

Pushing the Limit: Verified Performance-Optimal Causally-Consistent Database Transactions



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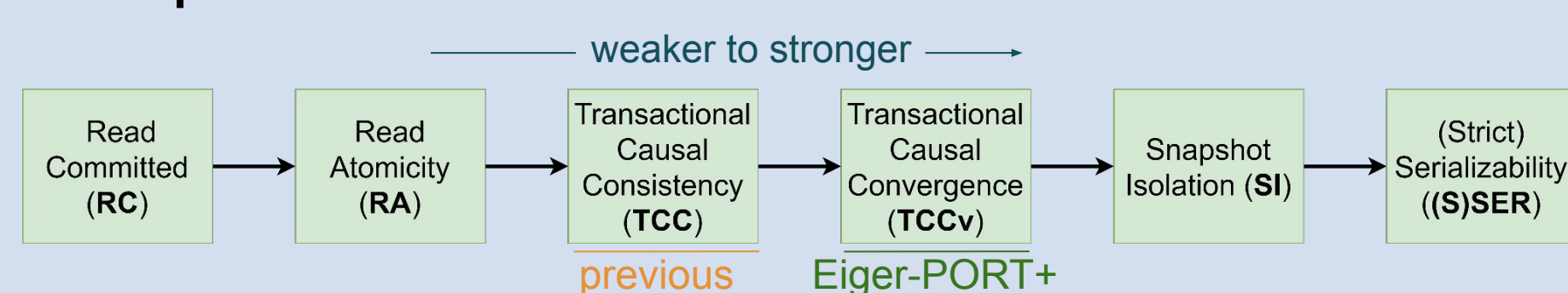


Project repository:



1 Motivation

- Distributed databases (key-value stores)
- Isolation** of concurrent transactions, realized by concurrency control protocols.
- Spectrum of isolation levels:



- Trade-off: isolation vs performance
- TCC previously conjectured to be the strongest achievable isolation level for performance-optimal reads in the presence of transactional writes.
- We **refute** the **conjecture** and push the limit to **TCCv** with our novel protocol **Eiger-PORT+**.
- Concurrency control protocols are highly complex and prone to design errors and isolation bugs. → Deductive **verification**

Contributions

❖ Eiger-PORT+

- Stronger isolation guarantee
- Superior performance

❖ Protocol verification in Isabelle/HOL

- Refinement
- Reduction

2 Abstract Model: Isolation Level

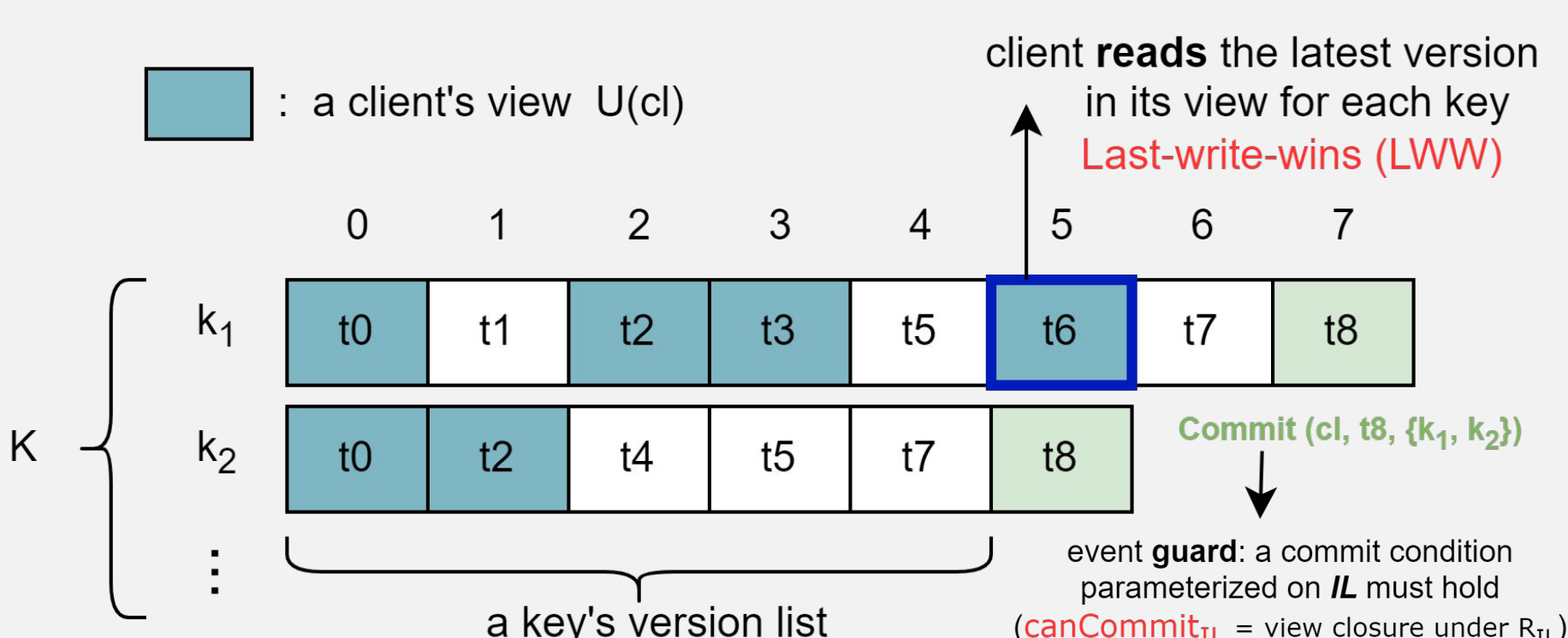
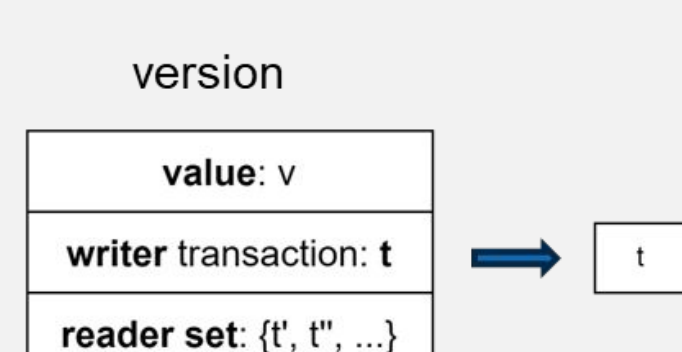
We specify isolation as an abstract model parameterized by an Isolation Level (IL).

