

# Discrete Fast Fourier Transform

Severin Davis

Teja Dharma

*High Performance Computing with Graphic Cards*

University of Stuttgart

# Contents

Motivation .....	3
How-to .....	3
Program structure.....	4
Reordering of the input.....	5
Mapping and implementation of interactions .....	5
Thread mapping.....	5
Interaction implementation .....	6
Computation and mapping of roots of unity .....	6
Parallelization and expected results .....	8
Parallelization .....	8
Root computation.....	8
FFT algorithm .....	8
Expected Results.....	8
Results .....	9
Sources .....	12
Appendix.....	13
FFT.cpp .....	13
FFT.cl .....	19
floatc.hpp.....	21

## Motivation

Fast Fourier Transforms are useful in a variety of applications, most notably signal processing. If signal processing must occur as fast as possible or with a very large input size, parallelization of a Fast Fourier Transform algorithm can reap large benefits in required processing time.

In this project, a Fast Fourier Transform algorithm is implemented. Initially it is implemented as single threaded, serially-executing program. This implementation is analyzed for candidates for parallelization which are then moved to an OpenCL kernel.

These two implementations are compared both in theoretical complexity as a function of input size, as well as real world execution speed on multiple platforms.

## How-to

Import the FFT project into eclipse.

After running the program, the results of the algorithm should be printed.

The number of elements in the input and the input is printed.

The FFT CPU results are printed.

The FFT GPU results are printed.

These results are automatically compared and checked for errors.

The speed test for the algorithm is run and the results are printed.

The input file is included in the deliverable. This file is named "input.txt". This can be modified to provide different inputs, as long as the format of the example is maintained.

The speed test input cannot be changed, it's purpose is solely to test the performance of the implementation with various input sizes. The results of these FFT executions are not returned. For the speed test, only the execution times are printed, both to the consoles and to "output.txt"

## Program structure

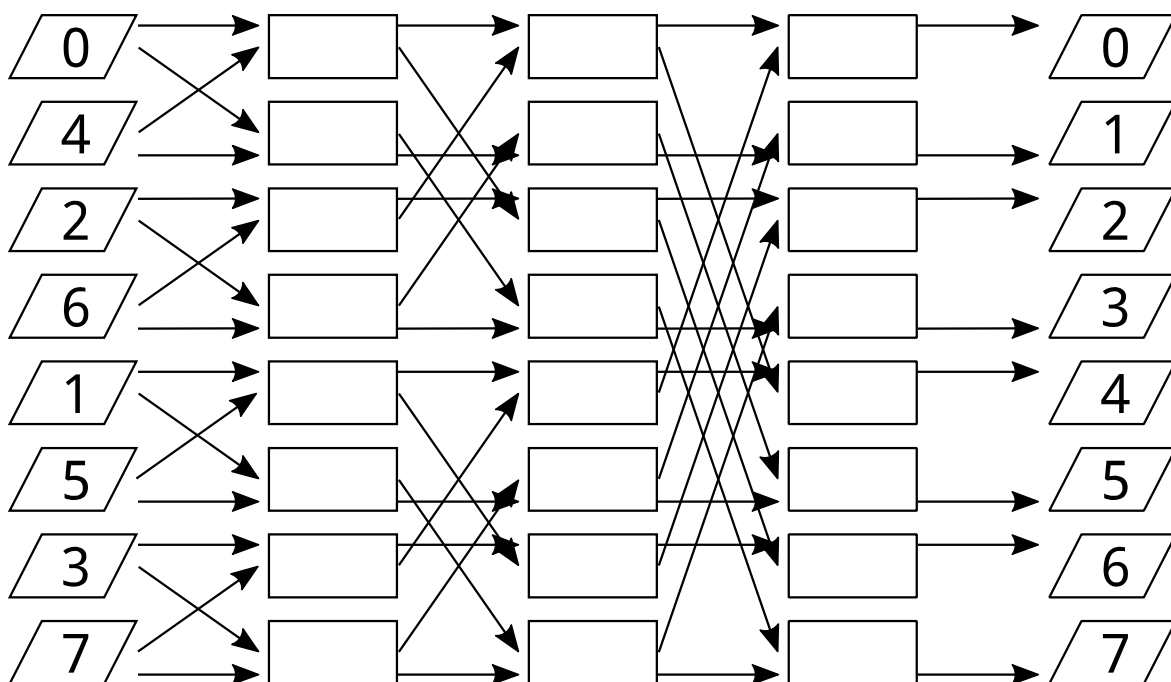
For this Fast Fourier Transform implementation, the Tukey-Cooley Algorithm was chosen. This algorithm works in stages in which two independent indices of the input array interact. The results of these interactions are placed back into the array and the algorithm is run in a chained manner, in which the results of previous stages are used in the current stage. Each interaction involves two indices and a computed root of unity. Additionally, the input must be placed in a specific order. As such, our implementation could be broken down into three sub problems, each with their unique challenges.

