Computation of Special Functions (Haskell)

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Special functions.

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2.1 Preamble

We start with the basic preamble.

```haskell
{-# Language BangPatterns #-}
{-# Language FlexibleContexts #-}
{-# Language FlexibleInstances #-}
{-# Language TypeFamilies #-}
-- {-# Language UndecidableInstances #-}
module Util where
import Data.Complex
```

2.2 Data Types

We start by defining a convenient type synonym for complex numbers over `Double`.

```haskell
type CDouble = Complex Double
```

Next, we define the `Value` typeclass which is useful for defining our special functions to work over both real (`Double`) values and over complex (`CDouble`) values with uniform implementations. This will also make it convenient for handling `Quad` values (later).

`class Value v` — Value

```haskell
class (Eq v, Floating v, Fractional v, Num v,
       Enum (RealKind v), Eq (RealKind v), Floating (RealKind v),
         Fractional (RealKind v), Num (RealKind v), Ord (RealKind v),
       Eq (ComplexKind v), Floating (ComplexKind v), Fractional (ComplexKind v),
         Num (ComplexKind v)
      ) => Value v where
  type RealKind v :: *
  type ComplexKind v :: *
  re :: v -> (RealKind v)
  im :: v -> (RealKind v)
  rabs :: v -> (RealKind v)
  is_inf :: v -> Bool
  is_nan :: v -> Bool
  fromDouble :: Double -> v
  fromReal :: (RealKind v) -> v
  toComplex :: v -> (ComplexKind v)
```

Both `Double` and `CDouble` are instances of the `Value` typeclass in the obvious ways.

`instance Value Double` — Value Double

```haskell
instance Value Double where
  type RealKind Double = Double
  type ComplexKind Double = CDouble
  re = id
  im = const 0
  rabs = abs
  is_inf = isInfinite
  is_nan = isNaN
```

instance Value Double (cont)   Value Double

```
fromDouble = id
fromReal = id
toComplex x = x :+ 0
```

instance Value CDouble   Value CDouble

```
instance Value CDouble where
  type RealKind CDouble = Double
  type ComplexKind CDouble = CDouble
  re = realPart
  im = imagPart
  rabs = realPart . abs
  is_inf z = (is_inf . re $ z) ∨ (is_inf . im $ z)
  is_nan z = (is_nan . re $ z) ∨ (is_nan . im $ z)
  fromDouble x = x :+ 0
  fromReal x = x :+ 0
  toComplex = id
```

TODO: add quad versions also

2.3 Helper functions

A convenient shortcut, as we often find ourselves converting indices (or other integral values) to our computation type.

```
(#) :: (Integral a, Num b) ⇒ a → b
(#) = fromIntegral
```

A version of `iterate` which passes along an index also (very useful for computing terms of a power-series, for example.)

ixiter i x f   ixiter

```
ixiter :: (Enum ix) ⇒ ix → a → (ix→a→a) → [a]
ixiter i x f = x:(ixiter (succ i) (f i x) f)
```

Computes the relative error in terms of decimal digits, handy for testing. Note that this fails when the exact value is zero.

$$\texttt{relerr e a} = \log_{10}\left|\frac{a-e}{e}\right|$$

```
relerr :: (Value v) ⇒ v → v → (RealKind v)
relerr !exact !approx = re $! logBase 10 (abs ((approx−exact)/exact))
```

2.4 Kahan summation

A useful tool is so-called Kahan summation, based on the observation that in floating-point arithmetic, one can ...

Here `kadd t s e k` is a single step of addition, adding a term to a sum+error and passing the updated sum+error to the continuation.

```
-- kadd value oldsum olderr ---> newsum newerr
kadd :: (Value v) ⇒ v → v → v → (v → v → a) → a
kadd t s e k =
  let y = t − e
      s' = s + y
      e' = (s' − s) − y
  in k s' e'
```

Here `ksum terms` sums a list with Kahan summation. The list is assumed to be (eventually) decreasing and the summation is terminated as soon as adding a term doesn't change the value. (Thus any zeros in the list will immediately terminate the sum.) This is typically used for power-series or asymptotic expansions.

`ksum terms` ksum

```
ksum :: (Value v) ⇒ [v] → v
ksum terms = k 0 0 terms
  where
    k !s !e [] = s
    k !s !e (t:terms) =
      let !y = t − e
          !s' = s + y
          !e' = (s' − s) − y
      in if s' == s
         then s
         else k s' e' terms
```

2.5 Continued fraction evaluation

This is Steed's algorithm for evaluation of a continued fraction

$$C = b_0 + a_1/(b_1 + a_2/(b_2 + a_3/(b_3 + \cdots)))$$

where $C_n = A_n/B_n$ is the partial evaluation up to $\dots a_n/b_n$. Here `steeds as bs` evaluates until $C_n = C_{n+1}$. TODO: describe the algorithm.

```
steeds :: (Value v) ⇒ [v] → [v] → v
steeds (a1:as) (b0:b1:bs) =
    let !c0 = b0
        !d1 = 1/b1
        !delc1 = a1*d1
        !c1 = c0 + delc1
    in recur c1 delc1 d1 as bs
    where recur !cn_1 !delcn_1 !dn_1 !(an:as) !(bn:bs) =
            let !dn = 1/(dn_1*an+bn)
                !delcn = (bn*dn − 1)*delcn_1
                !cn = cn_1 + delcn
            in if (cn == cn_1) ∨ is_nan cn then cn else (recur cn delcn dn as bs)
```

2.6 TO BE MOVED

```
sf_sqrt :: (Value v) ⇒ v → v
sf_sqrt = sqrt
```

3 Fibonacci Numbers

A silly approach to efficient computation of Fibonacci numbers

$$f_n = f_{n-1} + f_{n-2} \qquad f_0 = 0 \qquad f_1 = 1$$

The idea is to use the closed-form solution:

$$f_n = \frac{1}{\sqrt{5}}\left(\frac{1+\sqrt{5}}{2}\right)^n + \frac{-1}{\sqrt{5}}\left(\frac{1-\sqrt{5}}{2}\right)^n$$

and note that we can work in $\mathbb{Q}[\sqrt{5}]$ with terms of the form $a + b\sqrt{5}$ with $a, b \in \mathbb{Q}$ (notice that $\frac{1}{\sqrt{5}} = \frac{\sqrt{5}}{5}$.)

$$\begin{aligned}(a + b\sqrt{5}) + (c + d\sqrt{5}) &= (a + c) + (b + d)\sqrt{5} \\ (a + b\sqrt{5}) * (c + d\sqrt{5}) &= (ac + 5bd) + (ad + bc)\sqrt{5}\end{aligned}$$

