Compte rendu TP 2

Exercice 1: 1.select dept name, budget from department where budget in (select max(budget) from department) 2.select teacher_name, salary from teacher where salary >(select avg(salary) from teacher) 3.select teacher.teacher name, student.student name, count(*) as nombre cours from teacher, student, takes, teaches where takes.takes id = student.student id and teacher.teacher id= teaches.teaches id and takes.course_id =teaches.course_id and takes.sec_id = teaches.sec_id and takes.semester = teaches.semester and takes.takes vear = teaches.teaches vear GROUP BY teacher.teacher_name , student.student_name HAVING count (*) >= 2; 4.select S.teacher name, S.student name, S.nombre cours from (select teacher.teacher name as teacher name, student.student name as student name, count(*) as nombre_cours from teacher join teaches on teacher.teacher id = teaches.teaches id join takes on takes.course_id = teaches.course_id and takes.sec id = teaches.sec id and takes.semester = teaches.semester and takes.takes year = teaches.teaches year join student on takes.takes_id = student.student_id group by teacher.teacher_name, student.student_name) S where S.nombre cours \geq 2; 5. select student.student id, student.student name from student where student.student id not in (select takes.takes_id from takes where takes.takes_year < 2009);

6.select teacher name from teacher where teacher name LIKE 'E%';

```
7.select t1.teacher_name
from teacher t1
where 3 = (
       select count(distinct t2.salary)
       from teacher t2
       where t2.salary > t1.salary
);
8.select t1.teacher name, t1.salary
from teacher t1
where 2 \ge (
       select count(distinct t2.salary)
       from teacher t2
       where t2.salary < t1.salary
order by t1.salary desc;
9.select s.student_name
from student s
where (s.student_id) in (
       select takes.takes_id
       from takes
       where takes.semester = 'Fall'
       and takes.takes_year = 2009
);
10. select s.student_name
from student s
where ('Fall', 2009) = some (
       select takes.semester, takes.takes_year
       from takes
       where takes.takes_id = s.student_id
);
11.select s.student_name
from student s
natural inner join takes t
where t.semester = 'Fall'
 and t.takes_year = 2009;
12. SELECT student_name
FROM student
WHERE EXISTS ( SELECT *
```

```
FROM takes
WHERE takes.takes_id = student.student_id AND semester = 'Fall'
AND takes year = 2009;
13. select a.student name, b.student name
from (student natural inner join takes) as a,
      (student natural inner join takes) as b
where a.course id = b.course id
 and a.sec id = b.sec id
 and a.semester = b.semester
 and a.takes year = b.takes year
 and a.student id <> b.student id
 and a.student name < b.student name
group by a.student id, b.student id
having count(*) >= 1;
14. SELECT teacher.teacher_name , count (*)
from takes INNER JOIN teaches ON takes.course id = teaches.course id
             AND takes.sec id = teaches.sec id
             AND takes.semester = teaches.semester
             AND takes.takes_year = teaches.teaches_year
INNER JOIN teacher ON teaches.teaches id = teacher.teacher id
GROUP BY teacher.teacher name, teacher.teacher id ORDER BY count (*) DESC;
15. SELECT teacher.teacher_name, count(teaches.course_id)
FROM takes
INNER JOIN teaches ON takes.course id = teaches.course id
             AND takes.sec id = teaches.sec id
             AND takes.semester = teaches.semester
             AND takes_year = teaches.teaches year
RIGHT OUTER JOIN teacher ON teaches teaches id = teacher.teacher id
GROUP BY teacher.teacher name, teacher.teacher id
ORDER BY count(teaches.course_id) DESC;
16.WITH mytakes (id, course id, sec id, semester, year, grade) AS (
      SELECT takes id, course id, sec id, semester, takes year, grade
      FROM takes
      WHERE grade = 'A'
SELECT teacher.teacher name, count(mytakes.course id)
FROM mytakes
INNER JOIN teaches ON mytakes.course id = teaches.course id
             AND mytakes.sec id = teaches.sec id
             AND mytakes.semester = teaches.semester
             AND mytakes.year = teaches.teaches year
RIGHT OUTER JOIN teacher ON teaches.teaches id = teacher.teacher id
GROUP BY teacher.teacher name, teacher.teacher id
ORDER BY count(mytakes.course id) DESC;
```

```
17. SELECT teacher.teacher_name, student.student_name, count(*)
FROM teacher
INNER JOIN teaches ON teacher.teacher id = teaches.teaches id
INNER JOIN takes ON teaches.course id = takes.course id
               AND teaches.sec id = takes.sec id
               AND teaches.semester = takes.semester
               AND teaches teaches year = takes.takes year
INNER JOIN student ON takes.id = student.student id
GROUP BY teacher.teacher name, student.student name;
18. SELECT teacher teacher name, student student name, count(*)
FROM teacher
INNER JOIN teaches ON teacher.teacher id = teaches.teaches id
INNER JOIN takes ON teaches.course id = takes.course id
               AND teaches.sec id = takes.sec id
               AND teaches.semester = takes.semester
                AND teaches teaches year = takes.takes year
INNER JOIN student ON takes.id = student.student id
GROUP BY teacher.teacher_name, student.student_name;
Exercice 3:
1-A \rightarrow A (Réflexivité)
B → B (Réflexivité)
C → C (Réflexivité)
D → D (Réflexivité)
E → E (Réflexivité)
A \rightarrow B (Décomposition de A \rightarrow BC)
A \rightarrow C (Décomposition de A \rightarrow BC)
A \rightarrow D (Transitivité de A \rightarrow B et B \rightarrow D)
E → A (Donnée)
E \rightarrow B (Transitivité de E \rightarrow A et A \rightarrow B)
E \rightarrow C (Transitivité de E \rightarrow A et A \rightarrow C)
CD \rightarrow E (Donnée)
CD \rightarrow A (Transitivité de CD \rightarrow E et E \rightarrow A)
B \rightarrow D (Donnée)
A → BC (Donnée)
E \rightarrow BC (Transitivité de E \rightarrow A et A \rightarrow BC)
2-
a- B^+ = \{B, D, A, C, E\}
(AB)^+ = \{A, B, C, D, E\}
b- (AF)^+ = \{A, B, C, D, E, F\} => donc (AF) est une super clé.
3/a- Décomposition en R1(A,B,C) et R2(A,D,E) :
```

