

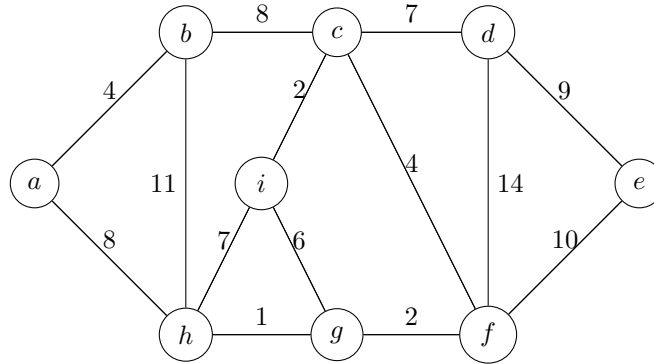
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## DSC 40B - Discussion 09

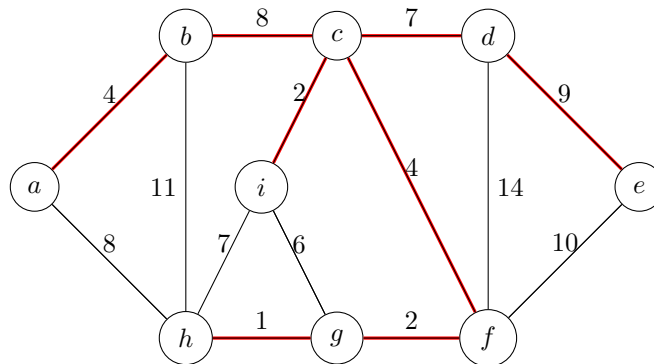
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### Problem 1.

Compute the minimum spanning tree for the following graph using Kruskal's algorithm. (Also compute the MST using Prim's algorithm and compare the results.)



**Solution:**



### Problem 2.

Suppose we are given both an undirected graph  $G$  with weighted edges and a minimum spanning tree  $T$  of  $G$ .

- a) Describe an efficient algorithm to update the minimum spanning tree when the weight of one edge  $e$  in  $T$  is decreased.

**Solution:** The minimum spanning tree of the updated graph would be  $T$ .

- b) Describe an efficient algorithm to update the minimum spanning tree when the weight of one edge  $e$  not in  $T$  is increased.

**Solution:** The minimum spanning tree of the updated graph would be  $T$ .

- c) Describe an efficient algorithm to update the minimum spanning tree when the weight of one edge  $e$  in  $T$  is increased.

**Solution:** Let  $e = (u,v)$  be the edge whose weight is increased. Remove the edge  $e$  from the minimum spanning tree  $T$ . This divides the tree  $T$  into two connected components. Let  $T_u$  be the component which contains  $u$  and  $T_v$  be the component which contains  $v$ . We can identify  $T_u$  and  $T_v$  by running a BFS with  $u$  and  $v$  as the sources. This takes time  $O(V + E)$ . While running BFS we can also label each node as 0 if it is a part of  $T_u$  and 1 if it is a part of  $T_v$ . We can now examine each edge  $e$  in the graph and find the minimum weight edge which connects a node labelled 0 to a node labelled 1. This takes time  $O(E)$ . We can then add this edge to  $T$  to get the minimum spanning tree of the updated graph. The total time complexity is  $O(V + E)$ .

- d) Describe an efficient algorithm to update the minimum spanning tree when the weight of one edge  $e$  not in  $T$  is decreased.

**Solution:** Let  $e = (u,v)$  be the edge whose weight is decreased. Add the edge  $e$  to the minimum spanning tree  $T$ . This would result in a cycle in  $T$ . We can identify the nodes and the edges on the cycle by running BFS with  $u$  or  $v$  as the source. This takes time  $O(V + E)$ . Find the maximum weight edge on the cycle and remove it from  $T$ . This takes time  $O(E)$ . The resultant tree would be a minimum spanning tree for the updated graph. The total time complexity is  $O(V + E)$ .