We use the `Rational` type to represent elements of $\mathbb{Q}$, which is a bit more than we actually need, as in the computations above the denominator of $\left(\frac{1\pm\sqrt{5}}{2}\right)^n$ is always, in fact, 1 or 2.

```
module Fibo (fibonacci) where
import Data.Ratio
data Q5 = Q5 Rational Rational
  deriving (Eq)
```

The number-theoretic norm $N(a + b\sqrt{5}) = a^2 - 5b^2$, though unused in our application.

```
norm (Q5 ra qa) = ra^2-5*qa^2
```

Human-friendly `Show` instantiation.

```
instance Show Q5 where
  show (Q5 ra qa) = (show ra)++"+"++(show qa)++"*sqrt(5)"
```

Implementation of the operations for typeclasses `Num` and `Fractional`. The `abs` and `signum` functions are unused, so we just give placeholder values.

```
instance Num Q5 where
  (Q5 ra qa) + (Q5 rb qb) = Q5 (ra+rb) (qa+qb)
  (Q5 ra qa) - (Q5 rb qb) = Q5 (ra-rb) (qa-qb)
  (Q5 ra qa) * (Q5 rb qb) = Q5 (ra*rb+5*qa*qb) (ra*qb+rb*qa)
  negate (Q5 ra qa) = Q5 (-ra) (-qa)
  abs a = Q5 (norm a) 0
  signum a@(Q5 ra qa) = if a==0 then 0 else Q5 (ra/(norm a)) (qa/(norm a))
  fromInteger n = Q5 (fromInteger n) 0

instance Fractional Q5 where
  recip a@(Q5 ra qa) = Q5 (ra/(norm a)) (-qa/(norm a))
  fromRational r = (Q5 r 0)
```

Finally, we define $\phi_\pm = \frac{1}{2}(1 \pm \sqrt{5})$ and $c_\pm = \pm\frac{1}{5}\sqrt{5}$ so that $f_n = c_+\phi_+^n + c_-\phi_-^n$. (We can shortcut and extract the value we want without actually computing the full expression.)

```
phip = Q5 (1%2) (1%2)
cp   = Q5 0     (1%5)
phim = Q5 (1%2) (-1%2)
cm   = Q5 0     (-1%5)
fibonacci' n = let (Q5 r q) = cp*phip^n + cm*phim^n in numerator r
fibonacci n = let (Q5 _ q) = phip^^n in numerator (2*q)
```

4 Numbers

4.1 Preamble

```
module Numbers where
import Data.Ratio
import qualified Fibo

fibonacci_number :: Int → Integer
fibonacci_number n = Fibo.fibonacci n

lucas_number :: Int → Integer
lucas_number = undefined

euler_number :: Int → Integer
euler_number = undefined

catalan_number :: Integer → Integer
catalan_number 0 = 1
catalan_number n = 2*(2*n−1)*(catalan_number (n−1))`div`(n+1)

bernoulli_number :: Int → Rational
bernoulli_number = undefined

tangent_number :: Int → Integer
tangent_number = undefined

triangular_number :: Integer → Integer
triangular_number n = n*(n+1)`div`2

factorial :: (Integral a) ⇒ a → a
factorial 0 = 1
factorial 1 = 1
factorial n = product [1..n]

binomial :: (Integral a) ⇒ a → a → a
binomial n k
    | k<0 = 0
    | n<0 = 0
    | k>n = 0
    | k==0 = 1
    | k==n = 1
    | k>n`div`2 = binomial n (n−k)
    | otherwise = (product [n−(k−1)..n]) `div` (product [1..k])
```

4.2 Stirling numbers

```
−− TODO: this is extremely inefficient approach
stirling_number_first_kind n k = s n k
  where s n k | k≤0 ∨ n≤0 = 0
        s n 1 = (−1)^(n−1)*(factorial (n−1))
        s n k = (s (n−1) (k−1)) − (n−1)*(s (n−1) k)

−− TODO: this is extremely inefficient approach
stirling_number_second_kind n k = s n k
  where s n k | k≤0 ∨ n≤0 = 0
        s n 1 = 1
        s n k = k*(s (n−1) k) + (s (n−1) (k−1))
```

5 Exponential & Logarithm

In this section, we implement the exponential function and logarithm function, as well as useful variations.

5.1 Preamble

We begin with a typical preamble.

module Exp

```
{-# Language BangPatterns #-}
{-# Language FlexibleInstances #-}
module Exp (
    sf_exp, sf_expn, sf_exp_m1, sf_exp_m1vx, sf_exp_men, sf_exp_menx,
    sf_log, sf_log_p1,
) where
import Numbers
import Util
```

5.2 Exponential

We start with implementation of the most basic special function, $exp(x)$ or e^x and variations thereof.

5.2.1 `sf_exp x`

For the exponential `sf_exp x` $= \exp(x)$ we use a simple series expansions

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{n!}$$

after first using the identity $e^{-x} = 1/e^x$ to ensure that the real part of the argument is positive. This avoids disastrous cancellation for negative arguments, (though note that for complex arguments this is not sufficient.) TODO: should do range-reduction first... TODO: maybe for complex, use explicit cis?

sf_exp x $= e^x$ sf_exp

```
sf_exp :: (Value v) => v -> v
sf_exp !x
  | is_inf x  = if (re x)<0 then 0 else (1/0)
  | is_nan x  = x
  | (re x)<0  = 1/(sf_exp (-x))
  | otherwise = ksum $ ixiter 1 1.0 $ \n t -> t*x/(#)n
```

5.2.2 `sf_exp_m1 x`

For numerical calculations, it is useful to have `sf_exp_m1 x` $= e^x - 1$ as explicitly calculating this expression will give poor results for x near 1. We use a series expansion for the calculation. Again for negative real part we reflect using $e^{-x} - 1 = -e^{-x}(e^x - 1)$. TODO: should do range-reduction first... TODO: maybe for complex, use explicit cis?

sf_exp_m1 x $= e^x - 1$ sf_exp_m1

```
sf_exp_m1 :: (Value v) ⇒ v → v
sf_exp_m1 !x
  | is_inf x   = if (re x)<0 then −1 else (1/0)
  | is_nan x   = x
  | (re x)<0   = −sf_exp x ∗ sf_exp_m1 (−x)
  | otherwise  = ksum $ ixiter 2 x $ λn t → t∗x/((#)n)
```

