

CPR E 431

## BASICS OF INFORMATION SYSTEM SECURITY

# Introduction to Cryptography Tools

Public Key Encryption



# Video Summary

- What is Public/Asymmetric Key Encryption
- Principles of Public-key Encryption
- Key Generation
- Public-key Encryption Assumptions
- Public-key Encryption Requirements



# Public-Key Encryption Structure

Publicly  
proposed by  
Diffie and  
Hellman in  
1976

Based on  
mathematical  
functions

## Asymmetric

- Uses two separate keys
- Public key and private key
- Public key is made public for others to use

Some form of  
protocol is  
needed for  
distribution

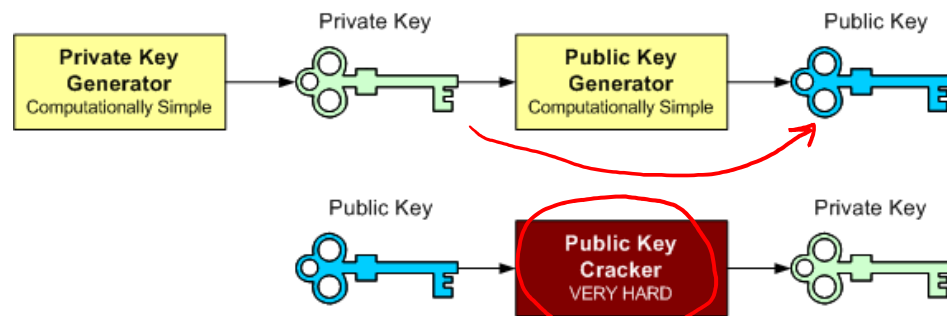
# Public-Key Encryption Structure

- Two different keys are used interchangeably to encrypt/decrypt the data
- The keys always come in pairs
- Each user is having two keys (one public and one private)

