

Computation of Special Functions (Haskell)

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Contents

1	Introduction	3
2	Utility	3
2.1	Preamble	3
2.2	Data Types	3
2.3	Helper functions	5
2.4	Kahan summation	5
2.5	Continued fraction evaluation	6
2.6	TO BE MOVED	6
3	Fibonacci Numbers	6
4	Numbers	7
4.1	Preamble	7
4.2	Stirling numbers	8
5	Exponential & Logarithm	8
5.1	Preamble	8
5.2	Exponential	9
5.2.1	<code>sf_exp x</code>	9
5.2.2	<code>sf_exp_m1 x</code>	9
5.2.3	<code>sf_exp_m1vx x</code>	10
5.2.4	<code>sf_exp_menx n x</code>	10
5.2.5	<code>sf_exp_men n x</code>	11
5.2.6	<code>sf_expn n x</code>	11
5.3	Logarithm	11
5.3.1	<code>sf_log x</code>	11
5.3.2	<code>sf_log_p1 x</code>	12
6	Gamma	12
6.1	Preamble	12
6.2	Misc	13
6.2.1	<code>euler_gamma</code>	13
6.2.2	<code>sf_beta a b</code>	13
6.3	Gamma	13
6.3.1	<code>sf_gamma z</code>	13
6.3.2	<code>*lngamma_asymp z</code>	13
6.3.3	<code>sf_invgamma z</code>	14
6.3.4	<code>sf_lngamma z</code>	14

6.3.5	<code>bernoulli_b n</code>	14
6.4	Digamma	15
6.4.1	<code>sf_digamma z</code>	15
7	Error function	16
7.1	Preamble	16
7.2	Error function	16
7.2.1	<code>sf_erf z</code>	17
7.2.2	<code>sf_erfc z</code>	17
7.3	Dawson's function	19
8	Exponential Integral	19
8.1	Preamble	19
8.2	Exponential integral Ei	19
8.2.1	<code>sf_expint_ei z</code>	19
8.3	Exponential integral E_n	20
8.3.1	<code>sf_expint_en n z</code>	20
9	AGM	22
9.1	Preamble	22
9.2	AGM	22
9.2.1	<code>sf_agm alpha beta</code>	22
9.2.2	<code>sf_agm' alpha beta</code>	22
10	Airy	23
10.1	Preamble	23
10.2	Ai	23
10.2.1	<code>sf_airy_ai z</code>	23
10.3	Bi	24
10.3.1	<code>sf_airy_bi z</code>	24
11	Riemann zeta function	24
11.1	Preamble	24
11.2	Zeta	24
11.2.1	<code>sf_zeta z</code>	24
11.2.2	<code>sf_zeta_m1 z</code>	25
12	Elliptic functions	25
12.1	Preamble	25
12.2	Elliptic integral of the first kind	26
12.2.1	<code>sf_elliptic_k k</code>	26
12.2.2	<code>sf_elliptic_f phi k</code>	26
12.3	Elliptic integral of the second kind	27
12.3.1	<code>sf_elliptic_e k</code>	27
12.3.2	<code>sf_elliptic_e_ic phi k</code>	28
12.4	Elliptic integral of the third kind	28
12.4.1	<code>sf_elliptic_pi c k</code>	28
12.4.2	<code>sf_elliptic_pi_ic phi c k</code>	29
12.5	Elliptic integral of Legendre's type	29
12.5.1	<code>sf_elliptic_d_ic phi k</code>	29
12.5.2	<code>sf_elliptic_d_ic phi k</code>	30
12.6	Burlisch's elliptic integrals	30
12.6.1	<code>sf_elliptic_cel kc p a b</code>	30
12.6.2	<code>sf_elliptic_el1 x kc</code>	30
12.6.3	<code>sf_elliptic_el2 x kc a b</code>	31

12.6.4	<code>sf_elliptic_el3 x kc p</code>	31
12.7	Symmetric elliptic integrals	31
12.7.1	<code>sf_elliptic_rc x y</code>	31
12.7.2	<code>sf_elliptic_rd x y z</code>	32
12.7.3	<code>sf_elliptic_rf x y z</code>	32
12.7.4	<code>sf_elliptic_rg x y z</code>	33
12.7.5	<code>sf_elliptic_rj x y z p</code>	33
13	Spence	34
13.1	Preamble	34
13.2	<code>sf_spence z</code>	35
14	Lommel functions	35
14.1	Preamble	35
14.2	First Lommel function	35
14.2.1	<code>sf_lommel_s mu nu z</code>	36
14.3	Second Lommel function	36
14.3.1	<code>sf_lommel_s2 mu nu z</code>	36

1 Introduction

Special functions.

2 Utility

2.1 Preamble

We start with the basic preamble.

```
{-# Language BangPatterns #-}
{-# Language FlexibleContexts #-}
{-# Language FlexibleInstances #-}
{-# Language ScopedTypeVariables #-}
{-# Language TypeFamilies #-}
-- {-# Language UndecidableSuperClasses #-}
-- {-# Language UndecidableInstances #-}
module Util where
import Data.Complex
```

2.2 Data Types

We start by defining a convenient type synonym for complex numbers over `Double`.

```
type CDouble = Complex Double
```

Next, we define the `Value` typeclass which is useful for defining our special functions to work over both real (`Double`) values and over complex (`CDouble`) values with uniform implementations. This will also make it convenient for handling `Quad` values (later).

```
class Value v
```

Value

```
class (Eq v, Floating v, Fractional v, Num v,
      Enum (RealKind v), Eq (RealKind v), Floating (RealKind v),
      Fractional (RealKind v), Num (RealKind v), Ord (RealKind v),
      Eq (ComplexKind v), Floating (ComplexKind v), Fractional (ComplexKind v),
      Num (ComplexKind v)
) => Value v where
type RealKind v :: *
type ComplexKind v :: *
pos_infty :: v
neg_infty :: v
nan :: v
re :: v -> (RealKind v)
im :: v -> (RealKind v)
rabs :: v -> (RealKind v)
is_inf :: v -> Bool
is_nan :: v -> Bool
is_real :: v -> Bool
fromDouble :: Double -> v
fromReal :: (RealKind v) -> v
toComplex :: v -> (ComplexKind v)
```

Both Double and CDouble are instances of the Value typeclass in the obvious ways.

```
instance Value Double
```

Value Double

```
instance Value Double where
type RealKind Double = Double
type ComplexKind Double = CDouble
pos_infty = 1.0/0.0
neg_infty = -1.0/0.0
nan = 0.0/0.0
re = id
im = const 0
rabs = abs
is_inf = isInfinite
is_nan = isNaN
is_real = True
fromDouble = id
fromReal = id
toComplex x = x :+ 0
```

```
instance Value CDouble
```

Value
CDouble

```
instance Value CDouble where
type RealKind CDouble = Double
type ComplexKind CDouble = CDouble
pos_infty = (1.0/0.0) :+ 0
neg_infty = (-1.0/0.0) :+ 0
nan = (0.0/0.0) :+ 0
re = realPart
```

```
instance Value CDouble (cont)
```

Value
CDouble

```
im = imagPart
rabs = realPart.abs
is_inf z = (is_inf.re$z) ∨ (is_inf.im$z)
is_nan z = (is_nan.re$z) ∨ (is_nan.im$z)
is_real _ = False
fromDouble x = x :+ 0
fromReal x = x :+ 0
toComplex = id
```

TODO: add quad versions also

2.3 Helper functions

A convenient shortcut, as we often find ourselves converting indices (or other integral values) to our computation type.

```
{-# INLINE (#) #-}
(#) :: (Integral a, Num b) ⇒ a → b
(#) = fromIntegral
```

A version of `iterate` which passes along an index also (very useful for computing terms of a power-series, for example.)

```
ixiter i x f
```

ixiter

```
{-# INLINE ixiter #-}
ixiter :: (Enum ix) ⇒ ix → a → (ix → a → a) → [a]
ixiter i x f = x:(ixiter (succ i) (f i x) f)
```

Computes the relative error in terms of decimal digits, handy for testing. Note that this fails when the exact value is zero.

$$\text{relerr } e \ a = \log_{10} \left| \frac{a - e}{e} \right|$$

```
relerr :: ∀ v. (Value v) ⇒ v → v → (RealKind v)
relerr !exact !approx = re $! logBase 10 (abs ((approx-exact)/exact))
```

2.4 Kahan summation

A useful tool is so-called Kahan summation, based on the observation that in floating-point arithmetic, one can ...

Here `kadd t s e k` is a single step of addition, adding a term to a sum+error and passing the updated sum+error to the continuation.

```
— kadd value oldsum olderr —> newsum newerr
{-# INLINE kadd #-}
{-# SPECIALISE kadd :: Double → Double → Double → (Double → Double → a) → a #-}
kadd :: (Value v) ⇒ v → v → v → (v → v → a) → a
kadd t s e k =
  let y = t - e
      s' = s + y
      e' = (s' - s) - y
  in k s' e'
```

Here `ksum terms` sums a list with Kahan summation. The list is assumed to be (eventually) decreasing and the summation is terminated as soon as adding a term doesn't change the value. (Thus any zeros in the list will immediately terminate the sum.) This is typically used for power-series or asymptotic expansions. (TODO: make generic over stopping condition)

`ksum terms`

`ksum`

```
{-# SPECIALISE ksum :: [Double] -> Double #-}
{-# SPECIALISE ksum' :: [Double] -> (Double -> Double -> a) -> a #-}
ksum :: (Value v) => [v] -> v
ksum terms = ksum' terms const

ksum' :: (Value v) => [v] -> (v -> v -> a) -> a
ksum' terms k = f 0 0 terms
  where
    f !s !e [] = k s e
    f !s !e (t:terms) =
      let !y = t - e
          !s' = s + y
          !e' = (s' - s) - y
      in if s' == s
          then k s' e'
          else f s' e' terms
```

