

JURONG PIONEER JUNIOR COLLEGE 2022 JC2 Preliminary Examination

COMPUTING Higher 2

Paper 1 (Written)

9569/01

19 September 2022

3 hours

Answer & Mark Scheme

(a) (i) Data Integrity	
 Data integrity is the accuracy, and completeness maintained over time and 	Total 1 mark
across formats [1m]. Preserving the data integrity is a constant process.	
(ii) Two threats:	Total 2 marks
Human error: For instance, accidentally deleting a row of data in a spreadsheet	Choose any 2 points
• Inconsistencies across format: For instance, a set of data in Microsoft Excel that	(without elaboration)
relies on cell referencing may not be accurate in a different format that doesn't allow	
those cells to be referenced	
• Collection error: For instance, data collected is inaccurate or lacking information,	
creating an incomplete picture of the subject	
	 Data integrity is the accuracy, and completeness maintained over time and across formats [1m]. Preserving the data integrity is a constant process. (ii) Two threats: Human error: For instance, accidentally deleting a row of data in a spreadsheet Inconsistencies across format: For instance, a set of data in Microsoft Excel that relies on cell referencing may not be accurate in a different format that doesn't allow those cells to be referenced Collection error: For instance, data collected is inaccurate or lacking information,

• External or internal cybersecurity or privacy breaches: For instance, someone	
hacks into your company's database with the intent to damage or steal information,	
or an internal employee damages data with malicious intent	
(iii)	
Data integrity refers to the completeness, consistency, accuracy and the validity of the	Total 2 marks
data [1m]. ie. when recorded, it is recorded exactly as the user intends, and when retrieved,	1 mark for each point
it is in the exact same state that it was recorded.	adequately mentioned.
Data security is the practice of preventing data from being accessed, altered, disclosed,	
or damaged without authorisation [1m]. The term also covers the methods used to do this.	
These include	
(iv)	
Some data may not be being used very often but it may still be useful or needed in the future.	Total 1 mark
In this case data can be archived. Archived data is copied to a suitable storage medium	
(perhaps DVDs or magnetic tape) then it is stored safely and securely. The original data	
is then deleted from the computer system. This is done to free up resources for new data.	
(b)	
Data validation is the process comparing information with a set of rules to ensure it is	Total 2 marks
sensible and reasonable [1m] enough before it is stored in the system.	

				e, hence it is unable to check for the	
C	correctness or acc				
(c)	Indiscriminate or	Total 2 marks			
	disclosure. NRIC	/FIN may be obta	ained and used	for illegal activities such as identity	
		•			
	theft and fraud	i as they can	be used to t	unlock large amounts of personal	
	information rela	ting to the indiv	vidual [1m]. Th	ne NRIC/FIN is also a permanent and	
	irreplaceable ide	entifier specific	to an individua	al [1m].	
(-1) (· >				
(a) C	Organisations can co	onsider to:			T
	 develop, imp 	Total 2 marks			
	• develop, lilip	noment and reg	jaiaily lotion	• •	
	• • •	_			
	them abreast	and updated. [1r	m]		
	them abreast	and updated. [1r	m]	professinals to help them store and	
	them abreast	and updated. [1r	m]		
(e)	them abreast employ author management	and updated. [1r	m]		
(e)	them abreast employ author	and updated. [1r	m] data specialist		
	them abreast employ author management	and updated. [1rrise third party of the data. [1m]	m] data specialist		
	them abreast employ author management (i)	and updated. [1rrise third party of the data. [1m] NRIC 3 rightnesses	m] data specialist		

C = [11 - (46 MC	1m – Obtain check digit = 9	
Check digit of ten	: A is <u>9</u>	

 Birthday in ddmm
 integers
 Digit

 3
 7
 6
 5
 4
 3
 2
 1
 ← digit position

 7
 1
 2
 8
 9
 6
 8-digit TenantID

Product sum = (1*8) + (7*7) + (1*6) + (2*5) + (8*4) + (9*3) + (6*2)= 8 + 49 + 6 + 10 + 32 + 27 + 12 = 144

C = [11 - (144 MOD 11)] MOD 11 = [11 - 1] MOD 11 = 10

Check digit = X (using X to represent 10 in the 8 digit number string code)

Therefore, the check digit the tenantID of B = 1712896x

(f)

Birthday in <i>ddmm</i>			NRIC 3 rightmost integers		Check Digit			
8	7	6	5	4	3	2	1	← digit position
0	9	0	8	9	9	5	<u>3</u>	8-digit TenantID

Product sum = (0*8) + (9*7) + (0*6) + (8*5) + (9*4) + (9*3) + (5*2) + (3*1)

1m – Obtain checkdigit = 10 1m – present check digit represented as by a symbol

Total 2 marks

stated.