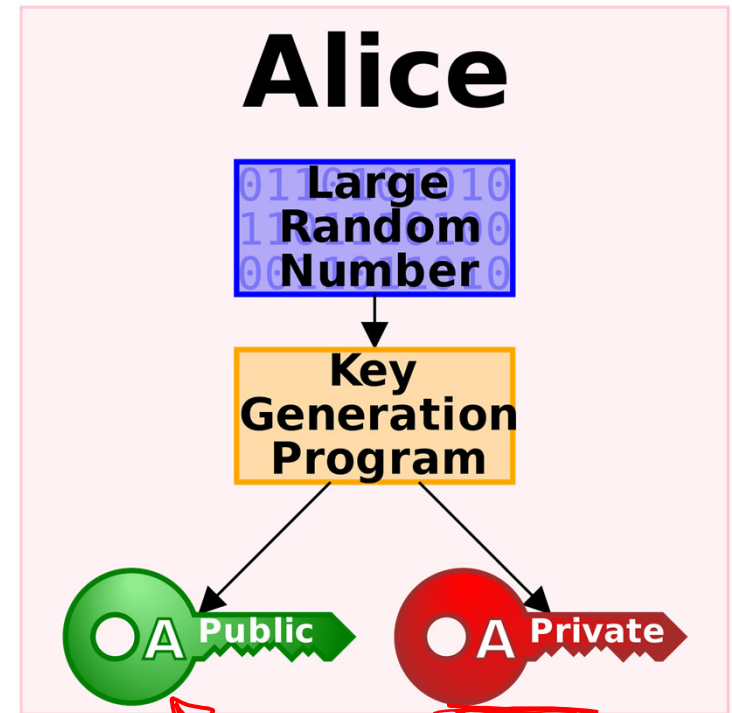
$PU_A$   $PR_A$



# Keys Generation



Source: <https://docs.huihoo.com/globus/gt3-tutorial/ch10s03.html>

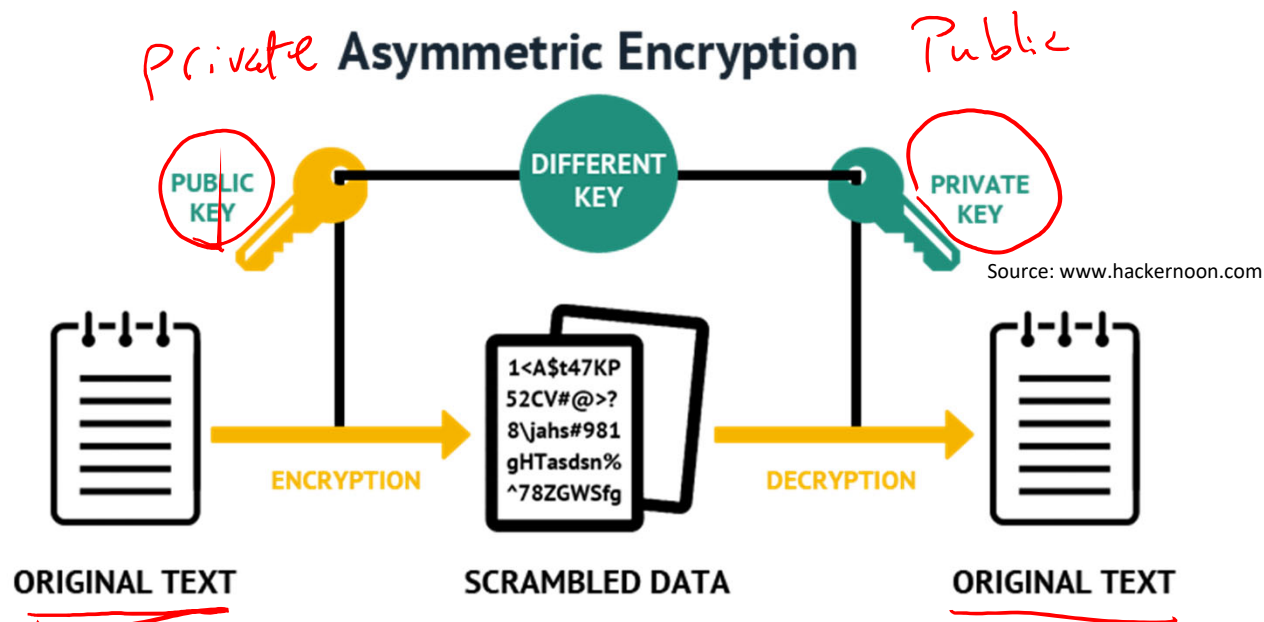


Source: [www.wikipedia.org](http://www.wikipedia.org)

