

ComS 311
Recitation 3, 2:00 Monday
Homework 6

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1) Recurrence:
 combo = null if remainder < 0
 combo = [] if remainder == 0
 combo = min(from(i = 0 → n) combo(set, remaining - set[i], list))

Algorithm 1 Find non-negative integers w1, ..., wn.

```

1: int[] comboIter(set, W){
2:
3: //Store sub problem answers here
4: int[] sub = new int[W+1]
5: *Assign each index in sub to infinity
6:
7: //Store the indices of the num used for each sum here
8: int[] indices = new int[W+1]
9:
10: //Sum of 0 will always be an empty set
11: sub[0] = 0
12:
13: for (int i = 1; i < W; i++) do
14:     for (int j = 0; j < set.length; j++) do
15:         int num = set[j]
16:
17:         //If the current number can fit
18:         if (num <= i) then
19:             //Replace current answer if it works better
20:             int current = sub[i]
21:             int new = sub[i - num] + 1
22:
23:             if (new < current) then
24:                 //Update our records
25:                 sub[i] = new
26:                 //Save the index of the number that was used
27:                 indices[i] = j
28:             end if
29:         end if
30:     end for
31: end for

```

```

32:
33: //If the last cell is unused, there is no answer
34: if (sub[W] == infinity) then
35:     return []
36: end if
37:
38:
39: //Otherwise use the stored indices to build a list of [w1, ..., wn]
40: int[] w = new int[set.length]
41: *Assign each index in w to 0
42:
43: //Backtrack through our lists
44: int i = W
45: while (i > 0) do
46:
47:     //Increment the use count (w) of each number used in the end
48:     int index = indices[i]
49:     w[index]++
50:
51:     //Go to the number used before this one
52:     i -= set[index]
53: end while
54:
55: return w
56: }
57:
58: _____ Runtime:  $O(n*W)$  _____

```

2) Recurrence:

```
func = true: spaceLeft == 0 && remainingSum == 0
func = false:(spaceLeft == 0 && remainingSum != 0) || index == U.length
func = func(U, spaceLeft, remainingSum, index+1) ||
      func(U, spaceLeft-1, remainingSum - U[index], index+1);
```

Algorithm 2 Test for subset of U of size k that adds to T.

```
1: boolean iterFunc(U, T, k){
2:
3: n = U.length;
4:
5: //Make a 2d array to map subsets
6: matrix[ ][ ] = new boolean[n][T+1];    //n, 0→T (not 1→T)
7:
8: //Make hashmap to track the lengths of each subset that adds to a sum
9: //key = current sum, value = array of subset lengths
10: Hashmap legths = new HashMap<Integer, List<Integer>>();
11:
12: //Set column 0 to true, as all sums == 0 use empty set
13: for int i = 0; i <= n; i++ do
14:     matrix[i][0] = true;
15: end for
16:
17: //For each number in the set
18: for int i = 0; i < n; i++ do
19:     int number = U[i];
20:
21:     //From 1→T
22:     for int sum = 1; sum <= T; sum++ do
23:
24:         //If this number is too big, grab the val above
25:         if number > sum then
26:             matrix[i][sum] = matrix[i-1][sum];
27:             continue;
28:         end if
29:
```

```

30:
31:     //Decide if # can be added to a prev subset to fit current sum
32:     //Use typical subset-sum lookback
33:     result1 = matrix[i - 1][sum - number];
34:
35:     if result then
36:         //We are adding this to the subset
37:         //Ex: if sum = 14, number = 9, and lengths@5 = [1, 3, 4],
38:         //lengths@14 will now = [2, 4, 5]
39:         /** This is a loop through an array **
40:         lengths@sum = lengths@(sum - number)++;
41:     end if
42:
43:     //If this number won't fit, we don't add it to the subsets, but
44:     //this sum may still be possible, so the matrix should reflect that
45:     result2 = result || matrix[i-1][sum];
46:     matrix[i][j] = result2;
47: end for
48: end for
49:
50:
51: if ! (matrix[n-1][T]) then
52:     return false;
53: end if
54:
55: //If there is a possible subset, check that there is one of length k
56: list = lengths@T
57:
58: for int i = 0; i < list.length; i++ do
59:     if(list[i]== k) return true;
60: end for
61: return false;
62: }
63:
64:
65: _____ Runtime:  $O(n \cdot T \cdot k)$  _____

```

3) Recurrence:

```
traverse = score if x==M && y==N
traverse = max(traverse(maze, M, N, x, y+1, score-2),
               traverse(maze, M, N, x+1, y, score-2),
               traverse(maze, M, N, x+1, y+1, score-3))
```

Algorithm 3 Maximize score in M x N maze

```
1: int iterTraverse(maze, M, N){
2:
3: //Go from left→right, top→bottom, [X][X]
4: //looking at the max of cells to left, top left, and top [X][O]
5:
6: for int x = 1; x <= M; x++ do
7:     for int y = 1; y <= N; y++ do
8:         //Find largest score for transitioning to this cell
9:         int score = max(maze[x-1][y] -2, maze[x-1][y-1] -3, maze[x][y-1] -2)
10:
11:         //Replace maze slot with score + maybe diamond,
12:         // as we no longer need it
13:         maze[x][y] += score
14:     end for
15: end for
16:
17: return maze[M][N]
18: }
19:
20:
21: ————— Runtime: O(M*N) —————
```