Les ensembles d'attributs sont :

- Attributs(R1) = {A, B, C}
- Attributs(R2) = {A, D, E}

L'intersection des attributs est :

• Attributs(R1) ∩ Attributs(R2) = {A}

Si A est une super clé, alors la décomposition est sans perte. **3/b-**

Les ensembles d'attributs sont :

- Attributs(R1) = {A, B, C}
- Attributs(R2) = {C, D, E}

L'intersection des attributs est :

Attributs(R1) ∩ Attributs(R2) = {C}

Si C est une super clé, alors la décomposition est sans perte.

Exercice 4:

```
1-def print_dependencies(dependencies):
    for dependency in dependencies:
        left, right = dependency
        left_item = ','.join(sorted(left))
        right_item = ','.join(sorted(right))
        print(f'{left_item}->{right_item}')
```

Execution:

2- def print_relations(relations):

```
for relation in relations:
if isinstance(relation, set):
print(f'R({", ".join(sorted(relation))})')
```

3- def powerSett(inputset):

subset_list = []

for i in range(len(inputset) + 1):

for subset in itertools.combinations(inputset, i):

subset_list.append(set(subset))

return subset_list

```
31 - def power_set(inputset):
32
        subset_list = []
         for i in range(len(inputset) + 1):
 33 *
 34 ₹
         for subset in itertools.combinations(inputset, i):
 35
                subset_list.append(set(subset))
 36
        return subset_list
 37
 38 example_set = {'A', 'B', 'C'}
 39 print("Power Set:", power_set(example_set))
Ln: 32, Col: 21
Run
         → Share  $ Command Line Arguments
    Power Set: [set(), {'C'}, {'A'}, {'B'}, {'C', 'A'}, {'C', 'B'}, {'A', 'B'}, {'C', 'A'}]
¥
    ** Process exited - Return Code: 0 **
>_ Press Enter to exit terminal
```

4- def closure(attributes, dependencies):

closure_set = set(attributes)

changed = True

while changed:

```
changed = False
for left, right in dependencies:
if left.issubset(closure_set) and not right.issubset(closure_set):
         closure_set.update(right)
         changed = True
return closure_set
         11 - def closure(attributes, dependencies):
                closure_set = set(attributes)
                changed = True
         13
                while changed:
         14 +
         15
                   changed = False
                    for left, right in dependencies:
         16 -
         17 -
                        if left.issubset(closure_set) and not right.issubset(closure_set):
                           closure_set.update(right)
         18
                           changed = True
         19
         50
               return closure_set
         51
         attributes = {'A'}
         print(attributes, "+ = ", closure(attributes, mydependencies))
         n: 53, Col: 24
                          $
                             Command Line Arguments
                 Share
           {'A'} + = {'H', 'C', 'A', 'B'}
```

 def determinates_functionally(F, alpha, beta): closure_alfa = closure(alpha, F) return beta.issubset(closure_alfa)

```
57 - def determinates_functionally(F, alpha, beta):
         closure_alfa = closure(alpha, F)
 58
 59
         return beta.issubset(closure_alfa)
 60
 61 alpha_att = {'A'}
62 beta_att = {'B','C','Z'}
 63 print(determinates_functionally(mydependencies,alpha_att, beta_att))
Ln: 62, Col: 23
Run
          Share
                        Command Line Arguments
    False
¥
    ** Process exited - Return Code: 0 **
   Press Enter to exit terminal
```

