

IOWA STATE UNIVERSITY

Department of Electrical and Computer Engineering

Lecture 21: Page Table

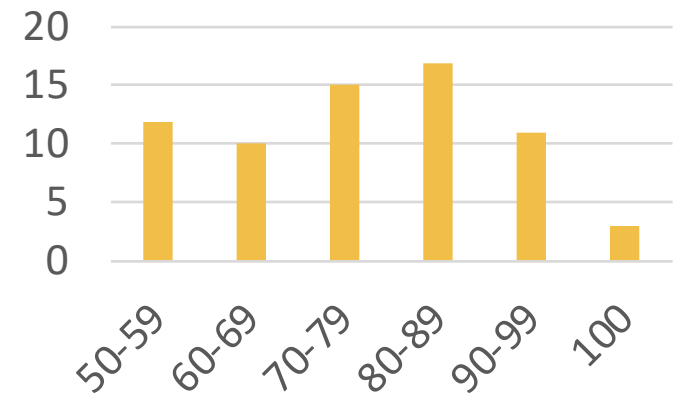


Agenda

- **Recap**
- **Page Table**
 - **Location of Page Table**
 - **Page Table Entry (PTE)**

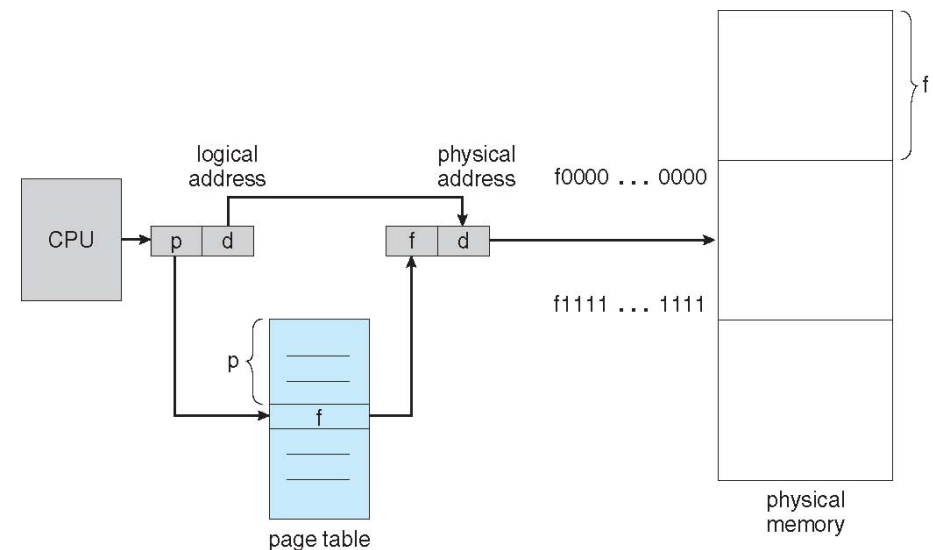
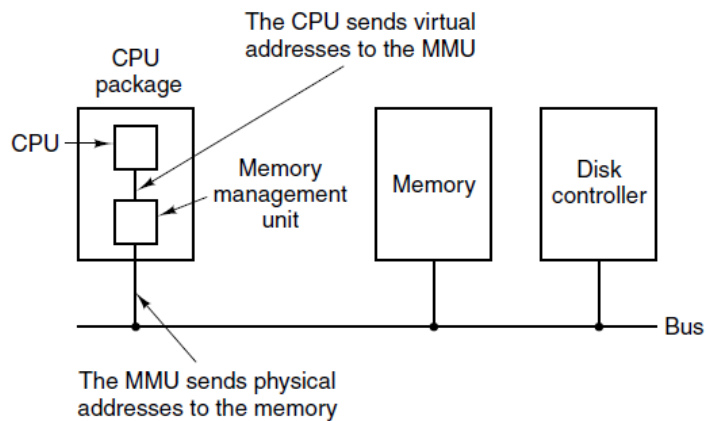
Recap

- Midterm 1 feedback
 - Observations
 - Top scorers
 - attended the class regularly
 - participated in Q&A actively
 - Those sent emails to me about class participation are above average in general
 - followed the instructions
 - Don't lose heart if you did not do well
 - score on paper does not necessarily reflect your potential
 - You may want to adjust your study method
 - Attend (more) class, ask (more) questions, ...
 - Welcome to visit my office hours



