# Computation of Special Functions (Haskell)

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# Contents

1	ntroduction	3
2	Jtility  1 Preamble 1 Data Types 2 Data Types 3 Helper functions 4 Kahan summation 5 Continued fraction evaluation 6 TO BE MOVED	3 3 5 5 6 6
3	Fibonacci Numbers	6
4	Numbers  .1 Preamble	<b>7</b> 7 8
5	Exponential & Logarithm  1 Preamble 2 Exponential 5.2.1 sf_exp x 5.2.2 sf_exp_m1 x 5.2.3 sf_exp_m1 x 5.2.4 sf_exp_menx n x 5.2.5 sf_exp_menx n x 5.2.6 sf_exp_men n x 5.2.6 sf_exp_n n x  3 Logarithm 5.3.1 sf_log x 5.3.2 sf_log_p1 x	8 9 9 9 10 11 11 11 11
6	Gamma	<b>12</b>
	1.1 Preamble 1.2 Misc 1.3 Gamma 1.4 Example 1.5 Section 1	12 12 12 13 13 13 13

			berno																				
	6.4	Digamn																					
		6.4.1	sf_di	gamm	ıa z											 							 14
-	1123	c ,																					1.
7		or funct Preamb																					16
	7.2	Error fu 7.2.1																					
			sf_er																				
		1.2.2	sf_er	10 2			• •		• •	•	 •	 •	•	 •	 •	 • •	•	 •	• •	•	 •	 ٠	 16
8	Exp	onentia	al Int	egra	ıl																		18
	-	Preamb		_												 		 					
		Expone																					
		-	sf_ex	_	-																		
	8.3	Expone	ential	$\frac{1}{1}$	ral I	$E_n$										 							 19
		8.3.1	sf_ex	pint	:_en	n z										 		 					 19
9	AGI																						20
		Preamb																					
	9.2																						
			sf_ag		-																		
		9.2.2	sf_ag	m'a	ιTph	a be	eta			•	 ٠	 •	•	 •	 •	 	•	 •		•	 ٠	 ٠	 2
10	Airy	i.r																					21
10		, Preamb	ole																				
		Ai																					
10.2	10.2	10.2.1																					
	10.3	Bi																					
		10.3.1																					
				J																			
11		mann z																					23
		Preamb																					
	11.2	Zeta .																					
		11.2.1	sf_ze	ta z	٠. :											 							
		11.2.2	sf_ze	$ta_m$	.1 z											 		 •		•			 23
10	T211:-	-+:- f																					2
14		p <b>tic fun</b> Preamb																					<b>2</b> 4
		Elliptic																					
	12.2	12.2.1	_																				
		12.2.1 $12.2.2$																					
	12.3	Elliptic																					
	12.0	12.3.1																					
		12.3.1 $12.3.2$																					
	12.4	Elliptic																					
	12.1	12.4.1																					
	12.5	Burlisch																					
	12.0	12.5.1																					
		12.5.2		_				_															
		12.5.2 $12.5.3$																					
		12.5.4		-																			
	12.6	Symme																					
	12.0	12.6.1																					
		12.6.2		_			-																
		12.6.3		_			-																

	12.6.4 sf_elliptic_rg x y z
13 Sp	pence
13	1 Preamble
13	2 sf_spence z
	ommel functions
	1 Preamble
14	2 First Lommel function
	14.2.1 sf_lommel_s mu nu z
14	3 Second Lommel function
	14.3.1 sf_lommel_s2 mu nu z

## 1 Introduction

Special functions.

# 2 Utility

## 2.1 Preamble

We start with the basic preamble.

```
{-# Language BangPatterns #-}
{-# Language FlexibleContexts #-}
{-# Language FlexibleInstances #-}
{-# Language TypeFamilies #-}
-- {-# Language UndecidableInstances #-}
module Util where
import Data.Complex
```

## 2.2 Data Types

We start by defining a convenient type synonym for complex numbers over Double.

```
type CDouble = Complex Double
```

Next, we define the Value typeclass which is useful for defining our special functions to work over both real (Double) values and over complex (CDouble) values with uniform implementations. This will also make it convenient for handling Quad values (later).

```
class Value v (cont)

re :: v \rightarrow (RealKind \ v)

im :: v \rightarrow (RealKind \ v)

rabs :: v \rightarrow (RealKind \ v)

is_inf :: v \rightarrow Bool

is_nan :: v \rightarrow Bool

is_real :: v \rightarrow Bool

fromDouble :: Double \rightarrow v

fromReal :: (RealKind \ v) \rightarrow v

toComplex :: v \rightarrow (ComplexKind \ v)
```

Value

Value Double

Value CDouble

Both Double and CDouble are instances of the Value typeclass in the obvious ways.

```
instance Value Double
instance Value Double where
  type RealKind Double = Double
  \mathbf{type} ComplexKind \mathbf{Double} = \mathbf{CDouble}
  pos_infty = 1.0/0.0
  neg_{infty} = -1.0/0.0
  nan = 0.0/0.0
  re = id
  im = const 0
  rabs = abs
  is_i = is_i = is_i
  is\_nan = isNaN
  is\_real \ \_ = \mathbf{True}
  from Double = id
  \mathrm{fromReal} = \mathbf{id}
  toComplex x = x :+ 0
```

```
instance Value CDouble
instance Value CDouble where
  \mathbf{type} RealKind CDouble = \mathbf{Double}
  type ComplexKind CDouble = CDouble
  pos_infty = (1.0/0.0) :+ 0
  neg_{infty} = (-1.0/0.0) :+ 0
  nan = (0.0/0.0) :+ 0
  \mathrm{re} = \mathbf{realPart}
  im = imagPart
  rabs = \mathbf{realPart.abs}
  is_inf z = (is_inf.re$z) \lor (is_inf.im$z)
  is_nan z = (is_nan.re$z) \lor (is_nan.im$z)
  is_real_{-} = False
  from Double x = x :+ 0
  from Real x = x :+ 0
  toComplex = id
```

TODO: add quad versions also

## 2.3 Helper functions

A convenient shortcut, as we often find ourselves converting indices (or other integral values) to our computation type.

```
(\#) :: (Integral a, Nm b) \Rightarrow a \rightarrow b (\#) = fromIntegral
```

A version of iterate which passes along an index also (very useful for computing terms of a power-series, for example.)

Computes the relative error in terms of decimal digits, handy for testing. Note that this fails when the exact value is zero.

$$\texttt{relerr e a} = \log_{10} \left| \frac{a-e}{e} \right|$$

```
relerr :: (Value v) \Rightarrow v \rightarrow v \rightarrow (RealKind v)
relerr !exact !approx = re $! logBase 10 (abs ((approx-exact)/exact))
```

## 2.4 Kahan summation

A useful tool is so-called Kahan summation, based on the observation that in floating-point arithmetic, one can ...

Here kadd t s e k is a single step of addition, adding a term to a sum+error and passing the updated sum+error to the continuation.

```
— kadd value oldsum olderr — newsum newerr kadd :: (Value v) \Rightarrow v \rightarrow v \rightarrow v \rightarrow (v \rightarrow v \rightarrow a) \rightarrow a kadd t s e k =

let y = t - e

s' = s + y

e' = (s' - s) - y

in k s' e'
```

Here ksum terms sums a list with Kahan summation. The list is assumed to be (eventually) decreasing and the summation is terminated as soon as adding a term doesn't change the value. (Thus any zeros in the list will immediately terminate the sum.) This is typically used for power-series or asymptotic expansions.

```
ksum terms

ksum :: (Value\ v) \Rightarrow [v] \rightarrow v
ksum terms = ksum' terms const

ksum' :: (Value\ v) \Rightarrow [v] \rightarrow (v \rightarrow v \rightarrow a) \rightarrow a
ksum' terms k = f 0 0 terms

where

f !s !e [] = k s e
f !s !e (t:terms) =
let !y = t - e
!s' = s + y
```

```
ksum terms (cont)
```

ksum

```
\begin{array}{c} !e'=(s'-s)-y\\ \textbf{in if } s'==s\\ \textbf{then } k\ s'\ e'\\ \textbf{else } f\ s'\ e'\ terms \end{array}
```

#### 2.5 Continued fraction evaluation

This is Steed's algorithm for evaluation of a continued fraction

$$C = b_0 + a_1/(b_1 + a_2/(b_2 + a_3/(b_3 + \cdots)))$$

where  $C_n = A_n/B_n$  is the partial evaluation up to ...  $a_n/b_n$ . Here steeds as bs evaluates until  $C_n = C_{n+1}$ . TODO: describe the algorithm.

```
steeds :: (Value\ v) \Rightarrow [v] \rightarrow [v] \rightarrow v

steeds (a1:as)\ (b0:b1:bs) =

let\ !c0 = b0

!d1 = 1/b1

!delc1 = a1*d1

!c1 = c0 + delc1

in recur c1 delc1 d1 as bs

where recur !cn_1 !delcn_1 !dn_1 !(an:as) !(bn:bs) =

let\ !dn = 1/(dn_1*an+bn)

!delcn = (bn*dn - 1)*delcn_1

!cn = cn_1 + delcn

in if (cn = cn_1) \lor is_nan\ cn\ then\ cn\ else\ (recur\ cn\ delcn\ dn\ as\ bs)
```

## 2.6 TO BE MOVED

```
sf\_sqrt :: (Value v) \Rightarrow v \rightarrow v
sf\_sqrt = sqrt
```

## 3 Fibonacci Numbers

A silly approach to efficient computation of Fibonacci numbers

$$f_n = f_{n-1} + f_{n-2} \qquad f_0 = 0 \qquad f_1 = 1$$

The idea is to use the closed-form solution:

$$f_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^n + \frac{-1}{\sqrt{5}} \left( \frac{1-\sqrt{5}}{2} \right)^n$$

and note that we can work in  $\mathbb{Q}[\sqrt{5}]$  with terms of the form  $a+b\sqrt{5}$  with  $a,b\in\mathbb{Q}$  (notice that  $\frac{1}{\sqrt{5}}=\frac{\sqrt{5}}{5}$ .)

$$(a+b\sqrt{5}) + (c+d\sqrt{5}) = (a+c) + (b+d)\sqrt{5}$$

$$(a+b\sqrt{5}) * (c+d\sqrt{5}) = (ac+5bd) + (ad+bc)\sqrt{5}$$

We use the Rational type to represent elements of  $\mathbb{Q}$ , which is a bit more than we actually need, as in the computations above the denominator of  $\left(\frac{1\pm\sqrt{5}}{2}\right)^n$  is always, in fact, 1 or 2.

```
module Fibo (fibonacci) where
import Data.Ratio
data Q5 = Q5 Rational Rational
deriving (Eq)
```

The number-theoretic norm  $N(a+b\sqrt{5})=a^2-5b^2$ , though unused in our application.

```
norm (Q5 ra qa) = ra^2-5*qa^2
```

Human-friendly Show instantiation.

