

## 2. Link State Routing

- ⊕ Each node maintains the three lists

H, D, C

- ⊕ Each node broadcasts its local link state information to all other nodes in the network

- Often via flooding

Protocol : local exchanged info globally via flooding

Each node has complete information about all link information in the network.

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## Link State Routing (Algorithm)

- ⊕ Each node will calculate its shortest paths to all other nodes upon receiving new link state information

- Dijkstra's Shortest-Path Routing Algorithm

- An iterative algorithm to find the shortest paths from a source node to all other nodes in the network

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## Dijkstra's Algorithm

### ⊕ Notations:

- s: source node
- N: the set of nodes whose shortest paths have already been found

### ⊕ Initialization Step

- $N = \{s\}$ ,  $D_{sj} = C_{sj}$  and  $H_{sj} = j$  for all j

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## Dijkstra's Algorithm

### ⊕ Notations:

- s: source node
- N: the set of nodes whose shortest paths have already been found

### ⊕ Step A: (Finding the next closest node i)

- Find node  $i \notin N$  such that  $D_{si} = \min D_{sj}$  for  $j \notin N$
- Add i to N
- If N contains all the nodes, Stop

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# Dijkstra's Algorithm

## Notations:

- s: source node
- N: the set of nodes whose shortest paths have already been found

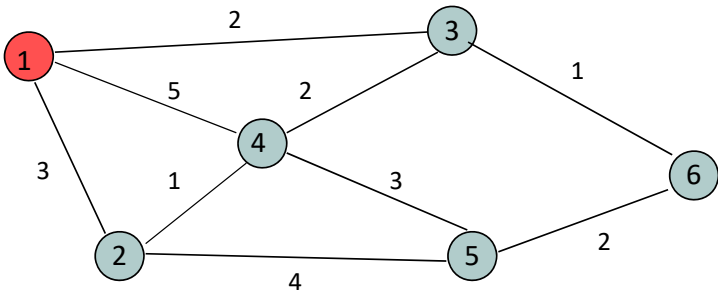
## Step B: (Updating minimum costs after node i is added to N)

- For each node  $j \notin N$ 
  - if  $(D_{si} + C_{ij}) < D_{sj}$  then  $D_{sj} = D_{si} + C_{ij}$  and  $H_{sj} = H_{si}$
- Go to Step A



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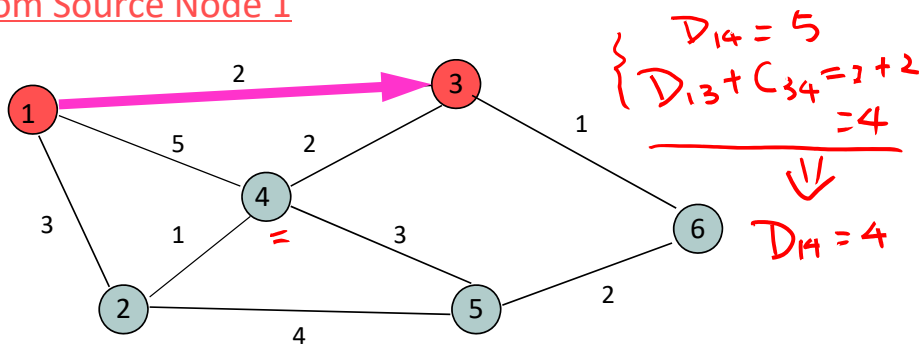
## Example: From Source Node 1



Iteration	N	H <sub>12</sub> ; D <sub>12</sub>	H <sub>13</sub> ; D <sub>13</sub>	H <sub>14</sub> ; D <sub>14</sub>	H <sub>15</sub> ; D <sub>15</sub>	H <sub>16</sub> ; D <sub>16</sub>
Initial	{1}	2; 3	3; 2	4; 5	5; ∞	6; ∞

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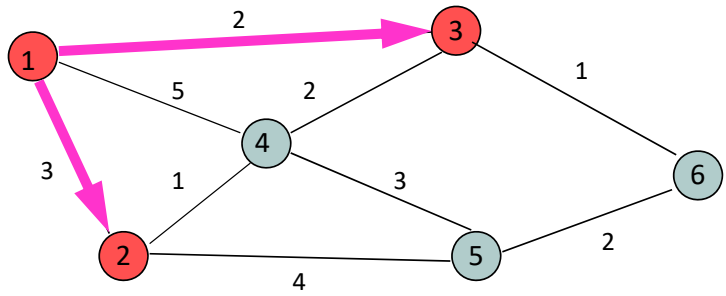
Example: From Source Node 1



Iteration	N	$H_{12}; D_{12}$	$H_{13}; D_{13}$	$H_{14}; D_{14}$	$H_{15}; D_{15}$	$H_{16}; D_{16}$
Initial	{1}	2; 3	3; 2	4; 5	5; $\infty$	6; $\infty$
1	{1,3}	2; 3	3; 2	3; 4	5; $\infty$	3; 3