Commented [DL1]: This part is 3 marks in the question paper?

	= 0 + 63 + 0 + 40 + 36 + 27 + 10 + 3 = 179 Since 179 MOD 11 <> 0 Therefore TenantID of value '09089953' is not a valid 8 digit number string	1m – obtain product sum = 179 1m – state 179 MOD 11 <>
		0, and conclude tenantID not valid.
(g)	 (i) A primary key ensures unique row identification [1m]. This results in faster sorting, searching, and querying operations [1m]. (ii) A foreign key creates a link between two tables [1m]. It maintains referential integrity [1m] between the referencing column and the referenced column. 	Total 2 marks Total 2 marks
(h)	<pre>Tenant(<u>tenantID</u>, Name, sex, DOB, email, Contact_No) Apartment(<u>apartmentID</u>, level, unit, occupied)</pre>	Total 2 marks 1m – correct PK underlined and seen in both tables. 1m – 2 other sensible fields
	(ii) RentalContract(tenantID, apartmentID, startDate, endDate, bookingDate) (by selecting as the composite keys will not be sufficiently unique enough to uniquely identify the records as they grow in size.)	Total 1 mark 1m – correct composite keys underlined.
(i)	ER diagram	Total 3 marks 1m – all entities correct

	TENANT	RENTAL	APARTMENT		1m - r/s between TENANT and RENTALCONTRACT 1m - r/s between APARTMENT and RENTALCONTRACT	
(j)	The foreign key Tenant TenantID of the of the ta The foreign key apartme apartmentID of the of the		Total 2 marks			
(k)	Data inconsistency happe files/locations within a dat Data redundancy occurs locations/files [1m] within	of itself [1m].	Total 2 marks			
(1)	increase the likelihood o	f the same data element of partial update [1m] when but not the others [1	re the data element (gets updated	Total 2 marks	