5.2.3 sf_exp_m1vx x

Similarly, it is useful to have the scaled variant **sf_exp_m1vx x** $= \frac{e^x-1}{x}$. In this case, we use a continued-fraction expansion

$$\frac{e^x-1}{x} = \frac{2}{2-x+}\,\frac{x^2/6}{1+}\,\frac{x^2/4\cdot 3\cdot 5}{1+}\,\frac{x^2/4\cdot 5\cdot 7}{1+}\,\frac{x^2/4\cdot 7\cdot 9}{1+}\cdots$$

For complex values, simple calculation is inaccurate (when $\Re z \sim 1$).

sf_exp_m1vx x $= \frac{e^x-1}{x}$ sf_exp_m1vx

```
sf_exp_m1vx :: (Value v) ⇒ v → v
sf_exp_m1vx !x
  | is_inf x = if (re x)<0 then 0 else (1/0)
  | is_nan x = x
  | rabs(x)>(1/2) = (sf_exp x − 1)/x −− inaccurate for some complex points
  | otherwise =
      let x2 = x^2
      in 2/(2 − x + x2/6/(1
          + x2/(4∗(2∗3−3)∗(2∗3−1))/(1
          + x2/(4∗(2∗4−3)∗(2∗4−1))/(1
          + x2/(4∗(2∗5−3)∗(2∗5−1))/(1
          + x2/(4∗(2∗6−3)∗(2∗6−1))/(1
          + x2/(4∗(2∗7−3)∗(2∗7−1))/(1
          + x2/(4∗(2∗8−3)∗(2∗8−1))/(1
          ))))))));
```

5.2.4 sf_exp_menx n x

Compute the scaled tail of series expansion of the exponential function.

$$\texttt{sf\_exp\_menx n x} = \frac{n!}{x^n}\left(e^z - \sum_{k=0}^{n-1}\frac{x^k}{k!}\right) = \frac{n!}{x^n}\sum_{k=n}^{\infty}\frac{x^k}{k!} = n!\sum_{k=0}^{\infty}\frac{x^k}{(k+n)!}$$

We use a continued fraction expansion and using the modified Lentz algorithm for evaluation.

```
sf_exp_menx :: (Value v) ⇒ Int → v → v
sf_exp_menx 0 z = sf_exp z
sf_exp_menx 1 z = sf_exp_m1vx z
sf_exp_menx n z
  | is_inf z   = if (re z)>0 then (1/0) else (0) −− TODO: verify
```

```
  | is_nan z  = z
  | otherwise = exp_menx__contfrac n z
  where
    !zeta = 1e-150
    !eps = 1e-16
    nz !z = if z==0 then zeta else z
    exp_menx__contfrac n z =
      let !fj = (#)$ n+1
          !cj = fj
          !dj = 0
          !j  = 1
      in lentz j dj cj fj
    lentz !j !dj !cj !fj =
      let !aj = if (odd j)
                then z*((#)$(j+1)`div`2)
                else -z*((#)$(n+(j`div`2)))
          bj = (#)$n+1+j
          !dj' = nz$ bj + aj*dj
          !cj' = nz$ bj + aj/cj
          !dji = 1/dj'
          !deltaj = cj'*dji
          !fj' = fj*deltaj
      in if (rabs(deltaj-1)<eps)
         then 1/(1-z/fj')
         else lentz (j+1) dji cj' fj'
```

5.2.5 `sf_exp_men n x`

This is the generalization of `sf_exp_m1 x`, giving the tail of the series expansion of the exponential function, for $n = 0, 1, \ldots$.

$$\texttt{sf\_exp\_men n z} = e^z - \sum_{k=0}^{n-1} \frac{z^k}{k!} = \sum_{k=n}^{\infty} \frac{z^k}{k!}$$

The special cases are: $n = 0$ gives $e^x =$ `sf_exp x` and $n = 1$ gives $e^x - 1 =$ `sf_exp_m1 x`. We compute this by calling the scaled version `sf_exp_menx` and rescaling back.

```
-- ($n=0, 1, 2, ∘ ..$)
sf_exp_men :: (Value v) => Int -> v -> v
sf_exp_men !n !x = (sf_exp_menx n x) * x^n / ((#)$factorial n)
```

5.2.6 `sf_expn n x`

```
-- Compute initial part of series for exponential, $\sum_(k=0)^n z^k/k!$
-- ($n=0,1,2,...$)
sf_expn :: (Value v) => Int -> v -> v
sf_expn n z
  | is_inf z  = if (re z)>0 then (1/0) else (if (odd n) then (-1/0) else (1/0))
  | is_nan z  = z
  | otherwise = expn__series n z
  where
    -- TODO: just call sf_exp when possible
    -- TODO: better handle large -ve values!
    expn__series :: (Value v) => Int -> v -> v
    expn__series n z = ksum $ take (n+1) $ ixiter 1 1.0 $ \k t -> t*z/(#)k
```

5.3 Logarithm

5.3.1 sf_log x

We simply use the built-in implementation (from the `Floating` typeclass).

```
sf_log :: (Value v) ⇒ v → v
sf_log = log
```

5.3.2 sf_log_p1 x

The accuracy preserving `sf_log_p1 x` $= \ln 1 + x$. For values close to zero, we use a power series expansion

$$\ln(1+x) = 2\sum_{n=0}^{\infty} \frac{(\frac{x}{x+2})^{2n+1}}{2n+1}$$

and otherwise just compute it directly.

```
sf_log_p1 :: (Value v) ⇒ v → v
sf_log_p1 !z
  | is_nan z = z
  | (rabs z)>0.25 = sf_log (1+z)
  | otherwise = series z
  where
    series z =
      let !r = z/(z+2)
          !zr2 = r^2
          !tterms = iterate (*zr2) (r*zr2)
          !terms = zipWith (λn t → t/((#)$2*n+1)) [1..] tterms
      in 2*(ksum (r:terms))
```

A simple continued fraction implementation for $\ln 1 + z$

$$\ln(1+z) = z/(1+z/(2+z/(3+4z/(4+4z/(5+9z/(6+9z/(7+\cdots)))))))$$

Though unused for now, it seems to have decent convergence properties.

```
ln_1_z_cf z = steeds (z:(ts 1)) [0..]
  where ts n = (n^2*z):(n^2*z):(ts (n+1))
```

6 Gamma

6.1 Preamble

A basic preamble.

```
module Gamma (
    euler_gamma,
    factorial,
    sf_beta,
    sf_gamma,
    sf_invgamma,
    sf_lngamma,
    sf_digamma,
    bernoulli_b,
    )
where
import Exp
import Numbers(factorial)
import Trig
import Util
```