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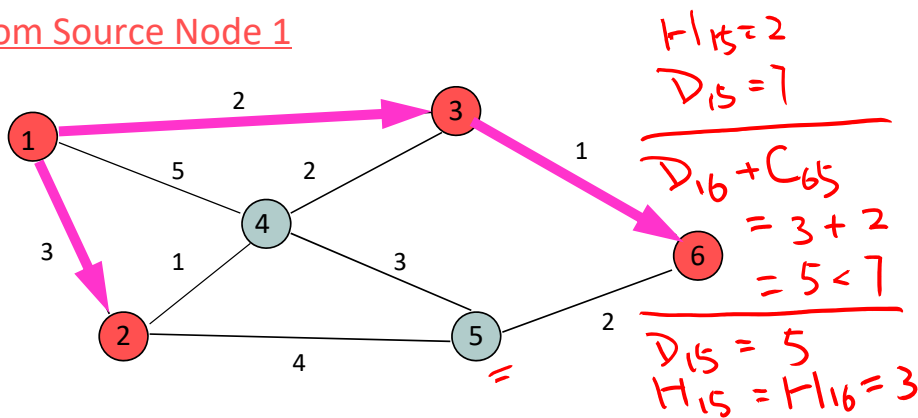
Example: From Source Node 1



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Initial	{1}	2; 3	3; 2	4; 5	5; $\infty$	6; $\infty$
1	{1,3}	2; 3	3; 2	3; 4	5; $\infty$	3; 3
2	{1,2,3}	2; 3	3; 2	3; 4	2; 7	3; 3

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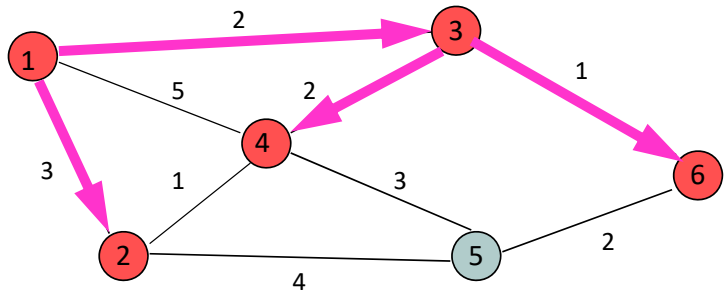
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Initial	{1}	2; 3	3; 2	4; 5	5; $\infty$	6; $\infty$
1	{1,3}	2; 3	3; 2	3; 4	5; $\infty$	3; 3
2	{1,2,3}	2; 3	3; 2	3; 4	2; 7	3; 3
3	{1,2,3,6}	2; 3	3; 2	3; 4	3; 5	3; 3

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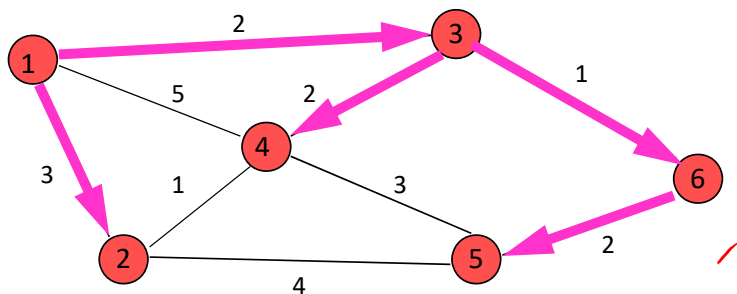
Example: From Source Node 1



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Initial	{1}	2; 3	3; 2	4; 5	5; $\infty$	6; $\infty$
1	{1,3}	2; 3	3; 2	3; 4	5; $\infty$	3; 3
2	{1,2,3}	2; 3	3; 2	3; 4	2; 7	3; 3
3	{1,2,3,6}	2; 3	3; 2	3; 4	3; 5	3; 3
4	{1,2,3,4,6}	2; 3	3; 2	3; 4	3; 5	3; 3

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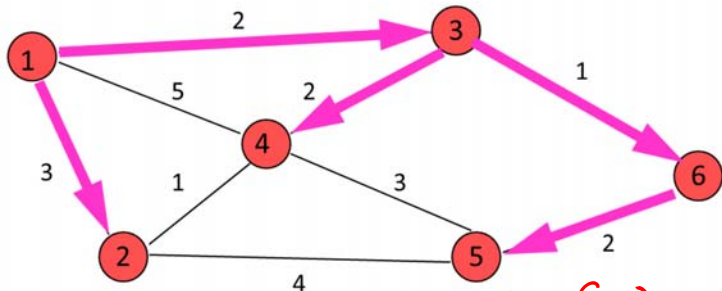
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Initial	{1}	2; 3	3; 2	4; 5	5; $\infty$	6; $\infty$
1	{1,3}	2; 3	3; 2	3; 4	5; $\infty$	3; 3
2	{1,2,3}	2; 3	3; 2	3; 4	2; 7	3; 3
3	{1,2,3,6}	2; 3	3; 2	3; 4	3; 5	3; 3
4	{1,2,3,4,6}	2; 3	3; 2	3; 4	3; 5	3; 3
5	{1,2,3,4,5,6}	2; 3	3; 2	3; 4	3; 5	3; 3

