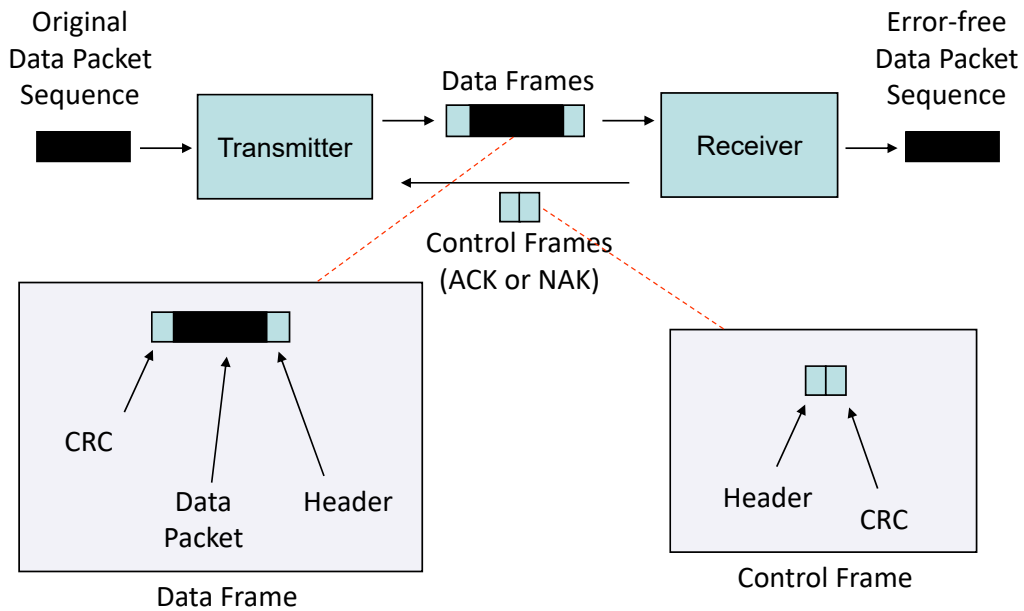


Automatic Repeat reQuest (ARQ)

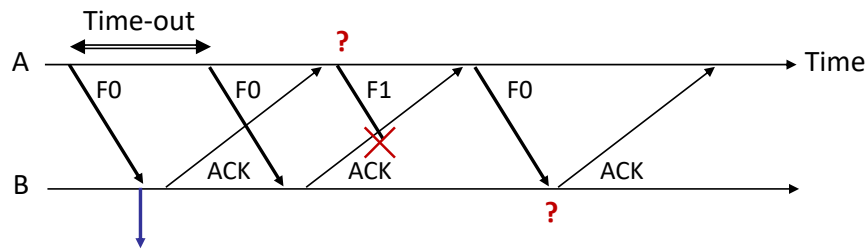


Design Goals of ARQ Protocols

✚ Three Design Goals of ARQ Protocols

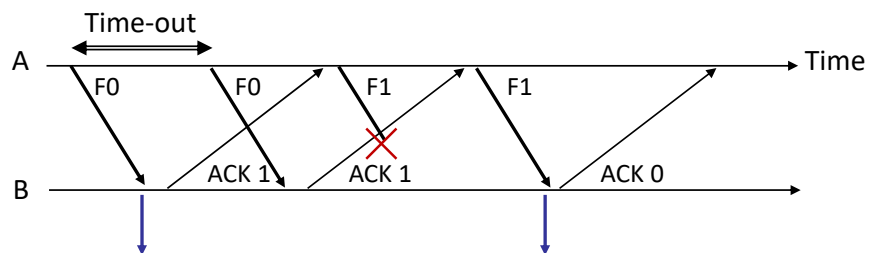
- ➡ **Goal #1:** to ensure that each packet is delivered **error-free**
- ➡ **Goal #2:** to ensure that each packet is delivered **exactly once without duplication**
- ➡ **Goal #3:** to ensure that packets are delivered **in order**

Problematic Scenario without Sequence Numbering



- ✦ **Problem:** A assumes that F1 was delivered error-free
B discards the new F0 by mistake
- ✦ **Solution:** Add frame sequence number (R_{next}) in the ACK header
 - R_{next} is the sequence number of **next frame** expected by the receiver
 - It acknowledges reception of all prior frames implicitly