6.2 Misc

6.2.1 `euler_gamma`

A constant for Euler's gamma:

$$\gamma = \lim_{n\to\infty}\left(\sum_{k=1}^{n}\frac{1}{n} - \ln n\right)$$

```
euler_gamma :: (Floating a) ⇒ a
euler_gamma = 0.5772156649015328606065120900824024310421593359399235988057672348848677267776646709369470632917467495
```

6.2.2 `sf_beta a b`

The Beta integral

$$B(a,b) = \int_0^1 t^{a-1}(1-t)^{b-1}\,dt = \frac{\Gamma(a)\Gamma(b)}{\Gamma(a+b)}$$

implemented in terms of log-gamma

$$\texttt{sf\_beta a b} = e^{\ln\Gamma(a)+\ln\Gamma(b)-\ln\Gamma(a+b)}$$

```
sf_beta :: (Value v) ⇒ v → v → v
sf_beta a b = sf_exp $ (sf_lngamma a) + (sf_lngamma b) − (sf_lngamma$a+b)
```

6.3 Gamma

The gamma function

$$\Gamma(z) = \int_0^\infty e^{-t}t^z\,\frac{dz}{z}$$

6.3.1 `sf_gamma z`

The gamma function implemented using the identity $\Gamma(z) = \frac{1}{z}\Gamma(z+1)$ to increase the real part of the argument to be > 15 and then using an asymptotic expansion for log-gamma, `lngamma_asymp`, to evaluate.

`sf_gamma x` $= \Gamma(x)$ `sf_gamma`

```
sf_gamma :: (Value v) ⇒ v → v
sf_gamma x =
  redup x 1 $ λ x' t → t ∗ (sf_exp (lngamma_asymp x'))
  where redup x t k
          | (re x)>15 = k x t
          | otherwise = redup (x+1) (t/x) k
```

6.3.2 `*lngamma_asymp z`

The asymptotic expansion for log-gamma

$$\ln\Gamma(z) \sim (z-\frac{1}{2})\ln z - z + \frac{1}{2}\ln(2\pi) + \sum_{k=1}^{\infty}\frac{B_{2k}}{2k(2k-1)z^{2k-1}}$$

where B_n is the n'th Bernoulli number.

```
lngamma_asymp :: (Value v) ⇒ v → v
lngamma_asymp z = (z − 1/2)∗(sf_log z) − z + (1/2)∗sf_log(2∗pi) + (ksum terms)
  where terms = [b2k/(2∗k∗(2∗k−1)∗z^(2∗k'−1)) | k'←[1..10], let k=(#)k', let b2k=bernoulli_b$2∗k']
```

6.3.3 sf_invgamma z

The inverse gamma function, `sf_invgamma z` $= \frac{1}{\Gamma(z)}$.

```
sf_invgamma :: (Value v) ⇒ v → v
sf_invgamma x =
  let (x',t) = redup x 1
      lngx = lngamma_asymp x'
  in t * (sf_exp$ −lngx)
  where redup x t
          | (re x)>15 = (x,t)
          | otherwise = redup (x+1) (t*x)
```

6.3.4 sf_lngamma z

The log-gamma function, `sf_lngamma z` $= \ln\Gamma(z)$.

```
sf_lngamma :: (Value v) ⇒ v → v
sf_lngamma x =
  let (x',t) = redup x 0
      lngx = lngamma_asymp x'
  in t + lngx
  where redup x t
          | (re x)>15 = (x,t)
          | otherwise = redup (x+1) (t−sf_log x)
```

6.3.5 bernoulli_b n

The Bernoulli numbers, B_n. A simple hard-coded table, for now. (Should be moved to Numbers module and general, cached, implementation done.)

```
bernoulli_b :: (Value v) ⇒ Int → v
bernoulli_b 1 = −1/2
bernoulli_b k | k`mod`2==1 = 0
bernoulli_b 0 = 1
bernoulli_b 2 = 1/6
bernoulli_b 4 = −1/30
bernoulli_b 6 = 1/42
bernoulli_b 8 = −1/30
bernoulli_b 10 = 5/66
bernoulli_b 12 = −691/2730
bernoulli_b 14 = 7/6
bernoulli_b 16 = −3617/510
bernoulli_b 18 = 43867/798
bernoulli_b 20 = −174611/330
bernoulli_b _ = undefined
```

Spouge's approximation to the gamma function

In tests, this gave disappointing results.

```
-- Spouge's approximation (a=17?)
spouge_approx :: (Value v) ⇒ Int → v → v
spouge_approx a z' =
  let z = z' − 1
      a' = (#)a
      res = (z+a')**(z+(1/2)) * sf_exp (−(z+a'))
      sm = fromDouble$sf_sqrt(2*pi)
      terms = [(spouge_c k a') / (z+k') | k←[1..(a−1)], let k' = (#)k]
```

```
    smm = sm + ksum terms
 in res*smm
 where
   spouge_c k a = ((if k`mod`2==0 then -1 else 1) / ((#) $ factorial (k-1)))
                   * (a-((#)k))**(((#)k)-1/2) * sf_exp(a-((#)k))
```

6.4 Digamma

The digamma function

$$\psi(z) = \frac{d}{dz} \ln \Gamma(z) = \frac{\Gamma'(z)}{\Gamma(z)}$$

6.4.1 `sf_digamma z`

We implement with a series expansion for $|z| <= 10$ and otherwise with an asymptotic expansion.

```
sf_digamma :: (Value v) => v -> v
--sf_digamma n | is_nonposint n = Inf
sf_digamma z | (rabs z)>10 = digamma__asympt z
             | otherwise   = digamma__series z
```

The series expansion is the following

$$\psi(z) = -\gamma - \frac{1}{z} + \sum_{k=1}^{\infty} \frac{z}{k(k+z)}$$

but with Euler-Maclaurin correction terms:

$$\psi(z) = -\gamma - \frac{1}{z} + \sum_{k=1}^{n} \frac{z}{k(k+z)} + (\ln \frac{k+z}{k} - \frac{z}{2k(k=z)} + \sum_{j=1}^{p} B_{2j}(k^{-2j} - (k+z)^{-2j})$$

```
digamma__series :: (Value v) => v -> v
digamma__series z =
 let res = -euler_gamma - (1/z)
     terms = map (\k->z/((#)k*(z+(#)k))) [1..]
     corrs = map (correction.(#)) [1..]
 in summer res res terms corrs
 where
   summer :: (Value v) => v -> v -> [v] -> [v] -> v
   summer res sum (t:terms) (c:corrs) =
     let sum' = sum + t
         res' = sum' + c
     in if res==res' then res
        else summer res' sum' terms corrs
   bn1 = bernoulli_b 2
   bn2 = bernoulli_b 4
   bn3 = bernoulli_b 6
   bn4 = bernoulli_b 8
   correction k =
     (sf_log$(k+z)/k) - z/2/(k*(k+z))
       + bn1*(k^^(-2) - (k+z)^^(-2))
       + bn2*(k^^(-4) - (k+z)^^(-4))
       + bn3*(k^^(-6) - (k+z)^^(-6))
       + bn4*(k^^(-8) - (k+z)^^(-8))
```

