

# Community Detection and Evaluation

## Chapter 3

# Community

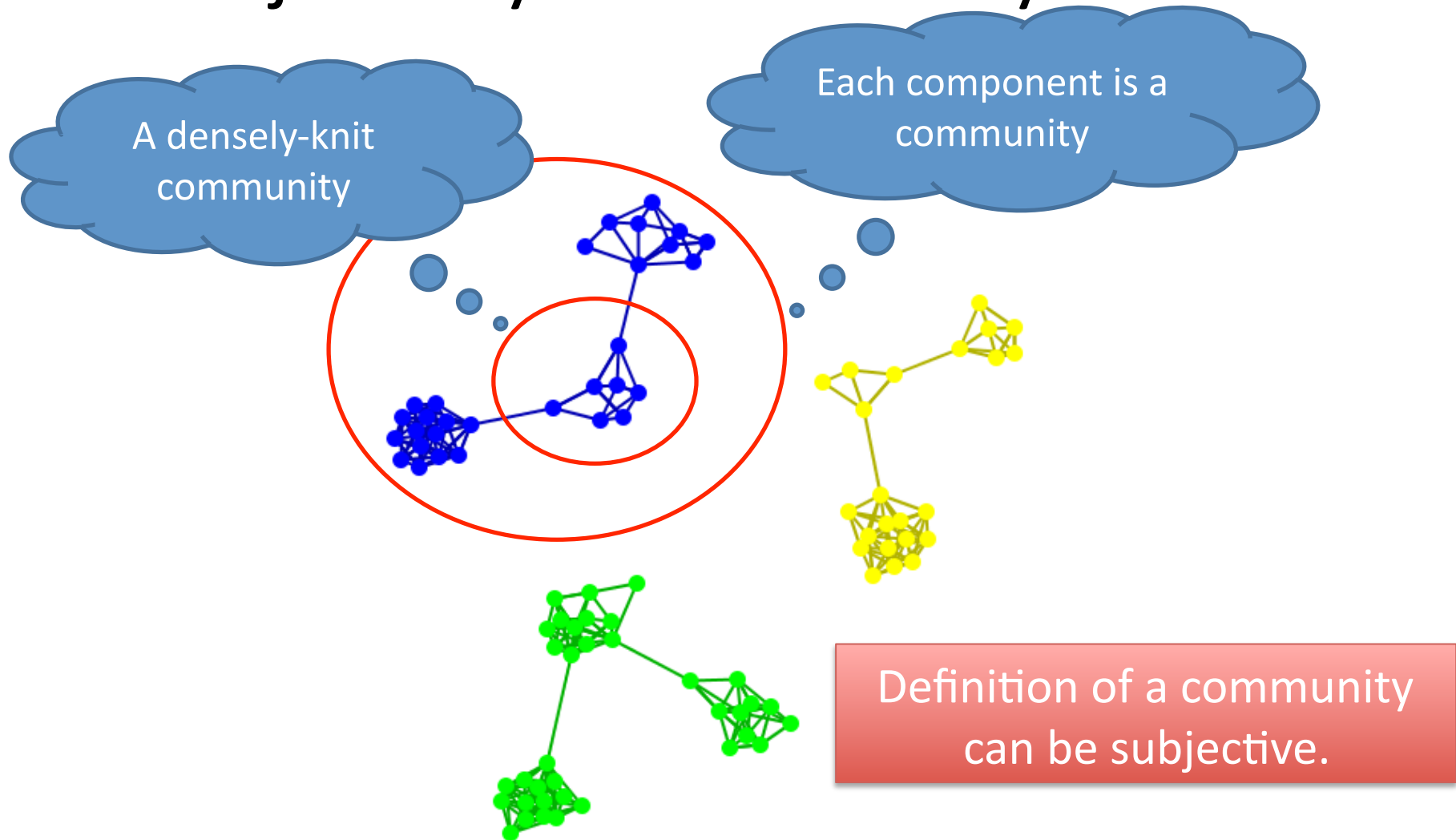
- **Community**: It is formed by individuals such that those within a group interact with each other more frequently than with those outside the group
  - a.k.a. **group**, **cluster**, **cohesive subgroup**, **module** in different contexts
- **Community detection**: discovering groups in a network where individuals' group memberships are not explicitly given
- **Why communities in social media?**
  - Human beings are social
  - Easy-to-use social media allows people to extend their social life in unprecedented ways
  - Difficult to meet friends in the physical world, but much easier to find friend online with similar interests
  - Interactions between nodes can help determine communities

# Communities in Social Media

- Two types of groups in social media
  - **Explicit Groups**: formed by user subscriptions
  - **Implicit Groups**: implicitly formed by social interactions
- Some social media sites allow people to join groups, is it necessary to extract groups based on network topology?
  - Not all sites provide community platform
  - Not all people want to make effort to join groups
  - Groups can change dynamically
- Network interaction provides rich information about the relationship between users
  - Can complement other kinds of information
  - Help network visualization and navigation
  - Provide basic information for other tasks

# **COMMUNITY DETECTION**

# Subjectivity of Community Definition



# Taxonomy of Community Criteria

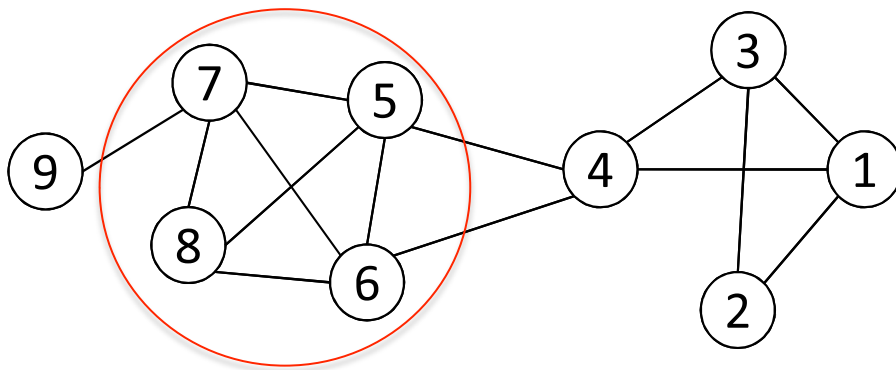
- Criteria vary depending on the tasks
- Roughly, community detection methods can be divided into 4 categories (not exclusive):
- **Node**-Centric Community
  - Each node in a group satisfies certain properties
- **Group**-Centric Community
  - Consider the connections within a group as a whole. The group has to satisfy certain properties without zooming into node-level
- **Network**-Centric Community
  - Partition the whole network into several disjoint sets
- **Hierarchy**-Centric Community
  - Construct a hierarchical structure of communities

# Node-Centric Community Detection

- Nodes satisfy different properties
  - Complete Mutuality
    - cliques
  - Reachability of members
    - k-clique, k-clan, k-club
  - Nodal degrees
    - k-plex, k-core
  - Relative frequency of Within-Outside Ties
    - LS sets, Lambda sets
- Commonly used in traditional social network analysis
- Here, we discuss some representative ones

# Complete Mutuality: Cliques

- **Clique**: a maximum complete subgraph in which all nodes are adjacent to each other



