# CprE 381: Computer Organization and Assembly Level Programming

Performance

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#### **Administrative**

- HW5 due Mar 4 (TONIGHT)
- Term Project
  - Test framework updated (v2)
    - Dmem.hex included
    - Added synthesis
  - Part 2 Posted → 10% of Final Course Grade

# **Using CPI**

Drawing on the previous equation:

$$Execution \ Time = \# \ Instructions \times \frac{Cycles}{Instruction} \times \frac{Seconds}{Cycle}$$

- To improve performance (i.e., reduce execution time)
  - Increase clock rate (decrease clock cycle time) OR
  - Decrease CPI OR
  - Reduce the number of instructions
- Designers balance cycle time against the number of cycles required
  - Improving one factor may make the other one worse...

## **Evaluating Performance**

- Performance best determined by running a real application
  - Use programs typical of expected workload
  - Or, typical of expected class of applications
    - Compilers/editors, scientific applications, graphics, etc.
- Small benchmarks
  - Nice for architects and designers
  - Easy to standardize
  - Can be abused
- SPEC (System Performance Evaluation Cooperative)
  - Companies have agreed on a set of real program and inputs
  - Valuable indicator of performance (and compiler technology)
  - Can still be abused

#### SPECINT2006 for Intel Core i7-920

Description	Name	Instruction Count x 10 <sup>9</sup>	CPI	Clock cycle time (seconds x 10 <sup>-9</sup> )	Execution Time (seconds)	Reference Time (seconds)	SPECratio
Interpreted string processing	perl	2252	0.60	0.376	508	9770	19.2
Block-sorting compression	bzip2	2390	0.70	0.376	629	9650	15.4
GNU C compiler	gcc	794	1.20	0.376	358	8050	22.5
Combinatorial optimization	mcf	221	2.66	0.376	221	9120	41.2
Go game (AI)	go	1274	1.10	0.376	527	10490	19.9
Search gene sequence	hmmer	2616	0.60	0.376	590	9330	15.8
Chess game (AI)	sjeng	1948	0.80	0.376	586	12100	20.7
Quantum computer simulation	libquantum	659	0.44	0.376	109	20720	190.0
Video compression	h264avc	3793	0.50	0.376	713	22130	31.0
Discrete event simulation library	omnetpp	367	2.10	0.376	290	6250	21.5
Games/path finding	astar	1250	1.00	0.376	470	7020	14.9
XML parsing	xalancbmk	1045	0.70	0.376	275	6900	25.1
Geometric mean	-	_	-	_	_	-	25.7

#### **Clock Rate != Performance**

- Mobile Intel Pentium 4 Vs Intel Pentium M
  - 2.4 GHz
     1.6 GHz
- Performance on Mobilemark with same memory and disk
  - Word, excel, photoshop, powerpoint, etc.
  - Pentium M takes only 15% longer
- What is the relative CPI?

#### **Clock Rate != Performance**

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- What is the relative CPI?
  - ExecTime = IC CPI/Clock rate
  - $IC \cdot CPIM/1.6 = 1.15 \cdot IC \cdot CPI4/2.4$
  - $CPI4/CPIM = 2.4/(1.15 \cdot 1.6) = 1.3$

#### **CPI Varies**

- Different instruction types require different numbers of cycles
- CPI is often reported for types of instructions

Clock Cycles = 
$$\sum_{i=1}^{n} (CPI_i \times IC_i)$$

 Where CPI<sub>i</sub> is the CPI for the type of instructions and IC<sub>i</sub> is the count of that type of instruction

Lec08.1.8

# **Computing CPI**

To compute the overall average CPI use

$$CPI = \sum_{i=1}^{n} \left( CPI_{i} \times \frac{Instruction Count_{i}}{Instruction Count} \right)$$

Frequency of Instruction Type

# **Computing CPI Example**

Instruction Type	СРІ	Frequency	CPI * Frequency
ALU	1	50%	0.5
Branch	2	20%	0.4
Load	2	20%	0.4
Store	2	10%	0.2

- Given this machine, the CPI is the sum of CPI × Frequency
- Average CPI is 0.5 + 0.4 + 0.4 + 0.2 = 1.5
- What fraction of the time for data transfer?

#### **CPI Question**

Instruction Type	СРІ	Frequency	CPI * Frequency
ALU	1	50%	0.5
Branch	2	20%	0.4
Load	2	20%	0.4
Store	2	10%	0.2

- What is the impact of the MIPS displacement based memory addressing mode?
- Assume 50% of MIPS loads and stores have a zero displacement

#### **CPU Time versus MIPS**

- Two different compilers are being tested for a 1 GHz. machine with three different classes of instructions:
  - Class A, Class B, and Class C,
  - Require 1, 2, and 3 cycles (respectively)
  - Both compilers are used to produce code for a large piece of software
- The first compiler's code uses 5 million Class A instructions, 1 million Class B instructions, and 1 million Class C instructions
- The second compiler's code uses 10 million Class A instructions, 1 million Class B instructions, and 1 million Class C instructions
- Which sequence will be faster according to MIPS?
- Which sequence will be faster according to execution time?

# Speedup

Speedup allows us to compare different CPUs or optimizations

$$Speedup = \frac{CPUtimeOld}{CPUtimeNew}$$

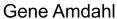
- Example
  - Original CPU takes 2sec to run a program
  - New CPU takes 1.5sec to run a program
  - Speedup = 1.33x or 33%

#### Amdahl's Law

 If an optimization improves a fraction f of execution time by a factor of a

$$Speedup = \frac{Told}{[(1-f)+f/a] \times Told} = \frac{1}{(1-f)+f/a}$$

- This formula is known as Amdahl's Law
- Lessons from
  - If f→100%, then speedup = a
  - If  $a\rightarrow \infty$ , the speedup = 1/(1-f)
- Summary
  - Make the common case fast
  - Watch out for the non-optimized/optimizable component





# Amdahl's Law Example

 Suppose a program runs in 100 seconds on a machine, with multiply responsible for 80 seconds of this time. How much do we have to improve the speed of multiplication if we want the program to run 4 times faster?

# In-class Assessment! Access Code: Nilla

Note: sharing access code to those outside of classroom or using access while outside of classroom is considered cheating

## Amdahl's Law Example

 Suppose a program runs in 100 seconds on a machine, with multiply responsible for 80 seconds of this time. How much do we have to improve the speed of multiplication if we want the program to run 4 times faster?

How about making it 5 times faster?

# Performance Summary

- Performance is specific to a particular program(s)
  - Total execution time is a consistent summary of performance
- For a given architecture performance increases come from:
  - Increases in clock rate (without adverse CPI affects)
  - Improvements in processor organization that lower CPI
  - Compiler enhancements that lower CPI and/or instruction count
  - Algorithm/Language choices that affect instruction count
- Pitfall: expecting improvement in one (limited) aspect of a machine's performance to affect the total performance

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