2.5 Continued fraction evaluation

This is Steed's algorithm for evaluation of a continued fraction

$$C = b_0 + a_1/(b_1 + a_2/(b_2 + a_3/(b_3 + \dots)))$$

where $C_n = A_n/B_n$ is the partial evaluation up to $\dots a_n/b_n$. Here `steeds as bs` evaluates until $C_n = C_{n+1}$. TODO: describe the algorithm.

```
steeds :: (Value v) => [v] -> [v] -> v
steeds (a1:as) (b0:b1:bs) =
  let !c0 = b0
      !d1 = 1/b1
      !delc1 = a1*d1
      !c1 = c0 + delc1
  in recur c1 delc1 d1 as bs
  where recur !cn_1 !delcn_1 !dn_1 !(an:as) !(bn:bs) =
    let !dn = 1/(dn_1*an+bn)
        !delcn = (bn*dn - 1)*delcn_1
        !cn = cn_1 + delcn
    in if (cn == cn_1) ∨ is_nan cn then cn else (recur cn delcn dn as bs)
```

2.6 TO BE MOVED

```
sf_sqrt :: (Value v) => v -> v
sf_sqrt = sqrt
```

3 Fibonacci Numbers

A silly approach to efficient computation of Fibonacci numbers

$$f_n = f_{n-1} + f_{n-2} \quad f_0 = 0 \quad f_1 = 1$$

The idea is to use the closed-form solution:

$$f_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^n + \frac{-1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)^n$$

and note that we can work in $\mathbb{Q}[\sqrt{5}]$ with terms of the form $a + b\sqrt{5}$ with $a, b \in \mathbb{Q}$ (notice that $\frac{1}{\sqrt{5}} = \frac{\sqrt{5}}{5}$.)

$$\begin{aligned} (a + b\sqrt{5}) + (c + d\sqrt{5}) &= (a + c) + (b + d)\sqrt{5} \\ (a + b\sqrt{5}) * (c + d\sqrt{5}) &= (ac + 5bd) + (ad + bc)\sqrt{5} \end{aligned}$$

We use the `Rational` type to represent elements of \mathbb{Q} , which is a bit more than we actually need, as in the computations above the denominator of $\left(\frac{1 \pm \sqrt{5}}{2}\right)^n$ is always, in fact, 1 or 2.

```
module Fibo (fibonacci) where
import Data.Ratio
data Q5 = Q5 Rational Rational
  deriving (Eq)
```

The number-theoretic norm $N(a + b\sqrt{5}) = a^2 - 5b^2$, though unused in our application.

```
norm (Q5 ra qa) = ra^2 - 5*qa^2
```

Human-friendly `Show` instantiation.

```
instance Show Q5 where
  show (Q5 ra qa) = (show ra)++"++"(show qa)++"*sqrt(5)"
```

Implementation of the operations for typeclasses `Num` and `Fractional`. The `abs` and `signum` functions are unused, so we just give placeholder values.

```
instance Num Q5 where
  (Q5 ra qa) + (Q5 rb qb) = Q5 (ra+rb) (qa+qb)
  (Q5 ra qa) - (Q5 rb qb) = Q5 (ra-rb) (qa-qb)
  (Q5 ra qa) * (Q5 rb qb) = Q5 (ra*rb+5*qa*qb) (ra*qb+rb*qa)
  negate (Q5 ra qa) = Q5 (-ra) (-qa)
  abs a = Q5 (norm a) 0
  signum a@(Q5 ra qa) = if a==0 then 0 else Q5 (ra/(norm a)) (qa/(norm a))
  fromInteger n = Q5 (fromInteger n) 0
```

```
instance Fractional Q5 where
  recip a@(Q5 ra qa) = Q5 (ra/(norm a)) (-qa/(norm a))
  fromRational r = (Q5 r 0)
```

Finally, we define $\phi_{\pm} = \frac{1}{2}(1 \pm \sqrt{5})$ and $c_{\pm} = \pm \frac{1}{5}\sqrt{5}$ so that $f_n = c_+ \phi_+^n + c_- \phi_-^n$. (We can shortcut and extract the value we want without actually computing the full expression.)

```
phip = Q5 (1%2) (1%2)
cp   = Q5 0      (1%5)
phim = Q5 (1%2) (-1%2)
cm   = Q5 0      (-1%5)
fibonacci' n = let (Q5 r q) = cp*phip^n + cm*phim^n in numerator r
fibonacci n = let (Q5 _ q) = phip^n in numerator (2*q)
```

4 Numbers

4.1 Preamble

```
module Numbers where
import Data.Ratio
import qualified Fibo
```

```

fibonacci_number :: Int → Integer
fibonacci_number n = Fibo.fibonacci n

lucas_number :: Int → Integer
lucas_number = undefined

euler_number :: Int → Integer
euler_number = undefined

catalan_number :: Integer → Integer
catalan_number 0 = 1
catalan_number n = 2*(2*n-1)*(catalan_number (n-1)) `div` (n+1)

bernoulli_number :: Int → Rational
bernoulli_number = undefined

tangent_number :: Int → Integer
tangent_number = undefined

triangular_number :: Integer → Integer
triangular_number n = n*(n+1) `div` 2

factorial :: (Integral a) ⇒ a → a
factorial 0 = 1
factorial 1 = 1
factorial n = product [1..n]

binomial :: (Integral a) ⇒ a → a → a
binomial n k
  | k < 0 = 0
  | n < 0 = 0
  | k > n = 0
  | k == 0 = 1
  | k == n = 1
  | k > n `div` 2 = binomial n (n-k)
  | otherwise = (product [n-(k-1)..n]) `div` (product [1..k])

```

4.2 Stirling numbers

— *TODO: this is extremely inefficient approach*

```

stirling_number_first_kind n k = s n k
  where s n k | k ≤ 0 ∨ n ≤ 0 = 0
             | s n 1 = (-1)^(n-1)*(factorial (n-1))
             | s n k = (s (n-1) (k-1)) - (n-1)*(s (n-1) k)

```

— *TODO: this is extremely inefficient approach*

```

stirling_number_second_kind n k = s n k
  where s n k | k ≤ 0 ∨ n ≤ 0 = 0
             | s n 1 = 1
             | s n k = k*(s (n-1) k) + (s (n-1) (k-1))

```

5 Exponential & Logarithm

In this section, we implement the exponential function and logarithm function, as well as useful variations.

5.1 Preamble

We begin with a typical preamble.


```
module Exp
```

```
{-# Language BangPatterns #-}  
{-# Language FlexibleInstances #-}  
module Exp (  
    sf_exp, sf_expn, sf_exp_m1, sf_exp_m1vx, sf_exp_men, sf_exp_menx,  
    sf_log, sf_log_p1,  
) where  
import Numbers  
import Util
```

5.2 Exponential

We start with implementation of the most basic special function, $\exp(x)$ or e^x and variations thereof.

5.2.1 sf_exp x

For the exponential **sf_exp** $x = \exp(x)$ we use a simple series expansions

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{n!}$$

after first using the identity $e^{-x} = 1/e^x$ to ensure that the real part of the argument is positive. This avoids disastrous cancellation for negative arguments, (though note that for complex arguments this is not sufficient.) TODO: should do range-reduction first... TODO: maybe for complex, use explicit cis?

```
sf_exp x = ex
```

```
sf_exp :: (Value v) => v -> v  
sf_exp !x  
| is_inf x  = if (re x)<0 then 0 else pos_infty  
| is_nan x  = x  
| (re x)<0   = 1/(sf_exp (-x))  
| otherwise = ksum $ ixiter 1 1.0 $ \n t -> t*x/(#)n
```

5.2.2 sf_exp_m1 x

For numerical calculations, it is useful to have **sf_exp_m1** $x = e^x - 1$ as explicitly calculating this expression will give poor results for x near 1. We use a series expansion for the calculation. Again for negative real part we reflect using $e^{-x} - 1 = -e^{-x}(e^x - 1)$. TODO: should do range-reduction first... TODO: maybe for complex, use explicit cis?

```
sf_exp_m1 x = ex - 1
```

```
sf_exp_m1 :: (Value v) => v -> v  
sf_exp_m1 !x  
| is_inf x  = if (re x)<0 then -1 else pos_infty  
| is_nan x  = x  
| (re x)<0   = -sf_exp x * sf_exp_m1 (-x)  
| otherwise = ksum $ ixiter 2 x $ \n t -> t*x/(#)n
```

5.2.3 sf_exp_mlvx x

Similarly, it is useful to have the scaled variant $\text{sf_exp_mlvx } x = \frac{e^x - 1}{x}$. In this case, we use a continued-fraction expansion

$$\frac{e^x - 1}{x} = \frac{2}{2 - x + \frac{x^2/6}{1 + \frac{x^2/4 \cdot 3 \cdot 5}{1 + \frac{x^2/4 \cdot 5 \cdot 7}{1 + \frac{x^2/4 \cdot 7 \cdot 9}{1 + \dots}}}}$$

For complex values, simple calculation is inaccurate (when $\Re z \sim 1$).

sf_exp_mlvx x = $\frac{e^x - 1}{x}$

```
sf_exp_mlvx :: (Value v) => v -> v
sf_exp_mlvx !x
| is_inf x = if (re x)<0 then 0 else pos_infty
| is_nan x = x
| rabs(x)>(1/2) = (sf_exp x - 1)/x — inaccurate for some complex points
| otherwise =
  let x2 = x^2
  in 2/(2 - x + x2/6/(1
    + x2/(4*(2*3-3)*(2*3-1))/(1
    + x2/(4*(2*4-3)*(2*4-1))/(1
    + x2/(4*(2*5-3)*(2*5-1))/(1
    + x2/(4*(2*6-3)*(2*6-1))/(1
    + x2/(4*(2*7-3)*(2*7-1))/(1
    + x2/(4*(2*8-3)*(2*8-1))/(1
    ))))));
```

5.2.4 sf_exp_menx n x

Compute the scaled tail of series expansion of the exponential function.