1. Reordering of the input
2. Mapping and implementation of interactions
3. Computation and mapping of roots of unity

Each of these is visible in the diagram of an example input size of 8 below.

The reordering of the input is visible on the input nodes at the left of the diagram.

The interacting index pairs, or stages, are shown as arrows. The boxes shown indicate the computations performed involving the roots of unity and the two inputs.



## Reordering of the input

Inputs need to be reordered before the algorithm can be run. Unfortunately, because this reordering does not operate on independent indices, this portion of the FFT implementation was not a candidate for parallelization. Instead, this was implemented as a recursive function that executes prior to any FFT execution.

```
std::vector<double> reorder_input(std::vector<double> input)
{
    //check base case
    if (input.size() <= 2)
        return input;

    //create even and odd index reorder vectors
    std::vector<double> even(input.size()/2);
    std::vector<double> odd(input.size()/2);

    //split even and odd indices
    for (unsigned int i = 0; i < even.size(); i++)
    {
        even.at(i) = input.at(i * 2);
        odd.at(i) = input.at((i * 2) + 1);
    }

    //recursive call to reorder even and odd index vectors
    even = reorder_input(even);
    odd = reorder_input(odd);

    //regroup returned vectors and return
    std::vector<double> result;
    result.reserve(input.size());
    result.insert(result.end(), even.begin(), even.end());
    result.insert(result.end(), odd.begin(), odd.end());

    return result;
}
```

## Mapping and implementation of interactions

The number of interactions occurring at each stage is half that of the input size. For example, in the given diagram, the sample size is 8. The left column of arrows, however, always shows two pairs of nodes interacting, creating a total of four interactions in that stage. Similarly, stages 2 and 4 also have 4 interacting pairs. As such, the number of threads spawned is half of the input size.

### Thread mapping

One major challenge to this approach was finding the corresponding mapping from the thread number to the nodes it operates on in a specific stage.

As an example, the diagram shows thread0 mapped to index 0 and index 1 of the input array. Similarly, thread 0 is mapped to index 0 and 2 and index 0 and 4 at

stages 2 and 4, respectively.

By analyzing the pattern, the following mapping was found to output an index, given the thread number and the current stage.

```
unsigned int thread_index_map(unsigned int thread_id, unsigned int stage)
{
    return (stage * 2 * (thread_id / stage)) + thread_id % stage;
}
```

Finally, after retrieving the one thread to index mapping, the other interacting index can be found simply by adding the current stage to the returned mapping.

### **Interaction implementation**

Once the correct mapping is found, the two data items at those indices interact with each other, as well as an additional root of unity.

This interaction is defined as:

$$\text{result}[\text{home\_index}] = \text{result}[\text{home\_index}] + (\text{result}[\text{target\_index}] * \text{root})$$
$$\text{result}[\text{target\_index}] = \text{result}[\text{home\_index}] - (\text{result}[\text{target\_index}] * \text{root})$$

Of interest is the fact that the root for a particular interaction remains constant. As such, the  $(\text{target\_index} * \text{root})$  can be computed only once, further optimizing for speed of execution.