Nodes 5, 6, 7 and 8 form a clique

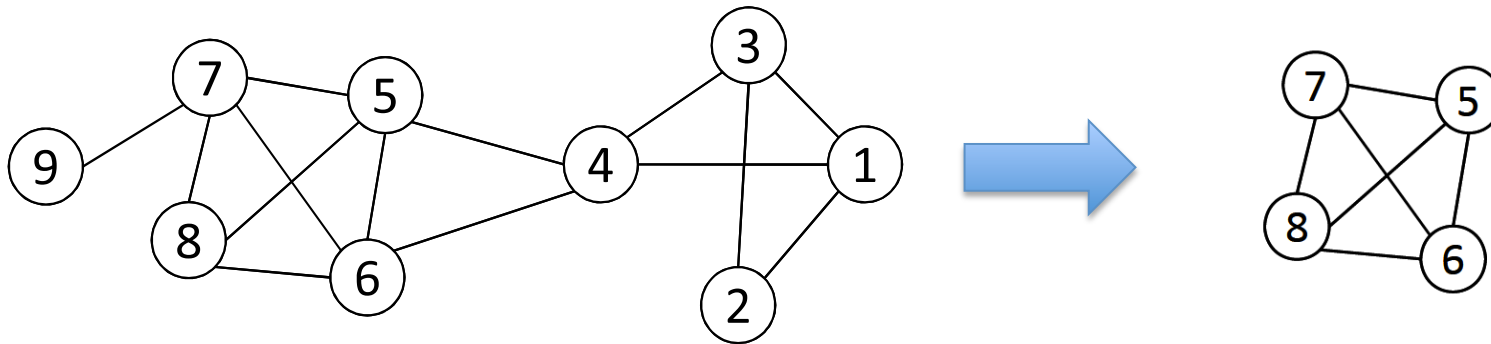
- NP-hard to find the maximum clique in a network
- Straightforward implementation to find cliques is very expensive in time complexity



# Finding the Maximum Clique

- In a clique of size  $k$ , each node maintains degree  $\geq k-1$
- Nodes with degree  $< k-1$  will not be included in the maximum clique
- Recursively apply the following **pruning** procedure
  - Sample a sub-network from the given network, and find a clique in the sub-network, say, by a greedy approach
  - Suppose the clique above is size  $k$ , in order to find out a *larger* clique, all nodes with degree  $\leq k-1$  should be removed.
- Repeat until the network is small enough
- Many nodes will be pruned as social media networks follow a power law distribution for node degrees

# Maximum Clique Example

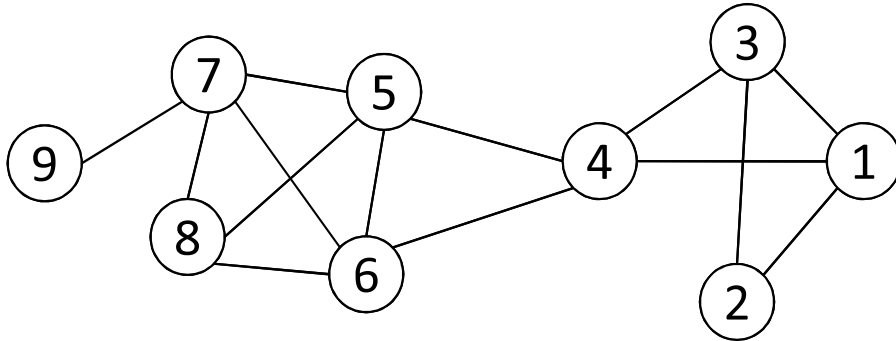


- Suppose we sample a sub-network with nodes {1-5} and find a clique {1, 2, 3} of size 3
- In order to find a clique  $>3$ , remove all nodes with degree  $\leq 3-1=2$ 
  - Remove nodes 2 and 9
  - Remove nodes 1 and 3
  - Remove node 4

# Clique Percolation Method (CPM)

- Clique is a very strict definition, unstable
- Normally use cliques as **a core or a seed** to find larger communities
- CPM is such a method to find **overlapping** communities
  - **Input**
    - A parameter  $k$ , and a network
  - **Procedure**
    - Find out all cliques of size  $k$  in a given network
    - Construct a clique graph. Two cliques are adjacent if they share  $k-1$  nodes
    - Each connected components in the clique graph form a community

# CPM Example



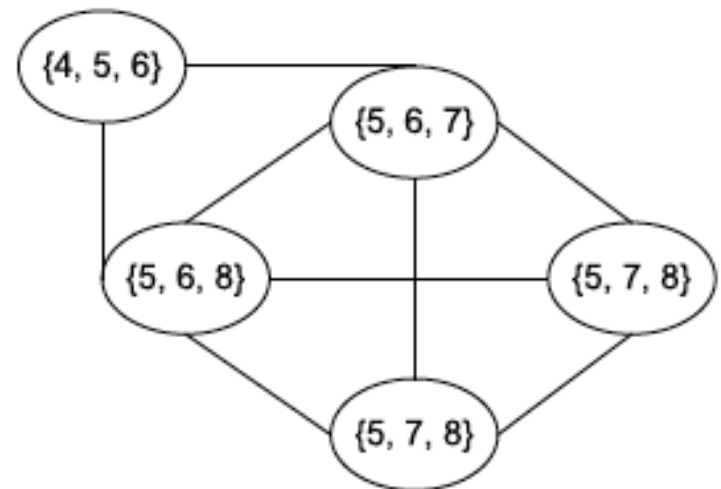
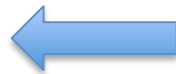
**Cliques of size 3:**

$\{1, 2, 3\}$ ,  $\{1, 3, 4\}$ ,  $\{4, 5, 6\}$ ,  
 $\{5, 6, 7\}$ ,  $\{5, 6, 8\}$ ,  $\{5, 7, 8\}$ ,  
 $\{6, 7, 8\}$



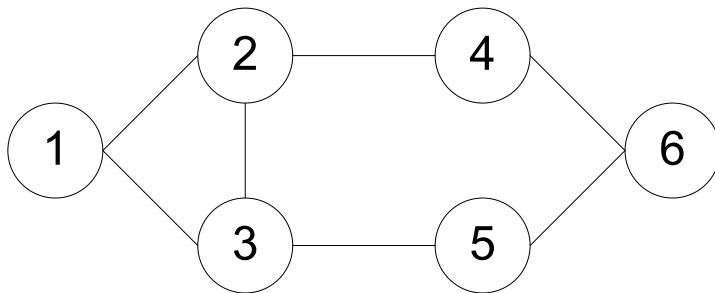
**Communities:**

$\{1, 2, 3, 4\}$   
 $\{4, 5, 6, 7, 8\}$



# Reachability : k-clique, k-club

- Any node in a group should be reachable in k hops
- **k-clique**: a maximal subgraph in which the largest geodesic distance between any nodes  $\leq k$
- **k-club**: a substructure of diameter  $\leq k$



Cliques: {1, 2, 3}

2-cliques: {1, 2, 3, 4, 5}, {2, 3, 4, 5, 6}

2-clubs: {1,2,3,4}, {1, 2, 3, 5}, {2, 3, 4, 5, 6}

- A k-clique might have diameter larger than k in the subgraph
- Commonly used in traditional SNA
- Often involves combinatorial optimization

# Group-Centric Community Detection:

## Density-Based Groups

- The group-centric criterion requires the whole group to satisfy a certain condition
  - E.g., the group density  $\geq$  a given threshold
- A subgraph  $G_s(V_s, E_s)$  is a  $\gamma$  – *dense* **quasi-clique** if

$$\frac{|E_s|}{|V_s|(|V_s| - 1)/2} \geq \gamma$$

- A similar strategy to that of cliques can be used
  - Sample a subgraph, and find a maximal  $\gamma$  – *dense* quasi-clique (say, of size  $k$ )
  - Remove nodes with degree  $< k\gamma$