#### instance Show Q5 where

```
show (Q5 ra qa) = (show ra)++"+"+"(show qa)++"**sqrt(5)"
```

Implementation of the operations for typeclasses Num and Fractional. The abs and signum functions are unused, so we just give placeholder values.

#### instance Num Q5 where

#### instance Fractional Q5 where

```
recip a@(Q5 ra qa) = Q5 (ra/(norm a)) (-qa/(norm a)) fromRational r = (Q5 \ r \ 0)
```

Finally, we define  $\phi_{\pm} = \frac{1}{2}(1 \pm \sqrt{5})$  and  $c_{\pm} = \pm \frac{1}{5}\sqrt{5}$  so that  $f_n = c_+\phi_+^n + c_-\phi_-^n$ . (We can shortcut and extract the value we want without actually computing the full expression.)

```
\begin{array}{lll} phip = Q5 \ (1\%2) \ (1\%2) \\ cp & = Q5 \ 0 & (1\%5) \\ phim = Q5 \ (1\%2) \ (-1\%2) \\ cm & = Q5 \ 0 & (-1\%5) \\ fibonacci 'n = \mathbf{let} \ (Q5 \ r \ q) = cp*phip^n + cm*phim^n \ \mathbf{in} \ \mathbf{numerator} \ r \\ fibonacci \ n = \mathbf{let} \ (Q5 \ _q) = phip^n \ \mathbf{in} \ \mathbf{numerator} \ (2*q) \end{array}
```

## 4 Numbers

#### 4.1 Preamble

```
module Numbers where
import Data.Ratio
import qualified Fibo

fibonacci_number :: Int → Integer
fibonacci_number n = Fibo.fibonacci n

lucas_number :: Int → Integer
lucas_number = undefined

euler_number :: Int → Integer
euler_number = undefined

catalan_number :: Integer → Integer
catalan_number 0 = 1
catalan_number n = 2*(2*n-1)*(catalan_number (n-1))*div*(n+1)
```

```
bernoulli_number :: Int \rightarrow Rational
bernoulli_number = undefined
tangent\_number :: Int \rightarrow Integer
tangent\_number = undefined
triangular\_number :: Integer \rightarrow Integer
triangular_number n = n*(n+1)'div'2
factorial :: (Integral a) \Rightarrow a \rightarrow a
factorial 0 = 1
factorial 1 = 1
factorial n = product [1..n]
binomial :: (Integral a) \Rightarrow a \rightarrow a \rightarrow a
binomial n k
      k < 0 = 0
      n < 0 = 0
      k > n = 0
      k=0 = 1
      k=n=1
      k > n' div' 2 = binomial n (n-k)
      otherwise = (product [n-(k-1)..n]) 'div' (product [1..k])
4.2
       Stirling numbers
— TODO: this is extremely inefficient approach
stirling_number_first_kind n k = s n k
  where s n k | k \le 0 \lor n \le 0 = 0
         s n 1 = (-1)^{(n-1)}*(factorial (n-1))
         s n k = (s (n-1) (k-1)) - (n-1)*(s (n-1) k)
— TODO: this is extremely inefficient approach
stirling\_number\_second\_kind n k = s n k
  where s n k | k \le 0 \lor n \le 0 = 0
         s n 1 = 1
         s n k = k*(s (n-1) k) + (s (n-1) (k-1))
```

# 5 Exponential & Logarithm

In this section, we implement the exponential function and logarithm function, as well as useful variations.

#### 5.1 Preamble

We begin with a typical preamble.

```
form in the state of the s
```

## 5.2 Exponential

We start with implementation of the most basic special function, exp(x) or  $e^x$  and variations thereof.

#### 5.2.1 sf\_exp x

For the exponential  $sf_{exp} = exp(x)$  we use a simple series expansions

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{n!}$$

after first using the identity  $e^{-x} = 1/e^x$  to ensure that the real part of the argument is positive. This avoids disastrous cancellation for negative arguments, (though note that for complex arguments this is not sufficient.) TODO: should do range-reduction first... TODO: maybe for complex, use explicit cis?

```
\begin{array}{l} \textbf{sf\_exp} \ \textbf{x} = e^x \\ \\ \textbf{sf\_exp} \ :: \ (\text{Value v}) \Rightarrow \textbf{v} \rightarrow \textbf{v} \\ \\ \textbf{sf\_exp} \ !\textbf{x} \\ | \ \textbf{is\_inf} \ \textbf{x} \ = \textbf{if} \ (\textbf{re} \ \textbf{x}) < 0 \ \textbf{then} \ 0 \ \textbf{else} \ \textbf{pos\_infty} \\ | \ \textbf{is\_nan} \ \textbf{x} \ = \textbf{x} \\ | \ (\textbf{re} \ \textbf{x}) < 0 \ = 1/(\textbf{sf\_exp} \ (-\textbf{x})) \\ | \ \textbf{otherwise} = \textbf{ksum} \ \$ \ \textbf{ixiter} \ 1 \ 1.0 \ \$ \ \lambda \textbf{n} \ \textbf{t} \rightarrow \textbf{t*x}/(\#) \textbf{n} \\ \end{array}
```

#### 5.2.2 sf\_exp\_m1 x

For numerical calculations, it is useful to have  $sf_{exp_m1} = e^x - 1$  as explicitly calculating this expression will give poor results for x near 1. We use a series expansion for the calculation. Again for negative real part we reflect using  $e^{-x} - 1 = -e^{-x}(e^x - 1)$ . TODO: should do range-reduction first... TODO: maybe for complex, use explicit cis?

```
\begin{array}{l} \textbf{sf\_exp\_m1} & \textbf{x} = e^x - 1 \\ \\ \textbf{sf\_exp\_m1} & :: & (\text{Value v}) \Rightarrow \textbf{v} \rightarrow \textbf{v} \\ \textbf{sf\_exp\_m1} & !\textbf{x} \\ | & \textbf{is\_inf x} & = \textbf{if (re x)} < 0 \textbf{ then } - 1 \textbf{ else pos\_infty} \\ | & \textbf{is\_nan x} & = \textbf{x} \\ | & (\text{re x}) < 0 & = -\text{sf\_exp x} * \text{sf\_exp\_m1 (-x)} \\ | & \textbf{otherwise} & = \text{ksum \$ ixiter 2 x \$ } \lambda \textbf{n} \text{ t} \rightarrow \text{t*x/((\#)n)} \end{array}
```

#### 5.2.3 sf\_exp\_m1vx x

Similarly, it is useful to have the scaled variant  $sf_{exp_m1vx} = \frac{e^x - 1}{x}$ . In this case, we use a continued-fraction expansion

$$\frac{e^x - 1}{x} = \frac{2}{2 - x + 1} \frac{x^2 / 6}{1 + 1} \frac{x^2 / 4 \cdot 3 \cdot 5}{1 + 1} \frac{x^2 / 4 \cdot 5 \cdot 7}{1 + 1} \frac{x^2 / 4 \cdot 7 \cdot 9}{1 + 1} \cdots$$

For complex values, simple calculation is inaccurate (when  $\Re z \sim 1$ ).

```
sf_exp_m1vx x = \frac{e^x-1}{x}
                                                                                                                sf_exp_m1vx
sf_exp_m1vx :: (Value v) \Rightarrow v \rightarrow v
sf_exp_m1vx !x
    is_inf x = if (re x)<0 then 0 else pos_infty
    is\_nan\ x=x
    rabs(x)>(1/2) = (sf_exp x - 1)/x - inaccurate for some complex points
    otherwise =
       let x2 = x^2
       in 2/(2 - x + x^2/6/(1 + x^2))
           + x2/(4*(2*3-3)*(2*3-1))/(1
           + x2/(4*(2*4-3)*(2*4-1))/(1
           + x2/(4*(2*5-3)*(2*5-1))/(1
           + x2/(4*(2*6-3)*(2*6-1))/(1
           + x2/(4*(2*7-3)*(2*7-1))/(1
           + x2/(4*(2*8-3)*(2*8-1))/(1
           ))))))));
```

#### 5.2.4 sf\_exp\_menx n x

Compute the scaled tail of series expansion of the exponential function.