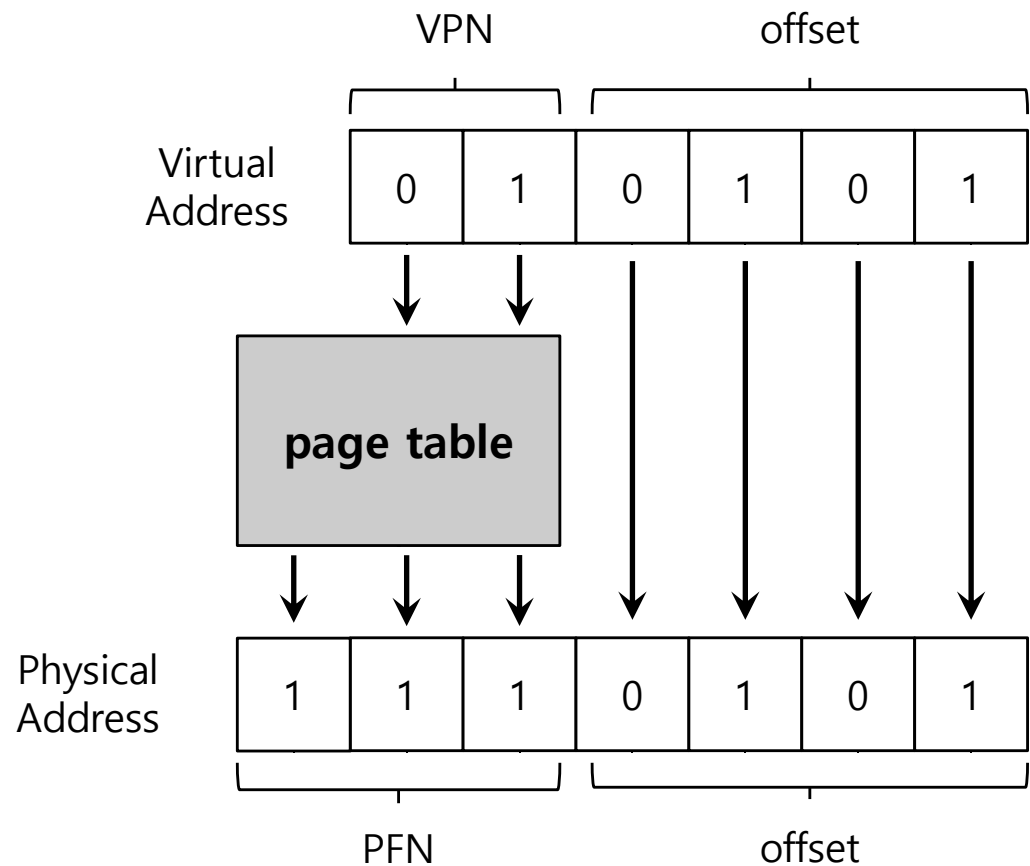
Recap

- Address Translation via **Page Table**
 - Each virtual address is divided into two parts:
 - VPN: virtual page number (p)
 - used as an index into the **page table**
 - Offset: offset within the page (d)



Recap

- Address Translation via **Page Table**
 - Example: the virtual address 21 in 64-byte address space
 - $64 = 2^6$
 - 21: 0x010101