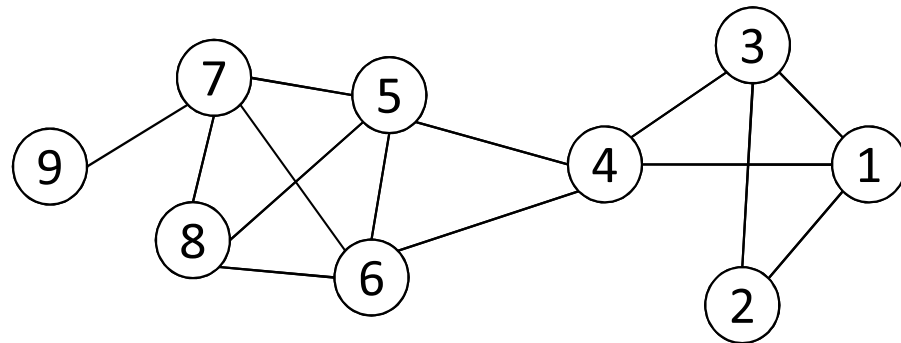
# Network-Centric Community Detection

- Network-centric criterion needs to consider the connections within a network globally
- Goal: partition nodes of a network into disjoint sets
- Approaches:
  - Clustering based on vertex similarity
  - Latent space models
  - Block model approximation
  - Spectral clustering
  - Modularity maximization

# Clustering based on Vertex Similarity

- Apply k-means or similarity-based clustering to nodes
- Vertex similarity is defined in terms of **the similarity of their neighborhood**
- **Structural equivalence**: two nodes are structurally equivalent iff they are connecting to the same set of actors

Nodes 1 and 3 are  
structurally equivalent;  
So are nodes 5 and 7.

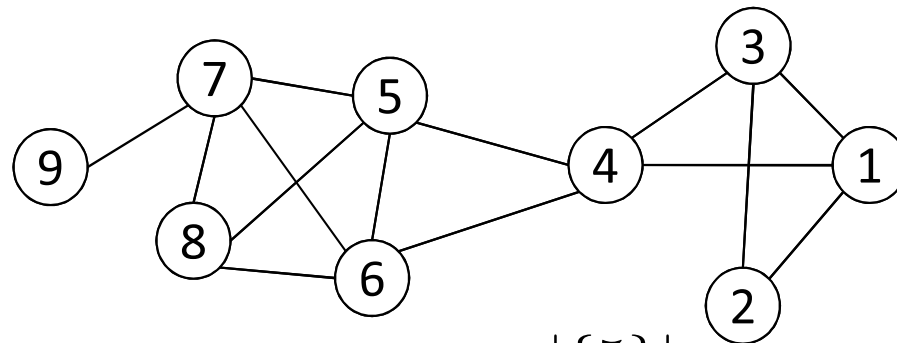


- Structural equivalence is too restrict for practical use.



# Vertex Similarity

- Jaccard Similarity  $Jaccard(v_i, v_j) = \frac{|N_i \cup N_j|}{|N_i \cap N_j|}$
- Cosine similarity  $cosine(v_i, v_j) = \frac{\sum_k A_{ik} A_{jk}}{\sqrt{A_{is}^2 \cdot \sum_t A_{jt}^2}}$



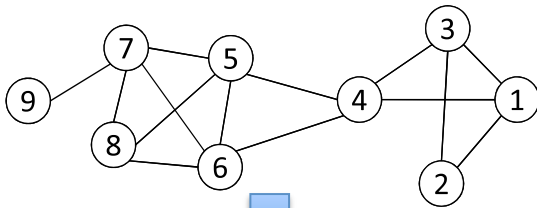
$$Jaccard(4, 6) = \frac{|\{5\}|}{|\{1, 3, 4, 5, 6, 7, 8\}|} = \frac{1}{7}$$

$$cosine(4, 6) = \frac{1}{\sqrt{4 \cdot 4}} = \frac{1}{4}$$

# Latent Space Models

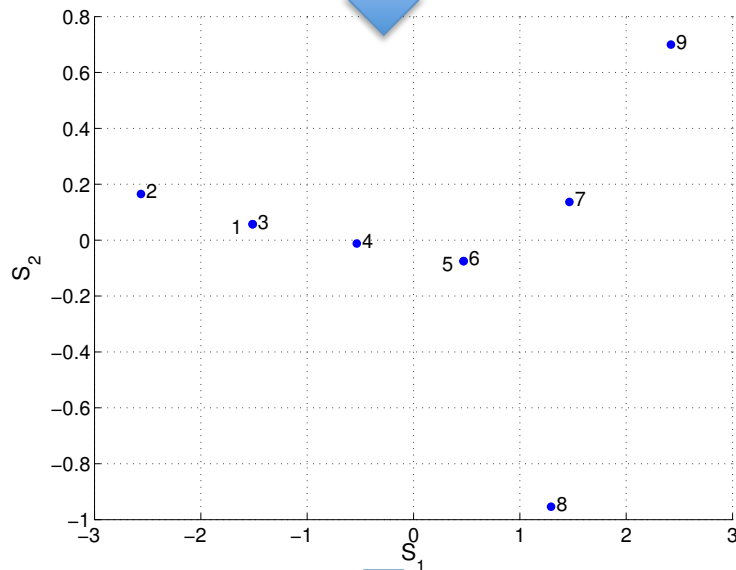
- Map nodes into a low-dimensional space such that the proximity between nodes based on network connectivity is preserved in the new space, then apply k-means clustering
- Multi-dimensional scaling (MDS)
  - Given a network, construct a proximity matrix  $P$  representing the pairwise distance between nodes (e.g., geodesic distance)
  - Let  $S \in R^{n \times \ell}$  denote the coordinates of nodes in the low-dimensional space
$$SS^T \approx -\frac{1}{2}\left(I - \frac{1}{n}\mathbf{1}\mathbf{1}^T\right)(P \circ P)\left(I - \frac{1}{n}\mathbf{1}\mathbf{1}^T\right) = \tilde{P}$$
  - Objective function:  $\min \|SS^T - \tilde{P}\|_F^2$
  - Solution:  $S = V\Lambda^{\frac{1}{2}}$
  - $V$  is the top  $\ell$  eigenvectors of  $\tilde{P}$ , and  $\Lambda$  is a diagonal matrix of top eigenvalues  $\Lambda = \text{diag}(\lambda_1, \lambda_2, \dots, \lambda_\ell)$

# MDS Example



geodesic  
distance

$$P = \begin{bmatrix} 0 & 1 & 1 & 1 & 2 & 2 & 3 & 3 & 4 \\ 1 & 0 & 1 & 2 & 3 & 3 & 4 & 4 & 5 \\ 1 & 1 & 0 & 1 & 2 & 2 & 3 & 3 & 4 \\ 1 & 2 & 1 & 0 & 1 & 1 & 2 & 2 & 3 \\ 2 & 3 & 2 & 1 & 0 & 1 & 1 & 1 & 2 \\ 2 & 3 & 2 & 1 & 1 & 0 & 1 & 1 & 2 \\ 3 & 4 & 3 & 2 & 1 & 1 & 0 & 1 & 1 \\ 3 & 4 & 3 & 2 & 1 & 1 & 1 & 0 & 2 \\ 4 & 5 & 4 & 3 & 2 & 2 & 1 & 2 & 0 \end{bmatrix}$$



$$\tilde{P} = \begin{bmatrix} 2.46 & 3.96 & 1.96 & 0.85 & -0.65 & -0.65 & -2.21 & -2.04 & -3.65 \\ 3.96 & 6.46 & 3.96 & 1.35 & -1.15 & -1.15 & -3.71 & -3.54 & -6.15 \\ 1.96 & 3.96 & 2.46 & 0.85 & -0.65 & -0.65 & -2.21 & -2.04 & -3.65 \\ 0.85 & 1.35 & 0.85 & 0.23 & -0.27 & -0.27 & -0.82 & -0.65 & -1.27 \\ -0.65 & -1.15 & -0.65 & -0.27 & 0.23 & -0.27 & 0.68 & 0.85 & 1.23 \\ -0.65 & -1.15 & -0.65 & -0.27 & -0.27 & 0.23 & 0.68 & 0.85 & 1.23 \\ -2.21 & -3.71 & -2.21 & -0.82 & 0.68 & 0.68 & 2.12 & 1.79 & 3.68 \\ -2.04 & -3.54 & -2.04 & -0.65 & 0.85 & 0.85 & 1.79 & 2.46 & 2.35 \\ -3.65 & -6.15 & -3.65 & -1.27 & 1.23 & 1.23 & 3.68 & 2.35 & 6.23 \end{bmatrix}$$