$$\texttt{sf\_exp\_menx n } \texttt{x} = \frac{n!}{x^n} \left( e^z - \sum_{k=0}^{n-1} \frac{x^k}{k!} \right) = \frac{n!}{x^n} \sum_{k=n}^{\infty} \frac{x^k}{k!} = n! \sum_{k=0}^{\infty} \frac{x^k}{(k+n)!}$$

We use a continued fraction expansion and using the modified Lentz algorithm for evaluation.

```
sf_exp_menx :: (Value v) \Rightarrow Int \rightarrow v \rightarrow v
sf_{exp_menx} 0 z = sf_{exp} z
sf_exp_menx 1 z = sf_exp_m1vx z
sf_exp_menx n z
  | is_inf z = if (re z)>0 then pos_infty else (0) — TODO: verify
    is_nan z = z
   otherwise = exp_menx_contfrac n z
 where
    !zeta = 1e-150
    ! eps = 1e-16
    nz ! z = if z = 0 then zeta else z
    exp_menx_contfrac n z =
      let ! fj = (#)$ n+1
          !cj = fj
          !dj = 0
          !j = 1
      in lentz j dj cj fj
    lentz !j !dj !cj !fj =
      let !aj = if (odd j)
                 then z*((\#)$(j+1)'div'2)
                 else -z*((\#)\$(n+(j'div'2)))
           bj = (\#) n+1+j
           !dj' = nz  bj + aj*dj
           !cj' = nz bj + aj/cj
           ! dji = 1/dj
           !deltaj = cj '*dji
           !fj' = fj*deltaj
```

```
in if (rabs(deltaj-1)<eps)
  then 1/(1-z/fj')
  else lentz (j+1) dji cj' fj'</pre>
```

#### 5.2.5 sf\_exp\_men n x

This is the generalization of  $sf_{exp_m1}$  x, giving the tail of the series expansion of the exponential function, for  $n = 0, 1, \ldots$ 

$$\texttt{sf\_exp\_men n } \texttt{z} = e^z - \sum_{k=0}^{n-1} \frac{z^k}{k!} = \sum_{k=n}^{\infty} \frac{z^k}{k!}$$

The special cases are: n = 0 gives  $e^x = \mathtt{sf\_exp} \ \mathtt{x}$  and n = 1 gives  $e^x - 1 = \mathtt{sf\_exp\_m1} \ \mathtt{x}$ . We compute this by calling the scaled version  $\mathtt{sf\_exp\_menx}$  and rescaling back.

```
— ($n=0, 1, 2, \circ ...$)

sf_exp_men :: (Value v) \Rightarrow Int \rightarrow v \rightarrow v

sf_exp_men !n !x = (sf_exp_menx n x) * x^n / ((#)$factorial n)
```

#### 5.2.6 sf\_expn n x

```
— Compute initial part of series for exponential, \lambda \sum_{k=0}^{\infty} \frac{1}{2} \cdot \frac{1}{2} \cdot
```

## 5.3 Logarithm

#### 5.3.1 sf\_log x

We simply use the built-in implementation (from the Floating typeclass).

```
sf_{\log} :: (Value \ v) \Rightarrow v \rightarrow v
sf_{\log} = \log
```

#### 5.3.2 sf\_log\_p1 x

The accuracy preserving  $sf_{log_p1} x = \ln 1 + x$ . For values close to zero, we use a power series expansion

$$\ln(1+x) = 2\sum_{n=0}^{\infty} \frac{\left(\frac{x}{x+2}\right)^{2n+1}}{2n+1}$$

and otherwise just compute it directly.

```
\begin{array}{l} sf\_log\_p1 \ :: \ (Value \ v) \ \Rightarrow \ v \ \rightarrow \ v \\ sf\_log\_p1 \ !z \\ | \ is\_nan \ z = z \\ | \ (rabs \ z) > 0.25 = sf\_log \ (1+z) \\ | \ \textbf{otherwise} = series \ z \\ \textbf{where} \\ series \ z = \end{array}
```

```
\begin{array}{ll} \textbf{let} & !\texttt{r} = \texttt{z}/(\texttt{z}+2) \\ & !\texttt{zr2} = \texttt{r}^2 \\ & !\texttt{tterms} = \textbf{iterate} \ (*\texttt{zr2}) \ (\texttt{r}*\texttt{zr2}) \\ & !\texttt{terms} = \textbf{zipWith} \ (\lambda \texttt{n} \ \texttt{t} \ \rightarrow \ \texttt{t}/((\#)\$2*\texttt{n}+1)) \ [1..] \ \texttt{tterms} \\ \textbf{in} \ 2*(\texttt{ksum} \ (\texttt{r}:\texttt{terms})) \end{array}
```

A simple continued fraction implementation for  $\ln 1 + z$ 

$$\ln(1+z) = z/(1+z/(2+z/(3+4z/(4+4z/(5+9z/(6+9z/(7+\cdots)))))))$$

Though unused for now, it seems to have decent convergence properties.

```
ln_1_z_cf z = steeds (z:(ts 1)) [0..]
where ts n = (n^2*z):(n^2*z):(ts (n+1))
```

## 6 Gamma

## 6.1 Preamble

A basic preamble.

```
module Gamma (
euler_gamma,
factorial,
sf_beta,
sf_gamma,
sf_invgamma,
sf_lngamma,
sf_digamma,
bernoulli_b,
)
where
import Exp
import Numbers(factorial)
import Trig
import Util
```

#### 6.2 Misc

#### 6.2.1 euler\_gamma

A constant for Euler's gamma:

$$\gamma = \lim_{n \to \infty} \left( \sum_{k=1}^{n} \frac{1}{n} - \ln n \right)$$

euler\_gamma :: (Floating a)  $\Rightarrow$  a euler\_gamma = 0.577215664901532860606512090082402431042159335939923598805767234884867726777664670936947063291746749

#### 6.2.2 sf\_beta a b

The Beta integral

$$B(a,b) = \int_0^1 t^{a-1} (1-t)^{b-1} dt = \frac{\Gamma(a)\Gamma(b)}{\Gamma(a+b)}$$

implemented in terms of log-gamma

$${\tt sf\_beta}$$
 a  ${\tt b} = e^{\ln\Gamma(a) + \ln\Gamma(b) - \ln\Gamma(a+b)}$ 

```
sf_beta :: (Value v) \Rightarrow v \rightarrow v \rightarrow v sf_beta a b = sf_exp \ (sf_lngamma a) + (sf_lngamma b) - (sf_lngamma$a+b)
```

## 6.3 Gamma

The gamma function

$$\Gamma(z) = \int_0^\infty e^{-t} t^z \, \frac{dz}{z}$$

#### 6.3.1 sf\_gamma z

The gamma function implemented using the identity  $\Gamma(z) = \frac{1}{z}\Gamma(z+1)$  to increase the real part of the argument to be > 15 and then using an asymptotic expansion for log-gamma, lngamma\_asymp, to evaluate.

sf\_gamma

```
\begin{array}{l} \textbf{sf\_gamma} \quad \textbf{x} = \Gamma(x) \\ \\ \textbf{sf\_gamma} \quad \vdots \quad (\text{Value } \textbf{v}) \, \Rightarrow \, \textbf{v} \, \rightarrow \, \textbf{v} \\ \\ \textbf{sf\_gamma} \quad \textbf{x} = \\ \\ \text{redup } \textbf{x} \; 1 \, \$ \; \lambda \; \textbf{x'} \; \textbf{t} \, \rightarrow \, \textbf{t} \; * \; (\textbf{sf\_exp} \; (\text{lngamma\_asymp } \textbf{x'})) \\ \\ \textbf{where} \; \text{redup } \textbf{x} \; \textbf{t} \; k \\ \\ | \; (\text{re } \textbf{x}) > 15 = k \; \textbf{x} \; \textbf{t} \\ | \; \textbf{otherwise} = \text{redup} \; (\textbf{x} + 1) \; (\textbf{t} / \textbf{x}) \; k \\ \end{array}
```

#### 6.3.2 \*lngamma\_asymp z

The asymptotic expansion for log-gamma

$$\ln \Gamma(z) \sim (z - \frac{1}{2}) \ln z - z + \frac{1}{2} \ln(2\pi) + \sum_{k=1}^{\infty} \frac{B_{2k}}{2k(2k-1)z^{2k-1}}$$

where  $B_n$  is the *n*'th Bernoulli number.

#### 6.3.3 sf\_invgamma z

The inverse gamma function,  $sf_{invgamma} z = \frac{1}{\Gamma(z)}$ .

```
\begin{array}{l} \text{sf\_invgamma} \ :: \ (\text{Value } v) \ \Rightarrow \ v \ \rightarrow \ v \\ \text{sf\_invgamma} \ x = \\ \textbf{let} \ (x',t) = \text{redup } x \ 1 \\ \text{lngx} = \text{lngamma\_asymp } x' \\ \textbf{in} \ t \ * \ (\text{sf\_exp\$-lngx}) \\ \textbf{where } \text{redup } x \ t \\ \text{|} \ (\text{re } x) > 15 = (x,t) \\ \text{|} \ \textbf{otherwise} = \text{redup } (x\!+\!1) \ (t\!*\!x) \end{array}
```

## 6.3.4 sf\_lngamma z

where redup x t

The log-gamma function, sf\_lngamma  $\mathbf{z} = \ln \Gamma(z)$ . sf\_lngamma :: (Value v)  $\Rightarrow$  v  $\rightarrow$  v sf\_lngamma x = let (x',t) = redup x 0 lngx = lngamma\_asymp x' in t + lngx

```
| (re x)>15 = (x,t)
| otherwise = redup (x+1) (t-sf_log x)
```

#### 6.3.5 bernoulli\_b n

The Bernoulli numbers,  $B_n$ . A simple hard-coded table, for now. (Should be moved to Numbers module and general, cached, implementation done.)

```
bernoulli_b :: (Value v) \Rightarrow Int \rightarrow v bernoulli_b 1 = -1/2 bernoulli_b k | k'mod 2==1 = 0 bernoulli_b 0 = 1 bernoulli_b 2 = 1/6 bernoulli_b 4 = -1/30 bernoulli_b 6 = 1/42 bernoulli_b 8 = -1/30 bernoulli_b 10 = 5/66 bernoulli_b 12 = -691/2730 bernoulli_b 14 = 7/6 bernoulli_b 16 = -3617/510 bernoulli_b 18 = 43867/798 bernoulli_b 20 = -174611/330 bernoulli_b = undefined
```

## Spouge's approximation to the gamma function

In tests, this gave disappointing results.