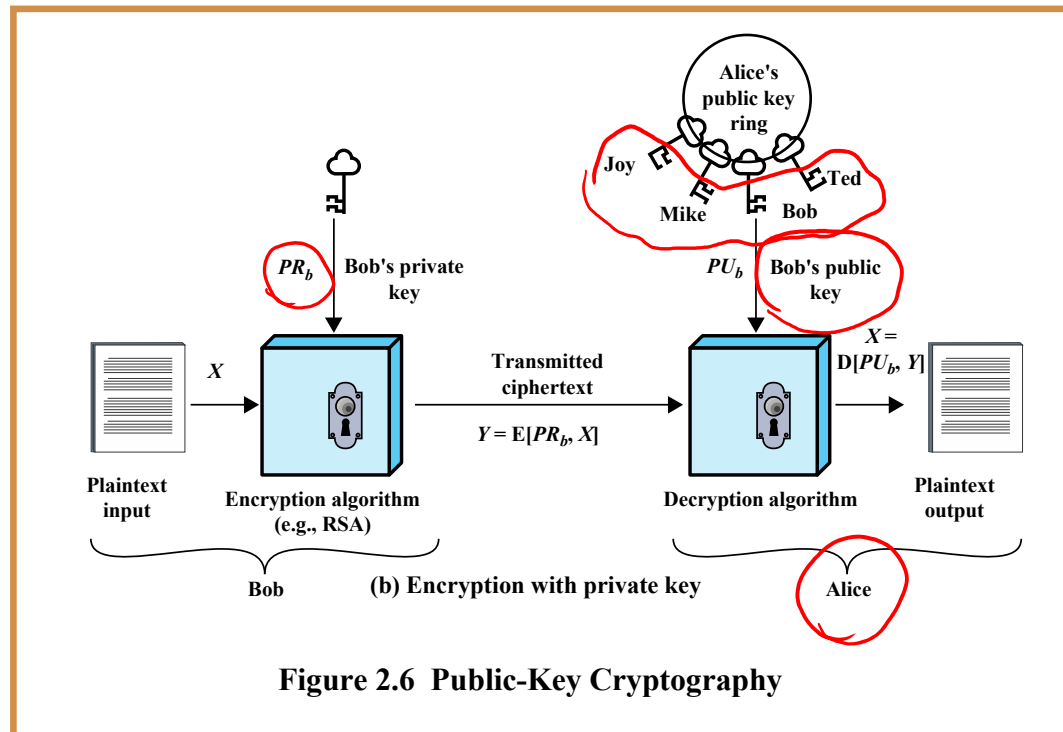
# Public and Private Keys

If you encrypted with public key  
then  
decrypt with private key

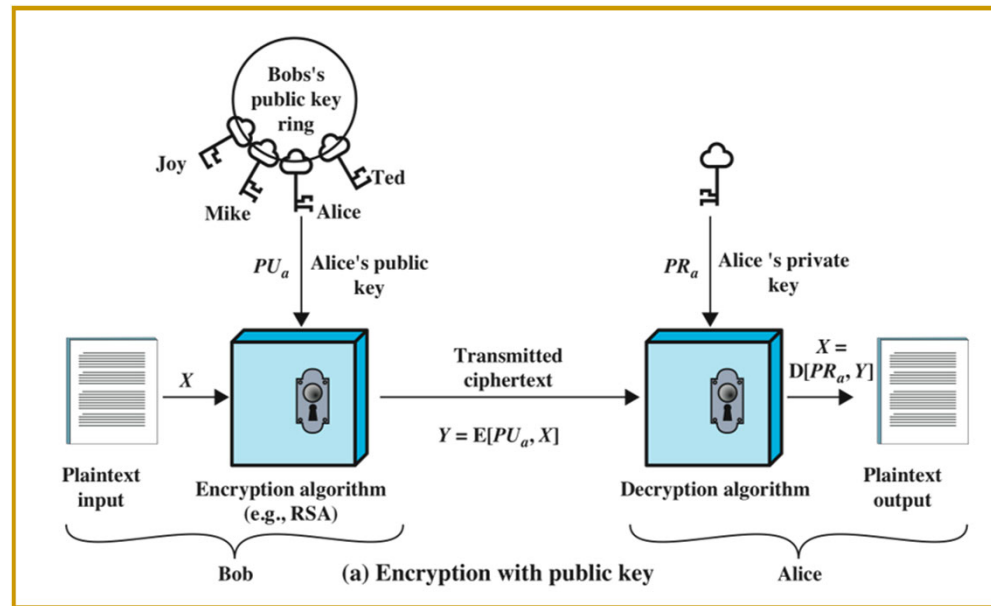
If you encrypted with private key  
then  
decrypt with public key







- User encrypts data using his or her own private key
- Anyone who knows the corresponding public key will be able to decrypt the message



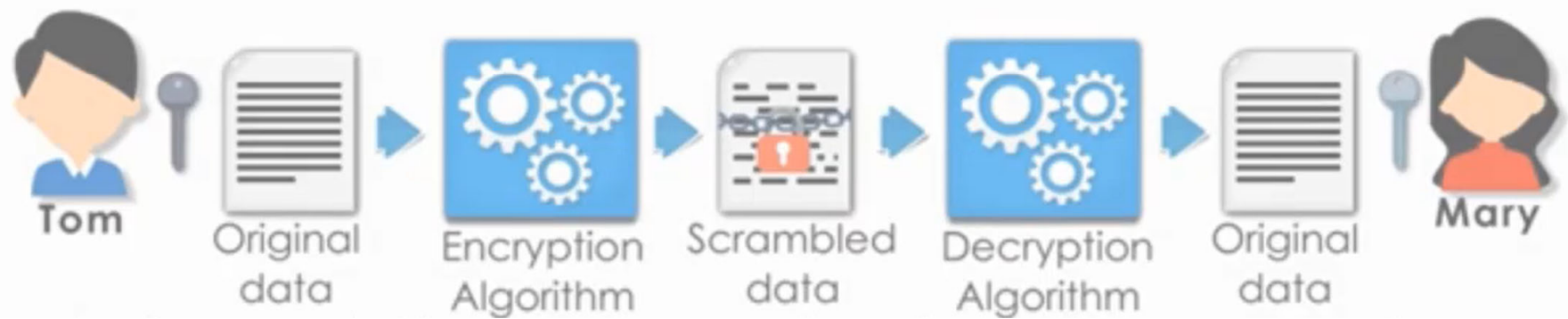
- User encrypts data using a certain public key
- Anyone who knows the corresponding private key will be able to decrypt the message





Tom wants to send an encrypted message to Mary





## Assumptions: Public Key Encryption

- ▶ There is a pair of keys, public ( $PU$ ) and private ( $PR$ ). One key from the pair is used for encryption, the other is used for decryption. Each entity has their own pair, e.g. ( $PU_A$ ,  $PR_A$ ).
- ▶ Encrypting a plaintext message,  $M$ , with a key, produces ciphertext  $C$ , e.g.  $C$  =  $E$ ( $PU_A$ ,  $M$ ).
- ▶ Decrypting ciphertext with the correct key will produce the original plaintext. The decrypter will be able to recognise that the plaintext is correct (and therefore the key is correct). E.g.  $M$  =  $D$ ( $PR_A$ ,  $C$ ).