$$V = \begin{bmatrix} -0.33 & 0.05 \\ -0.55 & 0.14 \\ -0.33 & 0.05 \\ -0.11 & -0.01 \\ 0.10 & -0.06 \\ 0.10 & -0.06 \\ 0.32 & 0.11 \\ 0.28 & -0.79 \\ 0.52 & 0.58 \end{bmatrix}, \quad \Lambda = \begin{bmatrix} 21.56 & 0 \\ 0 & 1.46 \end{bmatrix}, \quad S = V\Lambda^{1/2} = \begin{bmatrix} -1.51 & 0.06 \\ -2.56 & 0.17 \\ -1.51 & 0.06 \\ -0.53 & -0.01 \\ 0.47 & -0.08 \\ 0.47 & -0.08 \\ 1.47 & 0.14 \\ 1.29 & -0.95 \\ 2.42 & 0.70 \end{bmatrix}$$

Two communities:  
{1, 2, 3, 4} and {5, 6, 7, 8, 9}

# Block Models

Table 3.1: Adjacency Matrix									
-	1	1	1	0	0	0	0	0	0
1	-	1	0	0	0	0	0	0	0
1	1	-	1	0	0	0	0	0	0
1	0	1	-	1	1	0	0	0	0
0	0	0	1	-	1	1	1	0	0
0	0	0	1	1	-	1	1	0	0
0	0	0	0	1	1	-	1	1	0
0	0	0	0	1	1	1	-	0	0
0	0	0	0	0	0	1	0	-	0

$$\min \|A - S\Sigma S^T\|_F^2$$

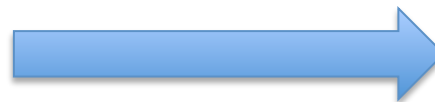


Table 3.2: Ideal Block Structure									
1	1	1	1	0	0	0	0	0	0
1	1	1	1	0	0	0	0	0	0
1	1	1	1	0	0	0	0	0	0
1	1	1	1	0	0	0	0	0	0
0	0	0	0	1	1	1	1	1	1
0	0	0	0	1	1	1	1	1	1
0	0	0	0	1	1	1	1	1	1
0	0	0	0	1	1	1	1	1	1
0	0	0	0	1	1	1	1	1	1
0	0	0	0	1	1	1	1	1	1

- $S$  is the community indicator matrix
- Relax  $S$  to be numerical values, then the optimal solution corresponds to the **top eigenvectors** of  $A$

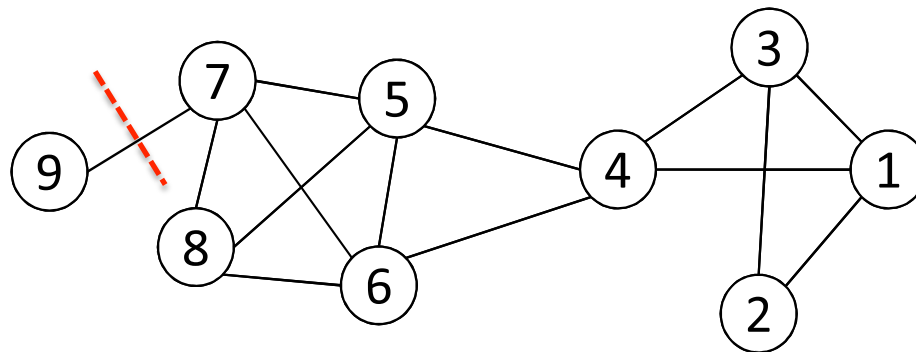
$$S = \begin{bmatrix} 0.20 & -0.52 \\ 0.11 & -0.43 \\ 0.20 & -0.52 \\ 0.38 & -0.30 \\ 0.47 & 0.15 \\ 0.47 & 0.15 \\ 0.41 & 0.28 \\ 0.38 & 0.24 \\ 0.12 & 0.11 \end{bmatrix}, \Sigma = \begin{bmatrix} 3.5 & 0 \\ 0 & 2.4 \end{bmatrix}.$$



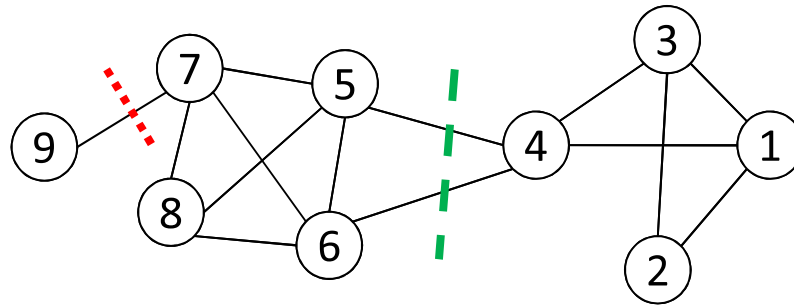
Two communities:  
 $\{1, 2, 3, 4\}$  and  $\{5, 6, 7, 8, 9\}$

# Cut

- Most interactions are within group whereas interactions between groups are few
- community detection → minimum cut problem
- **Cut**: A partition of vertices of a graph into two disjoint sets
- **Minimum cut problem**: find a graph partition such that the number of edges between the two sets is minimized



# Ratio Cut & Normalized Cut



- **Minimum cut often** returns an imbalanced partition, with one set being a singleton
- Change the objective function to consider community size

$$\text{Ratio Cut}(\pi) = \frac{1}{k} \sum_{i=1}^k \frac{\text{cut}(C_i, \bar{C}_i)}{|C_i|},$$

$$\text{Normalized Cut}(\pi) = \frac{1}{k} \sum_{i=1}^k \frac{\text{cut}(C_i, \bar{C}_i)}{\text{vol}(C_i)}$$

$C_i$ : a community

$|C_i|$ : number of nodes in  $C_i$

$\text{vol}(C_i)$ : sum of degrees in  $C_i$

# Ratio Cut & Normalized Cut Example

For partition in red:  $\pi_1$

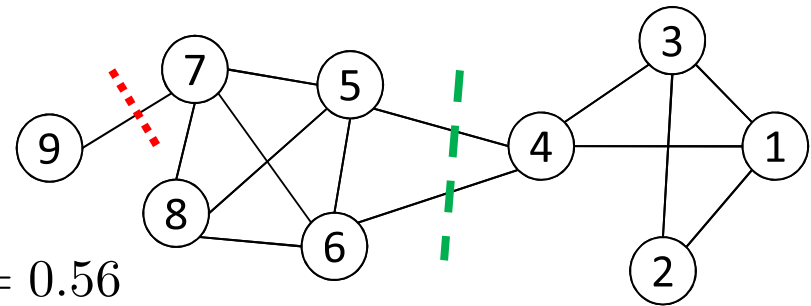
$$\text{Ratio Cut}(\pi_1) = \frac{1}{2} \left( \frac{1}{1} + \frac{1}{8} \right) = 9/16 = 0.56$$

$$\text{Normalized Cut}(\pi_1) = \frac{1}{2} \left( \frac{1}{1} + \frac{1}{27} \right) = 14/27 = 0.52$$

For partition in green:  $\pi_2$

$$\text{Ratio Cut}(\pi_2) = \frac{1}{2} \left( \frac{2}{4} + \frac{2}{5} \right) = 9/20 = 0.45 < \text{Ratio Cut}(\pi_1)$$

$$\text{Normalized Cut}(\pi_2) = \frac{1}{2} \left( \frac{2}{12} + \frac{2}{16} \right) = 7/48 = 0.15 < \text{Normalized Cut}(\pi_1)$$