## 6.4 Digamma

The digamma function

$$\psi(z) = \frac{d}{dz} \ln \Gamma(z) = \frac{\Gamma'(z)}{\Gamma(z)}$$

#### 6.4.1 sf\_digamma z

We implement with a series expansion for  $|z| \le 10$  and otherwise with an asymptotic expansion.

```
sf_digamma :: (Value v) \Rightarrow v \rightarrow v

—sf_digamma n | is_nonposint n = Inf

sf_digamma z | (rabs z)>10 = digamma_asympt z

| otherwise = digamma_series z
```

The series expansion is the following

$$\psi(z) = -\gamma - \frac{1}{z} + \sum_{k=1}^{\infty} \frac{z}{k(k+z)}$$

but with Euler-Maclaurin correction terms:

$$\psi(z) = -\gamma - \frac{1}{z} + \sum_{k=1}^{n} \frac{z}{k(k+z)} + \left(\ln \frac{k+z}{k} - \frac{z}{2k(k=z)} + \sum_{j=1}^{p} B_{2j}(k^{-2j} - (k+z)^{-2j})\right)$$

```
digamma\_series :: (Value v) \Rightarrow v \rightarrow v
digamma\_series z =
  let res = -\text{euler\_gamma} - (1/z)
       terms = map (\lambda k \rightarrow z/((\#)k*(z+(\#)k))) [1..]
       corrs = map (correction.(#)) [1..]
  in summer res res terms corrs
    summer :: (Value v) \Rightarrow v \rightarrow v \rightarrow [v] \rightarrow [v] \rightarrow v
    summer res sum (t:terms) (c:corrs) =
       let sum' = sum + t
           res' = sum' + c
       in if res=res' then res
           else summer res' sum' terms corrs
    bn1 = bernoulli_b 2
    bn2 = bernoulli_b 4
     bn3 = bernoulli_b 6
     bn4 = bernoulli_b 8
    correction k =
       (sf_log_k(k+z)/k) - z/2/(k*(k+z))
         + bn1*(k^{\hat{}}(-2) - (k+z)^{\hat{}}(-2))
         + bn2*(k^{\hat{}}(-4) - (k+z)^{\hat{}}(-4))
         + bn3*(k^{(-6)} - (k+z)^{(-6)})
         + bn4*(k^{\hat{}}(-8) - (k+z)^{\hat{}}(-8))
```

The asymptotic expansion (valid for  $|argz| < \pi$ ) is the following

$$\psi(z) \sim \ln z - \frac{1}{2z} + \sum_{k=1}^{\infty} \frac{B_{2k}}{2kz^{2k}}$$

Note that our implementation will fail if the bernoulli\_b table is exceeded. If  $\Re z < \frac{1}{2}$  then we use the reflection identity to ensure  $\Re z \geq \frac{1}{2}$ :

$$\psi(z) - \psi(1-z) = \frac{-\pi}{\tan(\pi z)}$$

```
digamma_asympt :: (Value v) \Rightarrow v \rightarrow v digamma_asympt z 

| (re z)<0.5 = compute (1-z) - \frac{pi}{(sf_{tan}(pi*z))} + (sf_{log}(1-z)) - \frac{1}{(2*(1-z))} | otherwise = compute z (sf_{log}(z) - \frac{1}{(2*z)}) where compute z res = let z_2 = z^(-2) z_2 terms = zipWith (2z_1 - z_2) + (2z_1 - z_2) terms = zipWith (2z_1 - z_2) + (2z_1 - z_2) in sumit res res terms sumit res of (z_1 - z_2) + (z_2 - z_2) let res' = res - t in if res=res' (z_1 - z_2) + (z_2 - z_2) (rabs ot) then res else sumit res' t terms
```

## 7 Error function

#### 7.1 Preamble

```
{-# Language BangPatterns #-}
module Erf (
    sf_erf,
    sf_erfc,
) where
import Exp
import Util
```

#### 7.2 Error function

#### 7.2.1 sf\_erf z

The error function  $sf_erf z = erf z$  where

$$\operatorname{erf} z = \frac{2}{\sqrt{\pi}} \int_{-\infty}^{z} e^{-x^{2}} \, dx$$

For  $\Re z < -1$ , we transform via erf  $z = -\operatorname{erf}(-z)$  and for |z| < 1 we use the power-series expansion, otherwise we use erf  $z = 1 - \operatorname{erfc} z$ . (TODO: this implementation is not perfect, but workable for now.)

```
sf_{erf} :: (Value v) \Rightarrow v \rightarrow v

sf_{erf} z

| (re z) < (-1) = -sf_{erf}(-z)

| (rabs z) < 1 = erf_{erf}(-z)

| otherwise = 1 - sf_{erf}(-z)
```

## 7.2.2 sf\_erfc z

The complementary error-function  $sf_{erfc} z = erfc z$  where

$$\operatorname{erfc} z = 1 - \operatorname{erf} z = \frac{2}{\sqrt{\pi}} \int_{z}^{\infty} e^{-x^{2}} dx$$

For  $\Re z < -1$  we transform via erfc  $z = 2 - \operatorname{erf}(-z)$  and if |z| < 1 then we use erfc  $z = 1 - \operatorname{erf} z$ . Finally, if |z| < 10 we use a continued-fraction expansion and an asymptotic expansion otherwise. (TODO: there are a few issues with this implementation: For pure imaginary values and for extremely large values it seems to hang.)

```
— infinite loop when (re z)=0 sf_erfc :: (Value v) ⇒ v → v sf_erfc z 

| (re z)<(-1) = 2-(sf_erfc (-z)) 

| (rabs z)<1 = 1-(sf_erf z) 

| (rabs z)<10 = erfc_cf_pos1 z 

| otherwise = erfc_asymp_pos z — TODO: hangs for very large input
```

#### erf\_series z

The series expansion for erf z:

erf 
$$z = \frac{2}{\sqrt{\pi}} \sum_{n=0}^{\infty} \frac{(-)^n z^{2n+1}}{n!(2n+1)}$$

There is an alternative expansion erf  $z = \frac{2}{\sqrt{\pi}}e^{-z^2}\sum_{n=0}^{\infty}\frac{2^nz^{2n+1}}{1\cdot 3\cdots (2n+1)}$ , but we don't use it here. (TODO: why not?)

```
erf_series z = 
let z2 = z^2
  rts = ixiter 1 z $ \lambdan t \rightarrow (-t)*z2/(\#)n
  terms = zipWith (\lambda n t \rightarrowt/(\#)(2*n+1)) [0..] rts
  in (2/sf_sqrt pi) * ksum terms
```

#### \*sf\_erf z

This asymptotic expansion for erfc z is valid as  $z \to +\infty$ :

erfc 
$$z \sim \frac{e^{-z^2}}{\sqrt{\pi}} \sum_{n=0}^{\infty} (-)^n \frac{(1/2)_m}{z^{2m+1}}$$

where the Pochhammer symbol  $(1/2)_m$  is given by:

$$\left(\frac{1}{2}\right)_m = \frac{1 \cdot 3 \cdot 5 \cdots (2m-1)}{2^m} = \frac{(2m)!}{m! 2^{2m}}$$

TODO: correct the asymptotic term checking (not smallest but pre-smallest term).

```
erfc_asymp_pos z = 

let z2 = z^2

iz2 = 1/2/z2

terms = ixiter 1 (1/z) $ \lambda n t \to (-t*iz2)*(\#)(2*n-1)

tterms = tk terms

in (sf_exp(-z2))/(sqrt pi) * ksum tterms

where tk (a:b:cs) = if (rabs a) < (rabs b) then [a] else a:(tk$b:cs)
```

#### \*erfc\_cf\_pos1 z

A continued-fraction expansion for erfc z:

$$\sqrt{\pi}e^{z^2}\operatorname{erfc} z = \frac{z}{z^2 + 1} \frac{1/2}{1 + z^2 + 1} \frac{3/2}{1 + \cdots}$$

```
erfc_cf_pos1 z =
let z2 = z^2
as = z: (map fromDouble [1/2,1..])
bs = 0: cycle [z2,1]
cf = steeds as bs
in sf_exp(-z2) / (sqrt pi) * cf
```

## \*erfc\_cf\_pos1 z

This is an alternative continued-fraction expansion.

$$\sqrt{\pi}e^{z^2}\operatorname{erfc} z = \frac{2z}{2z^2 + 1 - 2z^2 + 5 - 2z^2 + 9 - \cdots} \cdot \cdots$$

Unused for now.

```
erfc_cf_pos2 z = 

let z2 = z^2 

as = (2*z):(map (\lambdan \rightarrow (#)$ -(2*n+1)*(2*n+2)) [0..]) 

bs = 0:(map (\lambdan \rightarrow 2*z2+(#)4*n+1) [0..]) 

cf = steeds as bs 

in sf_exp(-z2) / (sqrt pi) * cf
```