$$\text{sf_exp_menx } n \ x = \frac{n!}{x^n} \left(e^x - \sum_{k=0}^{n-1} \frac{x^k}{k!} \right) = \frac{n!}{x^n} \sum_{k=n}^{\infty} \frac{x^k}{k!} = n! \sum_{k=0}^{\infty} \frac{x^k}{(k+n)!}$$

We use a continued fraction expansion and using the modified Lentz algorithm for evaluation.

sf_exp_menx n z

```
sf_exp_menx :: (Value v) => Int -> v -> v
sf_exp_menx 0 z = sf_exp z
sf_exp_menx 1 z = sf_exp_mlvx z
sf_exp_menx n z
| is_inf z = if (re z)>0 then pos_infty else (0) — TODO: verify
| is_nan z = z
| otherwise = exp_menx__contfrac n z
where
  !zeta = 1e-150
  !eps = 1e-16
  nz !z = if z==0 then zeta else z
  exp_menx__contfrac n z =
    let !fj = (#)$ n+1
    !cj = fj
```

sf_exp_menx n z (cont)

```

!dj = 0
!j = 1
in lentz j dj cj fj
lentz !j !dj !cj !fj =
let !aj = if (odd j)
    then z*((#)$(j+1)'div'2)
    else -z*((#)$(n+(j'div'2)))
bj = (#)$n+1+j
!dj' = nz$ bj + aj*dj
!cj' = nz$ bj + aj/cj
!dji = 1/dj'
!deltaj = cj'*dji
!fj' = fj*deltaj
in if (rabs(deltaj-1)<eps)
then 1/(1-z/fj')
else lentz (j+1) dji cj' fj'

```

5.2.5 sf_exp_men n x

This is the generalization of `sf_exp_m1 x`, giving the tail of the series expansion of the exponential function, for $n = 0, 1, \dots$

$$\text{sf_exp_men } n \ z = e^z - \sum_{k=0}^{n-1} \frac{z^k}{k!} = \sum_{k=n}^{\infty} \frac{z^k}{k!}$$

The special cases are: $n = 0$ gives $e^x = \text{sf_exp } x$ and $n = 1$ gives $e^x - 1 = \text{sf_exp_m1 } x$. We compute this by calling the scaled version `sf_exp_menx` and rescaling back.

```

— ($n=0, 1, 2, \dots$)
sf_exp_men :: (Value v) => Int -> v -> v
sf_exp_men !n !x = (sf_exp_menx n x) * x^n / ((#)$factorial n)

```

5.2.6 sf_expn n x

```

— Compute initial part of series for exponential, $lambda sum_{k=0}^n z^k/k!$
— ($n=0,1,2,\dots$)
sf_expn :: (Value v) => Int -> v -> v
sf_expn n z
| is_inf z = if (re z)>0 then (1/0) else (if (odd n) then (-1/0) else (1/0))
| is_nan z = z
| otherwise = expn__series n z
where
— TODO: just call sf_exp when possible
— TODO: better handle large -ve values!
expn__series :: (Value v) => Int -> v -> v
expn__series n z = ksum $ take (n+1) $ ixiter 1 1.0 $ \k t -> t*z/(#)k

```

5.3 Logarithm

5.3.1 sf_log x

We simply use the built-in implementation (from the `Floating` typeclass).

```

sf_log :: (Value v) => v -> v
sf_log = log

```

5.3.2 sf_log_p1 x

The accuracy preserving `sf_log_p1 x = ln 1 + x`. For values close to zero, we use a power series expansion

$$\ln(1+x) = 2 \sum_{n=0}^{\infty} \frac{\left(\frac{x}{x+2}\right)^{2n+1}}{2n+1}$$

and otherwise just compute it directly.

```
sf_log_p1 z = ln z + 1

sf_log_p1 :: (Value v) => v -> v
sf_log_p1 !z
  | is_nan z = z
  | (rabs z)>0.25 = sf_log (1+z)
  | otherwise = series z
where
  series z =
    let !r = z/(z+2)
        !zr2 = r^2
        !tterms = iterate (*zr2) (r*zr2)
        !terms = zipWith (\n t -> t/((#)s2*n+1)) [1..] tterms
    in 2*(ksum (r:terms))
```

A simple continued fraction implementation for $\ln 1 + z$

$$\ln(1+z) = z/(1 + z/(2 + z/(3 + 4z/(4 + 4z/(5 + 9z/(6 + 9z/(7 + \dots)))))))$$

Though unused for now, it seems to have decent convergence properties.

```
ln_1_z_cf z = steeds (z:(ts 1)) [0..]
  where ts n = (n^2*z):(n^2*z):(ts (n+1))
```

6 Gamma

6.1 Preamble

A basic preamble.

```
module Gamma (
  euler_gamma,
  factorial,
  sf_beta,
  sf_gamma,
  sf_invgamma,
  sf_lngamma,
  sf_digamma,
  bernoulli_b,
)
where
import Exp
import Numbers(factorial)
import Trig
import Util
```

6.2 Misc

6.2.1 euler_gamma

A constant for Euler's gamma:

$$\gamma = \lim_{n \rightarrow \infty} \left(\sum_{k=1}^n \frac{1}{k} - \ln n \right)$$

`euler_gamma :: (Floating a) => a`

`euler_gamma = 0.577215664901532860606512090082402431042159335939923598805767234884867726777664670936947063291746749`

6.2.2 sf_beta a b

The Beta integral

$$B(a, b) = \int_0^1 t^{a-1} (1-t)^{b-1} dt = \frac{\Gamma(a)\Gamma(b)}{\Gamma(a+b)}$$

implemented in terms of log-gamma

$$\text{sf_beta } a \ b = e^{\ln \Gamma(a) + \ln \Gamma(b) - \ln \Gamma(a+b)}$$

`sf_beta :: (Value v) => v -> v -> v`

`sf_beta a b = sf_exp $ (sf_lngamma a) + (sf_lngamma b) - (sf_lngamma$a+b)`

6.3 Gamma

The gamma function

$$\Gamma(z) = \int_0^\infty e^{-t} t^z \frac{dz}{z}$$

6.3.1 sf_gamma z

The gamma function implemented using the identity $\Gamma(z) = \frac{1}{z}\Gamma(z+1)$ to increase the real part of the argument to be > 15 and then using an asymptotic expansion for log-gamma, `lngamma_asymp`, to evaluate.

`sf_gamma x = $\Gamma(x)$`

`sf_gamma`

```
sf_gamma :: (Value v) => v -> v
sf_gamma x =
  redup x 1 $ \ x' t -> t * (sf_exp (lngamma_asymp x'))
  where redup x t k
        | (re x)>15 = k x t
        | otherwise = redup (x+1) (t/x) k
```

6.3.2 *lngamma_asymp z

The asymptotic expansion for log-gamma

$$\ln \Gamma(z) \sim \left(z - \frac{1}{2}\right) \ln z - z + \frac{1}{2} \ln(2\pi) + \sum_{k=1}^{\infty} \frac{B_{2k}}{2k(2k-1)z^{2k-1}}$$

where B_n is the n 'th Bernoulli number.

`lngamma_asymp :: (Value v) => v -> v`

`lngamma_asymp z = (z - 1/2)*(sf_log z) - z + (1/2)*sf_log(2*pi) + (ksum terms)`

`where terms = [b2k/(2*k*(2*k-1)*z^(2*k'-1)) | k' <- [1..10], let k=(#)k', let b2k=bernoulli_b$2*k']`

6.3.3 sf_invgamma z

The inverse gamma function, `sf_invgamma z` = $\frac{1}{\Gamma(z)}$.

```
sf_invgamma :: (Value v) => v -> v
sf_invgamma x =
  let (x', t) = redup x 1
      lngx = lngamma_asymp x'
  in t * (sf_exp$ -lngx)
  where redup x t
        | (re x)>15 = (x, t)
        | otherwise = redup (x+1) (t*x)
```

6.3.4 sf_lngamma z

The log-gamma function, `sf_lngamma z` = $\ln \Gamma(z)$.

```
sf_lngamma :: (Value v) => v -> v
sf_lngamma x =
  let (x', t) = redup x 0
      lngx = lngamma_asymp x'
  in t + lngx
  where redup x t
        | (re x)>15 = (x, t)
        | otherwise = redup (x+1) (t-sf_log x)
```

6.3.5 bernoulli_b n

The Bernoulli numbers, B_n . A simple hard-coded table, for now. (Should be moved to Numbers module and general, cached, implementation done.)

```
bernoulli_b :: (Value v) => Int -> v
bernoulli_b 1 = -1/2
bernoulli_b k | k `mod` 2 == 1 = 0
bernoulli_b 0 = 1
bernoulli_b 2 = 1/6
bernoulli_b 4 = -1/30
bernoulli_b 6 = 1/42
bernoulli_b 8 = -1/30
bernoulli_b 10 = 5/66
bernoulli_b 12 = -691/2730
bernoulli_b 14 = 7/6
bernoulli_b 16 = -3617/510
bernoulli_b 18 = 43867/798
bernoulli_b 20 = -174611/330
bernoulli_b _ = undefined
```

Spouge's approximation to the gamma function

In tests, this gave disappointing results.