This is implemented in the code shown:

```
c1_float2 pq = floatc_mul(result.at(target_index), roots.at(home_root));
c1_float2 top = floatc_add(pq, result.at(home_index));
c1_float2 bottom = floatc_sub(result.at(home_index), pq);

result.at(home_index) = top;
result.at(target_index) = bottom;
```

### **Computation and mapping of roots of unity**

The required degree of the roots of unity for this FFT implementation is directly related to the size of the input. For instance, an input of size 8 will require 8 roots of unity. As such, instead of computing the necessary root for that particular stage and thread, the required roots are computed at the beginning of the algorithm and are simply referenced later via an additional mapping function.

Additionally, because roots are symmetric around the unit circle, we can save time by simply computing half of them. For example, an input size of 8 will only require 4 roots to be computed. Coincidentally, since 4 threads are spawned for the same

input size, each thread is tasked with computing one root.

This computation is shown below:

```
for (unsigned int i = 0; i < roots.capacity(); i++)
{
    double arg = (2 * pi * (double)i) / ((double)result.capacity());
    roots.at(i).real = cos(arg);
    roots.at(i).imag = -1 * sin(arg);
}
```

Like the thread to index mapping function, threads need a mapping function to find the correct root to use, as shown below:

```
unsigned int thread_root_map(unsigned int thread_id, unsigned int estage,
unsigned int istage)
{
    return istage * (thread_id % estage);
}
```

In this function, estage counts up (1, 2, 4, 8, ...) while istage counts in reverse (... , 8, 4, 2, 1).

## Parallelization and expected results

Examination of the diagram of the FFT algorithm reveals that the number of stages of the algorithm is logarithmically related to the size of the input. Similarly, the number of operations performed within each stage is linearly related to the size of the input. Additionally, the non-recurring cost of computing the roots is also linearly related to the size of the input.

Thus, the runtime complexity of the serially executing algorithm  $O(n+n*\log(n))$  where  $n$  is the size of the input. This reduces to just  $O(n*\log(n))$ .

Analyses of the algorithm reveals several two primary candidates for optimization:

1. Root computation
2. FFT algorithm

### Parallelization

#### Root computation

Parallelizing the root computation simply maps one root computation to each thread. Thus, the runtime complexity for this portion reduces from  $O(n)$  to  $O(1)$ .

#### FFT algorithm

The stages of the algorithm cannot be parallelized, as they rely on previous stages' results. However, the interactions within each stage are easily parallelized by mapping each interaction to its own thread. Thus, the runtime complexity for this portion reduces from  $O(n*\log(n))$  to  $O(\log(n))$ .

Overall this results in a runtime complexity reduction from  $O(n*\log(n))$  to  $O(\log(n))$ .

### Expected Results

Timing results from the implementation are expected to roughly follow these complexities.

However, there is an overhead cost associated with moving data the OpenCL device as well as an associated cost in starting up the computation on the device. This introduces an overhead and causes a breakeven point in the total run time. This means that for dataset inputs below a certain size, the OpenCL overhead would cause the theoretically faster OpenCL implementation to be slower.