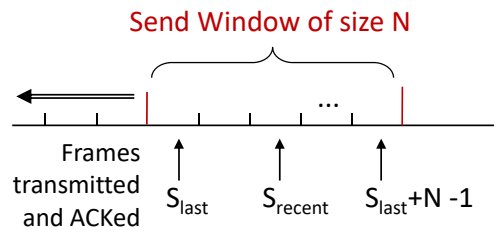
Correct Operation with Sequence Numbering



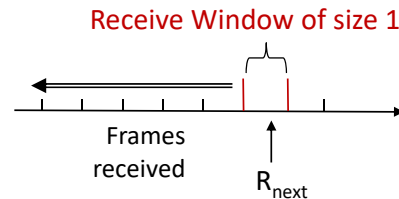
2. Go-Back-N ARQ (GBN)

- ⊕ GBN improves S&W by keeping the channel busy when the transmitter waits for acknowledgment from the receiver

- At the transmitter side:
Allow up to N
outstanding frames



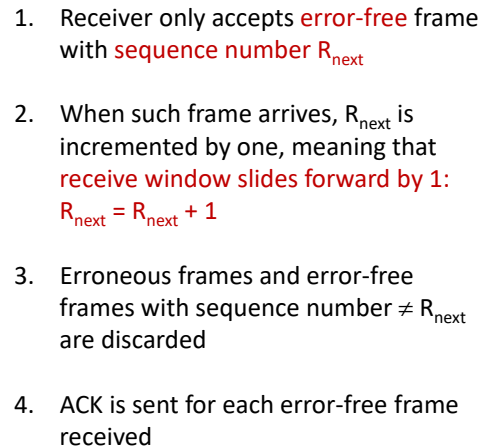
- At the receiver side:
Receive window of size 1



GBN Protocol

- ⊕ *Essential Components*: ACK, timeout, sequence numbering
 - ACK acknowledges reception of all prior frames implicitly
- ⊕ Upon timeout:
 - Frame in error and all subsequent frames are retransmitted
- ⊕ Needs m-bit sequence numbering to remove ambiguities
 - What is the minimum value for m?

Diagram illustrating a sliding window protocol. The timeline shows frames transmitted and acknowledged, with the current window spanning from S_{last} to $S_{last}+N-1$. The window size is N .



- Cpr E 489 -- D.Q.

3.22

- Example: 2-bit ($m = 2$) sequence numbering suffices for Go-Back-3 ($N = 3$) ARQ

[illegible]

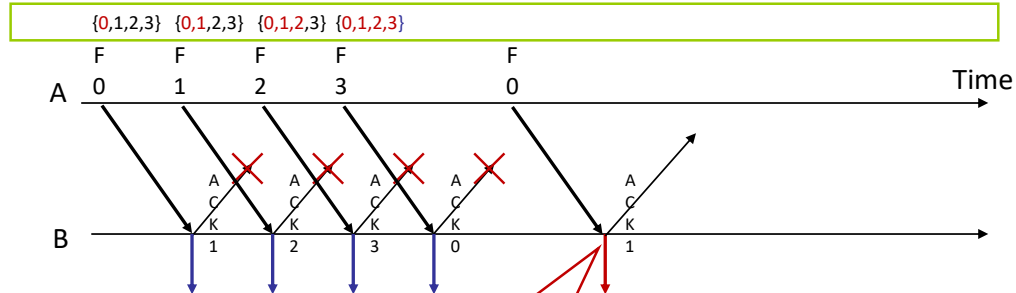
Cpr E 489 -- D.Q.

3.23

$$N + 1 \leq 2^m$$

- ✦ Example: 2-bit ($m = 2$) sequence numbering is **inadequate** for Go-Back-4 ($N = 4$) ARQ

send window



receive window

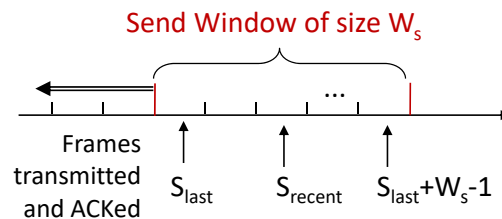


Receiver has $R_{next} = 0$, so it accepts the retransmitted old F0 as a new F0
 → **Wrong! Design goal #2 is violated!**

3. Selective Repeat ARQ (SR)

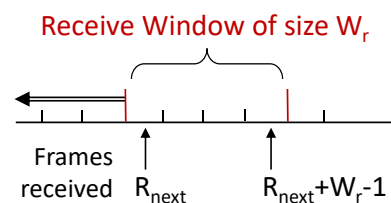
- ✦ SR improves upon GBN by buffering at the receiver side

- Allow up to W_s outstanding frames



- Allow a **receive window of size W_r** (>1)

