

Seth Jacobs  
Asymptotics drill

- 1) When doubling  $n$  in an algorithm which has a run time of  $n^2$  the algorithm gets slower by a factor of 4 each time you double  $n$ . If you run the algorithm after adding one to  $n$ , each time it is slower by adding  $(n*2) + 1$  to the total amount of operations.
- 2) When doubling  $n$  in an algorithm which has a run time of  $n^3$  the algorithm gets slower by a factor of 8 each time you double  $n$ . If you run the algorithm after adding one to  $n$ , each time it is slower by a factor  $3n^2 + 3n + 1$ .
- 3) When doubling  $n$  in an algorithm which has a run time of  $100n^2$  the algorithm gets slower by a factor of 4 each time you double  $n$ . If you run the algorithm after adding one to  $n$ , each time it is slower by a factor  $200n + 100$ .
- 4) When doubling  $n$  in an algorithm which has a run time of  $n \log n$  the algorithm gets slower by a factor of a little more than 2.5 each time you double  $n$ . If you run the algorithm after adding one to  $n$ , each time it is slower by adding 4 to the amount of total operations.
- 5) When doubling  $n$  in an algorithm which has a run time of  $2^n$  the algorithm gets slower by a factor of  $2^n$  each time you double  $n$ . If you run the algorithm after adding one to  $n$ , each time it is slower by a factor of 2.

If you have a computer that can run  $10^{10}$  operations per second, the largest input for an algorithm that has a run time of

- 1)  $n^2$  is 6,000,000
- 2)  $n^3$  is 33,019
- 3)  $100n^2$  is 600,000
- 4)  $n \log n$  is 1,290,000,000,000
- 5)  $2^n$  is 45
- 6)  $2^{2^n}$  is 22