## 8 Exponential Integral

## 8.1 Preamble

```
module ExpInt(
    sf_expint_ei,
    sf_expint_en,
)
where
import Exp
import Gamma
import Util
```

## 8.2 Exponential integral Ei

The exponential integral Ei z is defined for x < 0 by

$$\mathrm{Ei}(z) = -\int_{-x}^{\infty} \frac{e^{-t}}{t} \, dt$$

It can be defined

## 8.2.1 sf\_expint\_ei z

We give only an implementation for  $\Re z \geq 0$ . We use a series expansion for |z| < 40 and an asymptotic expansion otherwise.

sf\_expint\_ei

expint\_ei\_\_se

```
\begin{array}{lll} \mathbf{sf\_expint\_ei} & \mathbf{z} = \mathrm{Ei}(z) \\ \\ \mathbf{sf\_expint\_ei} & :: & (\mathrm{Value} \ \mathrm{v}) \Rightarrow \mathrm{v} \rightarrow \mathrm{v} \\ \\ \mathbf{sf\_expint\_ei} & z \\ | & (\mathrm{re} \ z) < 0.0 = (0/0) - (\mathit{NaN}) \\ | & z = 0.0 = (-1/0) - (-\mathit{Inf}) \\ | & (\mathrm{rabs} \ z) < 40 = \mathrm{expint\_ei\_series} \ z \\ | & \mathbf{otherwise} & = \mathrm{expint\_ei\_asymp} \ z \end{array}
```

The series expansion is given (for x > 0)

$$\mathrm{Ei}(x) = \gamma + \ln x + \sum_{n=1}^{\infty} \frac{x^n}{n!n}$$

We evaluate the addition of the two terms with the sum slightly differently when  $\Re z < 1/2$  to reduce floating-point cancellation error slightly.

```
expint_ei__series :: (Value v) \Rightarrow v \rightarrow v

expint_ei__series z =

let tterms = ixiter 2 z \Rightarrow \lambdan t \rightarrow t*z/(#)n

terms = zipWith (\lambda t n \rightarrowt/(#)n) tterms [1..]

res = ksum terms

in if (re z)<0.5

then sf_log(z * sf_exp(euler_gamma + res))

else res + sf_log(z) + euler_gamma
```

The asymptotic expansion as  $x \to +\infty$  is

$$\operatorname{Ei}(x) \sim \frac{e^x}{x} \sum_{n=0}^{\infty} \frac{n!}{x^n}$$

```
expint_ei__asymp z
```

expint\_ei\_\_as

sf\_expint\_en

```
expint_ei_asymp :: (Value v) \Rightarrow v \rightarrow v expint_ei_asymp z = let terms = tk $ ixiter 1 1.0 $ \lambdan t \rightarrow t/z*(#)n res = ksum terms in res * (sf_exp z) / z where tk (a:b:cs) = if (rabs a)<(rabs b) then [a] else a:(tk$b:cs)
```

## 8.3 Exponential integral $E_n$

The exponential integrals  $E_n(z)$  are defined as

$$E_n(z) = z^{n-1} \int_z^\infty \frac{e^{-t}}{t^n} dt$$

They satisfy the following relations:

$$E_0(z) = \frac{e^{-z}}{z}$$

$$E_{n+1}(z) = \int_z^{\infty} E_n(t) dt$$

And they can be expressed in terms of incomplete gamma functions:

$$E_n(z) = z^{n-1}\Gamma(1-n,z)$$

(which also gives a generalization for non-integer n).

#### 8.3.1 sf\_expint\_en n z

 $sf_{expint_n} = n z = E_n(z)$ 

We use this series expansion for  $E_1(z)$ :

$$E_1(z) = -\gamma - \ln z + \sum_{k=1}^{\infty} (-)^k \frac{z^k}{k!k}$$

(Note that this will not be good for large values of z.)

```
expint_en_{-1} :: (Value v) \Rightarrow v \rightarrow v
expint_en_1 z =
  let r0 = -euler\_gamma - (sf\_log z)
       tterms = ixiter 2 (z) \lambda k t \rightarrow -t*z/(\#)k
       terms = zipWith (\lambda t k \rightarrow t/(#)k) tterms [1..]
  in ksum (r0:terms)
-- assume n \ge 2, z \le 1
expint_en_series :: (Value v) \Rightarrow Int \rightarrow v \rightarrow v
expint_en_series n z =
  let n' = (\#)n
       res = (-(sf_log z) + (sf_log amma n')) * (-z)^(n-1)/(#)(factorial*n-1) + 1/(n'-1)
       terms' = ixiter 2 (-z) (\lambdam t \rightarrow -t*z/(#)m)
       terms = \max(\lambda(m,t) \to (-t)/(\#)(m-(n-1))) $ filter ((/=(n-1)) \circ fst) $ zip [1..] terms'
  in ksum (res:terms)
-- assume n \ge 2, z > 1
— modified Lentz algorithm
expint_en_contfrac :: (Value v) \Rightarrow Int \rightarrow v \rightarrow v
expint_en_-contfrac n z =
   \mathbf{let} \ \mathrm{fj} = \mathrm{zeta}
       cj = fj
       dj = 0
       j = 1
       n' = (\#)n
  in lentz j cj dj fj
  where
     zeta = 1e-100
     eps = 5e-16
     nz x = if x=0 then zeta else x
     lentz j cj dj fj =
       let aj = (\#)  $ if j=1 then 1 else -(j-1)*(n+j-2)
            bj = z + (\#)(n + 2*(j-1))
            dj' = nz \  bj + aj*dj
            cj' = nz   bj + aj/cj
            dji = 1/dj
            delta = cj '*dji
            fj\ '=fj*delta
       in if (rabs$delta−1)<eps
           then fj ' * sf_exp(-z)
           else lentz (j+1) cj 'dji fj '
```

## 9 AGM

## 9.1 Preamble

```
module AGM (

sf_agm,
sf_agm',
)
where
import Util
```

## 9.2 AGM

Gauss' arithmetic-geometric mean or AGM of two numbers is defined as the limit  $\operatorname{agm}(\alpha, \beta) = \lim_n \alpha_n = \lim_n \beta_n$  where we define

$$\alpha_{n+1} = \frac{\alpha_n + \beta_n}{2}$$

$$\beta_{n+1} = \sqrt{\alpha_n \cdot \beta_n}$$

(Note that we need real values to be positive for this to make sense.)

#### 9.2.1 sf\_agm alpha beta

Here we compute the AGM via the definition and return the full arrays of intermediate values ( $[\alpha_n], [\beta_n], [\gamma_n]$ ), where  $\gamma_n = \frac{\alpha_n - \beta_n}{2}$ . (The iteration converges quadratically so this is an efficient approach.)

```
 \begin{split} & \text{sf\_agm alpha beta} = \text{agm}(\alpha,\beta) \\ & \text{sf\_agm} \ :: \ (\text{Value } v) \ \Rightarrow v \rightarrow v \rightarrow ([v],[v],[v]) \\ & \text{sf\_agm alpha beta} = \text{agm [alpha] [beta] [alpha-beta]} \\ & \text{where agm as@}(a:\_) \ bs@(b:\_) \ cs@(c:\_) = \\ & \text{if } c \Longrightarrow \text{then } (as,bs,cs) \\ & \text{else let } a' = (a+b)/2 \\ & b' = \text{sf\_sqrt } (a*b) \\ & c' = (a-b)/2 \\ & \text{in if } c' \Longrightarrow \text{then } (as,bs,cs) \\ & \text{else agm } (a':as) \ (b':bs) \ (c':cs) \\ \end{split}
```

## 9.2.2 sf\_agm' alpha beta

Here we return simply the value  $sf_agm'$  a b = agm(a, b).

```
sf_agm' z = agm z

sf_agm' :: (Value v) \Rightarrow v \rightarrow v \rightarrow v sf_agm' alpha beta = agm alpha beta ((alpha-beta)/2)

— let (as, -, -) = sf_agm alpha beta in head as

where agm a b 0 = a

agm a b c =

let a' = (a+b)/2

b' = sf_sqrt (a*b)

c' = (a-b)/2

in agm a' b' c'
```

```
sf_agm_c0 :: (Value v) \Rightarrow v \rightarrow v \rightarrow v \rightarrow ([v],[v],[v]) sf_agm_c0 alpha beta c0 = undefined
```

## 10 Airy

The Airy functions Ai and Bi, standard solutions of the ode y'' - zy = 0.

