

# **CprE 381: Computer Organization and Assembly Level Programming**

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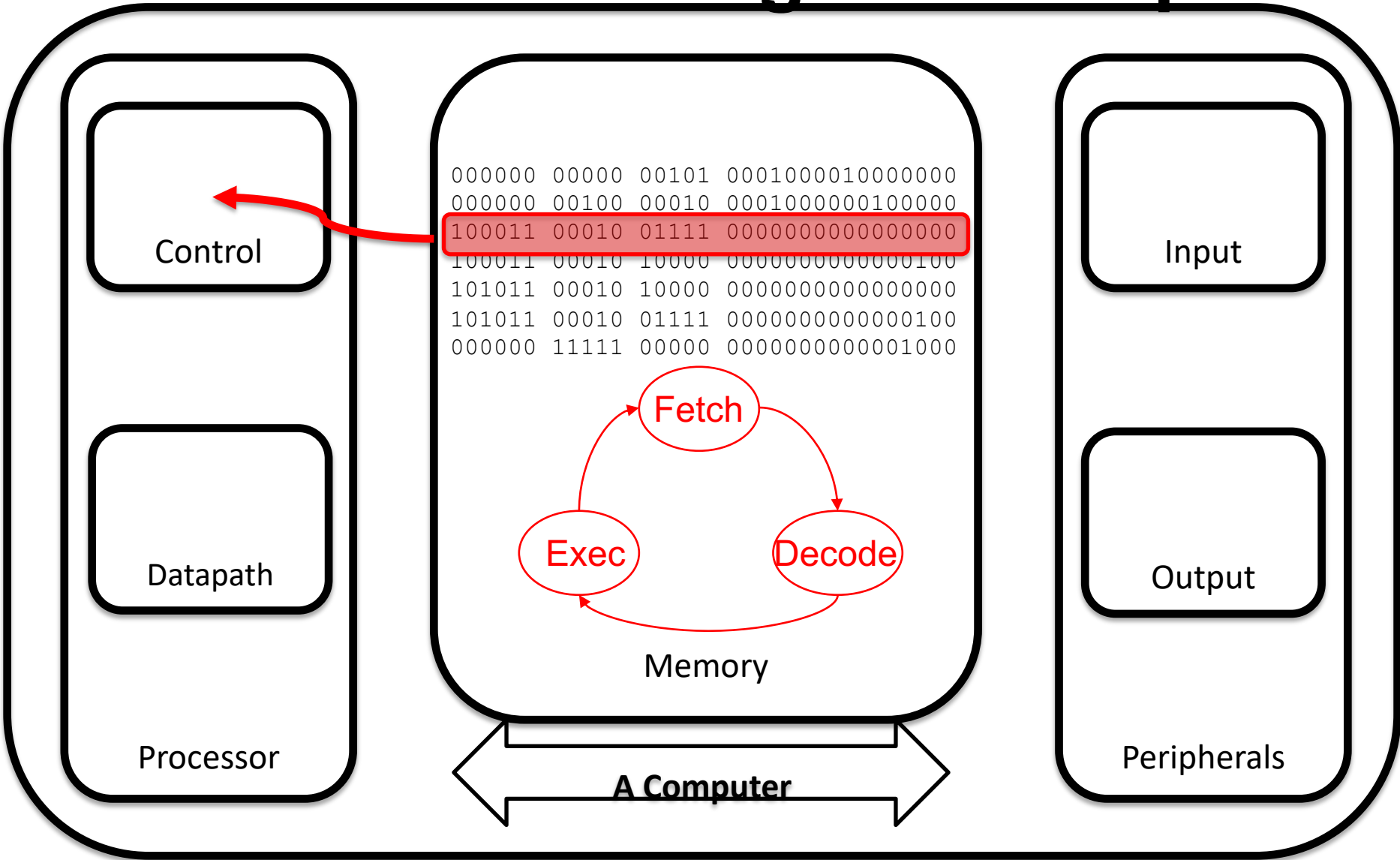
# Administrative