Both ratio cut and normalized cut prefer a balanced partition

# Spectral Clustering

- Both ratio cut and normalized cut can be reformulated as

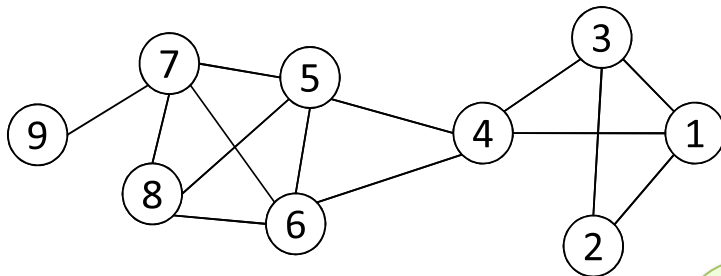
$$\min_{S \in \{0,1\}^{n \times k}} Tr(S^T \tilde{L} S)$$

- Where  $\tilde{L} = \begin{cases} D - A & \text{graph Laplacian for ratio cut} \\ I - D^{-1/2} A D^{-1/2} & \text{normalized graph Laplacian} \end{cases}$   
 $D = \text{diag}(d_1, d_2, \dots, d_n)$  A diagonal matrix of degrees

- Spectral relaxation:**  $\min_S Tr(S^T \tilde{L} S) \quad s.t. \quad S^T S = I_k$
- Optimal solution: top eigenvectors with the smallest eigenvalues



# Spectral Clustering Example



$$D = \text{diag}(3, 2, 3, 4, 4, 4, 4, 3, 1)$$

The 1<sup>st</sup> eigenvector means all nodes belong to the same cluster, no use

Two communities:  
 $\{1, 2, 3, 4\}$  and  $\{5, 6, 7, 8, 9\}$

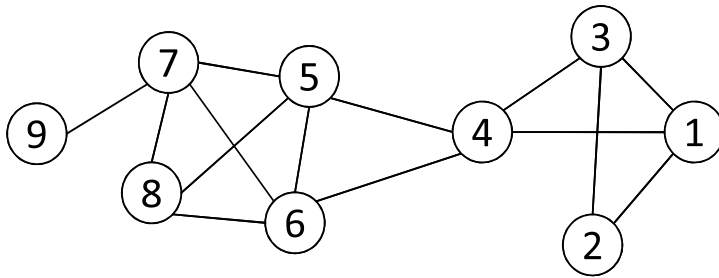
k-means



$$\tilde{L} = D - A = \begin{bmatrix} 3 & -1 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & 2 & -1 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & -1 & 3 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & -1 & 4 & -1 & -1 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 & 4 & -1 & -1 & -1 & 0 \\ 0 & 0 & 0 & -1 & -1 & 4 & -1 & -1 & 0 \\ 0 & 0 & 0 & 0 & -1 & -1 & 4 & -1 & -1 \\ 0 & 0 & 0 & 0 & -1 & -1 & -1 & 3 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & -1 & 0 & 1 \end{bmatrix} \rightarrow S = \begin{bmatrix} 0.33 & -0.38 \\ 0.33 & -0.48 \\ 0.33 & -0.38 \\ 0.33 & -0.12 \\ 0.33 & 0.16 \\ 0.33 & 0.16 \\ 0.33 & 0.30 \\ 0.33 & 0.24 \\ 0.33 & 0.51 \end{bmatrix}$$

# Modularity Maximization

- Modularity measures the strength of a community partition by taking into account the degree distribution
- Given a network with  $m$  edges, the expected number of edges between two nodes with  $d_i$  and  $d_j$  is  $d_i d_j / 2m$



The expected number of edges between nodes 1 and 2 is

$$3 \cdot 2 / (2 \cdot 14) = 3/14$$

- Strength of a community:  $\sum_{i \in C, j \in C} A_{ij} - d_i d_j / 2m$
- Modularity:  $Q = \frac{1}{2m} \sum_{\ell=1}^k \sum_{i \in C_\ell, j \in C_\ell} (A_{ij} - d_i d_j / 2m)$
- A larger value indicates a good community structure

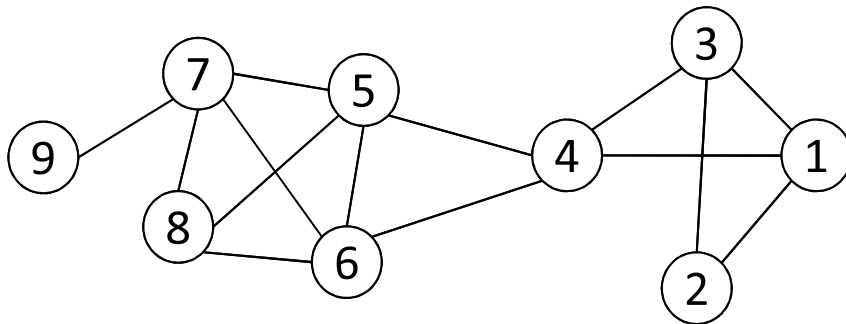
# Modularity Matrix

- Modularity matrix:  $B = A - \mathbf{d}\mathbf{d}^T/2m$  ( $B_{ij} = A_{ij} - d_i d_j / 2m$ )
- Similar to spectral clustering, Modularity maximization can be reformulated as

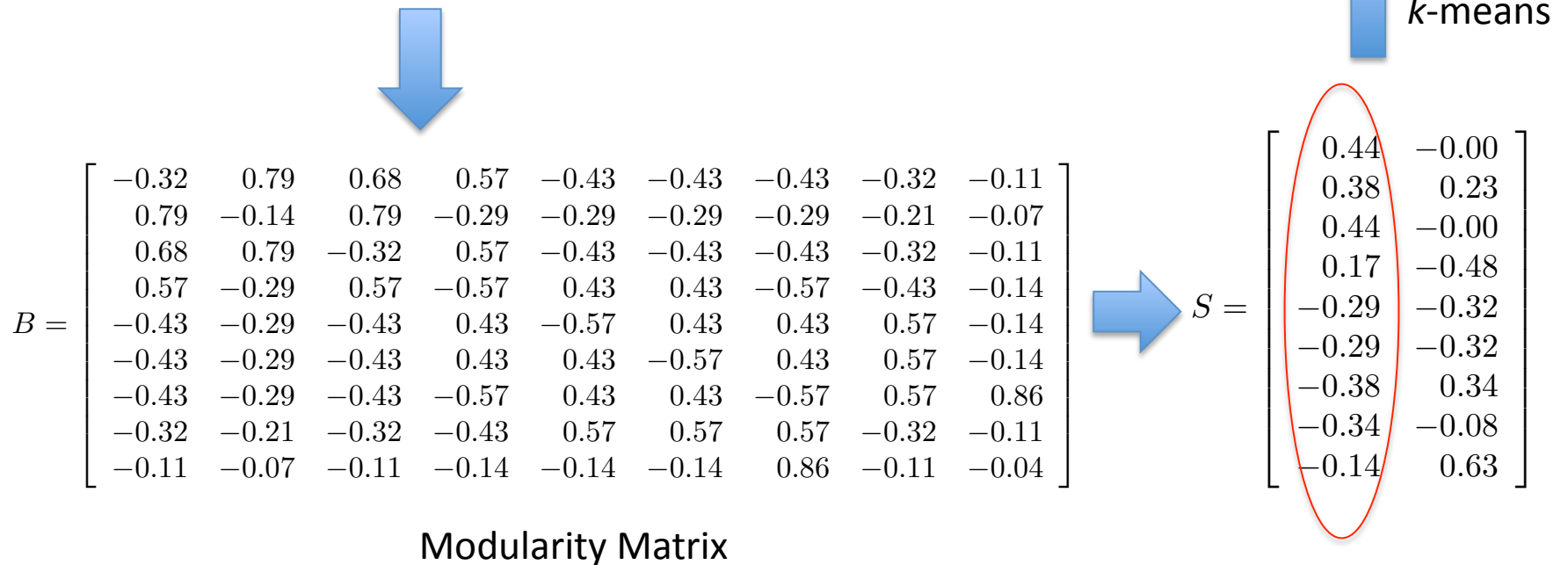
$$\max Q = \frac{1}{2m} \text{Tr}(S^T B S) \quad s.t. \quad S^T S = I_k$$