— *Spouge's approximation* ($a=17?$)

```
spouge_approx :: (Value v) => Int -> v -> v
spouge_approx a z' =
  let z = z' - 1
      a' = (#)a
      res = (z+a')**((z+(1/2)) * sf_exp (-(z+a')))
      sm = fromDouble$sf_sqrt(2*pi)
      terms = [(spouge_c k a') / (z+k') | k<-[1..(a-1)], let k' = (#)k]
```

```

    sum = sm + ksum terms
  in res*sum
where
  spouse_c k a = ((if k`mod`2==0 then -1 else 1) / ((#) $ factorial (k-1)))
                * (a-((#)k))*((#)k-1/2) * sf_exp(a-((#)k))

spouge :: (Value v) => Int -> v -> v
spouge a' z' =
  let z = z' - 1
      a = fromDouble$(#)a'
      — I don't quite understand why I can't do this:
      — q = fromReal $ (sf_sqrt(2*pi) :: (RealKind v))
      q = sf_sqrt(2*pi)
  in (z+a)**(z+1/2)*(sf_exp(-z-a))*(q + ksum (map (\k->(c a k)/(z+(#)k)) [1..(a'-1)]))
  where
    c :: (Value v) => v -> Int -> v
    c a k = let k' = (#)k
              sgn = if even k then -1 else 1
              in sgn*(a-k')**((k'-1/2)*(sf_exp(a-k')) / ((#)$factorial(k-1)))

```

6.4 Digamma

The digamma function

$$\psi(z) = \frac{d}{dz} \ln \Gamma(z) = \frac{\Gamma'(z)}{\Gamma(z)}$$

6.4.1 sf_digamma z

We implement with a series expansion for $|z| \leq 10$ and otherwise with an asymptotic expansion.

```

sf_digamma :: (Value v) => v -> v
—sf_digamma n | is_nonposint n = Inf
sf_digamma z | (rabs z)>10 = digamma_asympt z
              | otherwise = digamma_series z

```

The series expansion is the following

$$\psi(z) = -\gamma - \frac{1}{z} + \sum_{k=1}^{\infty} \frac{z}{k(k+z)}$$

but with Euler-Maclaurin correction terms:

$$\psi(z) = -\gamma - \frac{1}{z} + \sum_{k=1}^n \frac{z}{k(k+z)} + \left(\ln \frac{k+z}{k} - \frac{z}{2k(k+z)} + \sum_{j=1}^p B_{2j}(k^{-2j} - (k+z)^{-2j}) \right)$$

```

digamma_series :: (Value v) => v -> v
digamma_series z =
  let res = -euler_gamma - (1/z)
      terms = map (\k->z/((#)k*(z+(#)k))) [1..]
      corrs = map (correction.(#)) [1..]
  in summer res res terms corrs
where
  summer :: (Value v) => v -> v -> [v] -> [v] -> v
  summer res sum (t:terms) (c:corrs) =
    let sum' = sum + t
        res' = sum' + c
    in if res==res' then res
       else summer res' sum' terms corrs

```

```

bn1 = bernoulli_b 2
bn2 = bernoulli_b 4
bn3 = bernoulli_b 6
bn4 = bernoulli_b 8
correction k =
  (sf_log$(k+z)/k) - z/2/(k*(k+z))
  + bn1*(k-2 - (k+z)-2)
  + bn2*(k-4 - (k+z)-4)
  + bn3*(k-6 - (k+z)-6)
  + bn4*(k-8 - (k+z)-8)

```

The asymptotic expansion (valid for $|argz| < \pi$) is the following

$$\psi(z) \sim \ln z - \frac{1}{2z} + \sum_{k=1}^{\infty} \frac{B_{2k}}{2kz^{2k}}$$

Note that our implementation will fail if the `bernoulli_b` table is exceeded. If $\Re z < \frac{1}{2}$ then we use the reflection identity to ensure $\Re z \geq \frac{1}{2}$:

$$\psi(z) - \psi(1-z) = \frac{-\pi}{\tan(\pi z)}$$

```

digamma_asympt :: (Value v) => v -> v
digamma_asympt z
  | (re z)<0.5 = compute (1 - z) $ -pi/(sf_tan(pi*z)) + (sf_log(1-z)) - 1/(2*(1-z))
  | otherwise = compute z $ (sf_log z) - 1/(2*z)
  where
    compute z res =
      let z_2 = z-2
          zs = iterate (*z_2) z_2
          terms = zipWith (\n z2n -> z2n*(bernoulli_b(2*n+2))/(#)(2*n+2)) [0..] zs
      in sumit res res terms
    sumit res ot (t:terms) =
      let res' = res - t
      in if res==res' ∨ (rabs t)>(rabs ot)
         then res
         else sumit res' t terms

```

7 Error function

7.1 Preamble

```

{-# Language BangPatterns #-}
{-# Language BlockArguments #-}
module Erf (
  sf_erf,
  sf_erfc,
) where
import Exp
import Util

```

7.2 Error function

The error function is defined via

$$\operatorname{erf}(z) = \frac{2}{\sqrt{\pi}} \int_0^z e^{-x^2} dx \quad \operatorname{erf}(z)$$

and the complementary error function via

$$\operatorname{erfc}(z) = \frac{2}{\sqrt{\pi}} \int_z^\infty e^{-x^2} dx \quad \operatorname{erfc}(z)$$

Thus we have the relation $\operatorname{erf}(z) + \operatorname{erfc}(z) = 1$.

7.2.1 `sf_erf z`

The error function `sf_erf z` = $\operatorname{erf} z$ where

$$\operatorname{erf}(z) = \frac{2}{\sqrt{\pi}} \int_{-\infty}^z e^{-x^2} dx$$

For $\Re z < -1$, we transform via $\operatorname{erf}(z) = -\operatorname{erf}(-z)$ and for $|z| < 1$ we use the power-series expansion, otherwise we use $\operatorname{erf} z = 1 - \operatorname{erfc} z$. (TODO: this implementation is not perfect, but workable for now.)

`sf_erf z` = $\operatorname{erf}(z)$

```
sf_erf :: (Value v) => v -> v
sf_erf z
  | (re z)<(-1) = -sf_erf(-z)
  | (rabs z)<1  = erf_series z
  | otherwise  = 1 - sf_erfc z
```

7.2.2 `sf_erfc z`

The complementary error-function `sf_erfc z` = $\operatorname{erfc} z$ where

$$\operatorname{erfc} z = 1 - \operatorname{erf} z = \frac{2}{\sqrt{\pi}} \int_z^\infty e^{-x^2} dx$$

For $\Re z < -1$ we transform via $\operatorname{erfc} z = 2 - \operatorname{erf}(-z)$ and if $|z| < 1$ then we use $\operatorname{erfc} z = 1 - \operatorname{erf} z$. Finally, if $|z| < 10$ we use a continued-fraction expansion and an asymptotic expansion otherwise. (TODO: there are a few issues with this implementation: For pure imaginary values and for extremely large values it seems to hang.)

`sf_erfc z` = $\operatorname{erfc}(z)$

```
sf_erfc :: (Value v) => v -> v
sf_erfc z
  | (re z)<(-1) = 2-(sf_erfc (-z))
  | (rabs z)<1  = 1-(sf_erf z)
  | (rabs z)<10 = erfc_cf_pos1 z
  | otherwise  = erfc_asymp_pos z — TODO: hangs for very large input
```

`erf_series z`

The series expansion for $\operatorname{erf} z$:

$$\operatorname{erf} z = \frac{2}{\sqrt{\pi}} \sum_{n=0}^{\infty} \frac{(-)^n z^{2n+1}}{n!(2n+1)}$$

There is an alternative expansion $\operatorname{erf} z = \frac{2}{\sqrt{\pi}} e^{-z^2} \sum_{n=0}^{\infty} \frac{2^n z^{2n+1}}{1 \cdot 3 \cdots (2n+1)}$, but we don't use it here. (TODO: why not?)

```

erf_series z =
  let z2 = z^2
      rts = ixiter 1 z $ \n t → (-t)*z2/(#)n
      terms = zipWith (\n t → t/(#)(2*n+1)) [0..] rts
  in (2/sf_sqrt pi) * ksum terms

```

***sf_erf z**

This asymptotic expansion for $\operatorname{erfc} z$ is valid as $z \rightarrow +\infty$:

$$\operatorname{erfc} z \sim \frac{e^{-z^2}}{\sqrt{\pi}} \sum_{n=0}^{\infty} (-)^n \frac{(1/2)_m}{z^{2m+1}}$$

where the Pochhammer symbol $(1/2)_m$ is given by:

$$\left(\frac{1}{2}\right)_m = \frac{1 \cdot 3 \cdot 5 \cdots (2m-1)}{2^m} = \frac{(2m)!}{m!2^{2m}}$$

TODO: correct the asymptotic term checking (not smallest but pre-smallest term).

```

erfc_asymp_pos z =
  let z2 = z^2
      iz2 = 1/2/z2
      terms = ixiter 1 (1/z) $ \n t → (-t*iz2)*(#)(2*n-1)
      tterms = tk terms
  in (sf_exp (-z2))/(sqrt pi) * ksum tterms
  where tk (a:b:cs) = if (rabs a)<(rabs b) then [a] else a:(tk$b:cs)

```

***erfc_cf_pos1 z**

A continued-fraction expansion for $\operatorname{erfc} z$:

$$\sqrt{\pi}e^{z^2} \operatorname{erfc} z = \frac{z}{z^2 + 1} \frac{1/2}{1 + z^2} \frac{1}{z^2 + 1} \frac{3/2}{1 + z^2} \cdots$$

```

erfc_cf_pos1 z =
  let z2 = z^2
      as = z:(map fromDouble [1/2,1..])
      bs = 0:cycle [z2,1]
      cf = steeds as bs
  in sf_exp(-z2) / (sqrt pi) * cf

```

***erfc_cf_pos1 z**

This is an alternative continued-fraction expansion.

$$\sqrt{\pi}e^{z^2} \operatorname{erfc} z = \frac{2z}{2z^2 + 1} \frac{1 \cdot 2}{2z^2 + 5} \frac{3 \cdot 4}{2z^2 + 9} \cdots$$

Unused for now.

```

erfc_cf_pos2 z =
  let z2 = z^2
      as = (2*z):(map (\n→(#)$ -(2*n+1)*(2*n+2)) [0..])
      bs = 0:(map (\n→2*z2+(#)4*n+1) [0..])
      cf = steeds as bs
  in sf_exp(-z2) / (sqrt pi) * cf

```

7.3 Dawson’s function

Dawson’s function (or Dawson’s integral) is given by

$$D(z) = e^{-z^2} \int_0^z e^{t^2} dt = -\frac{i\sqrt{\pi}}{2} e^{-x^2} \operatorname{erf}(ix)$$