Array stores all its data elements continguously in order not to create space wasteages. However maintaining data elements stored continguosly requires high memory resources that is at the expense of performance. Inserting an item into a non-empty array at index location i in an array would first require moving all the existing data items from array locations i+1 onwards to the last data item in the array one location to the right [1m]. Similary after deleting a data item in location i of the array would require every data items from locations i+1 onwards to the last item forward by one location [1m] in order to ensure data items remain stored continguously. (ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable • start pointer to point to the new node.	2	(a)	(i)	
However maintaining data elements stored continguosly requires high memory resources that is at the expense of performance. Inserting an item into a non-empty array at index location i in an array would first require moving all the existing data items from array locations i+1 onwards to the last data item in the array one location to the right [1m]. Similary after deleting a data item in location i of the array would require every data items from locations i+1 onwards to the last item forward by one location [1m] in order to ensure data items remain stored continguously. (ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node/address of new node		(-)	• •	Total 3 marks
array at index location i in an array would first require moving all the existing data items from array locations i+1 onwards to the last data item in the array one location to the right [1m]. Similary after deleting a data item in location i of the array would require every data items from locations i+1 onwards to the last item forward by one location [1m] in order to ensure data items remain stored continguously. (ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF Start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node/address of new node.			However maintaining data elements stored continguosly requires high memory	
items from array locations i+1 onwards to the last data item in the array one location to the right [1m]. Similary after deleting a data item in location i of the array would require every data items from locations i+1 onwards to the last item forward by one location [1m] in order to ensure data items remain stored continguously. (ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node/address of new node.			resources that is at the expense of performance. Inserting an item into a non-empty	
location to the right [1m]. Similary after deleting a data item in location i of the array would require every data items from locations i+1 onwards to the last item forward by one location [1m] in order to ensure data items remain stored continguously. (ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF Start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node			array at index location i in an array would first require moving all the existing data	
Similary after deleting a data item in location i of the array would require every data items from locations i+1 onwards to the last item forward by one location [1m] in order to ensure data items remain stored continguously. (ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF Start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node/address of new node			items from array locations i+1 onwards to the last data item in the array one	
items from locations i+1 onwards to the last item forward by one location [1m] in order to ensure data items remain stored continguously. (ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF Start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node/address of new node			location to the right [1m].	
 (ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF Start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node/address of new node.			Similary after deleting a data item in location i of the array would require every data	
(ii) For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF Start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node			items from locations i+1 onwards to the last item forward by one location [1m] in	
For insertion and deletion of data item in a linked list, only the affected pointer values need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF Start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 1 mark Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node			order to ensure data items remain stored continguously.	
need to be updated [1m]. (b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node			(ii)	
(b) Inserting a node to the front of a linked list: • Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL • IF Start pointer is NULL THEN set start to point to new node/address of new node. • Else, store the value of start pointer into a temporary variable Total 3 marks 1m - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node			For insertion and deletion of data item in a linked list, only the affected pointer values	Total 1 mark
 (b) Inserting a node to the front of a linked list: Creation and initialisation of a new node that contains the data_pointer = newItem, and next_pointer = NULL IF Start pointer is NULL THEN set start to point to new node/address of new node. Else, store the value of start pointer into a temporary variable Im - initialise node with data and next pointer = NULL and if Linked List empty, set start to point to new node 			need to be updated [1m].	
 Creation and initialisation of a new node that contains the data_pointer =		4.		Total 3 marks
newItem, and next_pointer = NULL IF Start pointer is NULL THEN set start to point to new node/address of new node. Else, store the value of start pointer into a temporary variable data and next pointer = NULL and if Linked List empty, set start to point to new node		(a)		1m – initialise node with
 IF Start pointer is NULL THEN set start to point to new node/address of new node. Else, store the value of start pointer into a temporary variable NULL and if Linked List empty, set start to point to new node				data and next pointer =
Else, store the value of start pointer into a temporary variable new node			IF Start pointer is NULL THEN set start to point to new node/address of new	NULL and if Linked List
I new node				empty, set start to point to
• start pointer to point to the new node.				new node
next pointer of the new node to store the value of the temporary variable (next to			·	

point to the node that was previously the start node)

(c) Deleting a node from the rear of non-empty linked list:

If linked list contains **1 node** before deletion:

- Set start to NULL [1m]
- Deallocate and return the current node back to the main memory

If linked list contains more than 1 node before deletion:

- Using 2 pointers previous and current to traverse in tandem from the first node of the linked list until current reaches the last node (ie. next pointer of node points to null). [1m]
- Set next pointer of previous node to NULL. [1m]
- Deallocate and return the current node back to the main memory

 1m – if Linked List not empty, set start into temp var,

1m - set start to new node, and set next_pointer of new node to temp var.

Total 3 marks

1m – to consider if linked
list contains 1 node before
deletion and set start to null

1m – for traversing L from
start with prev and cur
until cur reaches the last
node.