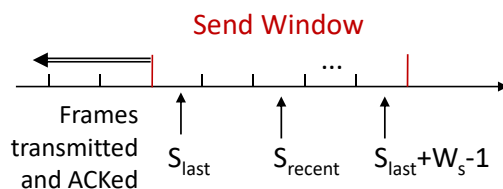
- Receiver buffers the error-free frames with sequence number $\in [R_{next} + 1, R_{next} + W_r - 1]$



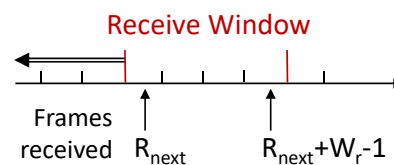
SR Protocol

- ✦ **Essential Components:** ACK, NAK, timeout, sequence numbering
 - NAK is sent when an error-free out-of-sequence frame is received
 - ACK is sent for all other error-free frames
 - Both ACK and NAK acknowledge reception of all prior frames
- ✦ Frame in error is retransmitted upon
 - Timeout or reception of NAK
 - **Only the frame in error is retransmitted**
- ✦ Needs m-bit sequence numbering to remove ambiguities
 - What is the minimum value for m?

SR Transmitter & Receiver



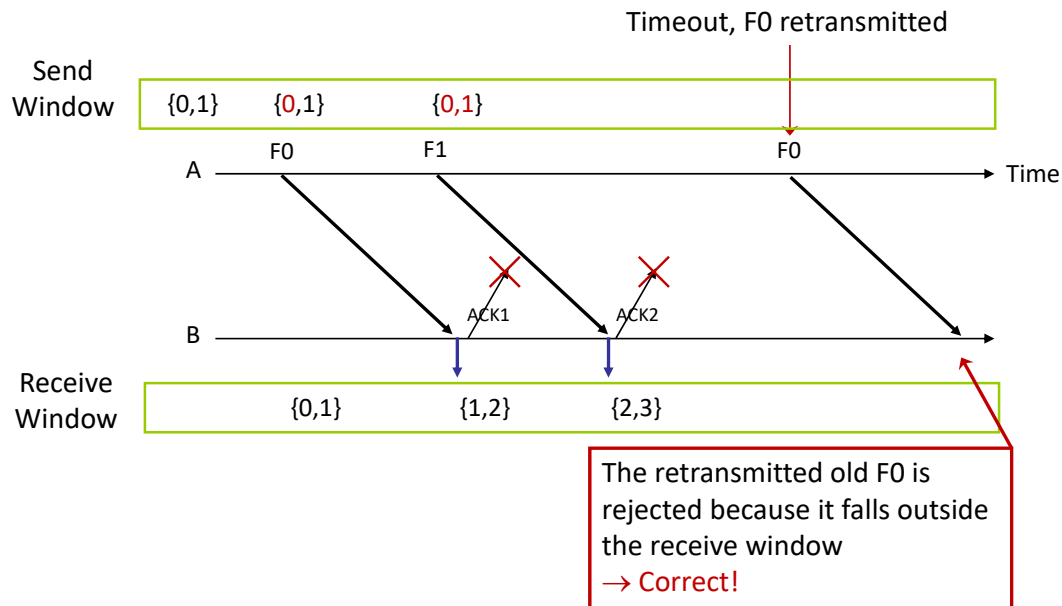
1. If an **error-free** ACK or NAK with $R_{\text{next}} \in [S_{\text{last}}, S_{\text{recent}}+1]$ arrives, **send window slides forward: $S_{\text{last}} = R_{\text{next}}$**
2. When timer for a frame expires or when a NAK arrives, transmitter **retransmits the corresponding frame only**



1. Receiver only accepts **error-free** frames with **sequence number $\in [R_{\text{next}}, R_{\text{next}}+W_r-1]$**
2. When frame with sequence number R_{next} arrives, R_{next} is incremented to a proper value \Rightarrow **receive window may slide forward by more than one**
3. Erroneous frames and error-free frames with sequence number $\notin [R_{\text{next}}, R_{\text{next}}+W_r-1]$ are discarded
4. NAK is sent when an error-free out-of-sequence frame is received
5. ACK is sent for all other error-free frames received

$$W_s + W_r \leq 2^m$$

- Example: 2-bit ($m = 2$) sequence numbering, $W_s = W_r = 2$



$$W_s + W_r \leq 2^m$$

- Example: Problem when $W_s + W_r > 2^m$
2-bit ($m = 2$) sequence numbering, $W_s = W_r = 3$