8 Exponential Integral

8.1 Preamble

```
module ExpInt(
  sf_expint_ei,
  sf_expint_en,
)
where
import Exp
import Gamma
import Util
```

8.2 Exponential integral Ei

The exponential integral $\operatorname{Ei} z$ is defined for $x < 0$ by

$$\operatorname{Ei}(z) = - \int_{-x}^{\infty} \frac{e^{-t}}{t} dt$$

It can be defined

8.2.1 sf_expint_ei z

We give only an implementation for $\Re z \geq 0$. We use a series expansion for $|z| < 40$ and an asymptotic expansion otherwise.

```
sf_expint_ei z = Ei(z)
```

sf_expint_ei

```
sf_expint_ei :: (Value v) => v -> v
sf_expint_ei z
  | (re z) < 0.0  = (0/0)  — (NaN)
  | z == 0.0     = (-1/0) — (-Inf)
  | (rabs z) < 40 = expint_ei_series z
  | otherwise    = expint_ei_asymp z
```

The series expansion is given (for $x > 0$)

$$\operatorname{Ei}(x) = \gamma + \ln x + \sum_{n=1}^{\infty} \frac{x^n}{n!n}$$

We evaluate the addition of the two terms with the sum slightly differently when $\Re z < 1/2$ to reduce floating-point cancellation error slightly.

```
expint_ei__series z
```

```
expint_ei__series :: (Value v) => v -> v
expint_ei__series z =
  let tterms = ixiter 2 z $ \n t -> t*z/(#)n
      terms = zipWith (\t n -> t/(#)n) tterms [1..]
      res = ksum terms
  in if (re z)<0.5
      then sf_log(z * sf_exp(euler_gamma + res))
      else res + sf_log(z) + euler_gamma
```

```
expint_ei__se
```

The asymptotic expansion as $x \rightarrow +\infty$ is

$$\text{Ei}(x) \sim \frac{e^x}{x} \sum_{n=0}^{\infty} \frac{n!}{x^n}$$

```
expint_ei__asymp z
```

```
expint_ei__asymp :: (Value v) => v -> v
expint_ei__asymp z =
  let terms = tk $ ixiter 1 1.0 $ \n t -> t/z*(#)n
      res = ksum terms
  in res * (sf_exp z) / z
  where tk (a:b:cs) = if (rabs a)<(rabs b) then [a] else a:(tk$b:cs)
```

```
expint_ei__as
```

8.3 Exponential integral E_n

The exponential integrals $E_n(z)$ are defined as

$$E_n(z) = z^{n-1} \int_z^{\infty} \frac{e^{-t}}{t^n} dt$$

They satisfy the following relations:

$$\begin{aligned} E_0(z) &= \frac{e^{-z}}{z} \\ E_{n+1}(z) &= \int_z^{\infty} E_n(t) dt \end{aligned}$$

And they can be expressed in terms of incomplete gamma functions:

$$E_n(z) = z^{n-1} \Gamma(1-n, z)$$

(which also gives a generalization for non-integer n).

8.3.1 sf_expint_en n z

```
sf_expint_en n z = En(z)
```

```
sf_expint_en
```

```
sf_expint_en :: (Value v) => Int -> v -> v
sf_expint_en n z | (re z)<0 = (0/0) — (NaN) TODO: confirm this
                  | z == 0   = (1/(#)(n-1)) — TODO: confirm this
sf_expint_en 0 z = sf_exp(-z) / z
sf_expint_en 1 z = expint_en_1 z
sf_expint_en n z | (rabs z) ≤ 1.0 = expint_en_series n z
                  | otherwise = expint_en_contfrac n z
```

We use this series expansion for $E_1(z)$:

$$E_1(z) = -\gamma - \ln z + \sum_{k=1}^{\infty} (-)^k \frac{z^k}{k!k}$$

(Note that this will not be good for large values of z .)

```
expint_en_1 :: (Value v) => v -> v
```

```
expint_en_1 z =
```

```
  let r0 = -euler_gamma - (sf_log z)
      tterms = ixiter 2 (z) $ \k t -> -t*z/(#)k
      terms = zipWith (\t k -> t/(#)k) tterms [1..]
  in ksum (r0:terms)
```

— *assume $n \geq 2$, $z \leq 1$*

```
expint_en_series :: (Value v) => Int -> v -> v
```

```
expint_en_series n z =
```

```
  let n' = (#)n
      res = -(sf_log z) + (sf_digamma n') * (-z)^(n-1)/(#)(factorial$ n-1) + 1/(n'-1)
      terms' = ixiter 2 (-z) (\m t -> -t*z/(#)m)
      terms = map (\(m,t) -> (-t)/(#)(m-(n-1))) $ filter ((/= (n-1)) o fst) $ zip [1..] terms'
  in ksum (res:terms)
```

— *assume $n \geq 2$, $z > 1$*

— *modified Lentz algorithm*

```
expint_en_contfrac :: (Value v) => Int -> v -> v
```

```
expint_en_contfrac n z =
```

```
  let fj = zeta
      cj = fj
      dj = 0
      j = 1
      n' = (#)n
  in lentz j cj dj fj
  where
    zeta = 1e-100
    eps = 5e-16
    nz x = if x==0 then zeta else x
    lentz j cj dj fj =
      let aj = (#) $ if j==1 then 1 else -(j-1)*(n+j-2)
          bj = z + (#)(n + 2*(j-1))
          dj' = nz $ bj + aj*dj
          cj' = nz $ bj + aj/cj
          dji = 1/dj'
          delta = cj'*dji
          fj' = fj*delta
      in if (rabs$delta-1)<eps
         then fj' * sf_exp(-z)
         else lentz (j+1) cj' dji fj'
```

9 AGM

9.1 Preamble

```
module AGM
```

```
  module AGM (
    sf_agm,
    sf_agm',
  )
  where
  import Util
```

9.2 AGM

Gauss' arithmetic-geometric mean or AGM of two numbers is defined as the limit $\text{agm}(\alpha, \beta) = \lim_n \alpha_n = \lim_n \beta_n$ where we define

$$\begin{aligned}\alpha_{n+1} &= \frac{\alpha_n + \beta_n}{2} \\ \beta_{n+1} &= \sqrt{\alpha_n \cdot \beta_n}\end{aligned}$$

(Note that we need real values to be positive for this to make sense.)

9.2.1 sf_agm alpha beta

Here we compute the AGM via the definition and return the full arrays of intermediate values $([\alpha_n], [\beta_n], [\gamma_n])$, where $\gamma_n = \frac{\alpha_n - \beta_n}{2}$. (The iteration converges quadratically so this is an efficient approach.)

```
sf_agm alpha beta = agm( $\alpha$ ,  $\beta$ )
```

sf_agm

```
sf_agm :: (Value v) => v -> v -> ([v], [v], [v])
sf_agm alpha beta = agm [alpha] [beta] [alpha-beta]
  where agm as@(a:_) bs@(b:_) cs@(c:_) =
    if c==0 then (as, bs, cs)
    else let a' = (a+b)/2
          b' = sf_sqrt (a*b)
          c' = (a-b)/2
          in if c'==c then (as, bs, cs)
             else agm (a':as) (b':bs) (c':cs)
```

9.2.2 sf_agm' alpha beta

Here we return simply the value $\text{sf_agm}' \ a \ b = \text{agm}(a, b)$.

```
sf_agm' z = agm z
```

```
sf_agm' :: (Value v) => v -> v -> v
sf_agm' alpha beta = agm alpha beta ((alpha-beta)/2)
  —let (as, -, -) = sf_agm alpha beta in head as
  where agm a b 0 = a
```

```
sf_agm' z = agm z (cont)
```

```

    agm a b c =
      let a' = (a+b)/2
          b' = sf_sqrt (a*b)
          c' = (a-b)/2
      in agm a' b' c'

```

```

sf_agm.c0 :: (Value v) => v -> v -> v -> ([v],[v],[v])
sf_agm.c0 alpha beta c0 = undefined

```

10 Airy

The Airy functions Ai and Bi , standard solutions of the ode $y'' - zy = 0$.

10.1 Preamble

A basic preamble.

```

module Airy (sf_airy_ai, sf_airy_bi) where
import Gamma
import Util

```

10.2 Ai

10.2.1 sf_airy_ai z

For now, just use a simple series expansion.

```

sf_airy_ai :: (Value v) => v -> v
sf_airy_ai z = airy_ai_series z

```

Initial conditions $Ai(0) = 3^{-2/3} \frac{1}{\Gamma(2/3)}$ and $Ai'(0) = -3^{-1/3} \frac{1}{\Gamma(1/3)}$

```

ai0 :: (Value v) => v
ai0 = 3**(-2/3)/sf_gamma(2/3)

```

```

ai'0 :: (Value v) => v
ai'0 = -3**(-1/3)/sf_gamma(1/3)

```

Series expansion, where $n!!! = \max(n, 1)$ for $n \leq 2$ and otherwise $n!!! = n \cdot (n-3)!!!$:

$$Ai(z) = Ai(0) \left(\sum_{n=0}^{\infty} \frac{(3n-2)!!!}{(3n)!} z^{3n} \right) + Ai'(0) \left(\frac{(3n-1)!!!}{(3n+1)!} z^{3n+1} \right)$$

```

airy_ai_series z =
  let z3 = z^3
      aiterms = ixiter 0 1 $ \n t -> t*z3*((#)$3*n+1)/((#)$3*n+1)*(3*n+2)*(3*n+3))
      ai'terms = ixiter 0 z $ \n t -> t*z3*((#)$3*n+2)/((#)$3*n+2)*(3*n+3)*(3*n+4))
  in ai0 * (ksum aiterms) + ai'0 * (ksum ai'terms)

```

10.3 Bi

10.3.1 sf_airy_bi z

For now, just use a simple series expansion.

```
sf_airy_bi :: (Value v) => v -> v
sf_airy_bi z = airy_bi_series z
```

Initial conditions $\text{Bi}(0) = 3^{-1/6} \frac{1}{\Gamma(2/3)}$ and $\text{Bi}'(0) = 3^{1/6} \frac{1}{\Gamma(1/3)}$

```
bi0 :: (Value v) => v
bi0 = 3**(-1/6)/sf_gamma(2/3)
```

```
bi'0 :: (Value v) => v
bi'0 = 3**(1/6)/sf_gamma(1/3)
```

Series expansion, where $n!!! = \max(n, 1)$ for $n \leq 2$ and otherwise $n!!! = n \cdot (n - 3)!!!$:

$$\text{Bi}(z) = \text{Bi}(0) \left(\sum_{n=0}^{\infty} \frac{(3n-2)!!!}{(3n)!} z^{3n} \right) + \text{Bi}'(0) \left(\frac{(3n-1)!!!}{(3n+1)!} z^{3n+1} \right)$$

```
airy_bi_series z =
  let z3 = z^3
      biterms = ixiter 0 1 $ \n t -> t*z3*((#)$3*n+1)/((#)$ (3*n+1)*(3*n+2)*(3*n+3))
      bi'terms = ixiter 0 z $ \n t -> t*z3*((#)$3*n+2)/((#)$ (3*n+2)*(3*n+3)*(3*n+4))
  in bi0 * (ksum biterms) + bi'0 * (ksum bi'terms)
```

11 Riemann zeta function

11.1 Preamble

```
{-# Language BangPatterns #-}
module Zeta (sf_zeta, sf_zeta_m1) where
import Gamma
import Trig
import Util
```

11.2 Zeta

The Riemann zeta function is defined by power series for $\Re z > 1$

$$\zeta(z) = \sum_{n=1}^{\infty} n^{-z}$$

and defined by analytic continuation elsewhere.