7.def IsSuperKey(dependencies, relation, K): return relation.issubset(closure(K,mydependencies))

```
relation= {\'A', 'B', 'C', 'H'}
             K={'A'}
            - def IsSuperKey(dependencies, relation, K):
                 return relation.issubset(closure(K,mydependencies))
             print(K , "is a superkey of ", relation, "= ", IsSuperKey(mydependencies, relation, K))
           65, Col: 29
                  ♦ Share
                           $ Command Line Arguments
            \{'A'\} is a superkey of \{'C', 'B', 'A', 'H'\} = True
            ** Process exited - Return Code: 0 **
            Press Enter to exit terminal
8. def IsCandidateKey (F, R, K):
         if not IsSuperKey(F,R,K):
         return False
         if not IsSuperKey (F, R, K): return False
         for A in K:
         K aux = set (K)
         K_aux.discard (A)
         if IsSuperKey (F, R, K_aux): return False
         return True
```

```
71 relation= {'A', 'B', 'C', 'H'}
72 K={'A'}
73 - def IsCandidateKey (F, R, K):
      if not IsSuperKey(F,R,K) :
      if not IsSuperKey (F , R , K ) : return False
76
      for A in K:
           K_aux = set (K)
           K_aux.discard (A)
           if IsSuperKey (F , R , K_aux ) : return False
80
       return True
82 print(IsCandidateKey(mydependencies, relation, K))
85
n: 80, Col: 37
                  $ Command Line Arguments
  True
```

```
relation= {'A', 'B', 'C', 'H'}
72 K={'A','B'}
73 - def IsCandidateKey (F, R, K):
      if not IsSuperKey(F,R,K) :
          return False
       if not IsSuperKey (F , R , K ) : return False
       for A in K :  
           K_aux = set (K)
           K_aux.discard (A)
          if IsSuperKey (F , R , K_aux ) : return False
  print(IsCandidateKey(mydependencies, relation, K))
: 72, Col: 10
        Share
                 $ Command Line Arguments
 Run
 False
```

```
9.def compute_all_candidate_keys(F, R):
    result = []
    for K in power_set(R):
    if is_candidate_key(F, R, K):
    result.append(K)
    return result
```

```
def compute_all_candidate_keys(F, R):
    result = []
    for K in power_set(R):
       if is_candidate_key(F, R, K):
            result.append(K)
    return result
R = {'A', 'B', 'C', 'D', 'E'}
F = [
    ( {'A', 'B'}, {'C'} ),
    ( {'C'}, {'D', 'E'} )
candidate_keys = compute_all_candidate_keys(F, R)
print("Clés candidates:", [sorted(k) for k in candidate_keys])
5, Col: 29
     Share
                   Command Line Arguments
Clés candidates: [['A', 'B']]
```

```
10. Cet opération se fait en deux étapes:
```

```
    Calculer toutes les sous-séquences de la relation :
        def calculate_all_subsets(r):
        subsets = [set()]
        for element in r:
            new_subsets = []
            for subset in subsets:
            new_subsets.append(subset.union({element}))
            subsets.extend(new_subsets)
```

return subsets

2. On va prendre les sous-séquences qui sont des super clés :

```
def calculate_all_superkeys (r,depandecy):
    subsets=calculate_all_subsets(r)
    result = []
    for s in subsets:
        if IsSuperKey(dependencies=mydependencies , relation = relation, K= s):
        result.append(s)
```

```
return result
     relation= {'A', 'B', 'C', 'H'}
    mydependencies = [
        [{'A'}, {'B'}], # A->B
         [{'A'}, {'C'}], # A->C
         [{'C', 'G'}, {'H'}], # CG ->H
         [{'C', 'G'}, {'I'}], # CG ->I
         [{'B'}, {'H'}] # B->H
    - def calculate_all_superkeys (r,depandecy):
       subsets=calculate_all_subsets(r)
       result = []
       for s in subsets:
        if IsSuperKey(dependencies=mydependencies , relation = relation, K= s):
            result.append(s)
       return result
     print(calculate all superkeys(r,mydependencies))
    2, Col: 27
         → Share  $ Command Line Arguments
   [{'A'}, {'B', 'A'}, {'A', 'C'}, {'B', 'A', 'C'}]
```

```
11.def find_one_candidate_key(R, F):
    all_subsets = power_set(R)
    all_subsets.sort(key=len)

for subset in all_subsets:
    if is_candidate_key(subset, R, F):
    return subset
```

return None

```
def find_one_candidate_key(R, F):
    all_subsets = power_set(R)
    all_subsets.sort(key=len)

for subset in all_subsets:
    if is_candidate_key(subset, R, F):
        return subset

return None
```

