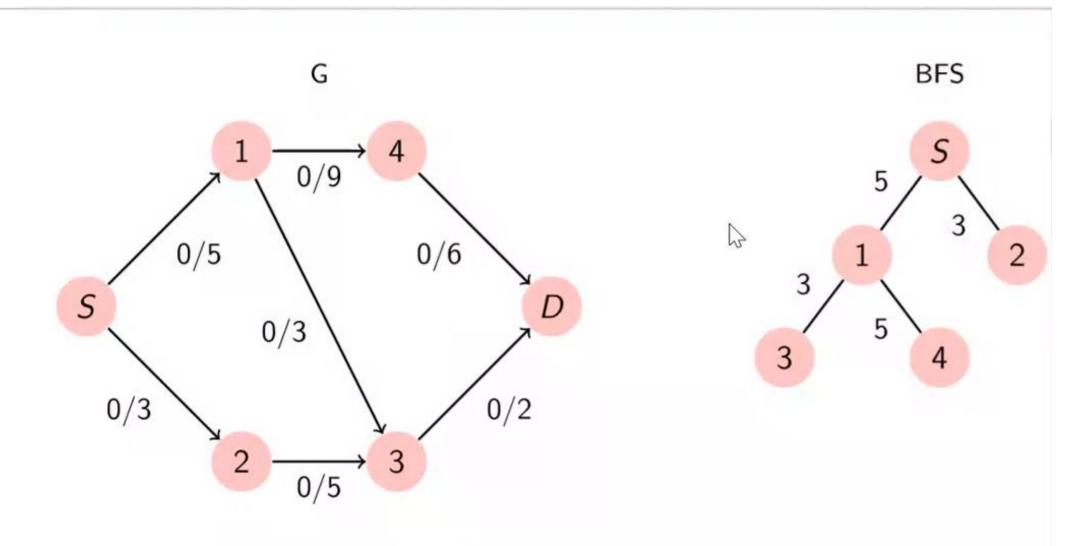


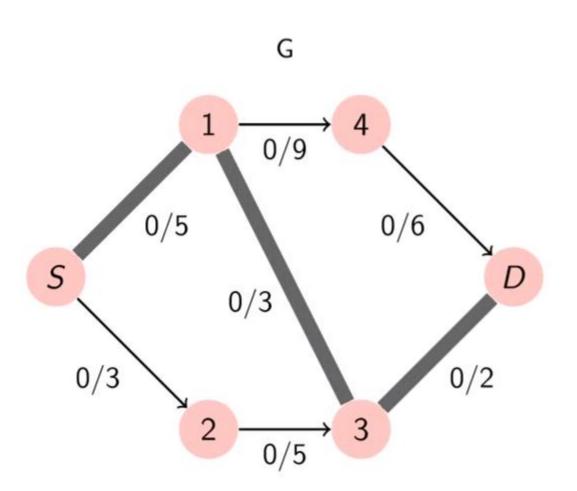
u	2	3	4	5	6	7	8
p[u]							
d[u]							

