

ComS 472

Homework 3

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- 6.1 -

- 1) 3 colors for (SA), 2 for (WA), 1 for (NT, Q, NSW, V), 3 for (T) \Rightarrow
 $3 * 2 * 1 * 1 * 1 * 1 * 3 = 18$ solutions
 - 2) 4:(SA), 3:(WA), 2:(NT, Q, NSW, V), 4:(T) $\Rightarrow 4 * 3 * 2 * 2 * 2 * 2 * 4 = 768$ solutions
 - 3) 2:(SA), 1:(WA), 0:(NT, Q, NSW, V), 2:(T) $\Rightarrow 2 * 1 * 0 * 0 * 0 * 0 * 2 = 0$ solutions
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- 6.6 -

- 1) Introduce variable D as a set of pair values (D_1, D_2) . $A = D_1, B = D_2$, and $C = D_1 + D_2$
 - 2) Constraints with more than 3 variables can be reduced in tiers, with $n-1$ variables reduced to n , n to $n-1$, ..., continuing until the set of constraints contains only binary constraints.
 - 3) Unary constraints can be completely eliminated by moving the effects of the constraint into the domain of the variable it is affecting.
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- 6.8 -

- 1) $A1 \rightarrow R$ ✓
- 2) $H \rightarrow R$ **X**, conflicts with $A1 \rightarrow$ backtrack
- 3) $H \rightarrow G$ ✓
- 4) $A4 \rightarrow R$ ✓
- 5) $F1 \rightarrow R$ ✓
- 6) $A2 \rightarrow R$ ✓
- 7) $F2 \rightarrow R$ ✓
- 8) $A3 \rightarrow R$ ✓
- 9) $T \rightarrow R$ **X**, conflicts with $F1 \rightarrow$ backtrack
- 9) $T \rightarrow G$ **X**, conflicts with $H \rightarrow$ backtrack
- 9) $T \rightarrow B$ ✓

Arcs:

$(SA \neq WA), (WA \neq SA), (SA \neq NT), (NT \neq SA), (SA \neq Q), (Q \neq SA),$
 $(SA \neq NSW), (NSW \neq SA), (SA \neq V), (V \neq SA), (WA \neq NT), (NT \neq WA),$
 $(NT \neq Q), (Q \neq NT), (Q \neq NSW), (NSW \neq Q), (NSW \neq V), (V \neq NSW)$

Initial Constraints:

$SA=\{RGB\}, WA=\{RGB\}, NT=\{RGB\}, Q=\{RGB\}, NSW=\{RGB\}, V=\{RGB\}, T=\{RGB\}$

Set $WA=green, V=red...$

1)

C: $SA=\{RGB\}, WA=\{G_ \}, NT=\{RGB\}, Q=\{RGB\}, NSW=\{RGB\}, V=\{R_ \}, T=\{RGB\}$

A: All arcs, $(SA \neq WA), (NT \neq WA), (SA \neq V), (NSW \neq V)$ at the end for conciseness.

Action: Check and dequeue all consistent arcs.

2)

C: $SA=\{RGB\}, WA=\{G_ \}, NT=\{RGB\}, Q=\{RGB\}, NSW=\{RGB\}, V=\{R_ \}, T=\{RGB\}$

A: $(SA \neq WA), (NT \neq WA), (SA \neq V), (NSW \neq V)$

Action: $(SA \neq WA)$ is inconsistent. Remove G from SA, queue arcs with SA on right.

3)

C: $SA=\{R_B\}, WA=\{G_ \}, NT=\{RGB\}, Q=\{RGB\}, NSW=\{RGB\}, V=\{R_ \}, T=\{RGB\}$

A: $(NT \neq WA), (SA \neq V), (NSW \neq V), (WA \neq SA), (NT \neq SA), (Q \neq SA),$
 $(NSW \neq SA), (V \neq SA)$

Action: $(NT \neq WA)$ is inconsistent. Remove G from NT, queue arcs with NT on right.

4)

C: $SA=\{R_B\}, WA=\{G_ \}, NT=\{R_B\}, Q=\{RGB\}, NSW=\{RGB\}, V=\{R_ \}, T=\{RGB\}$

A: $(SA \neq V), (NSW \neq V), (WA \neq SA), (NT \neq SA), (Q \neq SA), (NSW \neq SA),$
 $(V \neq SA), (WA \neq NT), (Q \neq NT)$

Action: $(SA \neq V)$ is inconsistent. Remove R from SA, queue arcs with SA on right.

5)

C: $SA=\{_B\}, WA=\{G_ \}, NT=\{R_B\}, Q=\{RGB\}, NSW=\{RGB\}, V=\{R_ \}, T=\{RGB\}$

A: $(NSW \neq V), (WA \neq SA), (NT \neq SA), (Q \neq SA), (NSW \neq SA), (V \neq SA),$
 $(WA \neq NT), (Q \neq NT)$

Action: $(NSW \neq V)$ is inconsistent. Remove R from NSW, queue arcs with NSW on right.

6)

C: $SA=\{_B\}, WA=\{G_ \}, NT=\{R_B\}, Q=\{RGB\}, NSW=\{_GB\}, V=\{R_ \}, T=\{RGB\}$

A: $(WA \neq SA), (NT \neq SA), (Q \neq SA), (NSW \neq SA), (V \neq SA), (WA \neq NT),$
 $(Q \neq NT), (SA \neq NSW), (Q \neq NSW), (V \neq NSW)$

Action: $(WA \neq SA)$ is consistent. Dequeue.