- Optimal solution: top eigenvectors of the modularity matrix
- Apply k-means to S as a post-processing step to obtain community partition

# Modularity Maximization Example

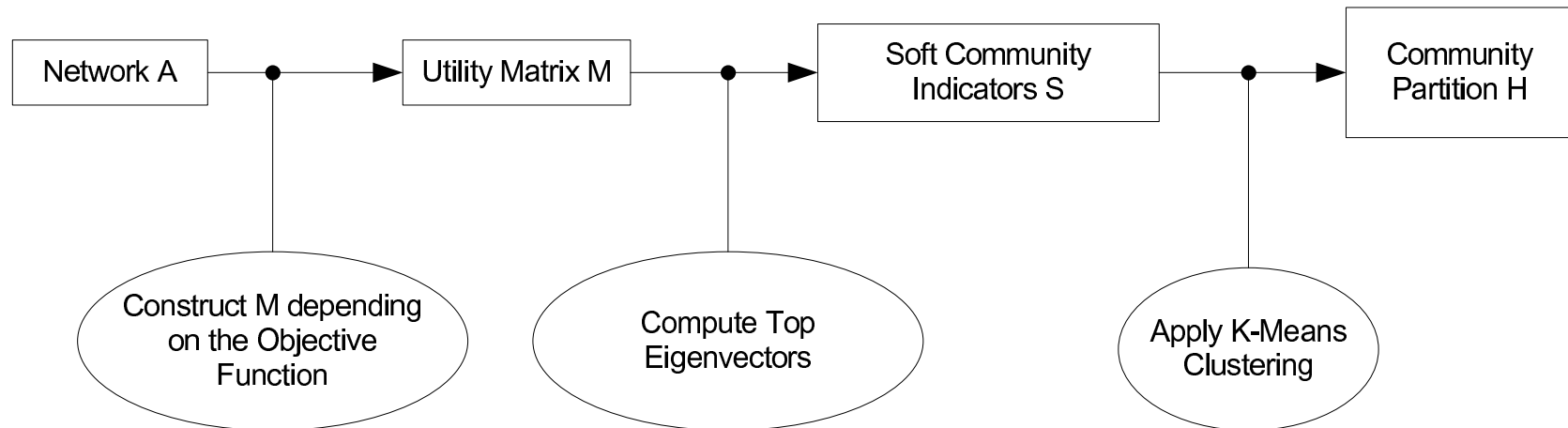


Two Communities:  
 $\{1, 2, 3, 4\}$  and  $\{5, 6, 7, 8, 9\}$



# A Unified View for Community Partition

- Latent space models, block models, spectral clustering, and modularity maximization can be unified as



$$\text{Utility Matrix } M = \begin{cases} \text{modified proximity matrix } \tilde{P} & \text{if latent space models} \\ \text{adjacency matrix } A & \text{if block models} \\ \text{graph Laplacian } \tilde{L} & \text{if spectral clustering} \\ \text{modularity maximization } B & \text{if modularity maximization} \end{cases}$$

# Hierarchy-Centric Community Detection

- Goal: build a hierarchical structure of communities based on network topology
- Allow the analysis of a network at different resolutions
- Representative approaches:
  - Divisive Hierarchical Clustering
  - Agglomerative Hierarchical clustering

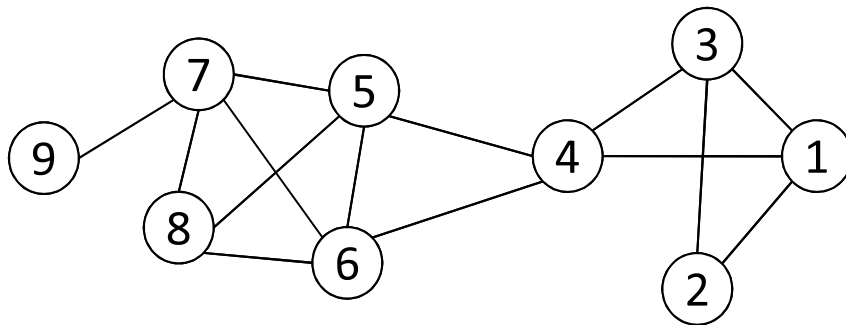
# Divisive Hierarchical Clustering

- Divisive clustering
  - Partition nodes into several sets
  - Each set is further divided into smaller ones
  - Network-centric partition can be applied for the partition
- One particular example: recursively remove the “weakest” tie
  - Find the edge with the least strength
  - Remove the edge and update the corresponding strength of each edge
- Recursively apply the above two steps until a network is discomposed into desired number of connected components.
- Each component forms a community

# Edge Betweenness

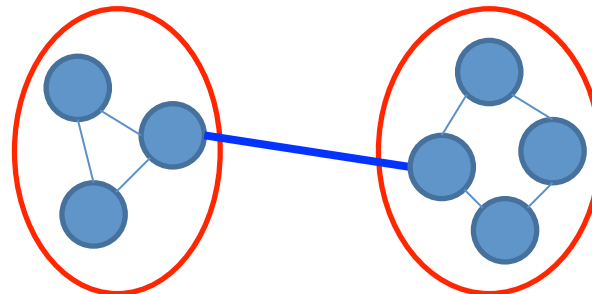
- The strength of a tie can be measured by **edge betweenness**
- **Edge betweenness**: the number of shortest paths that pass along with the edge

$$\text{edge-betweenness}(e) = \sum_{s < t} \frac{\sigma_{st}(e)}{\sigma_{s,t}}$$



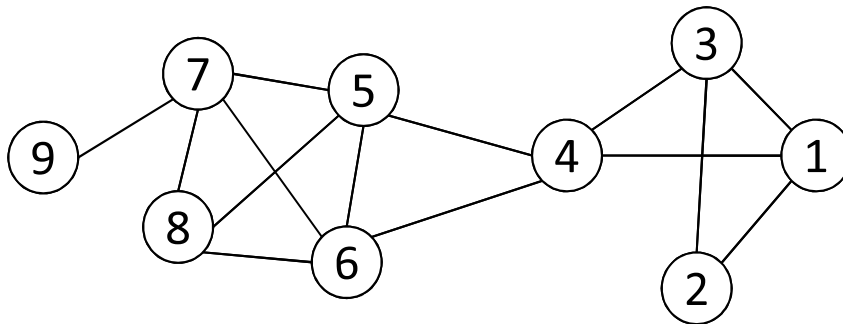
The edge betweenness of  $e(1, 2)$  is 4, as all the shortest paths from 2 to  $\{4, 5, 6, 7, 8, 9\}$  have to either pass  $e(1, 2)$  or  $e(2, 3)$ , and  $e(1, 2)$  is the shortest path between 1 and 2

- The edge with higher betweenness tends to be the bridge between two communities.