- HW1 posted (due Jan 28b – TODAY!)
  - Hard deadline – submit what you have several minutes before deadline (at least)
  - Typeset: in the future, generate figures with professional software (don't include snapshots of paper or whiteboard)
    - Visio, Powerpoint, etc. are all free
    - You will begin to lose points
  - Show your work!
    - HW intended to be formative → in order to get feedback, need to show process
    - Partial credit if answer is wrong
  - ABI vs ISA vs uArch → Who cares?
    - ISA is the interface between HW and SW
    - uArch (i.e., implementations) is what the HW designer can choose
    - ABI further constrains what SW will interoperate

# Administrative

- Labs
  - Prelab must be completed **prior** to start of lab
    - Starting with Lab 3, separate submission assignment – no points if not submitted by start of your lab
  - Submit what you have done before start of lab!
- Added Prof Office Hours:
  - M 10am (following lecture, walk back to office)
- TA Office Hours
  - You may go to any TA's office hours
  - You may demo at office hours, although I'd prefer you do so at the TA for your section
  - All office hours are located in 2050 Coover
  - **Ashraf: M 11am – Noon**
  - **Ryan: T 11am – Noon**
  - **Trent: W Noon – 1pm**
  - **Rohit: R 3pm – 4pm**

# Review: Stored Program Computer



# Review: MIPS Simple Arithmetic

Instruction	Example	Meaning	Comments
add	<code>add \$1,\$2,\$3</code>	$\$1 = \$2 + \$3$	3 operands; Overflow
subtract	<code>sub \$1,\$2,\$3</code>	$\$1 = \$2 - \$3$	3 operands; Overflow
add immediate	<code>addi \$1,\$2,100</code>	$\$1 = \$2 + 100$	+ constant; Overflow
add unsigned	<code>addu \$1,\$2,\$3</code>	$\$1 = \$2 + \$3$	3 operands; No overflow
sub unsigned	<code>subu \$1,\$2,\$3</code>	$\$1 = \$2 - \$3$	3 operands; No overflow
add imm unsign	<code>addiu \$1,\$2,100</code>	$\$1 = \$2 + 100$	+ constant; No overflow

- Your task: check out logical and shift instructions

# Review: MIPS Integer Load/Store

Instruction	Example	Meaning	Comments
store word	<b>sw</b> \$1, 8 (\$2)	<b>Mem</b> [ 8+\$2 ] =\$1	Store word
store half	<b>sh</b> \$1, 6 (\$2)	<b>Mem</b> [ 6+\$2 ] =\$1	Stores only lower 16b
store byte	<b>sb</b> \$1, 5 (\$2)	<b>Mem</b> [ 5+\$2 ] =\$1	Stores only lowest byte
load word	<b>lw</b> \$1, 8 (\$2)	<b>\$1</b> = <b>Mem</b> [ 8+\$2 ]	Load word
load halfword	<b>lh</b> \$1, 6 (\$2)	<b>\$1</b> = <b>Mem</b> [ 6+\$2 ]	Load half; sign extend
load half unsign	<b>lhu</b> \$1, 6 (\$2)	<b>\$1</b> = <b>Mem</b> [ 6+\$2 ]	Load half; zero extend
load byte	<b>lb</b> \$1, 5 (\$2)	<b>\$1</b> = <b>Mem</b> [ 5+\$2 ]	Load byte; sign extend
load byte unsign	<b>lbu</b> \$1, 5 (\$2)	<b>\$1</b> = <b>Mem</b> [ 5+\$2 ]	Load byte; zero extend

# More MIPS Control Flow

Instruction	Example	Meaning
jump	<code>j L</code>	<code>goto L</code>
jump register	<code>jr \$1</code>	<code>goto value in \$1</code>
jump and link	<code>jal L</code>	<code>goto L and set \$ra</code>
jump and link register	<code>jalr \$1</code>	<code>goto \$1 and set \$ra</code>
branch equal	<code>beq \$1,\$2,L</code>	<code>if (\$1 == \$2) goto L</code>
branch not equal	<code>bne \$1,\$2,L</code>	<code>if (\$1 != \$2) goto L</code>
branch less than 0	<code>bltz \$1,L</code>	<code>if (\$1 &lt; 0) goto L</code>
branch less than / eq 0	<code>blez \$1,L</code>	<code>if (\$1 &lt;= 0) goto L</code>
branch greater than 0	<code>bgtz \$1,L</code>	<code>if (\$1 &gt; 0) goto L</code>
branch greater than / eq 0	<code>bgez \$1,L</code>	<code>if (\$1 &gt;= 0) goto L</code>