# Requirements of Public-Key Cryptography

1. Computationally easy for  $B$  to generate pair  $(PU_b, PR_b)$
2. Computationally easy for  $A$ , knowing  $PU_b$  and message  $M$ , to generate ciphertext:

$$C = E(PU_b, M)$$

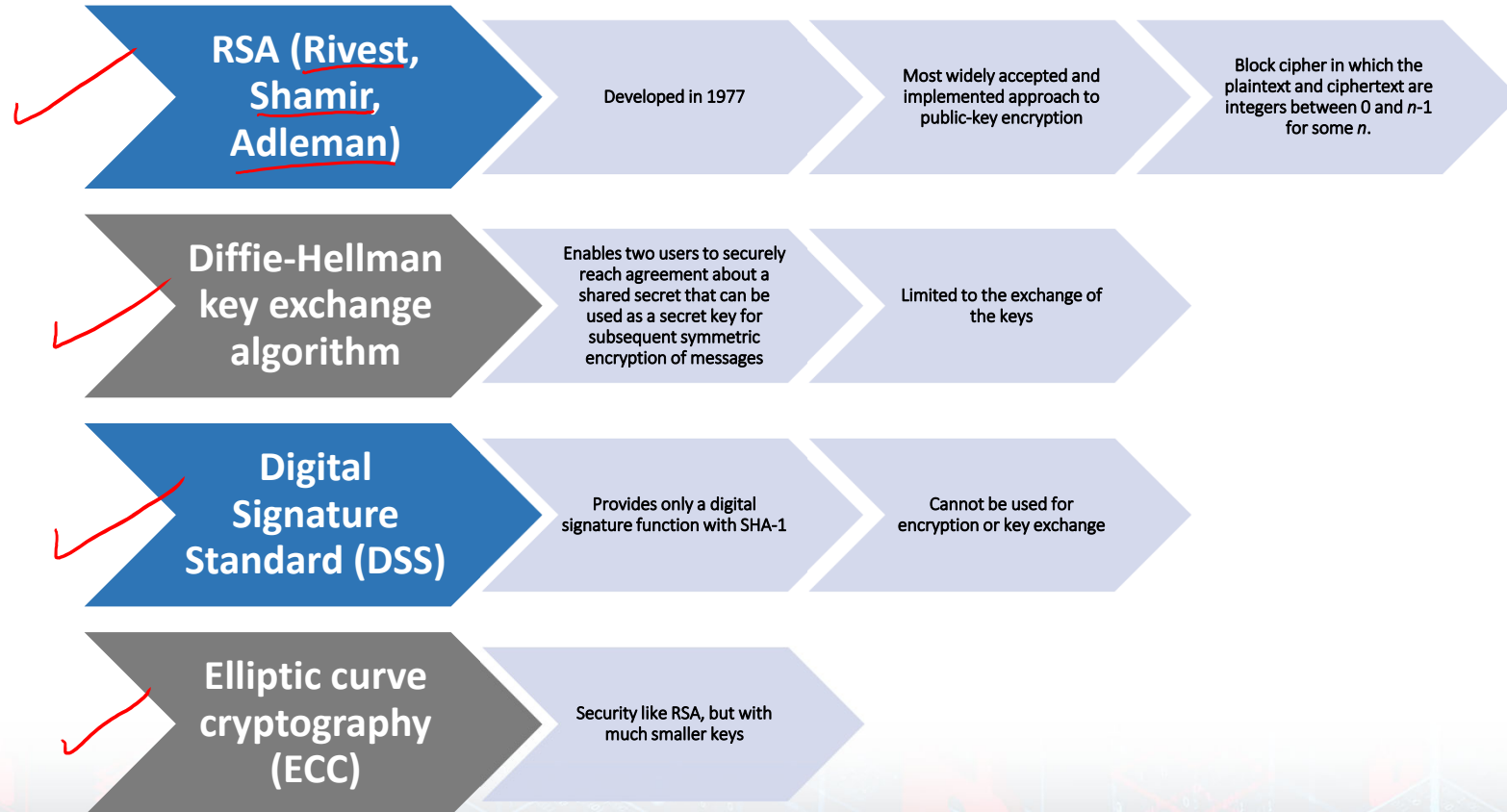
3. Computationally easy for  $B$  to decrypt ciphertext using  $PR_b$ :

$$M = D(PR_b, C) = D[PR_b, E(PU_b, M)]$$

4. Computationally infeasible for attacker, knowing  $PU_b$  and  $C$ , to determine  $PR_b$
5. Computationally infeasible for attacker, knowing  $PU_b$  and  $C$ , to determine  $M$
6. (Optional) Two keys can be applied in either order:

$$M = D[PU_b, E(PR_b, M)] = D[PR_b, E(PU_b, M)]$$

# Asymmetric Encryption Algorithms



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