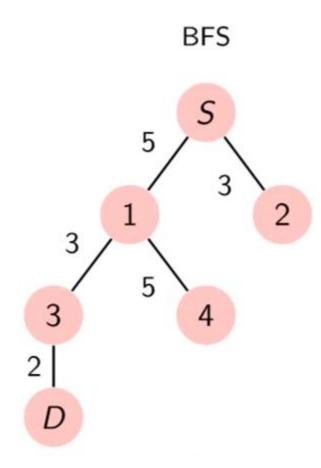


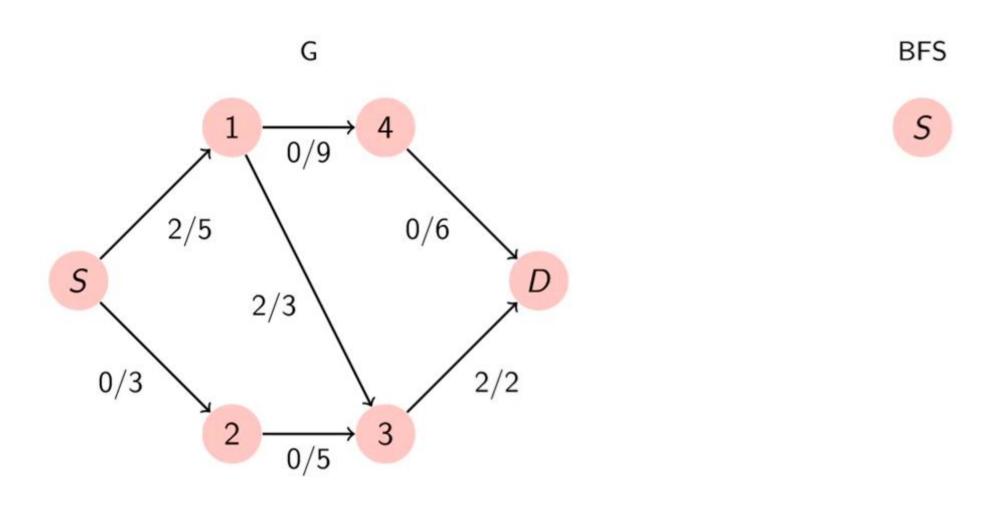


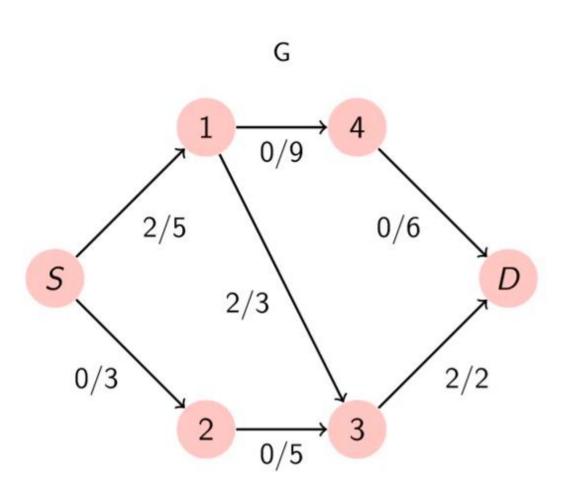


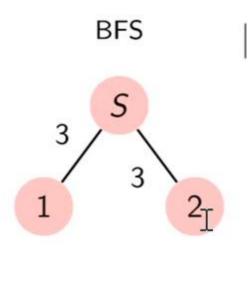


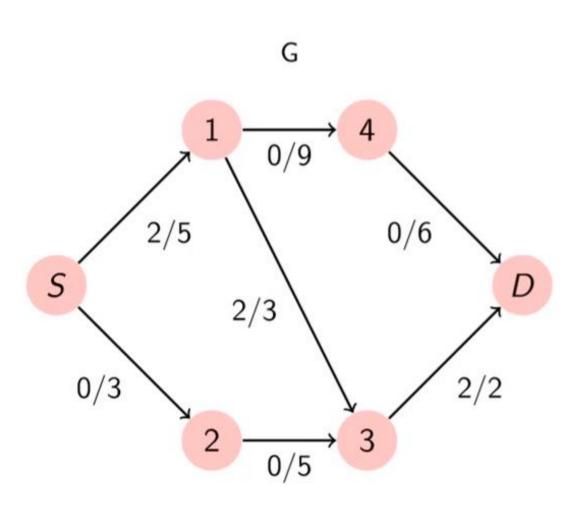


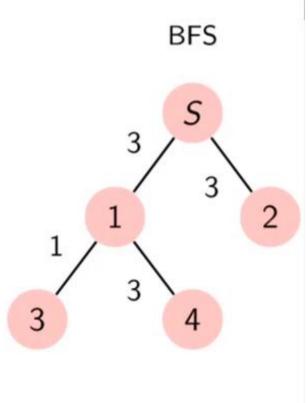


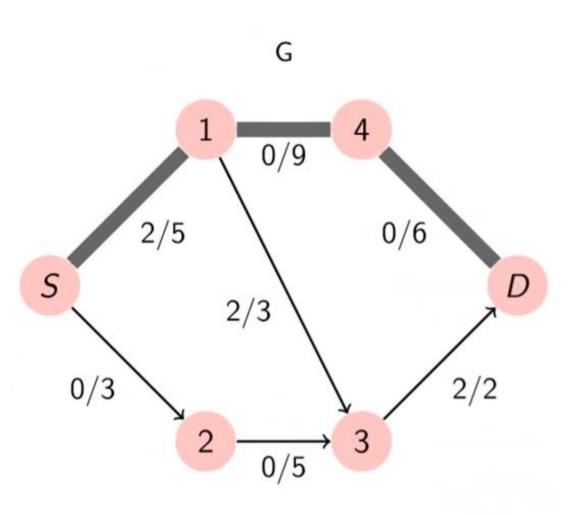


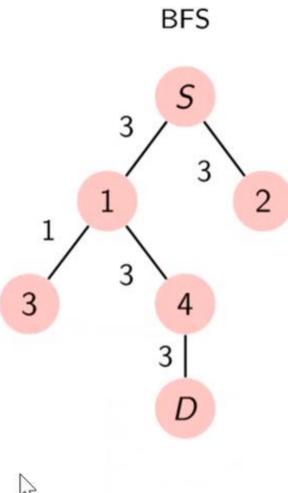


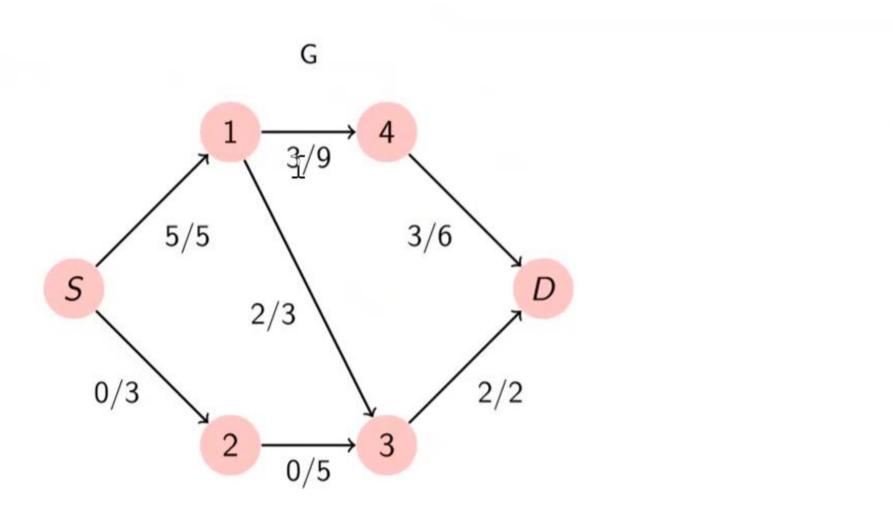












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