

Lecture 1

Introduction to Data Structure and Arrays
(Traversal, Insertion, Deletion, Searching)

Data Structure

- A particular way of storing and organizing data in a computer so that it can be used efficiently.
- They provide a means to manage large amounts of data efficiently, such as large databases.
- Data are simply values or set of values and Database is organized collection of data.

Abstract Data Type (ADT)

- Realization of a data type as a software component, i.e. focus on **what** can be done with the data, not **how** it is done.
- The interface of the ADT is defined in terms of a type and a set of operations on that type.
- Implementation details are hidden from the user of the ADT and protected from outside access, a concept referred to as encapsulation.

Study of Data Structure Includes

- Logical description of data structure.
- Implementation of data structure.
- Quantitative analysis of data structure, this include amount of memory, processing time

Classification of Data Structures

Data Structure

Primitive

(All data-types available in the language)

Non-primitive

Linear

Non-linear

Array

Stack

Queue

Linked List

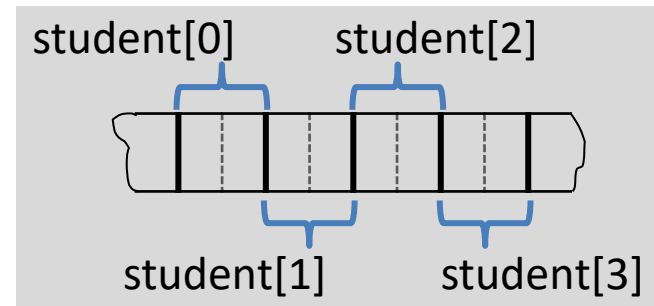
Trees

Graph

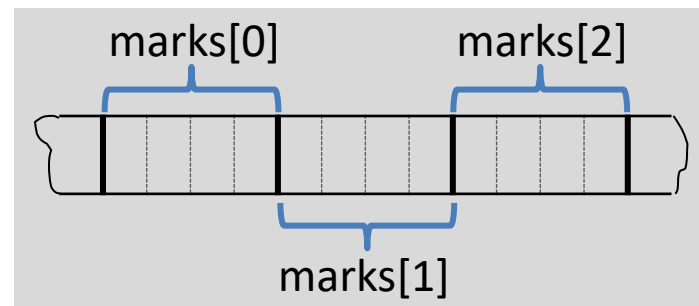
Array

Memory Storage – One Dimensional Array

```
int student[4];
```



```
float marks[3];
```



Memory Storage – Two Dimensional Array

```
int marks[3][5];
```

- Can be visualized in the form of a matrix as

| | Col 0 | Col 1 | Col 2 | Col 3 | Col 4 |
|-------|-------------|-------------|-------------|-------------|-------------|
| Row 0 | marks[0][0] | marks[0][1] | marks[0][2] | marks[0][3] | marks[0][4] |
| Row 1 | marks[1][0] | marks[1][1] | marks[1][2] | marks[1][3] | marks[1][4] |
| Row 2 | marks[2][0] | marks[2][1] | marks[2][2] | marks[2][3] | marks[2][4] |

Contd...

- Row-major order

| | | | | | | | | | | | | | | |
|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| (0,0) | (0,1) | (0,2) | (0,3) | (0,4) | (1,0) | (1,1) | (1,2) | (1,3) | (1,4) | (2,0) | (2,1) | (2,2) | (2,3) | (2,4) |
| Row0 | | | | | Row1 | | | | | Row2 | | | | |

- Column-major order

| | | | | | | | | | | | | | | |
|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| (0,0) | (1,0) | (2,0) | (0,1) | (1,1) | (2,1) | (0,2) | (1,2) | (2,2) | (0,3) | (1,3) | (2,3) | (0,4) | (1,4) | (2,4) |
| Col0 | | | Col1 | | | Col2 | | | Col3 | | | Col4 | | |

Array ADT

- The simplest but useful data structure.
- Assign single name to a homogeneous collection of instances of one abstract data type.
 - All array elements are of same type, so that a pre-defined equal amount of memory is allocated to each one of them.
- Individual elements in the collection have an associated index value that depends on array dimension.

Contd...

- One-dimensional and two-dimensional arrays are commonly used.
- Multi-dimensional arrays can also be defined.
- Usage:
 - Used frequently to store relatively permanent collections of data.
 - Not suitable if the size of the structure or the data in the structure are constantly changing.