11.2.1 sf_zeta z

Compute the Riemann zeta function $\text{sf_zeta } z = \zeta(z)$ where

```
sf_zeta z = ζ(z)
```

```
sf_zeta :: (Value v) => v -> v
sf_zeta z
  | z==1      = (1/0)
  | (re z)<0  = 2 * (2*pi)**(z-1) * (sf_sin$pi*z/2) * (sf_gamma$1-z) * (sf_zeta$1-z)
  | otherwise = zeta_series 1.0 z
```


11.2.2 sf.zeta.m1 z

For numerical purposes, it is useful to have `sf_zeta_m1 z = $\zeta(z) - 1$` .

```
sf_zeta_m1 z =  $\zeta(z) - 1$ 
```

```
sf_zeta_m1 :: (Value v) => v -> v
sf_zeta_m1 z
  | z==1      = (1/0)
  | (re z)<0  = 2 * (2*pi)**(z-1) * (sf_sin$pi*z/2) * (sf_gamma$1-z) * (sf_zeta$1-z) - 1
  — TODO:
  | otherwise = zeta_series 0.0 z
```

*zeta_series i z

We use the simple series expansion for $\zeta(z)$ with an Euler-Maclaurin correction:

$$\zeta(z) = \sum_{n=1}^N \frac{1}{n^z} + \sum_{k=1}^p \dots$$

```
zeta_series init z =
```

```
zeta_series :: (Value v) => v -> v -> v
zeta_series !init !z =
  let terms = map (\n->((#)n)**(-z)) [2..]
      corrs = map correction [2..]
  in summer terms corrs init 0.0 0.0
where
  —TODO: convert to use kahan summer
  summer !(t:ts) !(c:cs) !s !e !r =
    let !y = t + e
        !s' = s + y
        !e' = (s - s') + y
        !r' = s' + c + e'
    in if r==r' then r'
       else summer ts cs s' e' r'
  !zz1 = z/12
  !zz2 = z*(z+1)*(z+2)/720
  !zz3 = z*(z+1)*(z+2)*(z+3)*(z+4)/30240
  !zz4 = z*(z+1)*(z+2)*(z+3)*(z+4)*(z+5)*(z+6)/1209600
  !zz5 = z*(z+1)*(z+2)*(z+3)*(z+4)*(z+5)*(z+6)*(z+7)*(z+8)/239500800
  correction !n' =
    let n=(#)n'
    in n**(1-z)/(z-1) - n**(-z)/2
      + n**(-z-1)*zz1 - n**(-z-3)*zz2 + n**(-z-5)*zz3
      - n**(-z-7)*zz4 + n**(-z-9)*zz5
```

12 Elliptic functions

12.1 Preamble

```

{-# Language BangPatterns #-}
module Elliptic where
import AGM
import Exp
import Trig
import Util

    2-2/3

two23 :: Double
!two23 = 0.62996052494743658238

```

12.2 Elliptic integral of the first kind

Assume that $1 - \sin^2 \phi, 1 - k^2 \sin^2 \phi \in \mathbb{C} \setminus (-\infty, 0]$ except that one of them may be 0.

The elliptic integral of the first kind is defined by

$$F(\phi, k) = \int_0^\phi \frac{d\theta}{\sqrt{1 - k^2 \sin^2 \theta}} = \int_0^{\sin \phi} \frac{dt}{\sqrt{1 - t^2} \sqrt{1 - k^2 t^2}}$$

The complete integral is given by $\phi = \pi/2$:

$$K(k) = F(\pi/2, k) =$$

12.2.1 sf_elliptic_k k

Compute the complete elliptic integral of the first kind $K(k)$ To evaluate this, we use the AGM relation

$$K(k) = \frac{\pi}{2 \operatorname{agm}(1, k')} \quad \text{where } k' = \sqrt{1 - k^2} \quad K(k)$$

TODO: UNTESTED!

```

sf_elliptic_k k = K(k)

sf_elliptic_k :: Double → Double
sf_elliptic_k k =
    let an = sf_agm' 1.0 (sf_sqrt $ 1.0 - k^2)
    in pi / (2 * an)

```

12.2.2 sf_elliptic_f phi k

Compute the (incomplete) elliptic integral of the first kind $F(\phi, k)$. To evaluate, we use an ascending Landen transformation:

$$F(\phi, k) = \frac{2}{1 + k} F(\phi_2, k_2) \quad \text{where } k_2 = \frac{2\sqrt{k}}{1 + k} \text{ and } 2\phi_2 = \phi + \arcsin(k \sin \phi) \quad F(\phi, k)$$

Note that $0 < k < 1$ and $0 < \phi \leq \pi/2$ imply $k < k_2 < 1$ and $\phi_2 < \phi$. We iterate this transformation until we reach $k = 1$ and use the special case

$$F(\phi, 1) = \operatorname{gud}^{-1}(\phi)$$

(Where $\operatorname{gud}^{-1}(\phi)$ is the inverse Gudermannian function (TODO)). TODO: UNTESTED!

```

sf_elliptic.f phi k = F(φ, k)

sf_elliptic.f :: Double → Double → Double
sf_elliptic.f phi k
  | k==0 = phi
  | k==1 = sf_log((1 + (sf_sin phi)) / (1 - (sf_sin phi))) / 2
             — quad(@(t)(1/sqrt(1-k^2*sin(t)^2)), 0, phi)
  | phi==0 = 0
  | otherwise =
    ascending_landen phi k 1 $ λ phi' res' →
      res' * sf_log((1 + (sf_sin phi)) / (1 - (sf_sin phi))) / 2
  where
    ascending_landen phi k res kont =
      let k' = 2 * (sf_sqrt k) / (1 + k)
          phi' = (phi + (asin (k*(sin phi))))/2
          res' = res * 2/(1+k)
      in if k'==1 then kont phi' res
         else ascending_landen phi' k' res' kont
    —function res = agm_method(phi, k)
    — [an,bn,cn,phin] = sf_agm(1.0, sqrt(1 - k^2), phi, k);
    — res = phin(end) / (2^(length(phin)-1) * an(end));
    —endfunction

```

12.3 Elliptic integral of the second kind

Assume that $1 - \sin^2 \phi, 1 - k^2 \sin^2 \phi \in \mathbb{C} \setminus (-\infty, 0]$ except that one of them may be 0.

Legendre's (incomplete) elliptic integral of the second kind is defined via

$$E(\phi, k) = \int_0^\phi \sqrt{1 - k^2 \sin^2 \theta} d\theta = \int_0^{\sin \phi} \frac{\sqrt{1 - k^2 t^2}}{\sqrt{1 - t^2}} dt$$

The complete integral is

$$E(k) = E(\pi/2, k) =$$

12.3.1 sf_elliptic_e k

Compute the complete elliptic integral of the second kind $E(k)$. We evaluate this with an agm-based approach:

...

TODO: UNTESTED!

```

sf_elliptic.e k = E(k)

sf_elliptic.e :: Double → Double
sf_elliptic.e k =
  let phi = k
      (as,bs,cs') = sf_agm 1.0 (sf_sqrt (1.0 - k^2))
      cs = k:(tail.reverse$cs')
      res = foldl (-) 2 (map (λ(i,c)→2^(i-1)*c^2) (zip [1..] cs))
  in res * pi/(4*(head as))

```

12.3.2 sf_elliptic_e_ic phi k

Compute the incomplete elliptic integral of the second kind $E(\phi, k)$. We evaluate this with an ascending Landen transformation:

...