RQ: Au début, on a fait un tri pour au final avoir la clé la plus courte.

Exemple:

```
12. def est_en_BCNF(relation, dependencies):
    for left, right in dependencies:
    if not left:
        continue
    if not relation.issubset(closure(left, dependencies)):
        return False
        return True
```

```
131
132 ▼ myrelations = [
         {'A', 'B', 'C', 'G', 'H', 'I'},
133
         {'X', 'Y'}
134
135
136
137 → for r in myrelations:
138 -
         if est_en_BCNF(r, mydependencies):
              print(f"La relation R({', '.join(sorted(r))}) est en BCNF.")
139
         else:
140 *
              print(f"La relation R({', '.join(sorted(r))}) n'est pas en BCNF.")
141
Ln: 147, Col: 2
                       Command Line Arguments
Run
          Share
   La relation R(A, B, C, G, H, I) n'est pas en BCNF.
```

```
13.def est_en_BCNF_schema(schema_relations, dependances_fonctionnelles):
       def est en BCNF relation(relation, dependencies relation):
       for left, right in dependencies relation:
       if not left:
              continue
       if not relation.issubset(closure(left, dependencies relation)):
              return False
       return True
       for relation in schema_relations:
       dependances relation = []
       for dep_left, dep_right in dependances_fonctionnelles:
       if dep left.issubset(relation) and dep right.issubset(relation):
              dependances_relation.append([dep_left, dep_right])
       if not est en BCNF relation(relation, dependances relation):
       return False
       return True
```

```
0 - schema1_relations = [
       {'A', 'B', 'C'},
       {'D', 'E'}
72
73
74 - schema1_dependances = [
       [{'A'}, {'B'}],
76
       [{'C'}, {'A', 'B'}],
       [{'D'}, {'E'}]
78 ]
if est en BCNF schema(schema1 relations, schema1 dependances):
       print("Le schéma 1 est en BCNF.")
32 - else:
       print("Le schéma 1 n'est pas en BCNF.")
155, Col: 19
Run
       Share
                     Command Line Arguments
 Le schéma 1 n'est pas en BCNF.
```

```
14.def decomposition BCNF(relations, dependances fonctionnelles):
       result = list(relations) # Commencer avec les relations initiales
       i = 0
       while i < len(result):
       R = result[i]
       Fd R = [] # Dépendances fonctionnelles applicables à R
       for left, right in dependances_fonctionnelles:
       if left.issubset(R) and right.issubset(R):
               Fd R.append([left, right])
       if not est_en_BCNF_relation(R, Fd_R):
       # Trouver une dépendance fonctionnelle X -> Y dans Fd R qui viole BCNF
       violating_fd = None
       for left, right in Fd R:
               if not left:
               continue
               if not R.issubset(fermeture(left, Fd R)):
               violating fd = [left, right]
               break
       if violating fd:
               X, Y = violating fd
               R1 = X.union(Y)
               R2 = R.difference(Y)
               # Ajouter les nouvelles relations à la liste et supprimer l'ancienne
               if R1 not in result:
               result.append(R1)
```

```
if R2 and R2 not in result: # Ajouter R2 seulement s'il n'est pas vide
result.append(R2)
result.pop(i)
i = 0 # Recommencer la vérification depuis le début
else:
i += 1 # Aucune violation trouvée (cas improbable si la condition BCNF est
fausse)
else:
i += 1
```

return result

```
233 relations_initiales = [
235 - dependances_initiales = [
236
         [{'A'}, {'B', 'C'}]
237
238
relations_bcnf = decomposition_BCNF(relations_initiales, dependances_initiales)
print("Décomposition en BCNF:", [sorted(list(r)) for r in relations_bcnf])
241
242
243
244
n: 232, Col: 2
                   $

→ Share

                      Command Line Arguments
 Run
  Décomposition en BCNF: []
   ** Process exited - Return Code: 0 **
  Press Enter to exit terminal
```