Operations on Linear Data Structures

- Traversal
- Search
- Insertion
- Deletion
- Merging
- Sorting

Traversal

- Processing each element in the array.
- Example: Find minimum element in the array

Algorithm arrayMinElement(A,n)

Input: An array **A** containing **n** integers.

Output: The minimum element in **A**.

| | cost | time | |
|-------------------------|------|------------------------|-----------------------------|
| 1. min = 0 | c1 | 1 | → c1 |
| 2. for i = 1 to n-1 do | c2 | n | → c2 n |
| 3. if A[min] > A[i] | c3 | n - 1 | → c3 (n - 1) |
| 4. min = i | c4 | $\sum_{i=1}^{n-1} t_i$ | → c4 $\sum_{i=1}^{n-1} t_i$ |
| 5. return A[min] | c5 | 1 | → c5 |

Contd...

- $$T(n) = c1 + c2 n + c3 (n - 1) + c4 \sum_{i=1}^{n-1} t_i + c5$$
$$= (c1 - c3 + c5) + (c2 + c3) n + c4 \sum_{i=1}^{n-1} t_i$$
- Best case: $\Omega(n) \rightarrow$ summation evaluates to 0.
- Worst case: $O(n) \rightarrow$ summation evaluates to n .
- Average case: $\theta(n) \rightarrow$ summation evaluates to $n/2$.

Example – Traversal

1. //Determine smallest element in an array.
2. `#include<iostream>`
3. `using namespace std;`
4. `int arrayMinElement(int arr[], int n)`
5. `{`
6. `int min = 0;`
7. `for (int i = 1; i < n; i++)`
8. `{`
9. `if (arr[i] < arr[min])`
10. `min = i;`
11. `}`
12. `return arr[min]; }`

Contd...

```
13. int main()
14. {
15.     int i, n, a[10];
16.     cout << "Enter the array size (<=10): ";
17.     cin >> n;
18.     for (i = 0; i < n; i++)
19.     { cout << "Enter " << i+1 << " element: ";
20.         cin >> a[i];
21.     }
22.     cout << "\nSmallest element is " << arrayMinElement(a,n);
23.     return 0;
24. }
```


Insertion and Deletion

| | | | | | | | | | | |
|-----|---|---|---|---|---|---|---|---|---|---|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| a[] | 8 | 6 | 3 | 4 | 5 | | | | | |

- Insert 2 at index 1

| | | | | | | | | | | |
|-----|---|--------------|--------------|--------------|--------------|---|---|---|---|---|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| a[] | 8 | 6 | 3 | 4 | 5 | | | | | |
| | | 2 | 6 | 3 | 4 | 5 | | | | |

- Delete the value at index 2

| | | | | | | | | | | |
|-----|---|---|--------------|--------------|--------------|---|---|---|---|---|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| a[] | 8 | 2 | 6 | 3 | 4 | 5 | | | | |
| | | | 3 | 4 | 5 | | | | | |

Algorithm – Insertion

Algorithm `insertElement(A,n,num,indx)`

Input: An array **A** containing **n** integers and the number **num** to be inserted at index **indx**.

Output: Successful insertion of **num** at **indx**.

| | cost | time |
|---------------------------------|-------------|----------------|
| 1. for $i = n - 1$ to $indx$ do | c1 | $n - indx + 1$ |
| 2. $A[i + 1] = A[i]$ | c2 | $n - indx$ |
| 3. $A[indx] = num$ | c3 | 1 |
| 4. $n = n + 1$ | c4 | 1 |

Contd...

- $T(n) = c_1 (n - \text{indx} + 1) + c_2 (n - \text{indx}) + c_3 + c_4$
- Best case: $\Omega(1) \rightarrow \text{indx} = n - 1$
- Worst case: $O(n) \rightarrow \text{indx} = 0$
- Average case: $\theta(n) \rightarrow \text{indx} = n/2$

Algorithm – Deletion

Algorithm deleteElement(A,n,indx)

Input: An array **A** containing **n** integers and the index **indx** whose value is to be deleted.

Output: Deleted value stored initially at **indx**.

| | cost | time |
|-----------------------------|------|--------------|
| 1. temp = A[indx] | c1 | 1 |
| 2. for i = indx to n – 2 do | c2 | n – indx |
| 3. A[i] = A[i + 1] | c3 | n – indx – 1 |
| 4. n = n – 1 | c4 | 1 |
| 5. return temp | c5 | 1 |

Contd...