Agenda

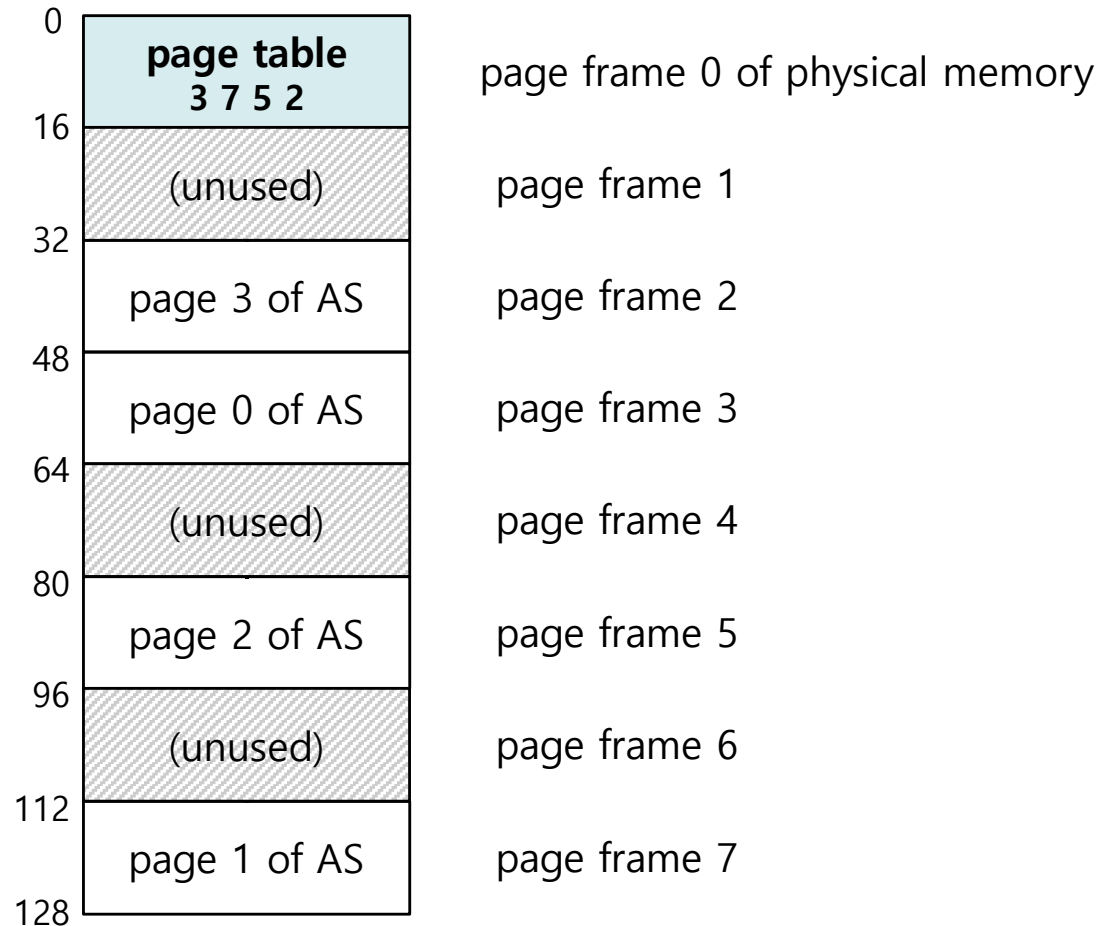
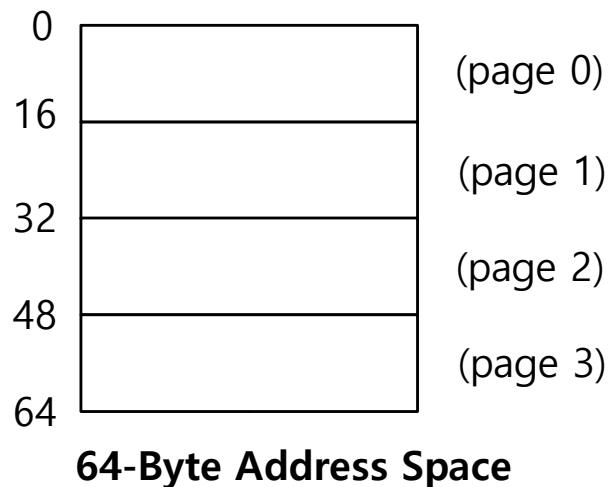
- ~~Recap~~
- Page Table
 - Location of Page Table
 - Page Table Entry (PTE)

Where Are Page Tables Stored?

- The page table for each process is stored in physical memory
 - accessible by kernel
 - Page-table base register (PTBR) points to the page table
 - Page-table length register (PTLR) indicates size of the page table

Where Are Page Tables Stored?

- The page table for each process is stored in physical memory



What Is in the Page Table?

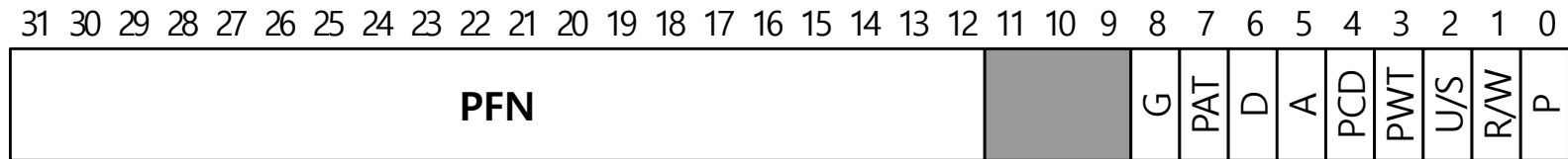
- The page table is just a **data structure** that is used to map the virtual address to physical address.
 - Simplest form: a linear page table, an array
- The OS **indexes** the array by VPN, and looks up the page-table entry
 - Each entry must contain the base address of each page in the physical memory
 - may also contain additional bits for other purposes (e.g., protection)

What Is in the Page Table?