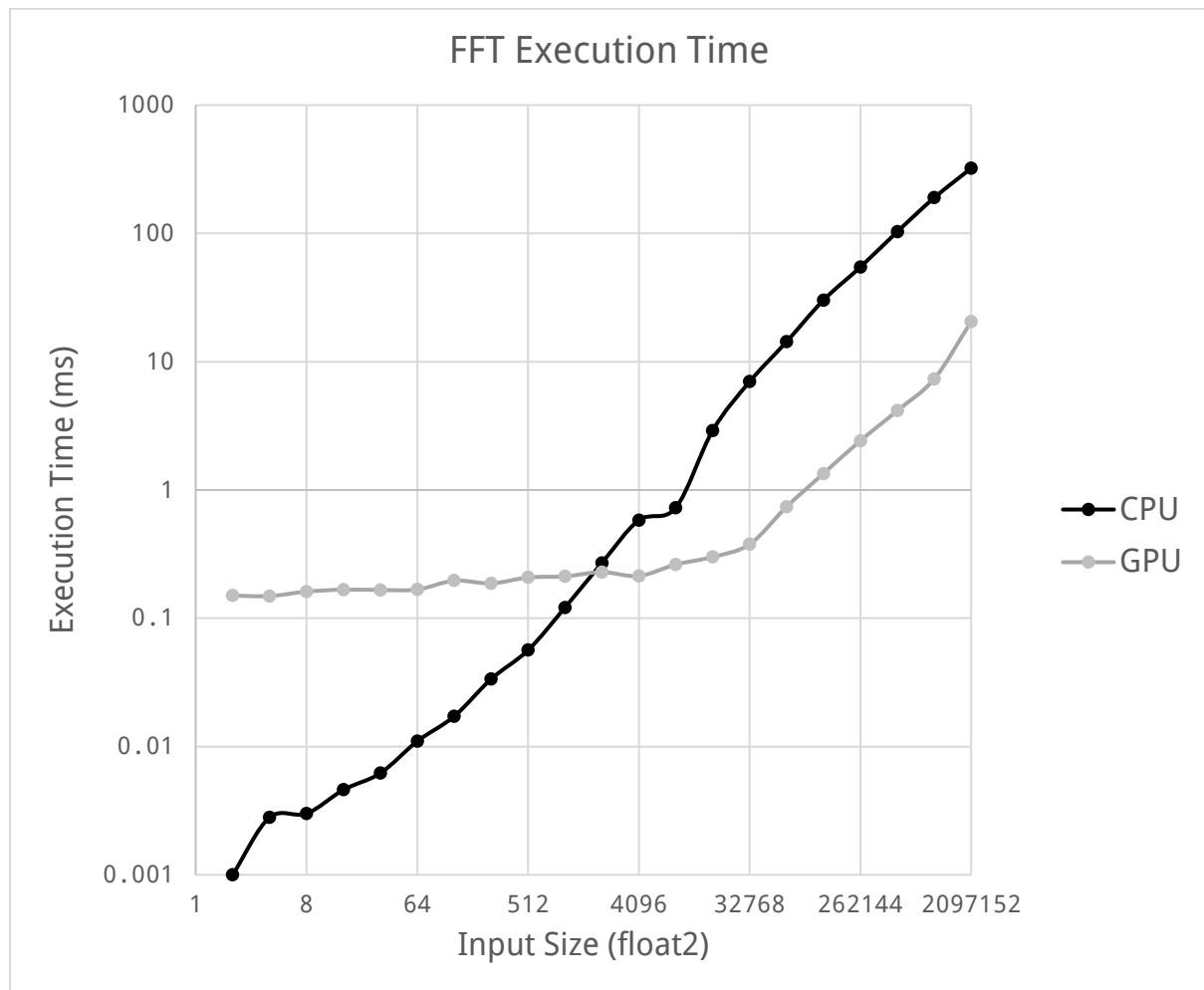


# Results

## Execution speed

The implementation was executed 5 times over a large range of input sizes. The timing results were averaged with each other to minimize spurious results.

These averaged results were graphed.



This graph clearly shows the breakeven point. This point is located at the intersections of the CPU and GPU lines. Thus, any dataset with a size of around 2048 or larger should be run on the GPU implementation as the execution time cost is lower. Sizes below this should not be run on the GPU as the GPU overhead kills destroys its parallelized execution improvements.



This graph more clearly shows the overhead cost of executing on a GPU. The long, negative slope line from 2 to about 32768 relates well to the increasing size of the input. This indicates that most of the execution time on the GPU seen in this range is the overhead. From 32768 to the end of the graph, this overhead appears to be amortized and we see the true execution time per element.

## Potential improvements

Since this implementation only computes the roots once at the beginning of the algorithm, the effects of the sine and cosine functions are negligible. However, in a different implementation where the roots are computed again when needed, using `native_sin` and `native_cos` functions may result in significant performance increases. The implementation used for this project assumed that a memory access of the stored roots would be faster than a repeated computation of each required root.

Additionally, as the size of the input and the required memory increases, the GPU

may be starved due to insufficient memory bandwidth. Without reducing the input size, increasing cache size, or compressing the individual elements, there is no easy way to solve this.