- $T(n) = c_1 + c_2 (n - \text{indx}) + c_3 (n - \text{indx} - 1) + c_4 + c_5$
- Best case: $\Omega(1) \rightarrow \text{indx} = n - 1$
- Worst case: $O(n) \rightarrow \text{indx} = 0$
- Average case: $\theta(n) \rightarrow \text{indx} = n/2$

Example – Insertion and Deletion

```
1. #include<stdio.h>
2. int n;
3. void insert(int a[], int num, int pos)
4. { for(int i = n-1; i >= pos; i--)          a[i+1] = a[i];
5.   a[pos] = num;
6.   n++; }
7. int deleteElement(int a[], int pos)
8. { int temp = a[pos];
9.   for(int i = pos; i <= n-2; i++)          a[i] = a[i+1];
10.  n--;
11.  return temp; }
12. void printArray(int a[])
13. { for (i = 0; i < n; i++)
14.   printf("%d ",a[i]); }
```

Contd...

```
15. int main()
16. { int i, a[10], num, pos;
17.   printf("Enter the size of an array (<10): ");
18.   scanf("%d",&n);
19.   for (i = 0; i < n; i++)
20.   {   printf("Enter element at index %d: ",i);
21.       scanf("%d",&a[i]);   }
22.   printf("\nEnter number to be inserted: ");
23.   scanf("%d",&num);
24.   printf("Enter the desired index: ");
25.   scanf("%d",&pos);
26.   insert(a,num,pos);
```

Contd...

```
27. printf("\nArray after insertion is...\n");
28. printArray(a);
29. printf("\nEnter the index whose value is to be deleted: ");
30. scanf("%d",&pos);
31. printf("\nThe deleted element is %d.", deleteElement(a,pos));
32. printf("\nArray after deletion is...\n");
33. printArray(a);
34. return 0;
35. }
```


Search

Find the location of the element with
a given value.

Linear Search

- Used if the array is unsorted.
- Example:

Search 7 in the following array

i → 0 → 1 → 2 → 3 → 4 → 5 → 6

| | | | | | | | | | | |
|-----|----|---|---|---|---|---|---|---|---|---|
| a[] | 10 | 5 | 1 | 6 | 2 | 9 | 7 | 8 | 3 | 4 |
|-----|----|---|---|---|---|---|---|---|---|---|

Found at
index 6

Search 11 in the following array

i → 0 → 1 → 2 → 3 → 4 → 5 → 6 → 7 → 8 → 9 → 10

| | | | | | | | | | | |
|-----|----|---|---|---|---|---|---|---|---|---|
| a[] | 10 | 5 | 1 | 6 | 2 | 9 | 7 | 8 | 3 | 4 |
|-----|----|---|---|---|---|---|---|---|---|---|

Not found

Contd...

Algorithm linearSearch(A,n,num)

Input: An array **A** containing **n** integers and number **num** to be searched.

Output: Index of **num** if found, otherwise -1.

| | cost | time |
|------------------------|---------------------------|---|
| 1. for i = 0 to n-1 do | c1 $\sum_{i=0}^n t_i$ | \rightarrow c1 $\sum_{i=0}^n t_i$ |
| 2. if A[i] == num | c2 $\sum_{i=0}^{n-1} t_i$ | \rightarrow c2 $\sum_{i=0}^{n-1} t_i$ |
| 3. return i | c3 1 | \rightarrow c3 |
| 4. return -1 | c4 1 | \rightarrow c4 |

Contd...

- $T(n) = c1 \sum_{i=0}^n t_i + c2 \sum_{i=0}^{n-1} t_i + (c3 + c4)$
- Best case: $\Omega(1) \rightarrow$ summation evaluates to 1.
- Worst case: $O(n) \rightarrow$ summation evaluates to n .
- Average case: $\theta(n) \rightarrow$ summation evaluates to $n/2$.