The asymptotic expansion (valid for $|argz| < \pi$) is the following

$$\psi(z) \sim \ln z - \frac{1}{2z} + \sum_{k=1}^{\infty} \frac{B_{2k}}{2kz^{2k}}$$

Note that our implementation will fail if the `bernoulli_b` table is exceeded. If $\Re z < \frac{1}{2}$ then we use the reflection identity to ensure $\Re z \geq \frac{1}{2}$:

$$\psi(z) - \psi(1-z) = \frac{-\pi}{\tan(\pi z)}$$

```
digamma__asympt :: (Value v) ⇒ v → v
digamma__asympt z
  | (re z)<0.5 = compute (1 − z) $ −pi/(sf_tan(pi∗z)) + (sf_log(1−z)) − 1/(2∗(1−z))
  | otherwise  = compute z $ (sf_log z) − 1/(2∗z)
   where
     compute z res =
       let z_2 = z^^(−2)
           zs = iterate (∗z_2) z_2
           terms = zipWith (λn z2n → z2n∗(bernoulli_b(2∗n+2))/(#)(2∗n+2)) [0..] zs
       in sumit res res terms
     sumit res ot (t:terms) =
       let res' = res − t
       in if res==res' ∨ (rabs t)>(rabs ot)
          then res
          else sumit res' t terms
```

7 Error function

7.1 Preamble

```
{−# Language BangPatterns #−}
module Erf (
     sf_erf,
     sf_erfc,
) where
import Exp
import Util
```

7.2 Error function

7.2.1 `sf_erf z`

The error function `sf_erf z` $= \operatorname{erf} z$ where

$$\operatorname{erf} z = \frac{2}{\sqrt{\pi}} \int_{-\infty}^{z} e^{-x^2}\, dx$$

For $\Re z < -1$, we transform via $\operatorname{erf} z = -\operatorname{erf}(-z)$ and for $|z| < 1$ we use the power-series expansion, otherwise we use $\operatorname{erf} z = 1 - \operatorname{erfc} z$. (TODO: this implementation is not perfect, but workable for now.)

```
sf_erf :: (Value v) ⇒ v → v
sf_erf z
  | (re z)<(−1) = −sf_erf(−z)
  | (rabs z)<1  = erf_series z
  | otherwise   = 1 − sf_erfc z
```

7.2.2 `sf_erfc z`

The complementary error-function `sf_erfc z` $= \operatorname{erfc} z$ where

$$\operatorname{erfc} z = 1 - \operatorname{erf} z = \frac{2}{\sqrt{\pi}} \int_{z}^{\infty} e^{-x^2}\, dx$$

For $\Re z < -1$ we transform via $\operatorname{erfc} z = 2 - \operatorname{erf}(-z)$ and if $|z| < 1$ then we use $\operatorname{erfc} z = 1 - \operatorname{erf} z$. Finally, if $|z| < 10$ we use a continued-fraction expansion and an asymptotic expansion otherwise. (TODO: there are a few issues with this implementation: For pure imaginary values and for extremely large values it seems to hang.)

```
-- infinite loop when (re z)==0
sf_erfc :: (Value v) => v -> v
sf_erfc z
  | (re z)<(-1) = 2-(sf_erfc (-z))
  | (rabs z)<1  = 1-(sf_erf z)
  | (rabs z)<10 = erfc_cf_pos1 z
  | otherwise   = erfc_asymp_pos z -- TODO: hangs for very large input
```

erf_series z

The series expansion for $\operatorname{erf} z$:

$$\operatorname{erf} z = \frac{2}{\sqrt{\pi}} \sum_{n=0}^{\infty} \frac{(-)^n z^{2n+1}}{n!(2n+1)}$$

There is an alternative expansion $\operatorname{erf} z = \frac{2}{\sqrt{\pi}} e^{-z^2} \sum_{n=0}^{\infty} \frac{2^n z^{2n+1}}{1\cdot 3\cdots(2n+1)}$, but we don't use it here. (TODO: why not?)

```
erf_series z =
  let z2 = z^2
      rts = ixiter 1 z $ \n t -> (-t)*z2/(#)n
      terms = zipWith (\ n t ->t/(#)(2*n+1)) [0..] rts
  in (2/sf_sqrt pi) * ksum terms
```

*sf_erf z

This asymptotic expansion for $\operatorname{erfc} z$ is valid as $z \to +\infty$:

$$\operatorname{erfc} z \sim \frac{e^{-z^2}}{\sqrt{\pi}} \sum_{n=0}^{\infty} (-)^n \frac{(1/2)_m}{z^{2m+1}}$$

where the Pochhammer symbol $(1/2)_m$ is given by:

$$\left(\frac{1}{2}\right)_m = \frac{1\cdot 3\cdot 5\cdots(2m-1)}{2^m} = \frac{(2m)!}{m!2^{2m}}$$

TODO: correct the asymptotic term checking (not smallest but pre-smallest term).

```
erfc_asymp_pos z =
  let z2 = z^2
      iz2 = 1/2/z2
      terms = ixiter 1 (1/z) $ \n t -> (-t*iz2)*(#)(2*n-1)
      tterms = tk terms
  in (sf_exp (-z2))/(sqrt pi) * ksum tterms
  where tk (a:b:cs) = if (rabs a)<(rabs b) then [a] else a:(tk$b:cs)
```

*erfc_cf_pos1 z

A continued-fraction expansion for $\operatorname{erfc} z$:

$$\sqrt{\pi} e^{z^2} \operatorname{erfc} z = \frac{z}{z^2+} \frac{1/2}{1+} \frac{1}{z^2+} \frac{3/2}{1+} \cdots$$

```
erfc_cf_pos1 z =
  let z2 = z^2
      as = z:(map fromDouble [1/2,1..])
      bs = 0:cycle [z2,1]
      cf = steeds as bs
  in sf_exp(-z2) / (sqrt pi) * cf
```