- Common Flags Of Page Table Entry (PTE)
 - **Valid Bit:** Indicating whether the particular translation is valid
 - **Protection Bit:** Indicating whether the page could be read from, written to, or executed from
 - **Present Bit:** Indicating whether this page is in physical memory or on disk(swapped out)
 - **Dirty Bit:** Indicating whether the page has been modified since it was brought into memory
 - **Reference Bit(Accessed Bit):** Indicating that a page has been accessed

What Is in the Page Table?

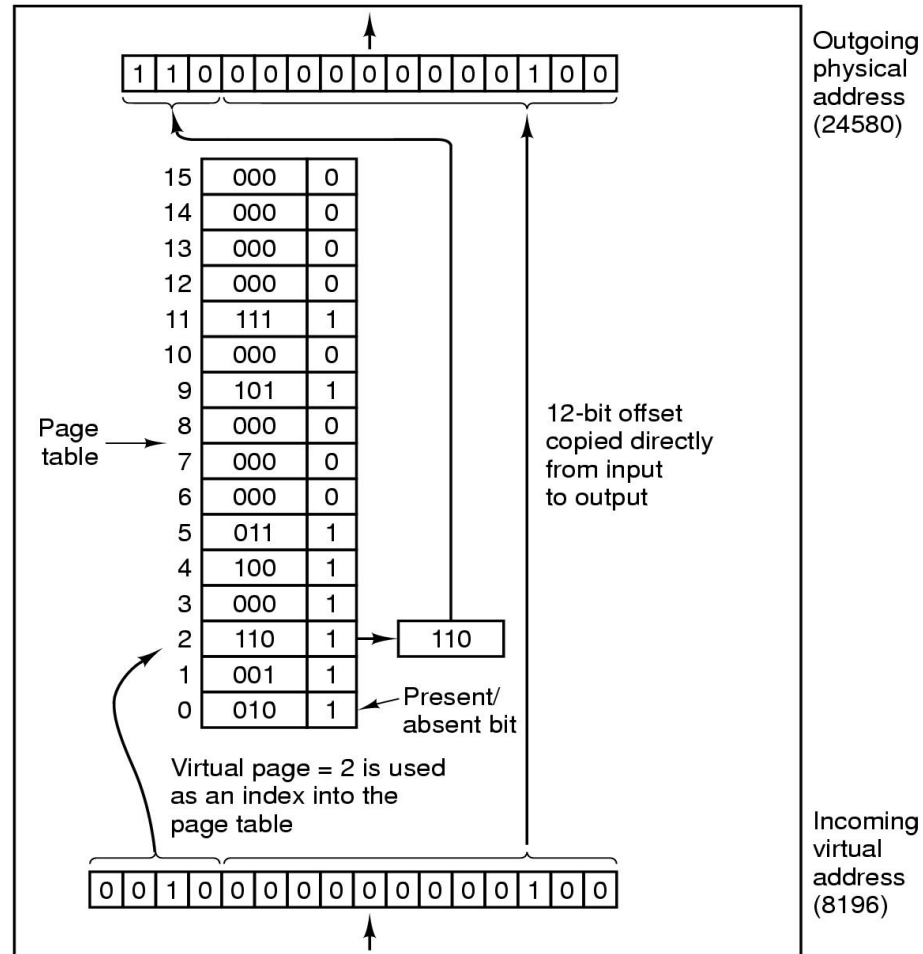
- Example: x86 Page Table Entry (PTE)



- P: present
- R/W: read/write bit
- U/S: supervisor
- A: accessed bit
- D: dirty bit
- PFN: the page frame number

What Is in the Page Table?