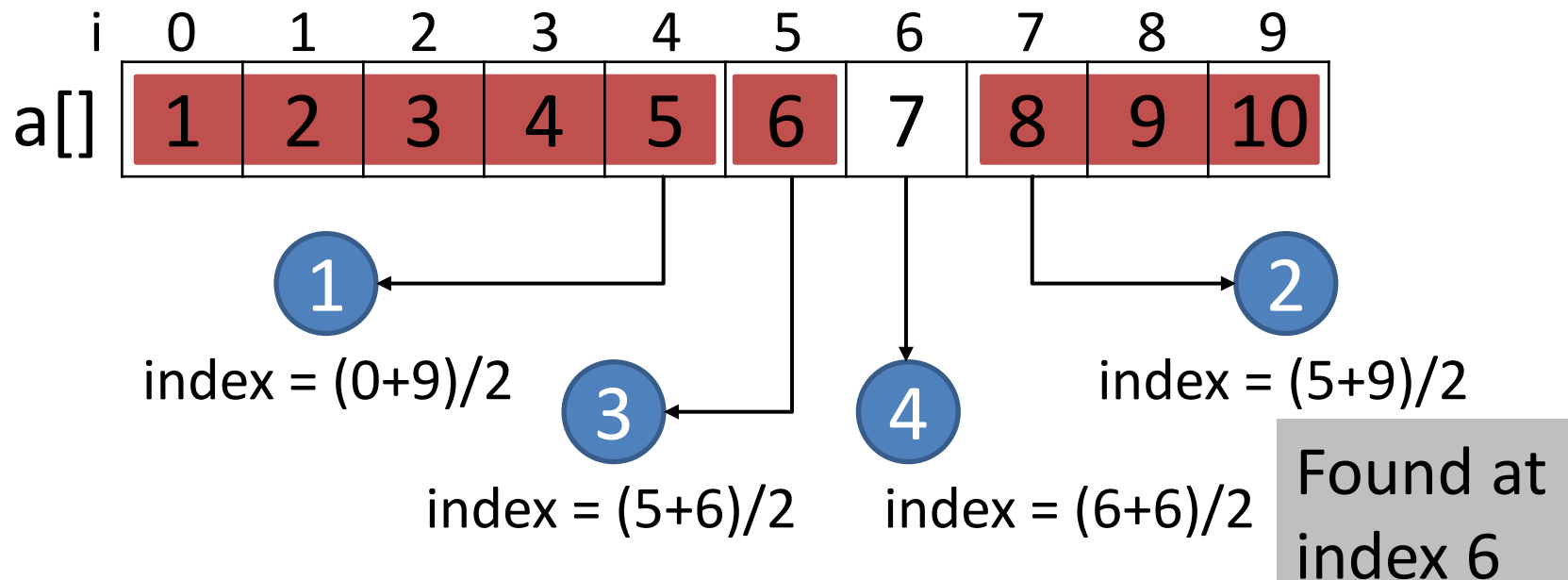
Example – Linear Search

```
1. #include<stdio.h>
2. int linearSearch(int a[], int n, int num)
3. { for (int i = 0; i < n; i++)
4.     if (a[i] == num) return i;
5.     return -1; }
6. int main()
7. { int i, n, a[10], num;
8.     printf("Enter the size of an array (<=10): ");    scanf("%d",&n);
9.     for (i = 0; i < n; i++)
10.    { printf("Enter element at index %d: ",i);          scanf("%d", &a[i]); }
11.    printf("\nEnter number to search: ");              scanf("%d",&num);
12.    int found = linearSearch(a,n,num);
13.    if(found != -1) printf("Element found at index %d",found);
14.    else             printf("Element not found.");
15.    return 0; }
```

Binary Searching

- Used if the array is sorted.
- Example:

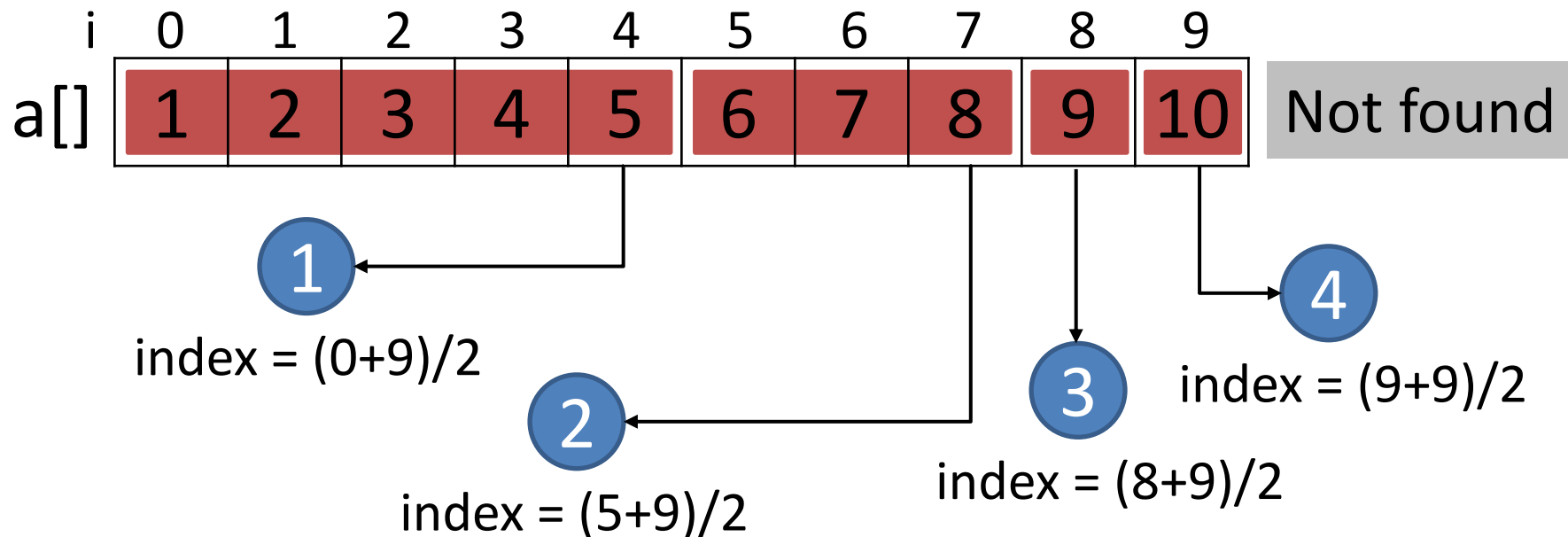
Search 7 in the following array



Contd...

- Used if the array is sorted.
- Example:

Search 11 in the following array



Contd...

Algorithm `binarySearch(A,n,num)`

Input: An array **A** containing **n** integers and number **num** to be searched.

Output: Index of **num** if found, otherwise -1.

1. `beg = 0, end = n-1`
2. `while beg <= end do`
3. `mid = (beg + end)/2`
4. `if A[mid] == num`
5. `return mid`
6. `else if A[mid] > num`
7. `end = mid - 1`
8. `else if A[mid] < num`
9. `beg = mid + 1`
10. `return -1`

Contd...

- In each iteration, the number of elements to be searched from gets reduced by half. This process continues till the number of elements (to be searched) reaches to 1.
 - 1st iteration – $n/2^0$
 - 2nd iteration – $n/2^1$
 - 3rd iteration – $n/2^2$
 - ...
 - $(m+1)^{\text{th}}$ iteration – $n/2^m$

Let the termination condition (i.e. single element to search from) reaches at the $(m+1)^{\text{th}}$ iteration, thus

$$1 = \frac{n}{2^m}$$

$$2^m = n$$

$$m \log_2 2 = \log_2 n$$

$$m = \lg n$$

\Rightarrow complexity is $O(\lg n)$

Example – Binary Search

```
1. #include<stdio.h>
2. int binarySearch(int arr[], int n, int num)
3. { int beg = 0, end = 9;
4.   while (beg <= end)
5.   { int mid = (beg + end)/2;
6.     if (arr[mid] == num)
7.       return mid;
8.     else if (arr[mid] > num)
9.       end = mid - 1;
10.    else if (arr[mid] < num)
11.      beg = mid + 1;
12.  }
13.  return -1;
14. }
```

Contd...

```
15. int main()
16. { int i, n, a[10], num;
17.   printf("Enter the size of an array (<=10): ");
18.   scanf("%d",&n);
19.   for (i = 0; i < n; i++)
20.   { printf("Enter element at index %d: ",i);
21.     scanf("%d", &a[i]); }
22.   printf("\nEnter number to search: ");
23.   scanf("%d",&num);
24.   int found = binarySearch(a,n,num);
25.   if(found != -1)
26.     printf("Element found at index %d",found);
27.   else
28.     printf("Element not found.");
29.   return 0; }
```



Thankyou