***erfc_cf_pos1 z**

This is an alternative continued-fraction expansion.

$$\sqrt{\pi}e^{z^2}\operatorname{erfc} z = \frac{2z}{2z^2+1-}\,\frac{1\cdot 2}{2z^2+5-}\,\frac{3\cdot 4}{2z^2+9-}\cdots$$

Unused for now.

```
erfc_cf_pos2 z =
  let z2 = z^2
      as = (2*z):(map (\n->(#)$ -(2*n+1)*(2*n+2)) [0..])
      bs = 0:(map (\n->2*z2+(#)4*n+1) [0..])
      cf = steeds as bs
  in sf_exp(-z2) / (sqrt pi) * cf
```

8 Exponential Integral

8.1 Preamble

```
module ExpInt(
  sf_expint_ei,
  sf_expint_en,
  )
where
import Exp
import Gamma
import Util
```

8.2 Exponential integral Ei

The exponential integral $\mathrm{Ei}\, z$ is defined for $x < 0$ by

$$\mathrm{Ei}(z) = -\int_{-x}^{\infty} \frac{e^{-t}}{t}\, dt$$

It can be defined

8.2.1 sf_expint_ei z

We give only an implementation for $\Re z \geq 0$. We use a series expansion for $|z| < 40$ and an asymptotic expansion otherwise.

sf_expint_ei z $= \mathrm{Ei}(z)$ sf_expint_ei

```
sf_expint_ei :: (Value v) => v -> v
sf_expint_ei z
  | (re z) < 0.0   = (0/0)  -- (NaN)
  | z == 0.0       = (-1/0) -- (-Inf)
  | (rabs z) < 40  = expint_ei__series z
  | otherwise      = expint_ei__asymp z
```

sf_expint_ei z = Ei(z) (cont)

sf_expint_ei

The series expansion is given (for $x > 0$)

$$\mathrm{Ei}(x) = \gamma + \ln x + \sum_{n=1}^{\infty} \frac{x^n}{n!n}$$

We evaluate the addition of the two terms with the sum slightly differently when $\Re z < 1/2$ to reduce floating-point cancellation error slightly.

expint_ei__series z

expint_ei__series

```
expint_ei__series :: (Value v) ⇒ v → v
expint_ei__series z =
  let tterms = ixiter 2 z $ λn t → t*z/(#)n
      terms = zipWith (λ t n →t/(#)n) tterms [1..]
      res = ksum terms
  in if (re z)<0.5
     then sf_log(z * sf_exp(euler_gamma + res))
     else res + sf_log(z) + euler_gamma
```

The asymptotic expansion as $x \to +\infty$ is

$$\mathrm{Ei}(x) \sim \frac{e^x}{x} \sum_{n=0}^{\infty} \frac{n!}{x^n}$$

expint_ei__asymp z

expint_ei__asymp

```
expint_ei__asymp :: (Value v) ⇒ v → v
expint_ei__asymp z =
  let terms = tk $ ixiter 1 1.0 $ λn t → t/z*(#)n
      res = ksum terms
  in res * (sf_exp z) / z
  where tk (a:b:cs) = if (rabs a)<(rabs b) then [a] else a:(tk$b:cs)
```

8.3 Exponential integral E_n

The exponential integrals $E_n(z)$ are defined as

$$E_n(z) = z^{n-1} \int_z^{\infty} \frac{e^{-t}}{t^n}\, dt$$

They satisfy the following relations:

$$\begin{aligned} E_0(z) &= \frac{e^{-z}}{z} \\ E_{n+1}(z) &= \int_z^{\infty} E_n(t)\, dt \end{aligned}$$

And they can be expressed in terms of incomplete gamma functions:

$$E_n(z) = z^{n-1}\Gamma(1-n, z)$$

(which also gives a generalization for non-integer n).

8.3.1 sf_expint_en n z

sf_expint_en n z $= E_n(z)$ sf_expint_en

```
sf_expint_en :: (Value v) ⇒ Int → v → v
sf_expint_en n z | (re z)<0 = (0/0) −− (NaN) TODO: confirm this
                 | z == 0   = (1/(#)(n−1)) −− TODO: confirm this
sf_expint_en 0 z = sf_exp(−z) / z
sf_expint_en 1 z = expint_en__1 z
sf_expint_en n z | (rabs z) ≤ 1.0 = expint_en__series n z
                 | otherwise = expint_en__contfrac n z
```

We use this series expansion for $E_1(z)$:

$$E_1(z) = -\gamma - \ln z + \sum_{k=1}^{\infty}(-)^k \frac{z^k}{k!k}$$

(Note that this will not be good for large values of z.)

```
expint_en__1 :: (Value v) ⇒ v → v
expint_en__1 z =
  let r0 = −euler_gamma − (sf_log z)
      tterms = ixiter 2 (z) $ λk t → −t*z/(#)k
      terms = zipWith (λ t k → t/(#)k) tterms [1..]
  in ksum (r0:terms)

−− assume n≥ 2, z≤1
expint_en__series :: (Value v) ⇒ Int → v → v
expint_en__series n z =
  let n' = (#)n
      res = (−(sf_log z) + (sf_digamma n')) * (−z)^(n−1)/(#)(factorial$n−1) + 1/(n'−1)
      terms' = ixiter 2 (−z) (λm t → −t*z/(#)m)
      terms = map (λ(m,t)→(−t)/(#)(m−(n−1))) $ filter ((/=(n−1)) ◦ fst) $ zip [1..] terms'
  in ksum (res:terms)

−− assume n≥ 2, z>1
−− modified Lentz algorithm
expint_en__contfrac :: (Value v) ⇒ Int → v → v
expint_en__contfrac n z =
  let fj = zeta
      cj = fj
      dj = 0
      j = 1
      n' = (#)n
  in lentz j cj dj fj
  where
    zeta = 1e−100
    eps = 5e−16
    nz x = if x==0 then zeta else x
    lentz j cj dj fj =
```

```
let aj = (#) $ if j==1 then 1 else -(j-1)*(n+j-2)
    bj = z + (#)(n + 2*(j-1))
    dj' = nz $ bj + aj*dj
    cj' = nz $ bj + aj/cj
    dji = 1/dj'
    delta = cj'*dji
    fj' = fj*delta
in if (rabs$delta-1)<eps
   then fj' * sf_exp(-z)
   else lentz (j+1) cj' dji fj'
```

9 AGM

9.1 Preamble

module AGM

```
module AGM (
    sf_agm,
    sf_agm',
    )
where
import Util
```

9.2 AGM

Gauss' arithmetic-geometric mean or AGM of two numbers is defined as the limit $\mathrm{agm}(\alpha,\beta) = \lim_n \alpha_n = \lim_n \beta_n$ where we define

$$\begin{aligned} \alpha_{n+1} &= \frac{\alpha_n+\beta_n}{2} \\ \beta_{n+1} &= \sqrt{\alpha_n\cdot\beta_n} \end{aligned}$$

(Note that we need real values to be positive for this to make sense.)