7) _____
C: $SA=\{_B\}$, $WA=\{_G_ \}$, $NT=\{R_B\}$, $Q=\{RGB\}$, $NSW=\{_GB\}$, $V=\{R_ \}$, $T=\{RGB\}$
A: $(NT \neq SA)$, $(Q \neq SA)$, $(NSW \neq SA)$, $(V \neq SA)$, $(WA \neq NT)$, $(Q \neq NT)$,
 $(SA \neq NSW)$, $(Q \neq NSW)$, $(V \neq NSW)$
Action: $(NT \neq SA)$ is inconsistent. Remove B from NT, queue arcs with NT on right.

8) _____
C: $SA=\{_B\}$, $WA=\{_G_ \}$, $NT=\{R_ \}$, $Q=\{RGB\}$, $NSW=\{_GB\}$, $V=\{R_ \}$, $T=\{RGB\}$
A: $(Q \neq SA)$, $(NSW \neq SA)$, $(V \neq SA)$, $(WA \neq NT)$, $(Q \neq NT)$, $(SA \neq NSW)$,
 $(Q \neq NSW)$, $(V \neq NSW)$, $(SA \neq NT)$
Action: $(Q \neq SA)$ is inconsistent. Remove B from Q, queue arcs with Q on right.

9) _____
C: $SA=\{_B\}$, $WA=\{_G_ \}$, $NT=\{R_ \}$, $Q=\{RG_ \}$, $NSW=\{_GB\}$, $V=\{R_ \}$, $T=\{RGB\}$
A: $(NSW \neq SA)$, $(V \neq SA)$, $(WA \neq NT)$, $(Q \neq NT)$, $(SA \neq NSW)$, $(Q \neq NSW)$,
 $(V \neq NSW)$, $(SA \neq NT)$, $(SA \neq Q)$, $(NT \neq Q)$, $(NSW \neq Q)$
Action: $(NSW \neq SA)$ is inconsistent. Remove B from NSW, queue arcs with NSW on right.

10) _____
C: $SA=\{_B\}$, $WA=\{_G_ \}$, $NT=\{R_ \}$, $Q=\{RG_ \}$, $NSW=\{_GB\}$, $V=\{R_ \}$, $T=\{RGB\}$
A: $(V \neq SA)$, $(WA \neq NT)$, $(Q \neq NT)$, $(SA \neq NSW)$, $(Q \neq NSW)$, $(V \neq NSW)$,
 $(SA \neq NT)$, $(SA \neq Q)$, $(NT \neq Q)$, $(NSW \neq Q)$
Action: $(V \neq SA)$ is consistent. Dequeue.

11) _____
C: $SA=\{_B\}$, $WA=\{_G_ \}$, $NT=\{R_ \}$, $Q=\{RG_ \}$, $NSW=\{_GB\}$, $V=\{R_ \}$, $T=\{RGB\}$
A: $(WA \neq NT)$, $(Q \neq NT)$, $(SA \neq NSW)$, $(Q \neq NSW)$, $(V \neq NSW)$, $(SA \neq NT)$,
 $(SA \neq Q)$, $(NT \neq Q)$, $(NSW \neq Q)$
Action: $(WA \neq NT)$ is consistent. Dequeue.

12) _____
C: $SA=\{_B\}$, $WA=\{_G_ \}$, $NT=\{R_ \}$, $Q=\{RG_ \}$, $NSW=\{_GB\}$, $V=\{R_ \}$, $T=\{RGB\}$
A: $(Q \neq NT)$, $(SA \neq NSW)$, $(Q \neq NSW)$, $(V \neq NSW)$, $(SA \neq NT)$, $(SA \neq Q)$,
 $(NT \neq Q)$, $(NSW \neq Q)$
Action: $(Q \neq NT)$ is inconsistent. Remove R from Q, queue arcs with Q on right.

13) _____
C: $SA=\{_B\}$, $WA=\{_G_ \}$, $NT=\{R_ \}$, $Q=\{_G_ \}$, $NSW=\{_GB\}$, $V=\{R_ \}$, $T=\{RGB\}$
A: $(SA \neq NSW)$, $(Q \neq NSW)$, $(V \neq NSW)$, $(SA \neq NT)$, $(SA \neq Q)$, $(NT \neq Q)$,
 $(NSW \neq Q)$
Action: $(SA \neq NSW)$ is inconsistent. Remove B from SA, queue arcs with SA on right.

This leaves SA with no possible values, revealing an inconsistency in the partial assignment.

- 6.20 -

Variables: S_1, S_2, \dots, S_n

Domain: Set of adjacent square pairs a domino can cover (s_1, s_2) .

Constraints: No two dominos can cover the same square.

Variables: D_1, D_2, \dots, D_n

Domain: Set of pairs (s, s') , with s = current square and s' = one of 4 adjacent squares.

Constraints: A square must only link to one other square. Square A can link to square B if square B can link to A.

A tree-structured constraint graph can be accomplished with a straight line of 6 squares.

A tree-structured graph requires there to be no loops. If the layout of the squares has no cycles, the resulting constraint graph will be tree-structured.

- 7.4 -

1) Correct: Because false will never be true, it doesn't matter what the second half of the equation is.

2) Incorrect: True is always true, but False will never be true, so by definition it is incorrect.

3)

- 7.6 -

- 7.7 -

- 7.15 -

- 7.16 -