## Sources

"Discrete Fourier Transform (DFT) Calculator - Fourier Series

Calculation", *Easycalculation.com*, 2017. [Online]. Available:

<https://www.easycalculation.com/engineering/mechanical/discrete-fourier-transform.php>.

P. Heckbert, Fourier Transforms and the Fast Fourier Transform (FFT) Algorithm.

1995. [Online]. Available: <https://www.cs.cmu.edu/afs/andrew/scs/cs/15-463/2001/pub/www/notes/fourier/fourier.pdf>.

S. Smith, The Scientist and Engineer's Guide to Digital Signal Processing. 2011.

Available: <http://www.dspguide.com/ch12/2.htm>

# Appendix

## FFT.cpp

```
#include <stdio.h>

#include <Core/Assert.hpp>
#include <Core/Time.hpp>
#include <OpenCL/cl-patched.hpp>
#include <OpenCL/Program.hpp>
#include <OpenCL/Event.hpp>
#include <OpenCL/Device.hpp>

#include <fstream>
#include <sstream>
#include <iostream>
#include <cmath>

#include "floatc.hpp"

using namespace Core;
using namespace std;

#define pi 3.14159265359f

int read_input(std::vector<cl_float2>* input, string filename);
void print_vector(std::vector<cl_float2>* input);
std::vector<cl_float2> reorder_input(std::vector<cl_float2> input);
unsigned int thread_index_map(unsigned int thread_id, unsigned int stage);
unsigned int thread_root_map(unsigned int thread_id, unsigned int estage, unsigned int istage);

void print_vector(std::vector<cl_float2>* input)
{
    cout << endl;
    for (unsigned int i = 0; i < input->size(); i++)
    {
        cout << floatc_to_string(input->at(i)) << endl;
    }
}

int read_input(std::vector<cl_float2>* input, string filename)
{
    input->clear();
    //read file
    ifstream inputFile;

    inputFile.open(filename.c_str(), ifstream::in); //input.csv
    //verify that file is open
    if (inputFile.is_open())
    {
        //log success
    }
    else
    {
        //log fail
        return -1;
    }

    string inputLine;

    //count number of lines (samples) in file
    while (getline(inputFile, inputLine))
    {
        cl_float2 sample;
        sample.y = 0;
        try
        {
            float line = ::atof(inputLine.c_str());
            sample.x = line;
        }
        catch (...)
    }
```

```

        {
            return -1;
        }
        input->push_back(sample);
    }

    //check sample size for usability
    //TODO
    if (input->size() < 2)
    {
        return -1;
    }

    unsigned int powertwo = 2;

    //find smallest power of two larger than the input vector
    while (powertwo < input->size())
        powertwo = powertwo * 2;

    //pad input vector with 0s until it's a power of two size
    while (input->size() != powertwo)
    {
        cl_float2 sample;
        sample.x = 0;
        sample.y = 0;
        input->push_back(sample);
    }
    return true;
}

std::vector<cl_float2> reorder_input(std::vector<cl_float2> input)
{
    //check base case
    if (input.size() <= 2)
        return input;

    //create even and odd index reorder vectors
    std::vector<cl_float2> even(input.size()/2);
    std::vector<cl_float2> odd(input.size()/2);

    //split even and odd indices
    for (unsigned int i = 0; i < even.size(); i++)
    {
        even.at(i) = input.at(i * 2);
        odd.at(i) = input.at((i * 2) + 1);
    }

    //recursive call to reorder even and odd index vectors
    even = reorder_input(even);
    odd = reorder_input(odd);

    //regroup returned vectors and return
    std::vector<cl_float2> result;
    result.reserve(input.size());
    result.insert(result.end(), even.begin(), even.end());
    result.insert(result.end(), odd.begin(), odd.end());

    return result;
}

std::vector<cl_float2> fft_cpu(std::vector<cl_float2> input)
{
    std::vector<cl_float2> result(input);
    std::vector<cl_float2> roots;
    roots.reserve(input.size()/2);

    //populate root vector with empty complex numbers
    for (unsigned int i = 0; i < roots.capacity(); i++)
    {
        cl_float2 root;
        root.x = 0;
        root.y = 0;
    }
}