9.2.1 sf_agm alpha beta

Here we compute the AGM via the definition and return the full arrays of intermediate values $([\alpha_n],[\beta_n],[\gamma_n])$, where $\gamma_n = \frac{\alpha_n-\beta_n}{2}$. (The iteration converges quadratically so this is an efficient approach.)

sf_agm alpha beta $= \mathrm{agm}(\alpha,\beta)$ sf_agm

```
sf_agm :: (Value v) => v -> v -> ([v],[v],[v])
sf_agm alpha beta = agm [alpha] [beta] [alpha-beta]
  where agm as@(a:_) bs@(b:_) cs@(c:_) =
          if c==0 then (as,bs,cs)
          else let a' = (a+b)/2
                   b' = sf_sqrt (a*b)
                   c' = (a-b)/2
               in if c'==c then (as,bs,cs)
                  else agm (a':as) (b':bs) (c':cs)
```

9.2.2 `sf_agm' alpha beta`

Here we return simply the value `sf_agm' a b` $= \text{agm}(a, b)$.

`sf_agm' z` $= \text{agm}\, z$

```
sf_agm' :: (Value v) => v -> v -> v
sf_agm' alpha beta = agm alpha beta ((alpha-beta)/2)
  --let (as,_,_) = sf_agm alpha beta in head as
  where agm a b 0 = a
        agm a b c =
          let a' = (a+b)/2
              b' = sf_sqrt (a*b)
              c' = (a-b)/2
          in agm a' b' c'
```

```
sf_agm_c0 :: (Value v) => v -> v -> v -> ([v],[v],[v])
sf_agm_c0 alpha beta c0 = undefined
```

10 Airy

The Airy functions Ai and Bi, standard solutions of the ode $y'' - zy = 0$.

10.1 Preamble

A basic preamble.

```
module Airy (sf_airy_ai, sf_airy_bi) where
import Gamma
import Util
```

10.2 Ai

10.2.1 `sf_airy_ai z`

For now, just use a simple series expansion.

```
sf_airy_ai :: (Value v) => v -> v
sf_airy_ai z = airy_ai_series z
```

Initial conditions $\text{Ai}(0) = 3^{-2/3}\frac{1}{\Gamma(2/3)}$ and $\text{Ai}'(0) = -3^{-1/3}\frac{1}{\Gamma(1/3)}$

```
ai0 :: (Value v) => v
ai0 = 3**(-2/3)/sf_gamma(2/3)

ai'0 :: (Value v) => v
ai'0 = -3**(-1/3)/sf_gamma(1/3)
```

Series expansion, where $n!!! = \max(n, 1)$ for $n \le 2$ and otherwise $n!!! = n \cdot (n-3)!!!$:

$$\text{Ai}(z) = \text{Ai}(0)\left(\sum_{n=0}^{\infty} \frac{(3n-2)!!!}{(3n)!} z^{3n}\right) + \text{Ai}'(0)\left(\frac{(3n-1)!!!}{(3n+1)!} z^{3n+1}\right)$$

```
airy_ai_series z =
    let z3 = z^3
        aiterms  = ixiter 0 1 $ \n t -> t*z3*((#)$3*n+1)/((#)$(3*n+1)*(3*n+2)*(3*n+3))
        ai'terms = ixiter 0 z $ \n t -> t*z3*((#)$3*n+2)/((#)$(3*n+2)*(3*n+3)*(3*n+4))
    in ai0 * (ksum aiterms) + ai'0 * (ksum ai'terms)
```

10.3 Bi

10.3.1 `sf_airy_bi z`

For now, just use a simple series expansion.

```
sf_airy_bi :: (Value v) => v -> v
sf_airy_bi z = airy_bi_series z
```

Initial conditions $\mathrm{Bi}(0) = 3^{-1/6}\frac{1}{\Gamma(2/3)}$ and $\mathrm{Bi}'(0) = 3^{1/6}\frac{1}{\Gamma(1/3)}$

```
bi0 :: (Value v) => v
bi0 = 3**(-1/6)/sf_gamma(2/3)

bi'0 :: (Value v) => v
bi'0 = 3**(1/6)/sf_gamma(1/3)
```

Series expansion, where $n!!! = \max(n, 1)$ for $n \leq 2$ and otherwise $n!!! = n \cdot (n-3)!!!$:

$$\mathrm{Bi}(z) = \mathrm{Bi}(0)\left(\sum_{n=0}^{\infty}\frac{(3n-2)!!!}{(3n)!}z^{3n}\right) + \mathrm{Bi}'(0)\left(\frac{(3n-1)!!!}{(3n+1)!}z^{3n+1}\right)$$

```
airy_bi_series z =
    let z3 = z^3
        biterms  = ixiter 0 1 $ \n t -> t*z3*((#)$3*n+1)/((#)$(3*n+1)*(3*n+2)*(3*n+3))
        bi'terms = ixiter 0 z $ \n t -> t*z3*((#)$3*n+2)/((#)$(3*n+2)*(3*n+3)*(3*n+4))
    in bi0 * (ksum biterms) + bi'0 * (ksum bi'terms)
```

11 Riemann zeta function

11.1 Preamble

```
{-# Language BangPatterns #-}
module Zeta (
  sf_zeta,
  sf_zeta_m1,
) where
import Gamma
import Trig
import Util
```

11.2 Zeta

11.2.1 `sf_zeta z`

Compute the Riemann zeta function `sf_zeta z` $= \zeta(z)$ where

$$\zeta(z) = \sum_{n=1}^{\infty} n^{-z}$$

(for $\Re z > 1$ and defined by analytic continuation elsewhere).

```
sf_zeta :: (Value v) ⇒ v → v
sf_zeta z
  | z==1       = (1/0)
  | (re z)<0   = 2 * (2*pi)**(z-1) * (sf_sin $ pi*z/2) * (sf_gamma $ 1-z) * (sf_zeta $ 1-z)
  | otherwise  = zeta_series 1.0 z
```