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Example: After Algorithm Terminates, Information at Node 1



dest	$H_{ij}$	$D_{ij}$	$C_{ij}$
1	1	0	0
2	2	3	3
3	3	2	2
4	3	4 ←	5 ←
5	3	5	$\infty$
6	3	3	$\infty$

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## Reaction to Link Failure

- ⊕ If a link is broken,
  - ➡ Affected nodes set link cost to infinity & flood the network with update packets
  - ➡ All nodes immediately update their link database & re-calculate their shortest paths
  - ➡ Recovery is very quick
  
- ⊕ Link State Routing is NOT loop-free
  - ➡ Due to delay in link state information propagation

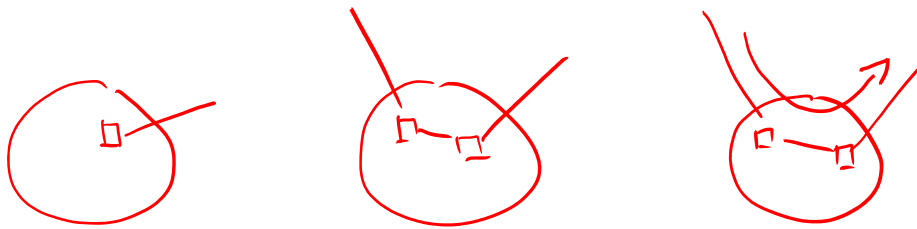
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## Internet Routing Protocols

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## Autonomous Systems

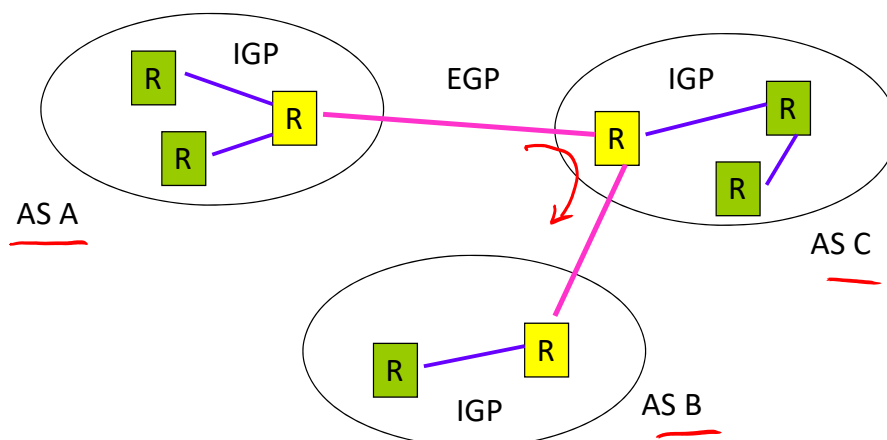
- ✦ **Autonomous System (AS)** is defined as a set of routers or networks administered by a single organization
  - ➡ **Stub AS:** has only a single connection to the outside world
  - ➡ **Multi-homed AS:** has multiple connections to the outside world, but refuses to carry transit traffic
  - ➡ **Transit AS:** has multiple connections to the outside world, and can carry both transit and local traffic



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## Inter and Intra AS Routing

- ✦ **IGP (Interior Gateway Protocol):** routing within an AS
  - ➡ RIP, OSPF
- ✦ **EGP (Exterior Gateway Protocol):** routing between ASs
  - ➡ BGPv4



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## RIP (Routing Information Protocol)

- ✦ Distance Vector Routing Protocol
  - ✦ Split Horizon with Poisoned Reverse
- ✦ Runs on top of UDP, port # 520, “**routed**” BSD Unix program
- ✦ Routing Metric: number of hops
- ✦ Max number of hops is limited to 15
  - ✦ Suitable for small networks (local area environments)
  - ✦ 16 is reserved to represent infinity
  - ✦ Small number helps to limit the Counting-to-Infinity Problem

Routing Loop

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## OSPF (Open Shortest Path First)

- ✦ Link State Routing Protocol
- ✦ OSPF runs directly over IP
  - ✦ Value in the protocol field of IP headers: 89
- ✦ OSPF typically converges faster than RIP when there is a failure in the network

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## BGP (Border Gateway Protocol)

### ⊕ Path Vector Routing Protocol

- ➡ Avoid routing loops

### ⊕ BGP

- ➡ Is a reachability protocol
- ➡ Uses TCP to send updates
  - Reliable transmission
  - Allow incremental updates
- ➡ Allows for policy routing
  - Path selection by policy rather than path optimality