TODO: UNTESTED! (Note: could try direct quadrature of the integral, also there is an AGM-based method).

sf_elliptic_e_ic phi k = $E(\phi, k)$

```
sf_elliptic_e_ic :: Double → Double → Double
sf_elliptic_e_ic phi k
  | k==1 = sf_sin phi
  | k==0 = phi
  | otherwise = ascending_landen phi k
where
  ascending_landen phi 1 = sin phi
  ascending_landen phi k =
    let !k' = 2*(sf_sqrt k) / (k+1)
        !phi' = (phi + (sf_asin (k*(sf_sin phi))))/2
    in (1+k)*(ascending_landen phi' k') + (1-k)*(sf_elliptic_f phi' k') - k*(sf_sin phi)
```

12.4 Elliptic integral of the third kind

We define Legendre's (incomplete) elliptic integral of the third kind via

$$\Pi(\phi, \alpha^2, k) = \int_0^\phi \frac{d\theta}{\sqrt{1-k^2 \sin^2 \theta} (1-\alpha^2 \sin^2 \theta)} = \int_0^{\sin \phi} \frac{dt}{\sqrt{1-t^2} \sqrt{1-k^2 t^2} (1-\alpha^2 t^2)}$$

The complete integral of the third kind is given by

$$\Pi(\alpha^2, k) = \Pi(\pi/2, \alpha^2, k) =$$

12.4.1 sf_elliptic_pi c k

Compute the complete elliptic integral of the third kind ($c = \alpha^2$ in DLMF notation) for real values only $0 < k < 1, 0 < c < 1$. Uses agm-based approach. (Could also try numerical quadrature `quad(@(t) (1.0/(1-c*sf_sin(t)^2)/sqrt(1-t^2))`).
 TODO: mostly untested

sf_elliptic_pi c k = $\Pi(c, k)$

```
sf_elliptic_pi :: Double → Double → Double
sf_elliptic_pi c k = complete_agm k c
where
  — -λinfy < k^2 < 1
  — -λinfy < c < 1
  complete_agm k c =
    let (ans, gns, _) = sf_agm 1 (sf_sqrt (1.0-k^2))
        pn1 = sf_sqrt (1-c)
        qn1 = 1
        an1 = last ans
        gn1 = last gns
    in en1 = (pn1^2 - an1*gn1) / (pn1^2 + an1*gn1)
```

```
sf_elliptic_pi c k =  $\Pi(c, k)$  (cont)
```

```
in iter pn1 en1 (reverse ans) (reverse gns) [qn1]

iter pnml enml [an] [gn] qns = pi/(4*an) * (2 + c/(1-c)*(ksum qns))
iter pnml enml (anml:an:ans) (gnml:gn:gns) (qnml:qns) =
  let pn = (pnml^2 + anml*gnml)/(2*pnml)
      en = (pn^2 - an*gn) / (pn^2 + an*gn)
      qn = qnml * enml/2
  in iter pn en (an:ans) (gn:gns) (qn:qnml:qns)
```

12.4.2 sf_elliptic_pi_ic phi c k

```
sf_elliptic_pi_ic phi c k =  $\Pi(\phi, c, k)$ 
```

```
sf_elliptic_pi_ic :: Double → Double → Double → Double
sf_elliptic_pi_ic 0 c k = 0.0
sf_elliptic_pi_ic phi c k = gauss_transform k c phi
where
  gauss_transform k c phi =
    if (sf_sqrt (1-k^2))==1
    then let cp=sf_sqrt(1-c)
          in sf_atan(cp*(sf_tan phi)) / cp
    else if (1-k^2/c)==0 — special case else rho below is zero...
    then ((sf_elliptic_e_ic phi k) - c*(sf_cos phi)*(sf_sin phi)
          / sqrt(1-c*(sf_sin phi)^2))/(1-c)
    else let kp = sf_sqrt (1-k^2)
          k' = (1 - kp) / (1 + kp)
          delta = sf_sqrt(1-k^2*(sf_sin phi)^2)
          psi' = sf_asin((1+kp)*(sf_sin phi) / (1+delta))
          rho = sf_sqrt(1 - (k^2/c))
          c' = c*(1+rho)^2/(1+kp)^2
          xi = (sf_csc phi)^2
          newgt = gauss_transform k' c' psi'
    in (4/(1+kp)*newgt + (rho-1)*(sf_elliptic_f phi k)
        - (sf_elliptic_rc (xi-1) (xi-c)))/rho
```

12.5 Elliptic integral of Legendre's type

The (incomplete) elliptic integral of Legendre's type is defined by

$$D(\phi, k) = \int_0^\phi \frac{\sin^2 \theta}{\sqrt{1 - k^2 \sin^2 \theta}} d\theta = \int_0^{\sin \phi} \frac{t^2}{\sqrt{1 - t^2} \sqrt{1 - k^2 t^2}} dt$$

This can be expressed as $D(\phi, k) = (F(\phi, k) - E(\phi, k))/k^2$.

The complete elliptic integral of Legendre's type is

$$D(k) = D(\pi/2, k) = (K(k) - E(k))/k^2$$

12.5.1 sf_elliptic_d_ic phi k

We simply reduce to $F(\phi, k)$ and $E(\phi, k)$.

```
sf_elliptic_d_ic phi k = D( $\phi, k$ )
```

```
sf_elliptic_d_ic :: Double → Double → Double
sf_elliptic_d_ic phi k = ((sf_elliptic_f phi k) - (sf_elliptic_e_ic phi k)) / (k^2)
```

12.5.2 sf_elliptic_d_ic phi k

We simply reduce to $K(k)$ and $E(k)$.

```
sf_elliptic_d k = D( $k$ )
```

```
sf_elliptic_d :: Double → Double
sf_elliptic_d k = ((sf_elliptic_k k) - (sf_elliptic_e k)) / (k^2)
```

12.6 Burlisch's elliptic integrals

DLMF: “Bulirschs integrals are linear combinations of Legendres integrals that are chosen to facilitate computational application of Bartkys transformation”

12.6.1 sf_elliptic_cel kc p a b

Compute Burlisch's elliptic integral where $p \neq 0$, $k_c \neq 0$.

$$cel(k_c, p, a, b) = \int_0^{\pi/2} \frac{a \cos^2 \theta + b \sin^2 \theta}{\cos^2 \theta + p \sin^2 \theta} \frac{1}{\sqrt{\cos^2 \theta + k_c^2 \sin^2 \theta}} d\theta \quad cel(k_c, p, a, b)$$

TODO: UNTESTED!

```
sf_elliptic_cel kc p a b = cel( $k_c, p, a, b$ )
```

```
sf_elliptic_cel :: Double → Double → Double → Double → Double
sf_elliptic_cel kc p a b = a * (sf_elliptic_rf 0 (kc^2) 1) + (b-p*a)/3 *
(sf_elliptic_rj 0 (kc^2) 1 p)
```

12.6.2 sf_elliptic_el1 x kc

Compute Burlisch's elliptic integral

$$el_1(x, k_c) =$$

TODO: UNTESTED!

```
sf_elliptic_el1 k kc = el_1( $x, k_c$ )
```

```
sf_elliptic_el1 :: Double → Double → Double
sf_elliptic_el1 x kc =
  —sf_elliptic_f (atan x) (sf_sqrt(1-kc^2))
  let r = 1/x^2
  in sf_elliptic_rf r (r+kc^2) (r+1)
```

12.6.3 sf_elliptic_el2 x kc a b

Compute Burlisch's elliptic integral

$$el_2(x, k_c, a, b) = \int_0^{\arctan x} \frac{a + b \tan^2 \theta}{\sqrt{(1 + \tan^2 \theta)(1 + k_c^2 \tan^2 \theta)}} d\theta$$

TODO: UNTESTED!

```
sf_elliptic_el2 x kc a b = el2(x, kc, a, b)
```

```
sf_elliptic_el2 :: Double → Double → Double → Double → Double
sf_elliptic_el2 x kc a b =
  let r = 1/x^2
  in a * (sf_elliptic_el1 x kc) + (b-a)/3 * (sf_elliptic_rd r (r+kc^2) (r+1))
```

12.6.4 sf_elliptic_el3 x kc p

Compute the Burlisch's elliptic integral

$$el_3(x, k_c, p) = \int_0^{\arctan x} \frac{d\theta}{(\cos^2 \theta + p \sin^2 \theta) \sqrt{\cos^2 \theta + k_c^2 \sin^2 \theta}}$$

TODO: UNTESTED!

```
sf_elliptic_el3 x kc p = el3(x, kc, p)
```

```
sf_elliptic_el3 :: Double → Double → Double → Double
sf_elliptic_el3 x kc p =
  — sf_elliptic_pi(atan(x), 1-p, sf_sqrt(1-kc.^2));
  let r = 1/x^2
  in (sf_elliptic_el1 x kc) + (1-p)/3 * (sf_elliptic_rj r (r+kc^2) (r+1) (r+p))
```

12.7 Symmetric elliptic integrals

12.7.1 sf_elliptic_rc x y

Compute the symmetric elliptic integral $R_C(x, y)$ for real parameters. Let $x \in \mathbb{C} \setminus (-\infty, 0)$, $y \in \mathbb{C} \setminus \{0\}$, then we define

$$R_C(x, y) = \frac{1}{2} \int_0^\infty \frac{dt}{\sqrt{t+x}(t+y)}$$

(where the Cauchy principal value is taken if $y < 0$.) TODO: UNTESTED!

```
sf_elliptic_rc x y = R_C(x, y)
```

```
— x ≥ 0, y ≠ 0
sf_elliptic_rc :: Double → Double → Double
sf_elliptic_rc x y
  | 0 == x ∧ x < y = 1/sf_sqrt(y-x) * sf_acos(sf_sqrt(x/y))
  | 0 < x ∧ x < y = 1/sf_sqrt(y-x) * sf_atan(sf_sqrt((y-x)/x))
  | 0 < y ∧ y < x = 1/sf_sqrt(x-y) * sf_atanh(sf_sqrt((x-y)/x))
```

```
sf_elliptic_rc x y = R_C(x,y) (cont)
```

```

— = 1/sf_sqrt(x-y) * sf_log((sf_sqrt(x) + sf_sqrt(x-y))/sf_sqrt(y))
| y<0  ∧  0≤x = 1/sf_sqrt(x-y) * sf_log((sf_sqrt(x)+sf_sqrt(x-y))/sf_sqrt(-y))
— = 1/sf_sqrt(x-y) * sf_atanh(sf_sqrt(x/(x-y)))
— = sf_sqrt(x/(x-y)) * (sf_elliptic_rc (x-y) (-y))
| x == y      = 1/(sf_sqrt x)
| otherwise    = error "sf_elliptic_rc:_domain_error"