# Preview: C Code Example

Simple C procedure:  $\text{sum\_pow2} = 2^{b+c}$

```
1: int sum_pow2 (int b, int c)
2: {
3:     int pow2[8] = {1, 2, 4, 8, 16, 32, 64, 128};
4:     int a, ret;
5:     a = b + c;
6:     if (a < 8)
7:         ret = pow2[a];
8:     else
9:         ret = 0;
10:    return(ret);
11: }
```



# Equivalent MIPS Assembly

```
sum_pow2:                # $a0 = 1, $a1 = c
    addu $a0,$a0,$a1      # a = 2*c, $a0 = a
    slti $v0,$a0,8        # $v0 = 1 if a < 8
    beq $v0,$zero,End     # if $v0==0
    addiu
```

POW!

Exceed

Return:

```
    jr r0                # return sum_pow2
```

# Equivalent MIPS Assembly

```
sum_pow2:                # $a0 = b, $a1 = c
    addu $a0,$a0,$a1      # a = b + c, $a0 = a
    slti $v0,$a0,8        # $v0 = a < 8
    beq $v0,$zero,Exceed  # goto Exceed if $v0==0
    addiu $v1,$sp,8        # $v1 = pow2 address
    sll $v0,$a0,2          # $v0 = a*4
    addu $v0,$v0,$v1       # $v0 = pow2 + a*4
    lw $v0,0($v0)         # $v0 = pow2[a]
    j Return              # goto Return
```

Exceed:

```
    addu $v0,$zero,$zero  # $v0 = 0
```

Return:

```
    jr ra                # return sum_pow2
```

# Support for Simple Branches Only

- Notice there is no branch less than instruction for comparing two registers?
  - The reason is that such an instruction would be too complicated and might require a longer cycle time
  - Therefore, two conditionals that do not compare against zero take at least two instructions
    - First is a set
    - Second is a conditional branch
- We'll see this later as a design trade-off
  - Less time per instruction vs. fewer instructions
    - How do you decide what to do?
  - Other RISC ISAs made a different choice (e.g. HP's PA-RISC)

# MIPS Comparisons

Instruction	Example	Meaning	Comments
set less than	<code>slt \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than signed
set less than imm	<code>slti \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const signed
set less unsgn	<code>sltu \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than unsigned
slt imm unsgn	<code>sltiu \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const unsigned

- C

```
if (a < 8) goto Exceed
```

```
if (a <= 8) goto Exceed
```

```
if (8 < a) goto Exceed
```

```
if (8 <= a) goto Exceed
```

# MIPS Comparisons

Instruction	Example	Meaning	Comments
set less than	<code>slt \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than signed
set less than imm	<code>slti \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const signed
set less unsgn	<code>sltu \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than unsigned
slt imm unsgn	<code>sltiu \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const unsigned

- C

**if (a < 8) goto Exceed**

`slti $v0, $a0, 8                   # $v0 = $a0 < 8`

`bne $v0, $zero, Exceed # goto if $v0 != 0`

# MIPS Comparisons

Instruction	Example	Meaning	Comments
set less than	<code>slt \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than signed
set less than imm	<code>slti \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const signed
set less unsgn	<code>sltu \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than unsigned
slt imm unsgn	<code>sltiu \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const unsigned

- C

`if (a <= 8) goto Exceed`

`slti $v0, $a0, 9                   # $v0 = $a0 <= 8`

`bne $v0, $zero, Exceed # goto if $v0 != 0`

# MIPS Comparisons

Instruction	Example	Meaning	Comments
set less than	<code>slt \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than signed
set less than imm	<code>slti \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const signed
set less unsgn	<code>sltu \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than unsigned
slt imm unsgn	<code>sltiu \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const unsigned