#### 10.1 Preamble

A basic preamble.

```
\begin{array}{ll} \textbf{module} \ Airy \ (sf\_airy\_ai \,, \ sf\_airy\_bi) \ \textbf{where} \\ \textbf{import} \ Gamma \\ \textbf{import} \ Util \end{array}
```

#### 10.2 Ai

## 10.2.1 sf\_airy\_ai z

For now, just use a simple series expansion.

$$\begin{array}{l} sf\_airy\_ai \ :: \ (Value \ v) \ \Rightarrow \ v \ \rightarrow \ v \\ sf\_airy\_ai \ z = airy\_ai\_series \ z \end{array}$$

Initial conditions Ai(0) = 
$$3^{-2/3} \frac{1}{\Gamma(2/3)}$$
 and Ai'(0) =  $-3^{-1/3} \frac{1}{\Gamma(1/3)}$ 

ai0 :: (Value v) 
$$\Rightarrow$$
 v  
ai0 =  $3**(-2/3)/sf_gamma(2/3)$ 

ai'0 :: (Value v) 
$$\Rightarrow$$
 v  
ai'0 =  $-3**(-1/3)/sf_gamma(1/3)$ 

Series expansion, where  $n!!! = \max(n, 1)$  for  $n \leq 2$  and otherwise  $n!!! = n \cdot (n-3)!!!$ :

$$Ai(z) = Ai(0) \left( \sum_{n=0}^{\infty} \frac{(3n-2)!!!}{(3n)!} z^{3n} \right) + Ai'(0) \left( \frac{(3n-1)!!!}{(3n+1)!} z^{3n+1} \right)$$

#### 10.3 Bi

#### 10.3.1 sf\_airy\_bi z

For now, just use a simple series expansion.

```
sf_airy_bi :: (Value v) \Rightarrow v \rightarrow v sf_airy_bi z = airy_bi_series z
```

Initial conditions Bi(0) = 
$$3^{-1/6} \frac{1}{\Gamma(2/3)}$$
 and Bi'(0) =  $3^{1/6} \frac{1}{\Gamma(1/3)}$ 

bi0 :: (Value v) 
$$\Rightarrow$$
 v  
bi0 =  $3**(-1/6)/\text{sf\_gamma}(2/3)$ 

bi'0 :: (Value v) 
$$\Rightarrow$$
 v  
bi'0 =  $3**(1/6)$ /sf\_gamma(1/3)

Series expansion, where  $n!!! = \max(n, 1)$  for  $n \leq 2$  and otherwise  $n!!! = n \cdot (n-3)!!!$ 

$$Bi(z) = Bi(0) \left( \sum_{n=0}^{\infty} \frac{(3n-2)!!!}{(3n)!} z^{3n} \right) + Bi'(0) \left( \frac{(3n-1)!!!}{(3n+1)!} z^{3n+1} \right)$$

```
airy_bi_series z = 
let z3 = z^3 
biterms = ixiter 0 1 $ $\lambda n t \rightarrow t*z3*((\pmu)$3*n+1)/((\pmu)$(3*n+1)*(3*n+2)*(3*n+3)) 
bi'terms = ixiter 0 z $ $\lambda n t \rightarrow t*z3*((\pmu)$3*n+2)/((\pmu)$(3*n+2)*(3*n+3)*(3*n+4)) 
in bi0 * (ksum biterms) + bi'0 * (ksum bi'terms)
```

## 11 Riemann zeta function

## 11.1 Preamble

```
{-# Language BangPatterns #-}
module Zeta (
    sf_zeta,
    sf_zeta_m1,
) where
import Gamma
import Trig
import Util
```

#### 11.2 Zeta

#### 11.2.1 sf\_zeta z

Compute the Riemann zeta function  $sf_zeta z = \zeta(z)$  where

$$\zeta(z) = \sum_{n=1}^{\infty} n^{-z}$$

(for  $\Re z > 1$  and defined by analytic continuation elsewhere).

```
sf_zeta :: (Value v) \Rightarrow v \rightarrow v sf_zeta z | z=1 = (1/0) | (re z)<0 = 2 * (2*pi)**(z-1) * (sf_sin*pi*z/2) * (sf_gamma$1-z) * (sf_zeta$1-z) | otherwise = zeta_series 1.0 z
```

#### 11.2.2 sf\_zeta\_m1 z

For numerical purposes, it is useful to have  $sf_zeta_m1 z = \zeta(z) - 1$ .

#### \*zeta\_series i z

We use the simple series expansion for  $\zeta(z)$  with an Euler-Maclaurin correction:

$$\zeta(z) = \sum_{n=1}^{N} \frac{1}{n^z} + \sum_{k=1}^{p} \cdots$$

```
zeta_series :: (Value v) \Rightarrow v \rightarrow v \rightarrow v
zeta_series !init !z =
  let terms = map (\lambda n \rightarrow ((\#)n)**(-z)) [2..]
      corrs = map correction [2..]
 in summer terms corrs init 0.0 0.0
 where
     -TODO: make general "corrected" kahan_sum!
    summer !(t:ts) !(c:cs) !s !e !r =
      let y = t + e
          !s' = s + y
           !e' = (s - s') + y
          !r' = \dot{s}' + c + e'
      in if r=r' then r'
         else summer ts cs s' e' r'
    !zz1 = z/12
    |zz2 = z*(z+1)*(z+2)/720
    |zz3| = z*(z+1)*(z+2)*(z+3)*(z+4)/30240
    |zz4| = z*(z+1)*(z+2)*(z+3)*(z+4)*(z+5)*(z+6)/1209600
    |zz5| = z*(z+1)*(z+2)*(z+3)*(z+4)*(z+5)*(z+6)*(z+7)*(z+8)/239500800
    correction !n' =
      let n=(#)n'
      in n**(1-z)/(z-1) - n**(-z)/2
         + n**(-z-1)*zz1 - n**(-z-3)*zz2 + n**(-z-5)*zz3
         - n**(-z-7)*zz4 + n**(-z-9)*zz5
```

## 12 Elliptic functions

#### 12.1 Preamble

```
{-# Language BangPatterns #-} module Elliptic where import AGM import Exp import Trig import Util 2^{-2/3}two23 :: Double !two23 = 0.62996052494743658238
```

## 12.2 Elliptic integral of the first kind

Assume that  $1 - \sin^2 \phi$ ,  $1 - k^2 \sin^2 \phi \in \mathbb{C} \setminus (-\infty, 0]$  except that one of them may be 0. The elliptic integral of the first kind is defined by

$$F(\phi, k) = \int_0^{\phi} \frac{d\theta}{\sqrt{1 - k^2 \sin^2 \theta}} = \int_0^{\sin \phi} \frac{dt}{\sqrt{1 - t^2} \sqrt{1 - k^2 t^2}}$$

The complete integral is given by  $\phi = \pi/2$ :

$$K(k) = F(\pi/2, k) =$$

## 12.2.1 sf\_elliptic\_k k

Compute the complete elliptic integral of the first kind K(k) To evaluate this, we use the AGM relation

$$K(k) = \frac{\pi}{2 \operatorname{agm}(1, k')}$$
 where  $k' = \sqrt{1 - k^2}$ 

TODO: UNTESTED!

sf\_elliptic\_

sf\_elliptic\_

#### 12.2.2 sf\_elliptic\_f phi k

Compute the (incomplete) elliptic integral of the first kind  $F(\phi, k)$ . To evaluate, we use an ascending Landen transformation:

$$F(\phi, k) = \frac{2}{1+k}F(\phi_2, k_2) \qquad \text{where } k_2 = \frac{2\sqrt{k}}{1+k} \text{ and } 2\phi_2 = \phi + \arcsin(k\sin\phi)$$

Note that 0 < k < 1 and  $0 < \phi \le \pi/2$  imply  $k < k_2 < 1$  and  $\phi_2 < \phi$ . We iterate this transformation until we reach k = 1 and use the special case

$$F(\phi, 1) = \operatorname{gud}^{-1}(\phi)$$

(Where  $\operatorname{gud}^{-1}(\phi)$  is the inverse Gudermannian function (TODO)). TODO: UNTESTED!

```
sf_elliptic_f phi k = F(\phi, k)
sf_elliptic_f :: Double \rightarrow Double \rightarrow Double
sf_elliptic_f phi k
  k=0 = phi
  k=1 = sf_{\log}((1 + (sf_{\sin} phi)) / (1 - (sf_{\sin} phi))) / 2
           -- quad(@(t)(1/sqrt(1-k^2*sin(t)^2)), 0, phi)
    phi=0 = 0
  | otherwise =
      ascending_landen phi k 1 $ \lambda phi' res' \rightarrow
        res' * sf_{\log}((1 + (sf_{\sin} phi)) / (1 - (sf_{\sin} phi))) / 2
  where
    ascending_landen phi k res kont =
      let k' = 2 * (sf\_sqrt k) / (1 + k)
          phi' = (phi + (asin (k*(sin phi))))/2
           res' = res * 2/(1+k)
      in if k'=1 then kont phi' res
         else ascending_landen phi' k' res' kont
    --function res = agm\_method(phi, k)
      -[an, bn, cn, phin] = sf_agm(1.0, sqrt(1 - k^2), phi, k);
    - res = phin(end) / (2^{(length(phin)-1)} * an(end));
```

## 12.3 Elliptic integral of the second kind

## 12.3.1 sf\_elliptic\_e k

Compute the complete elliptic integral of the second kind E(k). TODO: UNTESTED!

```
\begin{split} & \text{sf\_elliptic\_e } \mathbf{k} = E(k) \\ & \text{sf\_elliptic\_e } \mathbf{:: Double} \to \mathbf{Double} \\ & \text{sf\_elliptic\_e } \mathbf{k} = \\ & \text{let phi} = \mathbf{k} \\ & (as, bs, cs') = \text{sf\_agm } 1.0 \ (\mathbf{sqrt} \ (1.0 - \mathbf{k}^20)) \\ & cs = \mathbf{k} : (\mathbf{tail.reverse} \$ cs') \\ & res = \mathbf{foldl} \ (-) \ 2 \ (\mathbf{map} \ (\lambda(\mathbf{i} \ , c) \to 2^{\hat{}}(\mathbf{i} - 1) \ast c^2) \ (\mathbf{zip} \ [1..] \ cs)) \\ & \mathbf{in res} \ \ast \ \mathbf{pi} / (4 \ast (\mathbf{head} \ as)) \end{split}
```

#### 12.3.2 sf\_elliptic\_e\_ic phi k

Compute the incomplete elliptic integral of the second kind  $E(\phi, k)$  TODO: UNTESTED!