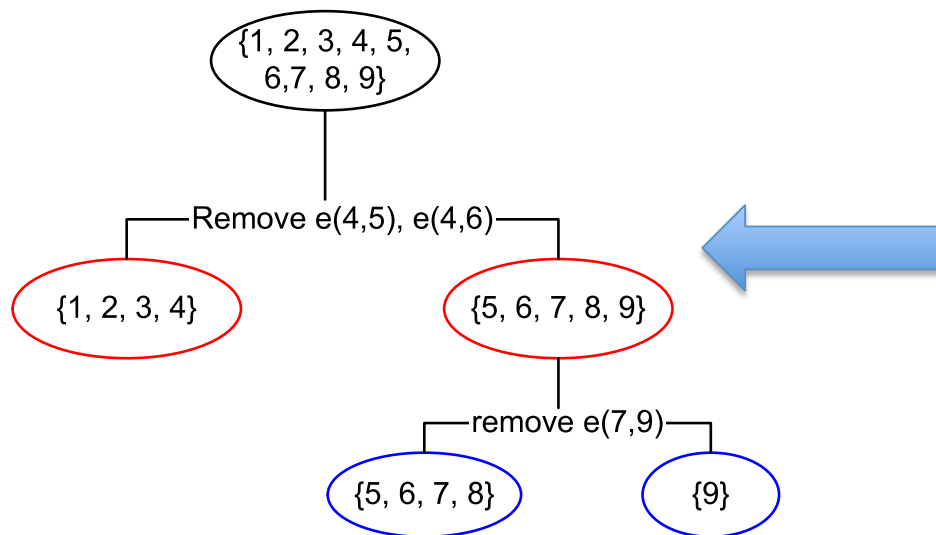
# Divisive clustering based on edge betweenness



Initial betweenness value

Table 3.3: Edge Betweenness

	1	2	3	4	5	6	7	8	9
1	0	4	1	9	0	0	0	0	0
2	4	0	4	0	0	0	0	0	0
3	1	4	0	9	0	0	0	0	0
4	9	0	9	0	10	10	0	0	0
5	0	0	0	10	0	1	6	3	0
6	0	0	0	10	1	0	6	3	0
7	0	0	0	0	6	6	0	2	8
8	0	0	0	0	3	3	2	0	0
9	0	0	0	0	0	0	8	0	0

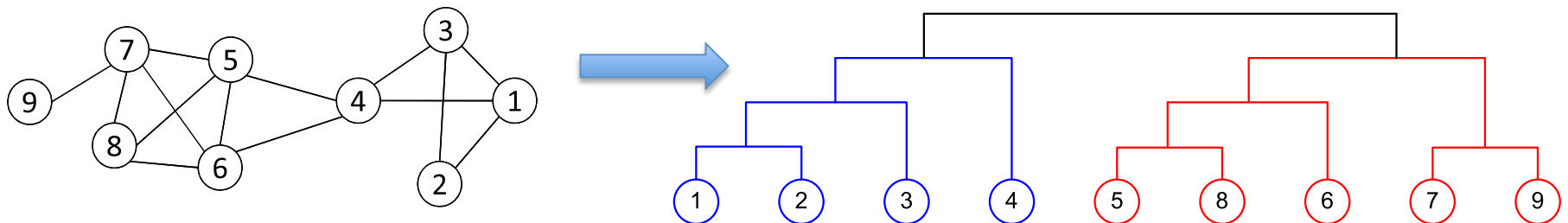


After remove  $e(4,5)$ , the betweenness of  $e(4, 6)$  becomes 20, which is the highest;

After remove  $e(4,6)$ , the edge  $e(7,9)$  has the highest betweenness value 4, and should be removed.

# Agglomerative Hierarchical Clustering

- Initialize each node as a community
- Merge communities successively into larger communities following a certain criterion
  - E.g., based on modularity increase



# Summary of Community Detection

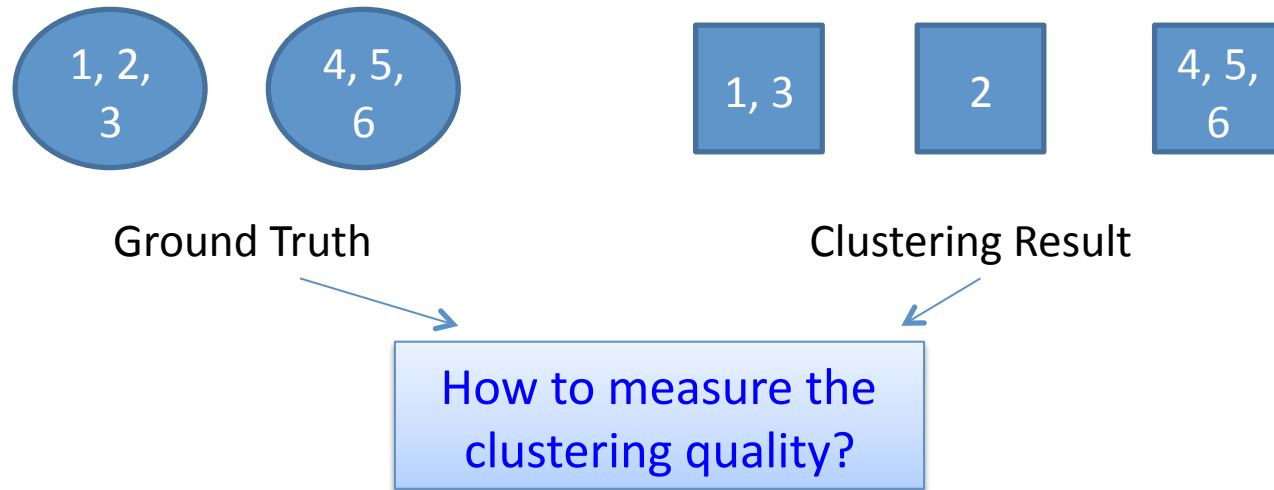
- **Node**-Centric Community Detection
  - *cliques, k-cliques, k-clubs*
- **Group**-Centric Community Detection
  - *quasi-cliques*
- **Network**-Centric Community Detection
  - *Clustering based on vertex similarity*
  - *Latent space models, block models, spectral clustering, modularity maximization*
- **Hierarchy**-Centric Community Detection
  - *Divisive clustering*
  - *Agglomerative clustering*

# **COMMUNITY EVALUATION**

# Evaluating Community Detection (1)

- For groups with clear definitions
  - E.g., Cliques, k-cliques, k-clubs, quasi-cliques
  - Verify whether extracted communities satisfy the definition
- For networks with ground truth information
  - Normalized mutual information
  - Accuracy of pairwise community memberships

# Measuring a Clustering Result



- The number of communities after grouping can be different from the ground truth
- No clear community correspondence between clustering result and the ground truth
- Normalized Mutual Information can be used

# Normalized Mutual Information

- **Entropy**: the information contained in a distribution

$$H(X) = \sum_{x \in X} p(x) \log p(x)$$

- **Mutual Information**: the shared information between two distributions

$$I(X; Y) = \sum_{y \in Y} \sum_{x \in X} p(x, y) \log \left( \frac{p(x, y)}{p_1(x)p_2(y)} \right)$$

- **Normalized Mutual Information** (between 0 and 1)

$$NMI(X; Y) = \frac{I(X; Y)}{\sqrt{H(X)H(Y)}}$$

- Consider a partition as a distribution (probability of one node falling into one community), we can compute the matching between two clusterings