- Address translation using a page table

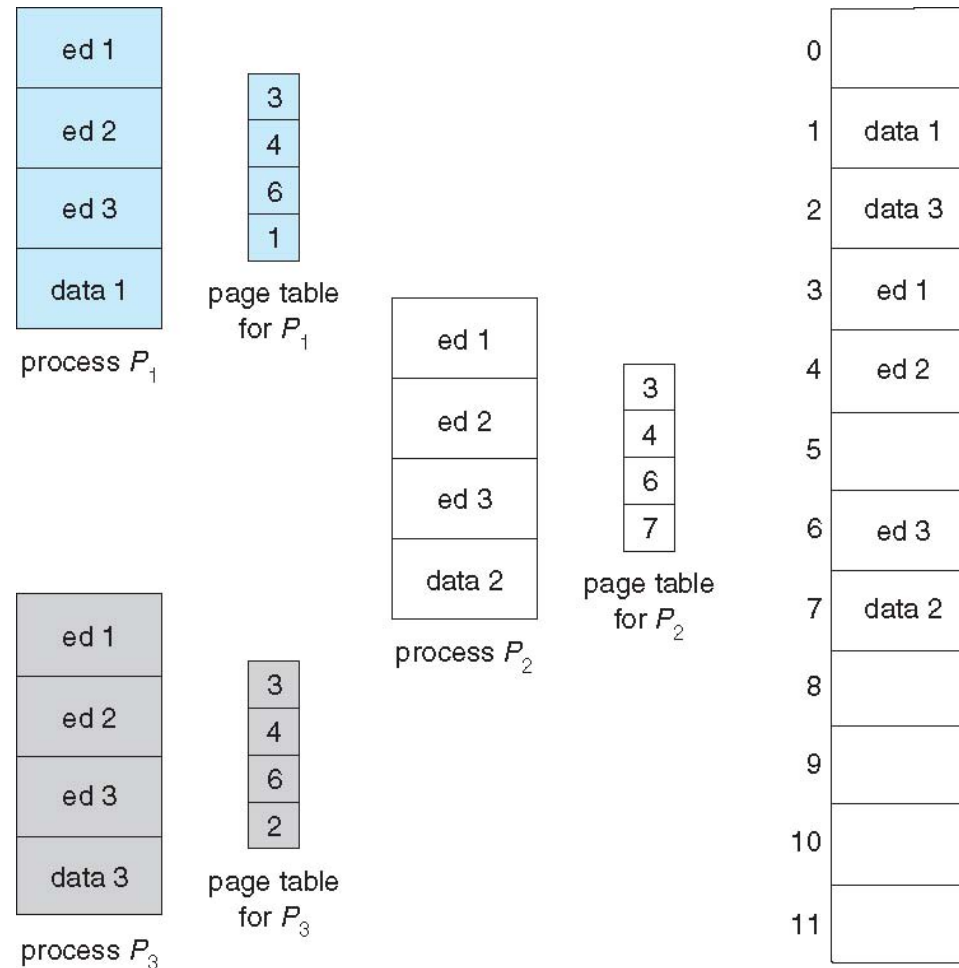


Shared Pages

- One page in the physical memory is mapped to the address spaces of multiple processes
 - Often used for code sharing
 - One copy of code shared among processes (i.e., text editors, compilers, window systems)
 - Similar to multiple threads sharing the same process space
 - Also useful for inter-process communication if sharing of read-write pages is allowed
 - Each process keeps a separate copy of the private code/data

Shared Pages

- Example of shared pages



What if requested page not in memory?

- **Page Fault**

- MMU reports a page fault to CPU
- CPU gives control to the OS (page fault handler routine)
- OS fetches a page from the disk
 - May need to evict an existing page from memory (page replacement policy)
- Instruction is restarted

Potential Issues of Paging

- Time
 - For every user memory access, one additional memory access is needed (for accessing page table in memory)
- Space
 - Page tables can get awfully large
 - 32-bit address space (4GB) with 4KB pages
 - 12 bits for offset within a page ($4K=2^{12}$)
 - 20 bits for VPN ($32 - 12 = 20$)
 - $4MB = 2^{20} \text{ entries} * 4 \text{ Bytes per page table entry}$
 - Each process needs to have a page table!

Agenda

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Questions?



*acknowledgement: slides include content from “Modern Operating Systems” by A. Tanenbaum, “Operating Systems Concepts” by A. Silberschatz etc., “Operating Systems: Three Easy Pieces” by R. Arpaci-Dusseau etc., and anonymous pictures from internet.