## 12.4 Elliptic integral of the third kind

## 12.4.1 sf\_elliptic\_pi c k

Compute the (in)complete elliptic integral of the third kind  $\Pi(c,k) = \Pi(\pi/2,c,k)$  or  $\Pi(\phi,c,k) = \int_0^\phi dt/((1-c\sin^2(t))\sqrt{(1-k^2\sin^2(t))})$  ( $c=\alpha^2$  in DLMF notation) for real values only  $0< k<1, \ 0< c<1$ . (Could also try numerical quadrature quad(@(t)(1.0/(1-c\*sf\_sin(t)^2)/sqrt(1.0 - k^2\*sf\_sin(t)^2)), 0, phi)). TODO: mostly untested

```
\begin{array}{l} \mathbf{sf\_elliptic\_pi} \ \ \mathbf{c} \ \ \mathbf{k} = \Pi(c,k) \\ \\ \mathbf{sf\_elliptic\_pi} \ \ \mathbf{c} \ \ \mathbf{bouble} \to \mathbf{Double} \to \mathbf{Double} \\ \mathbf{sf\_elliptic\_pi} \ \ \mathbf{c} \ \ \mathbf{k} = \mathbf{complete\_agm} \ \ \mathbf{k} \ \mathbf{c} \\ \mathbf{where} \\ \hline \qquad -\lambda infty < k^2 < 1 \\ \hline \qquad -\lambda infty < c < 1 \\ \mathbf{complete\_agm} \ \ \mathbf{k} \ \mathbf{c} = \\ \mathbf{let} \ \ (\mathbf{ans,gns,\_}) = \mathbf{sf\_agm} \ 1 \ \ (\mathbf{sf\_sqrt} \ \ (1.0 - \mathbf{k}^2)) \\ \mathbf{pn1} = \mathbf{sf\_sqrt} \ \ (1 - \mathbf{c}) \\ \mathbf{qn1} = 1 \\ \mathbf{an1} = \mathbf{last} \ \ \mathbf{ans} \end{array}
```

```
 \begin{array}{l} {\rm gn1} = {\bf last} \ {\rm gns} \\ {\rm en1} = ({\rm pn1^2}2 - {\rm an1*gn1}) \ / \ ({\rm pn1^2}2 + {\rm an1*gn1}) \\ {\bf in} \ {\rm iter} \ {\rm pn1} \ {\rm en1} \ ({\bf reverse} \ {\rm ans}) \ ({\bf reverse} \ {\rm gns}) \ [{\rm qn1}] \\ \\ {\rm iter} \ {\rm pmm1} \ {\rm em1} \ [{\rm an}] \ [{\rm gn}] \ {\rm qns} = {\bf pi}/(4*{\rm an}) \ * \ (2 + c/(1-c)*({\rm ksum} \ {\rm qns})) \\ {\rm iter} \ {\rm pmm1} \ {\rm emm1} \ ({\rm anm1:an:ans}) \ ({\rm gmm1:gn:gns}) \ ({\rm qmm1:qns}) = \\ \\ {\bf let} \ {\rm pn} = ({\rm pnm1^2}2 + {\rm anm1*gmn1})/(2*{\rm pmm1}) \\ {\rm en} = ({\rm pn^2}2 - {\rm an*gn}) \ / \ ({\rm pn^2}2 + {\rm an*gn}) \\ {\rm qn} = {\rm qnm1} \ * \ {\rm enm1/2} \\ \\ {\bf in} \ {\rm iter} \ {\rm pn} \ {\rm en} \ ({\rm an:ans}) \ ({\rm gn:gns}) \ ({\rm qn:qmm1:qns}) \\ \\ \end{array}
```

```
sf_{elliptic_pi_ic_phi} c k = \Pi(\phi, c, k)
sf_elliptic_pi_ic :: Double \rightarrow Double \rightarrow Double \rightarrow Double
sf_{elliptic_pi_i} = 0 c k = 0.0
sf_{elliptic_pi_ic} phi c k = gauss_{transform} k c phi
  where
    gauss\_transform k c phi =
      if (sf_sqrt (1-k^2))=1
      then let cp=sf\_sqrt(1-c)
            in sf_atan(cp*(sf_tan phi)) / cp
      else if (1-k^2/c)=0 - special case else rho below is zero...
      then ((sf_elliptic_e_ic_phi_k) - c*(sf_cos_phi)*(sf_sin_phi) / sqrt(1-c*(sf_sin_phi)^2))/(1-c)
      else let kp = sf\_sqrt (1-k^2)
                k' = (1 - kp) / (1 + kp)
                delta = sf_sqrt(1-k^2*(sf_sin phi)^2)
                psi' = sf_asin((1+kp)*(sf_sin phi) / (1+delta))
                rho = sf_sqrt(1 - (k^2/c))
                c' = c*(1+rho)^2/(1+kp)^2
                xi = (sf_csc phi)^2
                newgt = gauss_transform k' c' psi'
            in (4/(1+kp)*newgt + (rho-1)*(sf_elliptic_f phi k) - (sf_elliptic_rc (xi-1) (xi-c)))/rno
```

## 12.5 Burlisch's elliptic integrals

#### 12.5.1 sf\_elliptic\_cel kc p a b

Compute Burlisch's elliptic integral  $cel(k_c, p, a, b)$  TODO: UNTESTED!

```
\begin{split} & \texttt{sf\_elliptic\_cel kc p a b} = cel(k_c, p, a, b) \\ & \texttt{sf\_elliptic\_cel} :: \mathbf{Double} \to \mathbf{Double} \to \mathbf{Double} \to \mathbf{Double} \to \mathbf{Double} \to \mathbf{Double} \\ & \texttt{sf\_elliptic\_cel kc p a b} = \texttt{a} * (\texttt{sf\_elliptic\_rf 0 (kc^2) 1)} + (\texttt{b\_p*a})/3 * \\ & (\texttt{sf\_elliptic\_rj 0 (kc^2) 1 p)} \end{split}
```

sf\_elliptic\_

#### 12.5.2 sf\_elliptic\_el1 x kc

Compute Burlisch's elliptic integral  $el_1(x, k_c)$  TODO: UNTESTED!

```
\begin{array}{l} \textbf{sf\_elliptic\_el1} \ \ \textbf{k} \ \ \textbf{kc} = el_1(x,k_c) \\ \\ \textbf{sf\_elliptic\_el1} \ \ \vdots \ \ \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \\ \textbf{sf\_elliptic\_el1} \ \ \textbf{k} \ \ \textbf{kc} = \\ \\ --sf\_elliptic\_f \ (atan \ x) \ (sf\_sqrt(1-kc^2)) \\ \textbf{let} \ \ \textbf{r} = 1/x^2 \\ \textbf{in} \ \ \textbf{sf\_elliptic\_rf} \ \ \textbf{r} \ \ (\textbf{r+kc}^2) \ \ (\textbf{r+1}) \\ \end{array}
```

#### 12.5.3 sf\_elliptic\_el2 x kc a b

Compute Burlisch's elliptic integral  $el_2(x, k_c, a, b)$  TODO: UNTESTED!

```
\begin{array}{l} \textbf{sf\_elliptic\_el2} \ \textbf{x} \ \textbf{kc} \ \textbf{a} \ \textbf{b} = el_2(x,k_c,a,b) \\ \\ \textbf{sf\_elliptic\_el2} \ :: \ \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \\ \textbf{sf\_elliptic\_el2} \ x \ \textbf{kc} \ \textbf{a} \ \textbf{b} = \\ \textbf{let} \ \textbf{r} = 1/x^2 \\ \textbf{in} \ \textbf{a} \ * \ (\textbf{sf\_elliptic\_el1} \ x \ \textbf{kc}) \ + \ (\textbf{b-a})/3 \ * \ (\textbf{sf\_elliptic\_rd} \ \textbf{r} \ (\textbf{r+kc}^2) \ (\textbf{r+1})) \\ \end{array}
```

#### 12.5.4 sf\_elliptic\_el3 x kc p

Compute the Burlisch's elliptic integral  $el_3(x, k_c, p)$  TODO: UNTESTED!

```
\begin{array}{l} \textbf{sf\_elliptic\_el3} \ \textbf{x} \ \textbf{kc} \ \textbf{p} = el_3(x,k_c,p) \\ \\ \textbf{sf\_elliptic\_el3} \ :: \ \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \\ \textbf{sf\_elliptic\_el3} \ \textbf{x} \ \textbf{kc} \ \textbf{p} = \\ \\ \textbf{--} \ sf\_elliptic\_pi(atan(x), 1-p, sf\_sqrt(1-kc.^2));} \\ \textbf{let} \ \textbf{r} = 1/x^2 \\ \textbf{in} \ (\textbf{sf\_elliptic\_el1} \ \textbf{x} \ \textbf{kc}) + (1-p)/3 * (\textbf{sf\_elliptic\_rj} \ \textbf{r} \ (\textbf{r+kc}^2) \ (\textbf{r+1}) \ (\textbf{r+p})) \\ \end{array}
```

## 12.6 Symmetric elliptic integrals

#### 12.6.1 sf\_elliptic\_rc x y

Compute the symmetric elliptic integral  $R_C(x, y)$  for real parameters TODO: UNTESTED!