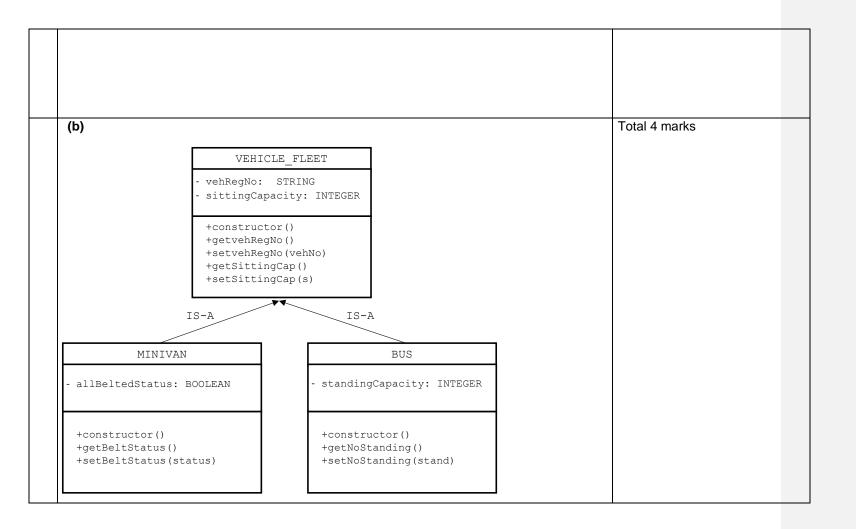
1m for update prev.next
= null

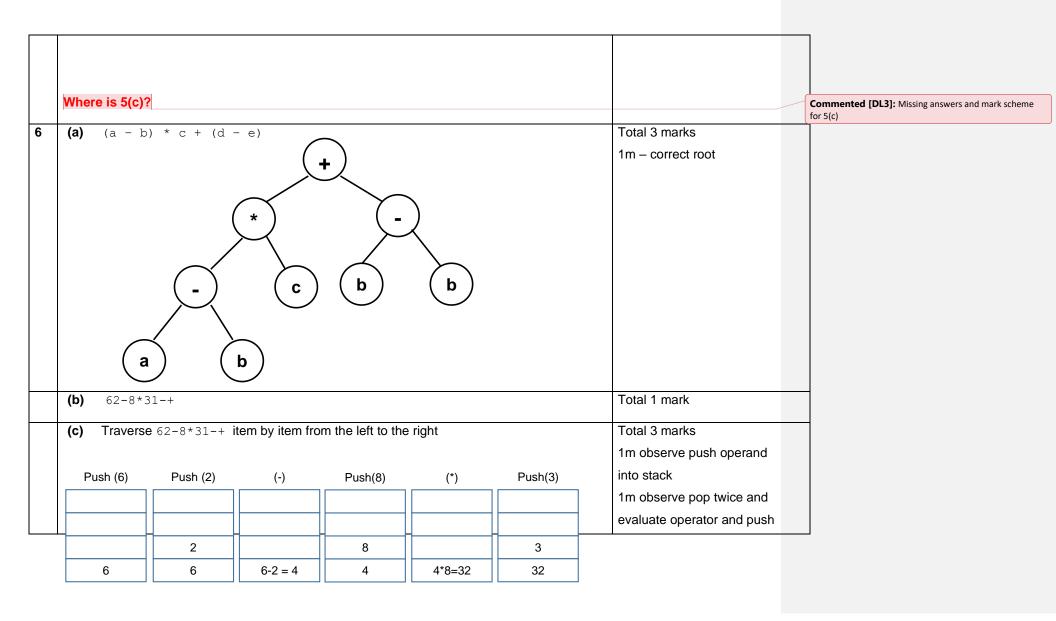
	(d)	Slower search/traversal time as it does not allow direct access to node or traversal needs to always start from the first node. (SR: No marks if student answers "does not allow direct access or trversal starts from first node as they do not answer the question directly)	Total 1 mark
	(e)	<pre>FUNCTION find(k) returns BOOLEAN current = start - 1m WHILE current <> NULL - 1m</pre>	Total 3 marks 1m – initialise current to start 1m – while loop to check if the end is reached. 1m – if statement to check if search key is found.
	(f)	<pre>(i) Create(Q) (ii) Insert(Q, data, Length(Q)) (iii) IF isEmpty(Q) = True THEN -1m</pre>	Total 1 mark Total 1 mark Total 3 marks
3	(a)		T
	1.	First chooses the first item in the array as the pivot .	Total 4 marks 1m for 1 point
	2.	$\textbf{Partitions} \text{ by } \underline{\textbf{moving}} \text{ the items in the data set } \textbf{about the pivot} \text{ items to smaller than pivot}$	Titi for a point
		to the left and larger to the right.	