```

```

        roots.push_back(root);
    }

    //calculate half of the roots
    //we do this to match the number of threads (size/2) that run on the OpenCL
implementation
    for (unsigned int i = 0; i < roots.capacity(); i++)
    {
        float arg = (2 * pi * (float)i) / ((float)result.capacity());
        roots.at(i).x = cos(arg);
        roots.at(i).y = -1 * sin(arg);
    }

#ifdef DEBUG
    cout << "unity roots";
    print_vector(&roots);
    cout << endl;
#endif

    //threads always operate on pairs of indices in each stage
    //so we start half as many threads as the input size
    unsigned int num_of_threads = (unsigned int)(result.size() / 2);

    //we maintain a power of two up counter and a power of two down counter to avoid exp
or log functions
    //here we loop through the stages which has a runtime complexity of log(n)
    unsigned int istage = num_of_threads;
    for (unsigned int estage = 1; estage <= num_of_threads; estage = estage * 2)
    {
#ifdef DEBUG
        cout << "STAGE " << estage << endl;
#endif
        //the inner loops simply executes our "threads".
        //this will be parallelized in the opencl implementation
        //runtime for this is n
        //total for fft serial execution is nlogn
        for (unsigned int thread_id = 0; thread_id < num_of_threads; thread_id++)
        {
            //we query for our home and target indeces
            //as well as our home and target root indeces
            unsigned int home_index = thread_index_map(thread_id, estage);
            unsigned int target_index = home_index + estage;
            unsigned int home_root = thread_root_map(thread_id, estage, istage);
#ifdef DEBUG
            cout << "t" << thread_id << ": " << "hi" << home_index << ", ti" <<
target_index << endl;
            cout << "t" << thread_id << ": " << "hr" << home_root << endl;
#endif
            cl_float2 pq = floatc_mul(result.at(target_index),
roots.at(home_root));
            cl_float2 top = floatc_add(pq, result.at(home_index));
            cl_float2 bottom = floatc_sub(result.at(home_index), pq);
            result.at(home_index) = top;
            result.at(target_index) = bottom;
        }
#ifdef DEBUG
        cout << endl;
#endif
        istage = istage / 2;
    }
    return result;
}

unsigned int thread_index_map(unsigned int thread_id, unsigned int stage)
{
    return (stage * 2 * (thread_id / stage)) + thread_id % stage;
}

unsigned int thread_root_map(unsigned int thread_id, unsigned int estage, unsigned int
istage)
{
    return istage * (thread_id % estage);
}