#### 12.6.2 sf\_elliptic\_rd x y z

```
Compute the symmetric elliptic integral R_D(x, y, z) TODO: UNTESTED!
```

```
 \begin{array}{l} ---x,y, \gg 0 \\ \text{sf_elliptic\_rd} & :: \  \, \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \\ \text{sf_elliptic\_rd} & x \ y \ z = \textbf{let} \ (x',s) = (\text{iter} \ x \ y \ z \ 0.0) \ \textbf{in} \ (x'**(-3/2) + s) \\ \textbf{where} \\ \text{iter} & x \ y \ z \ s = \\ & \textbf{let} \ \text{lam} = \textbf{sqrt}(x*y) + \textbf{sqrt}(y*z) + \textbf{sqrt}(z*x); \\ & s' = s + 3/\textbf{sqrt}(z)/(z+\text{lam}); \\ & x' = (x+\text{lam})*\text{two23} \\ & y' = (y+\text{lam})*\text{two23} \\ & z' = (z+\text{lam})*\text{two23} \\ & z' = (z+\text{lam})*\text{two23} \\ & \text{mu} = (x+y+z)/3; \\ & \text{eps} = \textbf{fold11} \ \textbf{max} \ (\textbf{map} \ (\lambda t \rightarrow \textbf{abs}(1-t/\text{mu})) \ [x,y,z]) \\ & \textbf{in} \ \textbf{if} \ \text{eps} < 1e-16 \ \lor \ [x,y,z] = [x',y',z'] \ \textbf{then} \ (x',s') \\ & \text{else} \ \text{iter} \ x' \ y' \ z' \ s' \\ \end{array}
```

#### 12.6.3 sf\_elliptic\_rf x y z

Compute the symmetric elliptic integral of the first kind  $R_F(x,y,z)$  TODO: UNTESTED!

```
 \begin{array}{lll} & -x,y, \gg 0 \\ & \text{sf\_elliptic\_rf} & :: \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \rightarrow \textbf{Double} \\ & \text{sf\_elliptic\_rf} & x \ y \ z = 1/(sf\_sqrt \ \$ \ iter \ x \ y \ z) \\ & \textbf{where} \\ & \text{iter} & x \ y \ z = \\ & \textbf{let} & \text{lam} = (sf\_sqrt \ \$ \ x*y) + (sf\_sqrt \ \$ \ y*z) + (sf\_sqrt \ \$ \ z*x) \\ & & \text{mu} = (x+y+z)/3 \\ & & \text{eps} = \textbf{foldl1 max} \ \$ \ \textbf{map} \ (\lambda a \rightarrow \textbf{abs}(1-a/mu)) \ [x,y,z] \\ & & x' = (x+lam)/4 \\ & y' = (y+lam)/4 \\ & z' = (z+lam)/4 \\ & \textbf{in} & \textbf{if} \ (\text{eps}<1e-16) \ \lor \ ([x,y,z]=[x',y',z']) \\ & \textbf{then} & x \\ & \textbf{else} & \text{iter} & x' \ y' \ z' \end{array}
```

#### 12.6.4 sf\_elliptic\_rg x y z

Compute the symmetric elliptic integral  $R_G(x, y, z)$  TODO: UNTESTED!

```
-- x, y, z > 0
{\tt sf\_elliptic\_rg} \; :: \; \textbf{Double} \, \rightarrow \, \textbf{Double} \, \rightarrow \, \textbf{Double} \, \rightarrow \, \textbf{Double} \, \rightarrow \, \textbf{Double}
sf_elliptic_rg x y z
     x>y = sf_elliptic_rg y x z
     x>z = sf_elliptic_rg z y x
     y>z = sf_elliptic_rg \times z y
    otherwise =
     let !a0 = sqrt (z-x)
           !c0 = \mathbf{sqrt} (y-x)
           !h0 = \mathbf{sqrt} z
           !t0 = \mathbf{sqrt} \ x
           !(an, tn, cn\_sum, hn\_sum) = iter 1 a0 t0 c0 (c0^2/2) h0 0
     in ((t0^2 + theta*cn.sum)*(sf_elliptic_rc (tn^2+theta*an^2) tn^2) + h0 + hn.sum)/2
     where
        theta = 1
        iter n an tn cn cn_sum hn hn_sum =
           let an' = (an + \mathbf{sqrt}(an^2 - cn^2))/2
                tn' = (tn + \mathbf{sqrt}(tn^2 + theta*cn^2))/2
```

```
\begin{array}{l} cn' = cn^2/(2*an')/2 \\ cn\_sum' = cn\_sum + 2^((\#)n-1)*cn'^2 \\ hn' = hn*tn'/sqrt(tn'^2+theta*cn'^2) \\ hn\_sum' = hn\_sum + 2^n*(hn' - hn) \\ n' = n + 1 \\ \textbf{in if } cn^2 = 0 \ \textbf{then } (an,tn,cn\_sum,hn\_sum) \\ \textbf{else } \text{iter } n' \ an' \ tn' \ cn' \ hn\_sum' \ hn' \ hn\_sum' \end{array}
```

#### 12.6.5 sf\_elliptic\_rj x y z p

Compute the symmetric elliptic integral  $R_J(x, y, z, p)$  TODO: UNTESTED!

```
---x, y, z>0
sf_elliptic_rj :: Double \rightarrow Double \rightarrow Double \rightarrow Double \rightarrow Double
sf_elliptic_rj x y z p =
  let (x', smm, scale) = iter x y z p 0.0 1.0
  in scale*x'**(-3/2) + smm
  where
    iter x y z p smm scale =
       let lam = sqrt(x*y) + sqrt(y*z) + sqrt(z*x)
            alpha = p*(\mathbf{sqrt}(x)+\mathbf{sqrt}(y)+\mathbf{sqrt}(z)) + \mathbf{sqrt}(x*y*z)
           beta = \mathbf{sqrt}(p)*(pHam)
           smm' = smm + (if (abs(1 - alpha^2/beta^2) < 5e-16)
                      — optimization to reduce external calls
                      scale*3/alpha;
                    else
                      scale*3*(sf_elliptic_rc (alpha^2) (beta^2))
           mu = (x+y+z+p)/4
           eps = foldl1 max (map (\lambda t \rightarrow abs(1-t/mu)) [x,y,z,p])
           x' = (x+lam)*two23/mu
           y' = (y+lam)*two23/mu
           z^{\,\prime} = (z \!\!+\!\! lam) \!*\! two 23 \!/\! mu
           p' = (p+lam)*two23/mu
           scale' = scale * (mu**(-3/2))
       in if eps<1e-16 \lor [x,y,z,p]=[x',y',z',p'] \lor smm'=smm
          then (x',smm', scale')
          else iter x' y' z' p' smm' scale'
```

# 13 Spence

Spence's integral for  $z \geq 0$  is

$$S(z) = -\int_{1}^{z} \frac{\ln t}{t - 1} dt = -\int_{0}^{z - 1} \frac{\ln(1 + u)}{z} dz$$

and we extend the function via analytic continuation. Spence's function S(z) is related to the dilogarithm function via  $S(z) = \text{Li}_2(1-z)$ .

#### 13.1 Preamble

```
module Spence (
sf_spence,
) where
import Exp
import Util
```

A useful constant pi2\_6 =  $\frac{\pi^2}{6}$ pi2\_6 :: (Value v)  $\Rightarrow$  v pi2\_6 = pi^2/6

## 13.2 sf\_spence z

Compute Spence's integral sf\_spence z = S(z). We use a variety of transformations to to allow efficient computation with a series.

(TODO: this code has not be solidly retested after conversion, especially verify complex.)

#### \*series z

The series expansion used for Spence's integral:

series 
$$\mathbf{z} = -\sum_{k=1}^{\infty} \frac{z^k}{k^2}$$

```
series z = let zk = iterate (*z) z terms = zipWith (\lambda t k \rightarrow -t/(#)k^2) zk [1..] in ksum terms
```

## 14 Lommel functions

## 14.1 Preamble

```
module Lommel (
    sf_lommel_s,
    sf_lommel_s2,
) where
import Util
    -TODO: These are completely untested!
```

## 14.2 First Lommel function

For  $\mu \pm \nu \neq \pm 1, \pm 3, \pm 5, \cdots$  we define the first Lommel function sf\_lommel\_s mu nu  $z = S_{\mu,\nu}(z)$  via series-expansion:

$$S_{\mu,\nu}(z) = \frac{z^{mu+1}}{(\mu+1)^2 - \nu^2} \sum_{k=0}^{\infty} t_k$$

where

$$t_0 = 1$$
  $t_k = t_{k-1} \frac{-z^2}{(\mu + 2k + 1)^2 - \nu^2}$ 

#### 14.2.1 sf\_lommel\_s mu nu z

```
sf_lommel_s mu nu z = let terms = ixiter 1 1.0 $ $\lambda$ k t $\rightarrow -t*z^2$ / ((mu+((#)$2*k+1))^2 - nu^2) res = ksum terms in res * z**(mu+1) / ((mu+1)^2 - nu^2)
```

## 14.3 Second Lommel function

For  $\mu \pm \nu \neq \pm 1, \pm 3, \pm 5, \cdots$  the second Lommel function sf\_lommel\_s2 mu nu  $z = s_{\mu,\nu}(z)$  is given via an asymptotic expansion:

$$s_{\mu,\nu}(z) \sim \sum_{k=0}^{\infty} u_k$$

where

$$u_0 = 1$$
  $u_k = u_{k-1} \frac{-(\mu - 2k + 1)^2 - \nu^2}{z^2}$ 

#### 14.3.1 sf\_lommel\_s2 mu nu z

```
sf_lommel_s2 mu nu z = let tterms = ixiter 1 1.0 $ $\lambda$ k t $\rightarrow -t*((mu-((\#)$2*k+1))^2 - nu^2) / z^2$ terms = tk tterms res = ksum terms in res where tk (a:b:cs) = if (rabs a)<(rabs b) then [a] else a:(tk$b:cs)
```