11.2.2 `sf_zeta_m1 z`

For numerical purposes, it is useful to have `sf_zeta_m1 z` $= \zeta(z) - 1$.

```
sf_zeta_m1 :: (Value v) ⇒ v → v
sf_zeta_m1 z
  | z==1       = (1/0)
  | (re z)<0   = 2 * (2*pi)**(z-1) * (sf_sin $ pi*z/2) * (sf_gamma $ 1-z) * (sf_zeta $ 1-z) - 1  -- TODO:
  | otherwise  = zeta_series 0.0 z
```

*`zeta_series i z`

We use the simple series expansion for $\zeta(z)$ with an Euler-Maclaurin correction:

$$\zeta(z) = \sum_{n=1}^{N} \frac{1}{n^z} + \sum_{k=1}^{p} \cdots$$

```
zeta_series :: (Value v) ⇒ v → v → v
zeta_series !init !z =
  let terms = map (\n→((#)n)**(-z)) [2..]
      corrs = map correction [2..]
  in summer terms corrs init 0.0 0.0
  where
    --TODO: make general "corrected" kahan_sum!
    summer !(t:ts) !(c:cs) !s !e !r =
      let !y = t + e
          !s' = s + y
          !e' = (s - s') + y
          !r' = s' + c + e'
      in if r==r' then r'
         else summer ts cs s' e' r'
    !zz1 = z/12
    !zz2 = z*(z+1)*(z+2)/720
    !zz3 = z*(z+1)*(z+2)*(z+3)*(z+4)/30240
    !zz4 = z*(z+1)*(z+2)*(z+3)*(z+4)*(z+5)*(z+6)/1209600
    !zz5 = z*(z+1)*(z+2)*(z+3)*(z+4)*(z+5)*(z+6)*(z+7)*(z+8)/239500800
    correction !n' =
      let n=(#)n'
      in n**(1-z)/(z-1) - n**(-z)/2
         + n**(-z-1)*zz1 - n**(-z-3)*zz2 + n**(-z-5)*zz3
         - n**(-z-7)*zz4 + n**(-z-9)*zz5
```

12 Spence

Spence's integral for $z \geq 0$ is

$$S(z) = -\int_1^z \frac{\ln t}{t-1}\,dt = -\int_0^{z-1} \frac{ln(1+u)}{z}\,dz$$

and we extend the function via analytic continuation. Spence's function $S(z)$ is related to the dilogarithm function via $S(z) = \mathrm{Li}_2(1-z)$.

12.1 Preamble

```
module Spence (
    sf_spence,
) where
import Exp
import Util
```

A useful constant `pi2_6` $= \frac{\pi^2}{6}$

```
pi2_6 :: (Value v) ⇒ v
pi2_6 = pi^2/6
```

12.2 `sf_spence z`

Compute Spence's integral `sf_spence z` $= S(z)$. We use a variety of transformations to to allow efficient computation with a series.

$$\begin{aligned}
\mathrm{Li}_2(z) + \mathrm{Li}_2(\frac{z}{z-1}) &= -\frac{1}{2}(\ln(1-z))^2 \qquad z \in \mathbb{C} \setminus [1,\infty) \\
\mathrm{Li}_2(z) + \mathrm{Li}_2(\frac{1}{z}) &= -\frac{\pi^2}{6} - \frac{1}{2}(\ln(-z))^2 \qquad z \in \mathbb{C} \setminus [0,\infty) \\
\mathrm{Li}_2(z) + \mathrm{Li}_2(1-z) &= \frac{\pi^2}{6} - \ln(z)\ln(1-z) \qquad 0 < z < 1
\end{aligned}$$

(TODO: this code has not be solidly retested after conversion, especially verify complex.)

```
sf_spence :: (Value v) ⇒ v → v
sf_spence z
  | is_nan z       = z
  | (re z)<0       = 0/0
  | z == 0         = pi2_6
  | (rabs z)<0.5 = (series z) + (pi2_6 − (sf_log z)*(sf_log (1−z)))
  | (rabs z)<1.0 = −(series (1−z))
  | (rabs z)<2.5 = (series ((z−1)/z)) − (sf_log z)^2/2
  | otherwise      = (series (1/(1−z))) − pi2_6 − (sf_log (z−1))^2/2
```

*`series z`

The series expansion used for Spence's integral:

$$\texttt{series z} = -\sum_{k=1}^{\infty} \frac{z^k}{k^2}$$

```
series z =
  let zk = iterate (*z) z
      terms = zipWith (λ t k → −t/(#)k^2) zk [1..]
  in ksum terms
```

13 Lommel functions

13.1 Preamble

```
module Lommel (
  sf_lommel_s,
  sf_lommel_s2,
) where
import Util
```

–TODO: These are completely untested!

13.2 First Lommel function

For $\mu \pm \nu \neq \pm 1, \pm 3, \pm 5, \cdots$ we define the first Lommel function `sf_lommel_s mu nu z` $= S_{\mu,\nu}(z)$ via series-expansion:

$$S_{\mu,\nu}(z) = \frac{z^{mu+1}}{(\mu+1)^2 - \nu^2} \sum_{k=0}^{\infty} t_k$$

where

$$t_0 = 1 \qquad t_k = t_{k-1} \frac{-z^2}{(\mu + 2k + 1)^2 - \nu^2}$$

13.2.1 `sf_lommel_s mu nu z`

```
sf_lommel_s mu nu z =
  let terms = ixiter 1 1.0 $ λ k t → −t∗z^2 / ((mu+((#)$2∗k+1))^2 − nu^2)
      res = ksum terms
  in res ∗ z∗∗(mu+1) / ((mu+1)^2 − nu^2)
```

13.3 Second Lommel function

For $\mu \pm \nu \neq \pm 1, \pm 3, \pm 5, \cdots$ the second Lommel function `sf_lommel_s2 mu nu z` $= s_{\mu,\nu}(z)$ is given via an asymptotic expansion:

$$s_{\mu,\nu}(z) \sim \sum_{k=0}^{\infty} u_k$$

where

$$u_0 = 1 \qquad u_k = u_{k-1} \frac{-(\mu - 2k + 1)^2 - \nu^2}{z^2}$$

13.3.1 `sf_lommel_s2 mu nu z`

```
sf_lommel_s2 mu nu z =
  let tterms = ixiter 1 1.0 $ λ k t → −t∗((mu−((#)$2∗k+1))^2 − nu^2) / z^2
      terms = tk tterms
      res = ksum terms
  in res
  where tk (a:b:cs) = if (rabs a)<(rabs b) then [a] else a:(tk$b:cs)
```