- C

`if (8 < a) goto Exceed`

`slti $v0, $a0, 9                   # $v0 = $a0 < 9`

`beq $v0, $zero, Exceed # goto if $v0 != 0`

# MIPS Comparisons

Instruction	Example	Meaning	Comments
set less than	<code>slt \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than signed
set less than or equal			
set less than or equal to unsigned			
set less than or equal to unsigned integer			

**In-class Assessment!**

**Access Code: <=|>?**

Note: sharing access code to those outside of classroom or using access while outside of classroom is considered cheating

`if (8 <= a) goto Exceed`



# MIPS Comparisons

Instruction	
set less than	slt
set less than imm	slti
set less unsgn	sltu
slt imm unsgn	sltiu



Comments
comp less than signed
comp w/const signed
comp less than unsigned
comp w/const unsigned

- C  
if (8 <= \$a0)  
{  
    slti \$v0, 8  
    beq \$v0, \$zero, then  
}

\$a0 < 8  
if \$v0 != 0

# MIPS Comparisons

Instruction	Example	Meaning	Comments
set less than	<code>slt \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than signed
set less than imm	<code>slti \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const signed
set less unsgn	<code>sltu \$1,\$2,\$3</code>	$\$1 = (\$2 < \$3)$	Comp less than unsigned
slt imm unsgn	<code>sltiu \$1,\$2,100</code>	$\$1 = (\$2 < 100)$	Comp w/const unsigned

- C

`if (8 <= a) goto Exceed`

`slti $v0, $a0, 8                   # $v0 = $a0 < 8`

`beq $v0, $zero, Exceed # goto if $v0 != 0`

# While Loops in C

- Consider a `while` loop

```
while (A[i] == k)
    i = i + j;
```
- MIPS assembly loop
- Assume `i=$s0`, `j=$s1`, `k=$s2`, `&A=$s3`

# While Loops in C

- Consider a `while` loop

```
while (A[i] == k)
    i = i + j;
```

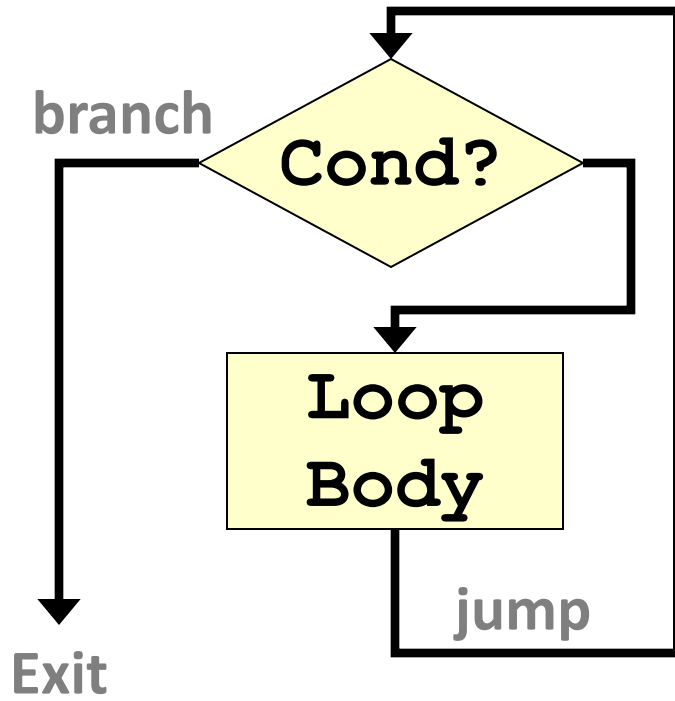
- MIPS assembly loop
- Assume `i=$s0`, `j=$s1`, `k=$s2`, `&A=$s3`

```
Loop: sll $t0, $s0, 2      # $t0 = 4 * i
      addu $t1, $t0, $s3   # $t1 = &(A[i])
      lw $t2, 0($t1)      # $t2 = A[i]
      bne $t2, $s2, Exit   # goto Exit if !=
      addu $s0, $s0, $s1   # i = i + j
      j Loop              # goto Loop
Exit:
```

