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17. Apr -
$$\frac{1}{\sqrt{1}}$$
 the binary is -0.01001

as the machine number is 8 bits long (1 bit for sign)

so the $\begin{bmatrix} \frac{1}{\sqrt{1}} \end{bmatrix}_{TF} = 1.0100010$
 $\begin{bmatrix} -\frac{1}{\sqrt{1}} \end{bmatrix}_{1}^{1} = 1.011101$
 $\begin{bmatrix} \frac{1}{\sqrt{1}} \end{bmatrix}_{1}^{2} = 1.011101$

for $\frac{92}{18}$ the binary is +0.1011101

so the $\begin{bmatrix} \frac{92}{18} \end{bmatrix}_{1}^{2} = 0.101101$
 $\begin{bmatrix} \frac{1}{\sqrt{2}} \end{bmatrix}_{1}^{2} = 0.100010$
 $\begin{bmatrix} \frac{1}{\sqrt{2}} \end{bmatrix}_{1}^{2} = 0.0100011$

for loo the binary is + 1100100

so the $\begin{bmatrix} 100 \end{bmatrix}_{1}^{2} = 0.011010$

for -73 the binary is -100100

so the $\begin{bmatrix} -73 \end{bmatrix}_{1}^{2} = 1.00101$
 $\begin{bmatrix} -73 \end{bmatrix}_{2}^{2} = 1.010110$
 $\begin{bmatrix} -73 \end{bmatrix}_{2}^{2} = 1.010110$
 $\begin{bmatrix} -73 \end{bmatrix}_{2}^{2} = 1.010111$

for - $\frac{7}{7}$ the binary is $\frac{1}{7} \times 2 = \frac{9}{7} \Rightarrow 1$
 $\begin{bmatrix} \frac{9}{7} - 1/2 \times 2 = \frac{9}{7} \Rightarrow 0 \\ \frac{7}{7} \times 2 = \frac{9}{7} \Rightarrow 1 \end{bmatrix}$ (the same as the first step)

so the binary is -0.100100

Thus the $\begin{bmatrix} -\frac{1}{7} \end{bmatrix}_{1}^{2} = 1.0110110$
 $\begin{bmatrix} -\frac{7}{7} \end{bmatrix}_{2}^{2} = 1.0110110$

27. Apr from $[x]_{TF}$ to $[x]_{S}$, we first get the invert, by bit lexcept sign bit) and add I at the end position. Thus, from $[x]_{S}$ to $[x]_{TF}$, we first minus I at the end position and then invert by bit lexcept sign bit) So for $[x_1]_{S} = 1.100$ $[x_1]_{TF} = 1.0100$ $x_1 = -\frac{1}{16}$ $[x_2]_{S} = 1.100$ $[x_3]_{TF} = 1.0111$ $x_2 = -\frac{1}{16}$ $[x_3]_{S} = 0.1110$ $[x_4]_{TF} = 0.110$ $[x_5]_{TF} = 0.110$

 $[X_4]_{5}' = 1.0000$ $[X_4]_{TF}$ no exist. $X_4 = -1$ (no express -1 as a decimal) $[X_4]_{5}' = 1,0101$ $[X_5]_{TF} = 1,1011$ $X_5 = -11$ $[X_6]_{5}' = 1,1101$ $[X_6]_{TF} = 1,0011$ $X_6 = -3$ $[X_7]_{5}' = 0,0111$ $[X_7]_{TF} = 0,0111$ $X_7 = -7$ $[X_8]_{5}' = 1,0000$ $[X_8]_{TF} = 1,10000$ $X_8 = -16$ (5 bits long if it only has 4 bits may produce loss, 1,0000)

3.T解: in for it is represented by IEEE 754 short in computer So there are 32 bits would be in FRI

So the finally code is

the greatest positive appear when the biased representation is 1111111 and the martissa is 0.0000001

the minimum positive appear when the biased representation is 0000001 and the mantissa is 0.1111111

the true form of the biased representation is -111111 the true form of the mantissa is 0.00000001

Thus the minimum positive is $0.00000001 \times 5^{-111111} = 5^{-8} \times 5^{-127}$ So the greatest negative is $-5^{-8} \times 5^{-127}$ [a+b]2' = [a]3' + [b]3'

(1) $[a]_{TF} = 1,01|$. $[b]_{TF} = 1,100$ Thus $[a]_{s} = 1,101$ $[b]_{s} = 1,100$ 90 1,101+1,100 = 101

Thus there is no overflow the result in true value is 1,111=-7

13) $[a]_{TF} = 1.1000 \quad [b]_{TF} = 1.1001$ $So [a]_{S} = 1.1000 \quad [b]_{S} = 1.0111$ Thus $\frac{1.1000}{10.1111} = [a+b]_{S} = 0.1111$

Thus there is an overflow

(x) $[a]_{7}F = 1,0101$. $[b]_{7}F = 0.2010$ Thus $[a]_{5} = 1,1011$ $[b]_{5} = 0,1010$

 $\frac{1,1011}{10,0101} \leftarrow [a+b]; = 0,0101$

Thus there is no overflow the result in true value is 0,0101 = 5

(4) $[a]_{TF} = 11,01000 \quad [b]_{TF} = 11,01001$ $[a]_{\lambda}' = 11,110000 \quad [b]_{\lambda}' = 11,10111$ Thus $\frac{11,10000}{+11,10111} = [a+b]_{\lambda}'$

Thus there is no overflow the result in true value is

1, 10001 = -1