# NMI

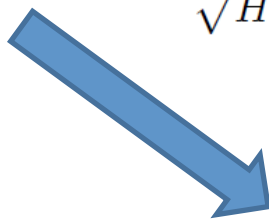
$$H(X) = \sum_{x \in X} p(x) \log p(x)$$



$$\left\{ \begin{array}{l} H(\pi^a) = \sum_h^{k^{(a)}} \frac{n_h^a}{n} \log\left(\frac{n_h^a}{n}\right) \\ H(\pi^b) = \sum_\ell^{k^{(b)}} \frac{n_\ell^b}{n} \log\left(\frac{n_\ell^b}{n}\right) \end{array} \right.$$

$$I(X; Y) = \sum_{y \in Y} \sum_{x \in X} p(x, y) \log \left( \frac{p(x, y)}{p_1(x)p_2(y)} \right) \Rightarrow I(\pi^a, \pi^b) = \sum_h \sum_\ell \frac{n_{h,\ell}}{n} \log \left( \frac{\frac{n_{h,\ell}}{n}}{\frac{n_h^a}{n} \frac{n_\ell^b}{n}} \right)$$

$$NMI(X; Y) = \frac{I(X; Y)}{\sqrt{H(X)H(Y)}}$$

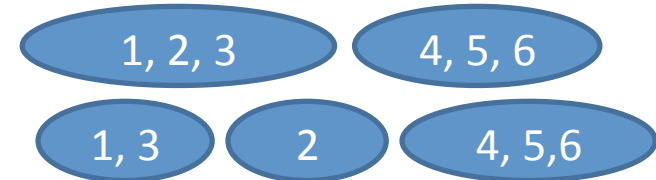


$$NMI(\pi^a, \pi^b) = \frac{\sum_{h=1}^{k^{(a)}} \sum_{\ell=1}^{k^{(b)}} n_{h,\ell} \log \left( \frac{n \cdot n_{h,\ell}}{n_h^{(a)} \cdot n_\ell^{(b)}} \right)}{\sqrt{\left( \sum_{h=1}^{k^{(a)}} n_h^{(a)} \log \frac{n_h^a}{n} \right) \left( \sum_{\ell=1}^{k^{(b)}} n_\ell^{(b)} \log \frac{n_\ell^b}{n} \right)}}$$



# NMI-Example

- Partition a: [1, 1, 1, 2, 2, 2]
- Partition b: [1, 2, 1, 3, 3, 3]



$n = 6$		$n_h^a$		$n_l^b$	$n_{h,l}$	$l=1$	$l=2$	$l=3$
$k^{(a)} = 2$	h=1	3	l=1	2	h=1	2	1	0
$k^{(b)} = 3$	h=2	3	l=2	1	h=2	0	0	3
			l=3	3				

$$NMI(\pi^a, \pi^b) = \frac{\sum_{h=1}^{k^{(a)}} \sum_{\ell=1}^{k^{(b)}} n_{h,\ell} \log \left( \frac{n \cdot n_{h,\ell}}{n_h^{(a)} \cdot n_\ell^{(b)}} \right)}{\sqrt{\left( \sum_{h=1}^{k^{(a)}} n_h^{(a)} \log \frac{n_h^{(a)}}{n} \right) \left( \sum_{\ell=1}^{k^{(b)}} n_\ell^{(b)} \log \frac{n_\ell^{(b)}}{n} \right)}} = 0.8278$$

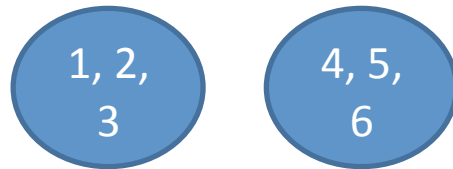
# Accuracy of Pairwise Community Memberships

- Consider all the possible pairs of nodes and check whether they reside in the same community
- An **error** occurs *if*
  - Two nodes belonging to the **same** community are assigned to **different** communities after clustering
  - Two nodes belonging to **different** communities are assigned to the **same** community
- Construct a **contingency table**

Clustering Result		Ground Truth	
		$C(v_i) = C(v_j)$	$C(v_i) \neq C(v_j)$
	$C(v_i) = C(v_j)$	a	b
	$C(v_i) \neq C(v_j)$	c	d

$$accuracy = \frac{a + d}{a + b + c + d} = \frac{a + d}{n(n - 1)/2}$$

# Accuracy Example



Ground Truth



Clustering Result

		Ground Truth	
		$C(v_i) = C(v_j)$	$C(v_i) \neq C(v_j)$
Clustering Result	$C(v_i) = C(v_j)$	4	0
	$C(v_i) \neq C(v_j)$	2	9

$$\text{Accuracy} = (4+9) / (4+2+9+0) = 13/15$$

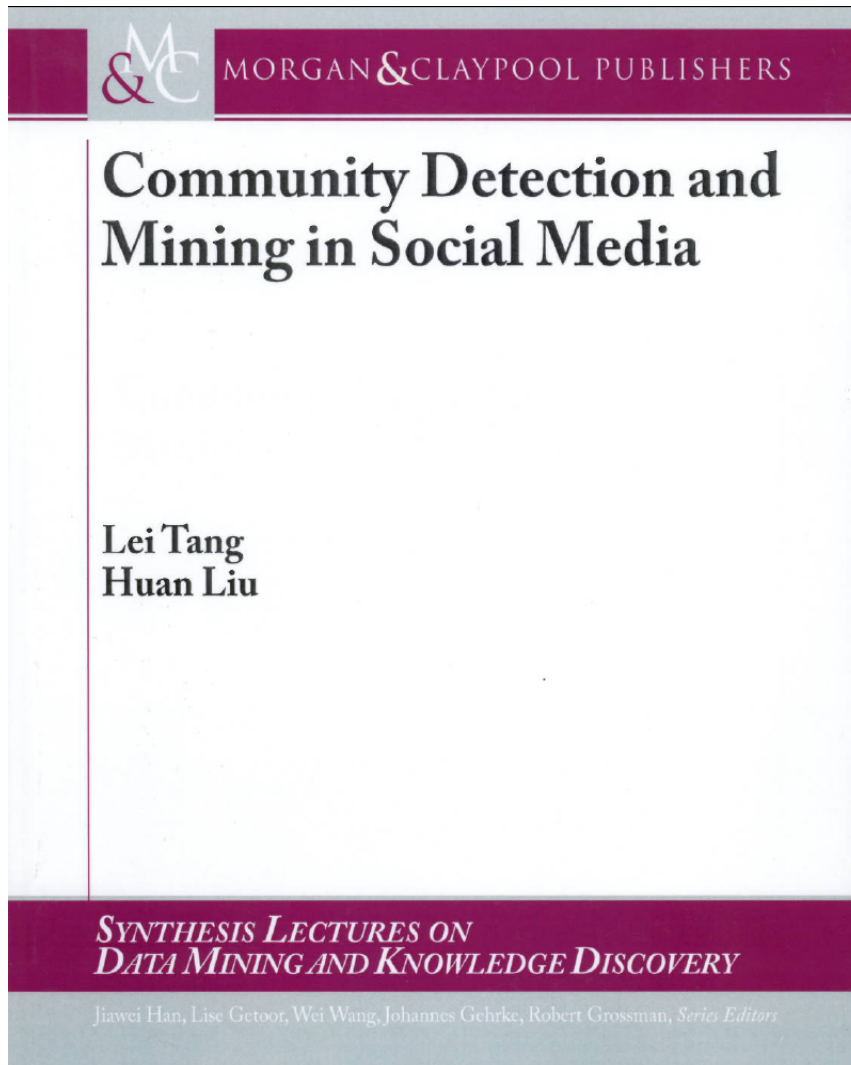
# Evaluation using Semantics

- For networks with semantics
  - Networks come with semantic or attribute information of nodes or connections
  - Human subjects can verify whether the extracted communities are coherent
- Evaluation is qualitative
- It is also intuitive and helps understand a community



# Evaluation without Ground Truth

- For networks without ground truth or semantic information
- This is the most common situation
- An option is to resort to cross-validation
  - Extract communities from a (training) network
  - Evaluate the quality of the community structure on a network constructed from a different date or based on a related type of interaction
- Quantitative evaluation functions
  - modularity
  - block model approximation error



## Book Available at

- [Morgan & claypool Publishers](#)
- [Amazon](#)

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