```

12.7.2 sf_elliptic_rd x y z

Compute the symmetric elliptic integral $R_D(x, y, z)$ TODO: UNTESTED!

```
sf_elliptic_rc x y z = R_D(x,y,z)
```

```

— x,y,z>0
sf_elliptic_rd :: Double → Double → Double → Double
sf_elliptic_rd x y z = let (x',s) = (iter x y z 0.0) in (x'**(-3/2) + s)
  where
    iter x y z s =
      let lam = sf_sqrt(x*y) + sf_sqrt(y*z) + sf_sqrt(z*x);
          s' = s + 3/sf_sqrt(z)/(z+lam);
          x' = (x+lam)*two23
          y' = (y+lam)*two23
          z' = (z+lam)*two23
          mu = (x+y+z)/3;
          eps = foldl1 max (map (λt→abs(1-t/mu)) [x,y,z])
      in if eps<2e-16 ∨ [x,y,z]==[x',y',z'] then (x',s')
        else iter x' y' z' s'

```

12.7.3 sf_elliptic_rf x y z

Compute the symmetric elliptic integral of the first kind

$$R_F(x, y, z) = \frac{1}{2} \int_0^\infty \frac{dt}{\sqrt{t+x}\sqrt{t+y}\sqrt{t+z}}$$

TODO: UNTESTED!

```
sf_elliptic_rf x y z = R_F(x,y,z)
```

```

— x,y,z>0
sf_elliptic_rf :: Double → Double → Double → Double
sf_elliptic_rf x y z = 1/(sf_sqrt $ iter x y z)
  where
    iter x y z =
      let lam = (sf_sqrt $ x*y) + (sf_sqrt $ y*z) + (sf_sqrt $ z*x)
          mu = (x+y+z)/3
          eps = foldl1 max $ map (λa→abs(1-a/mu)) [x,y,z]
          x' = (x+lam)/4
          y' = (y+lam)/4
          z' = (z+lam)/4

```



```
sf_elliptic_rf x y z = R_F(x,y,z) (cont)
```

```
in if (eps<1e-16) ∨ ([x,y,z]==[x',y',z'])
then x
else iter x' y' z'
```

12.7.4 sf_elliptic_rg x y z

Compute the symmetric elliptic integral

$$R_G(x,y,z) = \frac{1}{4\pi} \int_0^{2\pi} \int_0^\pi \sqrt{x \sin^2 \theta \cos^2 \phi + y \sin^2 \theta \sin^2 \phi + z \cos^2 \theta \sin \theta} d\theta d\phi$$

TODO: UNTESTED!

```
sf_elliptic_rg x y z = R_G(x,y,z)
```

```
— x,y,z>0
sf_elliptic_rg :: Double → Double → Double → Double
sf_elliptic_rg x y z
| x>y = sf_elliptic_rg y x z
| x>z = sf_elliptic_rg z y x
| y>z = sf_elliptic_rg x z y
| otherwise =
  let !a0 = sqrt (z-x)
      !c0 = sqrt (y-x)
      !h0 = sqrt z
      !t0 = sqrt x
      !(an,tn,cn_sum,hn_sum) = iter 1 a0 t0 c0 (c0^2/2) h0 0
  in ((t0^2 + theta*cn_sum)*(sf_elliptic_rc (tn^2+theta*an^2) tn^2) + h0 + hn_sum)/2
  where
    theta = 1
    iter n an tn cn cn_sum hn hn_sum =
      let an' = (an + sf_sqrt(an^2 - cn^2))/2
          tn' = (tn + sf_sqrt(tn^2 + theta*cn^2))/2
          cn' = cn^2/(2*an')/2
          cn_sum' = cn_sum + 2^((#)n-1)*cn'^2
          hn' = hn*tn'/sf_sqrt(tn'^2+theta*cn'^2)
          hn_sum' = hn_sum + 2^n*(hn' - hn)
          n' = n + 1
      in if cn'^2==0 then (an,tn,cn_sum,hn_sum)
          else iter n' an' tn' cn' hn_sum' hn' hn_sum'
```

12.7.5 sf_elliptic_rj x y z p

Compute the symmetric elliptic integral

$$R_J(x,y,z,p) = \frac{3}{2} \int_0^\infty \frac{dt}{\sqrt{t+x}\sqrt{t+y}\sqrt{t+z}(t+p)}$$

TODO: UNTESTED!

```

sf_elliptic_rj x y z p = R_J(x, y, z, p)

— x, y, z ≥ 0
sf_elliptic_rj :: Double → Double → Double → Double → Double
sf_elliptic_rj x y z p =
  let (x', smm, scale) = iter x y z p 0.0 1.0
  in scale * x' ** (-3/2) + smm
  where
    iter x y z p smm scale =
      let lam = sf_sqrt(x*y) + sf_sqrt(y*z) + sf_sqrt(z*x)
          alpha = p * (sf_sqrt(x) + sf_sqrt(y) + sf_sqrt(z)) + sf_sqrt(x*y*z)
          beta = sf_sqrt(p) * (p + lam)
          smm' = smm + (if (abs(1 - alpha^2/beta^2) < 5e-16)
                        then
                          — optimization to reduce external calls
                          scale * 3 / alpha;
                        else
                          scale * 3 * (sf_elliptic_rc (alpha^2) (beta^2))
                          )
          mu = (x + y + z + p) / 4
          eps = foldl1 max (map (\t → abs(1 - t / mu)) [x, y, z, p])
          x' = (x + lam) * two23 / mu
          y' = (y + lam) * two23 / mu
          z' = (z + lam) * two23 / mu
          p' = (p + lam) * two23 / mu
          scale' = scale * (mu ** (-3/2))
      in if eps < 1e-16 ∨ [x, y, z, p] == [x', y', z', p'] ∨ smm' == smm
         then (x', smm', scale')
         else iter x' y' z' p' smm' scale'

```

13 Spence

Spence's integral for $z \geq 0$ is

$$S(z) = - \int_1^z \frac{\ln t}{t-1} dt = - \int_0^{z-1} \frac{\ln(1+u)}{z} dz$$

and we extend the function via analytic continuation. Spence's function $S(z)$ is related to the dilogarithm function via $S(z) = \text{Li}_2(1-z)$.

13.1 Preamble

```

module Spence (sf_spence) where
import Exp
import Util

```

A useful constant $\text{pi2_6} = \frac{\pi^2}{6}$

```

pi2_6 :: (Value v) => v
pi2_6 = pi^2/6

```

13.2 sf_spence z

Compute Spence's integral **sf_spence** $z = S(z)$. We use a variety of transformations to to allow efficient computation with a series.

$$\begin{aligned} \operatorname{Li}_2(z) + \operatorname{Li}_2\left(\frac{z}{z-1}\right) &= -\frac{1}{2}(\ln(1-z))^2 & z \in \mathbb{C} \setminus [1, \infty) \\ \operatorname{Li}_2(z) + \operatorname{Li}_2\left(\frac{1}{z}\right) &= -\frac{\pi^2}{6} - \frac{1}{2}(\ln(-z))^2 & z \in \mathbb{C} \setminus [0, \infty) \\ \operatorname{Li}_2(z) + \operatorname{Li}_2(1-z) &= \frac{\pi^2}{6} - \ln(z) \ln(1-z) & 0 < z < 1 \end{aligned}$$

(TODO: this code has not be solidly retested after conversion, especially verify complex.)

```
sf_spence z = Li2(z)

sf_spence :: (Value v) => v -> v
sf_spence z
  | is_nan z      = z
  | (re z)<0      = 0/0
  | z == 0        = pi2_6
  | (rabs z)<0.5 = (series z) + (pi2_6 - (sf_log z)*(sf_log (1-z)))
  | (rabs z)<1.0 = -(series (1-z))
  | (rabs z)<2.5 = (series ((z-1)/z)) - (sf_log z)^2/2
  | otherwise    = (series (1/(1-z))) - pi2_6 - (sf_log (z-1))^2/2
```

*series z

The series expansion used for Spence's integral:

$$\text{series } z = - \sum_{k=1}^{\infty} \frac{z^k}{k^2}$$

```
series z =
  let zk = iterate (*z) z
      terms = zipWith (\ t k -> -t/(#)k^2) zk [1..]
  in ksum terms
```

14 Lommel functions

14.1 Preamble

```
module Lommel (
  sf_lommel_s,
  sf_lommel_s2,
) where
import Util
```

-TODO: These are completely untested!

14.2 First Lommel function

For $\mu \pm \nu \neq \pm 1, \pm 3, \pm 5, \dots$ we define the first Lommel function **sf_lommel_s** μ ν $z = S_{\mu,\nu}(z)$ via series-expansion:

$$S_{\mu,\nu}(z) = \frac{z^{\mu+1}}{(\mu+1)^2 - \nu^2} \sum_{k=0}^{\infty} t_k$$

where

$$t_0 = 1 \quad t_k = t_{k-1} \frac{-z^2}{(\mu + 2k + 1)^2 - \nu^2}$$

14.2.1 sf_lommel_s mu nu z

`sf_lommel_s mu nu z` = $S_{\mu,\nu}(z)$

```
sf_lommel_s mu nu z =
  let terms = ixiter 1 1.0 $ \ k t -> -t*z^2 / ((mu+((#)$2*k+1))^2 - nu^2)
      res = ksum terms
  in res * z**(mu+1) / ((mu+1)^2 - nu^2)
```

14.3 Second Lommel function

For $\mu \pm \nu \neq \pm 1, \pm 3, \pm 5, \dots$ the second Lommel function `sf_lommel_s2 mu nu z` = $s_{\mu,\nu}(z)$ is given via an asymptotic expansion:

$$s_{\mu,\nu}(z) \sim \sum_{k=0}^{\infty} u_k$$

where

$$u_0 = 1 \quad u_k = u_{k-1} \frac{-(\mu - 2k + 1)^2 - \nu^2}{z^2}$$

14.3.1 sf_lommel_s2 mu nu z

`sf_lommel_s2 mu nu z` = $s_{\mu,\nu}(z)$

```
sf_lommel_s2 mu nu z =
  let tterms = ixiter 1 1.0 $ \ k t -> -t*((mu-((#)$2*k+1))^2 - nu^2) / z^2
      terms = tk tterms
      res = ksum terms
  in res
  where tk (a:b:cs) = if (rabs a)<(rabs b) then [a] else a:(tk$b:cs)
```