```

```

float round(float x){
    float y = ceilf(x*100)/100;
    return y;
}

unsigned int error_calculator(std::vector<cl_float2> fromCPU, std::vector<cl_float2>
fromGPU)
{
    unsigned int errorCount = 0;
    for(size_t i = 0; i < fromCPU.size(); i++){
        float xdif = abs(fromCPU[i].x - fromGPU[i].x);
        float ydif = abs(fromCPU[i].y - fromGPU[i].y);
        if(xdif > 0.0001 || ydif > 0.0001)
            errorCount++;
    }
    return errorCount;
}

std::vector<cl_float2> opencl(std::vector<cl_float2> input, Core::TimeSpan *gpuTime)
{
    std::vector<cl_float2> output;
    // Create a context
    //cl::Context context(CL_DEVICE_TYPE_GPU)
    std::vector<cl::Platform> platforms;
    cl::Platform::get(&platforms);
    if (platforms.size() == 0) {
        std::cerr << "No platforms found" << std::endl;
        return output;
    }
    int platformId = 0;
    for (size_t i = 0; i < platforms.size(); i++) {
        //cout << platforms[i].getInfo<CL_PLATFORM_NAME>() << endl;
        platformId = i;
    }
    cl_context_properties prop[4] = { CL_CONTEXT_PLATFORM,
    (cl_context_properties)platforms[platformId](), 0, 0 };
    //std::cout << "Using platform '" <<
platforms[platformId].getInfo<CL_PLATFORM_NAME>() << "' from '" <<
platforms[platformId].getInfo<CL_PLATFORM_VENDOR>() << "'" << std::endl;
    cl::Context context(CL_DEVICE_TYPE_GPU, prop);

    // Get a device of the context
    //std::cout << "Context has " << context.getInfo<CL_CONTEXT_DEVICES>().size() << "
devices" << std::endl;
    cl::Device device = context.getInfo<CL_CONTEXT_DEVICES>()[0];
    std::vector<cl::Device> devices;
    devices.push_back(device);

    // Create a command queue
    cl::CommandQueue queue(context, device, CL_QUEUE_PROFILING_ENABLE);

    // Load the source code
    cl::Program program = OpenCL::loadProgramSource(context, "src/FFT.cl");
    // Compile the source code. This is similar to program.build(devices) but will print
more detailed error messages
    OpenCL::buildProgram(program, devices);

    std::size_t data_size = input.size()*sizeof(cl_float2);
    cl::Buffer d_data(context, CL_MEM_READ_WRITE, data_size);
    cl::Buffer d_roots(context, CL_MEM_READ_WRITE, (input.size()/2) *
sizeof(cl_float2));

    cl::Event copytoGPU;
    queue.enqueueWriteBuffer(d_data, true, 0 , data_size, input.data(), NULL,
&copytoGPU);

    cl::Event event;

    // Create a kernel object
    cl::Kernel kernel1(program, "fftKernel");
    kernel1.setArg(0, d_data);
    kernel1.setArg(1, d_roots);

```



```

        unsigned int WG_size = input.size()/2;
        if(WG_size > 256)
            WG_size = 256;

        queue.enqueueNDRangeKernel(kernel1, 0, input.size()/2, WG_size, 0, &event);

        cl::Event copytoCPU;
        queue.enqueueReadBuffer(d_data, true, 0, data_size, input.data(), NULL, &copytoCPU);
        queue.finish();

        Core::TimeSpan gpuExecTime = OpenCL::getElapsedTime(event);
        Core::TimeSpan copyTime = OpenCL::getElapsedTime(copytoGPU) +
OpenCL::getElapsedTime(copytoCPU);
        *gpuTime = gpuExecTime + copyTime;

        //cout << "Total time in GPU is " << gpuTime << endl;
        return input;
    }

void speed_test(unsigned int size)
{
    ofstream myfile;
    myfile.open ("output.txt");

    unsigned vector_size = 1;
    for(unsigned int i = 0; i <= size; i++)
    {
        vector_size = vector_size * 2;
        std::vector<cl_float2> test;
        for(unsigned int j = 0; j < vector_size; j++)
        {
            cl_float2 sample;
            sample.y = 0;
            sample.x = (float)(j+1);
            test.push_back(sample);
        }

        //CPU
        Core::TimeSpan cpuStart = Core::getCurrentTime();
        vector<cl_float2> result = fft_cpu(test);
        Core::TimeSpan cpuEnd = Core::getCurrentTime();
        Core::TimeSpan cpuTime = cpuEnd - cpuStart;
        cout << "CPU, " << vector_size << ", " << cpuTime << endl;

        //GPU
        Core::TimeSpan gpuTime = Core::getCurrentTime();
        vector<cl_float2> resultgpu = opencl(test, &gpuTime);
        cout << "GPU, " << vector_size << ", " << gpuTime << endl;

        myfile << "CPU, " << vector_size << ", " << cpuTime << ", GPU, " <<
vector_size << ", " << gpuTime << endl;
    }
    myfile.close();
}

int main()
{
    std::vector<cl_float2> h_input;

    if (read_input(&h_input, "input.txt"))
    {
        cout << "Input has " << h_input.size() << " samples." << endl;
    }
    else
    {
        cout << "Input invalid" << endl;
        return 1;
    }

    cout << "raw input";
    print_vector(&h_input);
    cout << endl;
}