- Basic block:
  - Maximal sequence of instructions without branches or branch targets

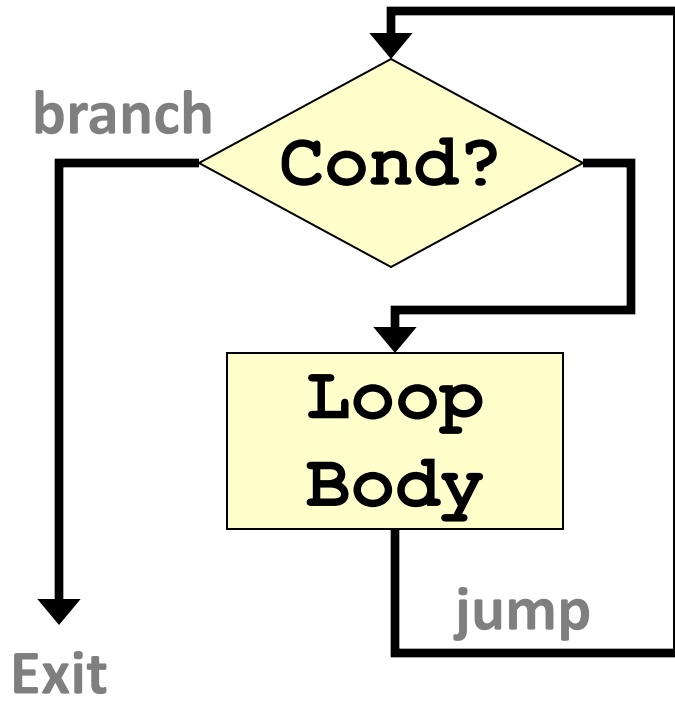
# Improve Loop Efficiency

- Code uses two branches per iteration:

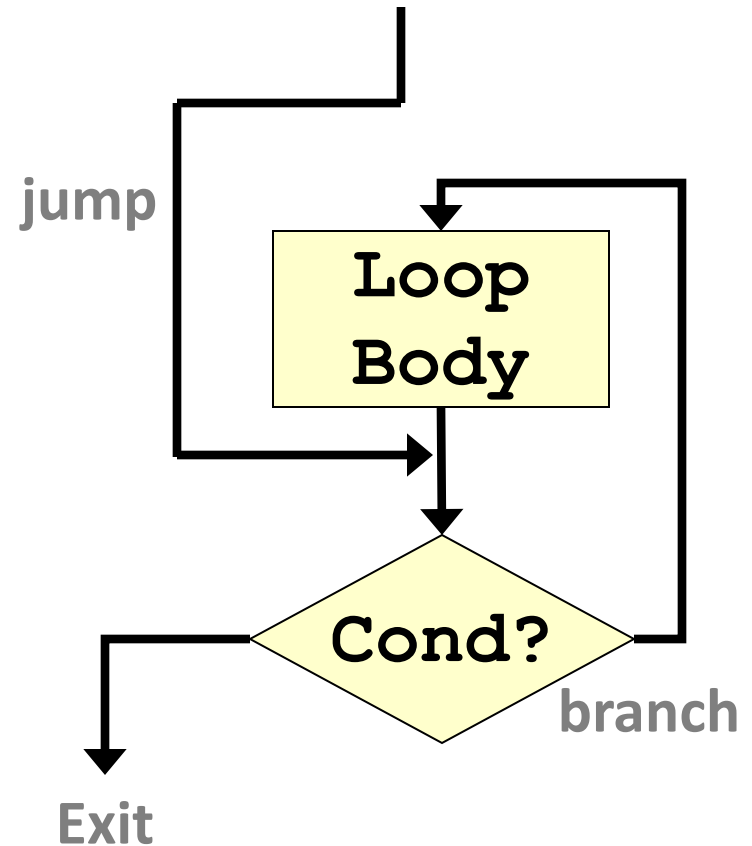


# Improve Loop Efficiency

- Code uses two branches per iteration:



- More efficient structure:



# Acknowledgments

- These slides contain material developed and copyright by:
  - Joe Zambreno (Iowa State)
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