	3	This wor	ıld result in the n	ivot be in a position that	splits/partitions the data set into 2		
	0.			Total and pooling it that			
		parts (al	so known as the	spilt point).			
	4.	. Recursiv	ely performs the	tasks 1, 2 and 3 on each	sub-array partitioned u the spilt point		
		until data	a until 1 item rema				
	(b)	Best cas	e time complexity	'=0(nlog₂n)		Total 1 mark	
	(c)	The wor	st case scenario	Total 2 marks			
		spilting	the array of n el	rays of size 1 and n-1 [1m].			
4	(a)	The three	programming con	structs are:		Total 2 marks	
		Sequence	, Selection, Iterati	on (Repetition)		3 correct 2 marks	
		(Do not a	ccept if students	answers IF-THEN-ELSE f	or selection, or LOOPS for iteration	2 correct 1 mark	
		without me	entioning the actu	al keywords).		0 mark otherwise	
	(b)	Sequence	and Selection			Total 2 marks	
						Each word 1 mark	
	(c)	Use mean	ingful names for i	Total 1 mark (exactly)			
	(d)	Line 5		Total 1 mark			
	(e)						
		n	Recursive call	Print			
		36	X(36)				

Commented [DL2]: How are the marks distributed?

			1			
	18	X(18)				
	4	X(4)				
	2	X(2)				
	1	X(1)	1			
			0			
			0			
			1			
			0			
			0			
	(f) Xisap	rocedure that conve	erts prints a denary number	to its binary equvalent	Total 1 mark	
5	(a) Object	t is an instance of a	class [1m].		Total 2 marks	
	Class	is a blueprint or tem	nplate from which objects a	e created [1m].		
		, , , , , , , , , , , , , , , , , , , ,	,			
1						





						evaluated result back into
						stack if end of postfix not
						reached.
						1m observe pop final result
	Push(1)	(-)	(+)	POP()		out from stack when the
					Result = POP()	postfix has reached the
	1				POP() from STACK since all operators and	end.
	3	3-1 = 2			operands in postfix	
	32	32	32+2=34	34	expression have been visited	
7.	domain r Step 2: When th If there is Step 3: If the loc the requ returns th	name server. e local domain rethis record, the all cache does not est to the root	name server rece local domain nar not have the record domain name so	vives the request me server directl rd, the local dom erver, and then	d sends the request to the local t, it first queries the local cache. It is preturns the result of the query. It is name server directly sends the root domain name server domain of the root) of the local	1 mark for 1 step except step 6.

	Step 4: The local server sends a request to the domain name server (Top Level Domain server)	
	returned in the previous step, and then the server that accepts the request queries its	
	own cache. If there is no such record, it returns the address of the relevant lower-level	
	domain name server (Authoritative Server).	
	Step 5: Repeat step 4 until you find the correct record.	
	Step 6: The local domain name server saves the returned results to the cache for the next use	
	and returns the results (the webpage ip address) to the client.	
8.	(a) The goal of network security is to achieve the following:	Total 3 marks
	Confidentiality - Ensure that private data that is disclosed or made available only to	1m for each point with
		correct reason.
	authorised persons or organisations and be protected from any unauthorised	
	access.	
	Integrity - to make sure that data is accurate and consistent, and not changed by	
	Integrity - to make sure that data is accurate and consistent , and not changed by unauthorised personnel.	
	unauthorised personnel.	

(b) In P2P network, the absence of a centralised server would mean that any device	Total 2 marks
connected to the requesting device can share a fragment of the resource to the requesting	1m – no centralised server,
	any device can be send.
device. Therefore, there is no way for any parties to control what content is being	1m – no centralised server,
$\textbf{transmitted} \ \text{from the senders and the receiver}, \ \text{which could carry risks and vulnerabilities relating}$	unable to have regulations
to data integrity, viruses, spyware, adware, and unwanted files.	on the data sent resulting in
	viruses, spyware, adware
	etc sent to the receiver.
(c)	Total 3 marks
A firewall is :	
1. a barrier between a trusted internal network and untrusted external network, such	
as the Internet [1m].	
2. can be a software program or a hardware device, or a combination of both that monitors	
and controls incoming and outgoing network traffic based on predetermined	
security rules [1m].	
3. Serves as a barrier between a device's internal network and the incoming traffic	
from external sources (such as the internet) with an intention to block malicious	

traffic like viruses and hackers, and they allow incoming traffic sent as a response	
to requests from internal hosts [1m].	