```

```

        h_input = reorder_input(h_input);
#ifdef DEBUG
        cout << "reordered input";
        print_vector(&h_input);
        cout << endl;
#endif

        Core::TimeSpan cpuStart = Core::getCurrentTime();
        vector<cl_float2> result = fft_cpu(h_input);
        Core::TimeSpan cpuEnd = Core::getCurrentTime();

        Core::TimeSpan cpuTime = cpuEnd - cpuStart;

        cout << "fft result";
        print_vector(&result);
        cout << endl;

        Core::TimeSpan gpuTime = Core::getCurrentTime();
        vector<cl_float2> resultgpu = opencl(h_input, &gpuTime);

        cout << "opencl result";
        print_vector(&resultgpu);
        cout << endl;

        cout << "Found errors in " << error_calculator(result,resultgpu) << " values" <<
endl;

        cout << "Total time in CPU is " << cpuTime << endl;
        cout << "Total time in GPU is " << gpuTime << endl;

        speed_test(20);
}

```

## FFT.cl

```
#ifndef __OPENCL_VERSION__
#include <OpenCL/OpenCLKernel.hpp> // Hack to make syntax highlighting in Eclipse work
#endif

#define pi 3.14159265359f

unsigned int thread_index_map(unsigned int thread_id, unsigned int stage)
{
    return (stage * 2 * (thread_id / stage)) + thread_id % stage;
}

unsigned int thread_root_map(unsigned int thread_id, unsigned int estage, unsigned int
istage)
{
    return istage * (thread_id % estage);
}

float2 floatc_add(float2 a, float2 b)
{
    float2 c;
    c.x = a.x + b.x;
    c.y = a.y + b.y;
    return c;
}

float2 floatc_sub(float2 a, float2 b)
{
    float2 c;
    c.x = a.x - b.x;
    c.y = a.y - b.y;
    return c;
}

float2 floatc_mul(float2 a, float2 b)
{
    float2 c;
    c.x = a.x*b.x - a.y*b.y;
    c.y = a.y*b.x + a.x*b.y;
    return c;
}

__kernel void fftKernel(__global float2* d_input, __global float2* d_roots) {

    size_t i = get_global_id(0);

    size_t count = get_global_size(0);

    //populate root array
    float arg = (2 * pi * (float)i) / ((float)count * 2);
    d_roots[i].x = cos(arg);
    d_roots[i].y = -1 * sin(arg);

    barrier(CLK_GLOBAL_MEM_FENCE);

    unsigned int istage = count;
    for (unsigned int estage = 1; estage <= count; estage = estage * 2)
    {
        //we query for our home and target indeces
        //as well as our home and target root indeces
        unsigned int home_index = thread_index_map(i, estage);
        unsigned int target_index = home_index + estage;
        unsigned int home_root = thread_root_map(i, estage, istage);

        float2 pq = floatc_mul(d_input[target_index], d_roots[home_root]);
        float2 top = floatc_add(pq, d_input[home_index]);
        float2 bottom = floatc_sub(d_input[home_index], pq);
        d_input[home_index] = top;
        d_input[target_index] = bottom;

        istage = istage / 2;

        barrier(CLK_GLOBAL_MEM_FENCE);
    }
}
```

}	}
---	---

## floatc.hpp

```
#ifndef FLOATC_H
#define FLOAT_H

#include <iostream>

using namespace std;

cl_float2 floatc_add(cl_float2 a, cl_float2 b)
{
    cl_float2 c;
    c.x = a.x + b.x;
    c.y = a.y + b.y;
    return c;
}

cl_float2 floatc_sub(cl_float2 a, cl_float2 b)
{
    cl_float2 c;
    c.x = a.x - b.x;
    c.y = a.y - b.y;
    return c;
}

cl_float2 floatc_mul(cl_float2 a, cl_float2 b)
{
    cl_float2 c;
    c.x = a.x*b.x - a.y*b.y;
    c.y = a.y*b.x + a.x*b.y;
    return c;
}

string floatc_to_string(cl_float2 a)
{
    stringstream ss;
    ss << a.x << " + " << a.y << "i";
    return ss.str();
}

#endif
```