

show that $(P \rightarrow (Q \rightarrow P)) \Rightarrow (\neg P \rightarrow (P \rightarrow Q))$

we have to prove that $(P \rightarrow (Q \rightarrow P)) \rightarrow (\neg P \rightarrow (P \rightarrow Q)) \equiv T$

$$\neg(\neg P \vee (\neg Q \vee P)) \vee (P \vee (\neg P \vee Q)) \text{ [conditional equivalence]}$$

$$(P \wedge (Q \wedge \neg P)) \vee (P \vee (\neg P \vee Q)) \text{ [De Morgan's law]}$$

$$[(P \wedge Q) \wedge (P \wedge \neg P)] \vee [(P \vee \neg P) \vee (P \vee Q)] \text{ [distributive law]}$$

$$[(P \wedge Q) \wedge F] \vee [T \vee (P \vee Q)] \text{ [complement laws]}$$

$$F \vee T \text{ [identity laws]}$$

$$\equiv T \text{ [identity laws]}$$

Establish the validity of following arguments using the rule of contradictions.

$$[(P \rightarrow Q) \wedge (\neg R \vee S) \wedge (P \vee R)] \rightarrow [Q \rightarrow S]$$

Given premises : $P \rightarrow Q$, $\neg R \rightarrow S$, $\neg P \rightarrow R$

Step no	Statement	Reason
1.	$\neg(Q \rightarrow S)$	Assume premises
2.	$P \rightarrow Q$	Rule P
3.	$\neg Q \rightarrow \neg P$	Rule T (2) $P \rightarrow Q \Leftrightarrow \neg Q \rightarrow \neg P$
4.	$\neg P \rightarrow R$	Rule T (2, 3) $P \rightarrow Q \wedge \neg Q \rightarrow \neg P \Rightarrow P \rightarrow R$ Rule P
5.	$\neg Q \rightarrow R$	Rule T (3, 4) chain rule ($P \rightarrow Q, Q \rightarrow R \Rightarrow P \rightarrow R$)
6.	$\neg R \rightarrow S$	Rule P
7.	$\neg R \rightarrow S$ $\neg R \rightarrow Q$	Rule T (5, 6) ($P \rightarrow Q \Leftrightarrow \neg Q \rightarrow \neg P$)
8.	$Q \rightarrow S$	Rule T (6, 7)
9.	$(Q \rightarrow S) \wedge \neg(Q \rightarrow S)$	contradiction Rule T (8, 1)
10.	$\equiv F$	

3) write each of the following in symbolic form by assuming the universal consist of literally everything

i) All men are giants

$M(x)$: x is a men

$G(x)$: x is giant

$$\therefore [M(x) \rightarrow G(x)]$$

ii) no men are giants.

if x is a men, then x is not giant

$$M(x) \rightarrow \neg G(x)$$

iii) some men are giants.

if x is men, $G(x)$ is giants.

$$\exists x (M(x) \wedge G(x))$$

iv) some men are not giants.

x is men, $G(x)$ is not giants

$$\exists x (M(x) \wedge \neg G(x))$$

prove by mathematical induction, that $2n+1 \leq 2^n$ for $n \geq 3$.

Given: $2n+1 \leq 2^n$ for $n \geq 3$.

1) let check $P(3)$ is true or not

$$P(n) = 2n+1 \leq 2^n$$

$$P(3) = 2(3) + 1 \leq 2^3$$

$$= 6 + 1 \leq 2^3$$

$$\therefore 7 \leq 8$$

$P(3)$ is true

2) $P(k)$ let assume the statement is true

for $n=k$

$$P(k) = 2k+1 \leq 2^k \quad [2k \leq 2^k - 1] \quad \text{--- (1)}$$

3) let us now try to establish that $P(k+1)$ is true

$$P(k+1) = 2(k+1) + 1 \leq 2^{k+1}$$

$$= 2k+2+1 \leq 2^k \cdot 2$$

from Equation (1)

$$2^k - 1 + 2 + 1 \leq 2^{k+1}$$

$$2^k + 2 \leq 2^{k+1}$$

or equal

2^{k+1} will be always greater than the $2^k + 2$

(4)

It is proved that $P(k+1)$ holds true, whenever the statement $P(k)$ is true.

thus, $2^{(2n+1)}$ is less or equal than 2^n is proved using the principles of mathematical induction