The abstract model's event system:

- States: (K, U) $\left\{ \begin{array}{l} K: \text{centralized multi-versioned key-value store} \\ U: \text{client views, capturing the distributed aspect} \end{array} \right.$

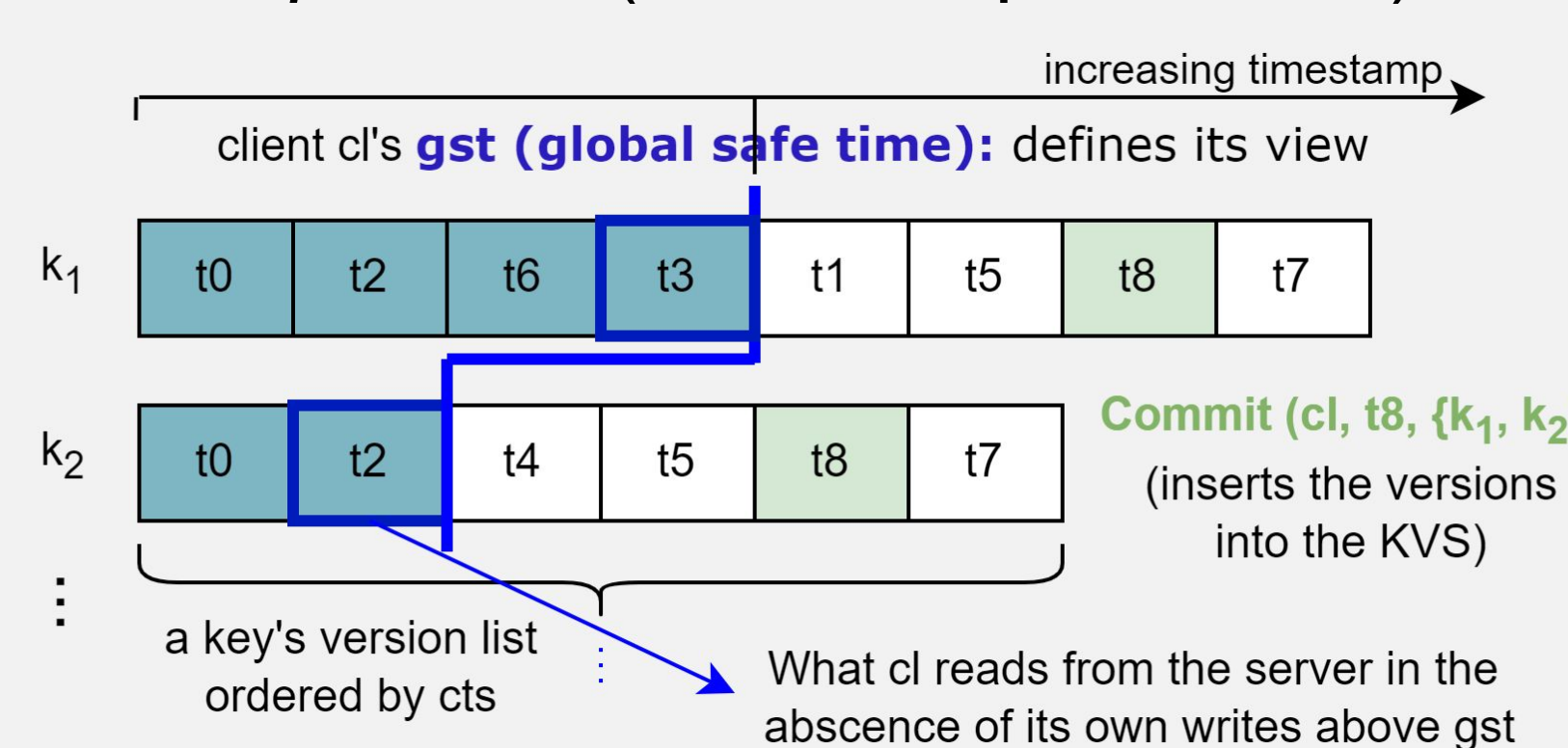
→ Events:

- Atomic Commit $(K, U) \xrightarrow{\text{commit}}_{IL} (K', U')$
- View Extension $(K, U) \xrightarrow{\text{view ext}}_{IL} (K, U')$

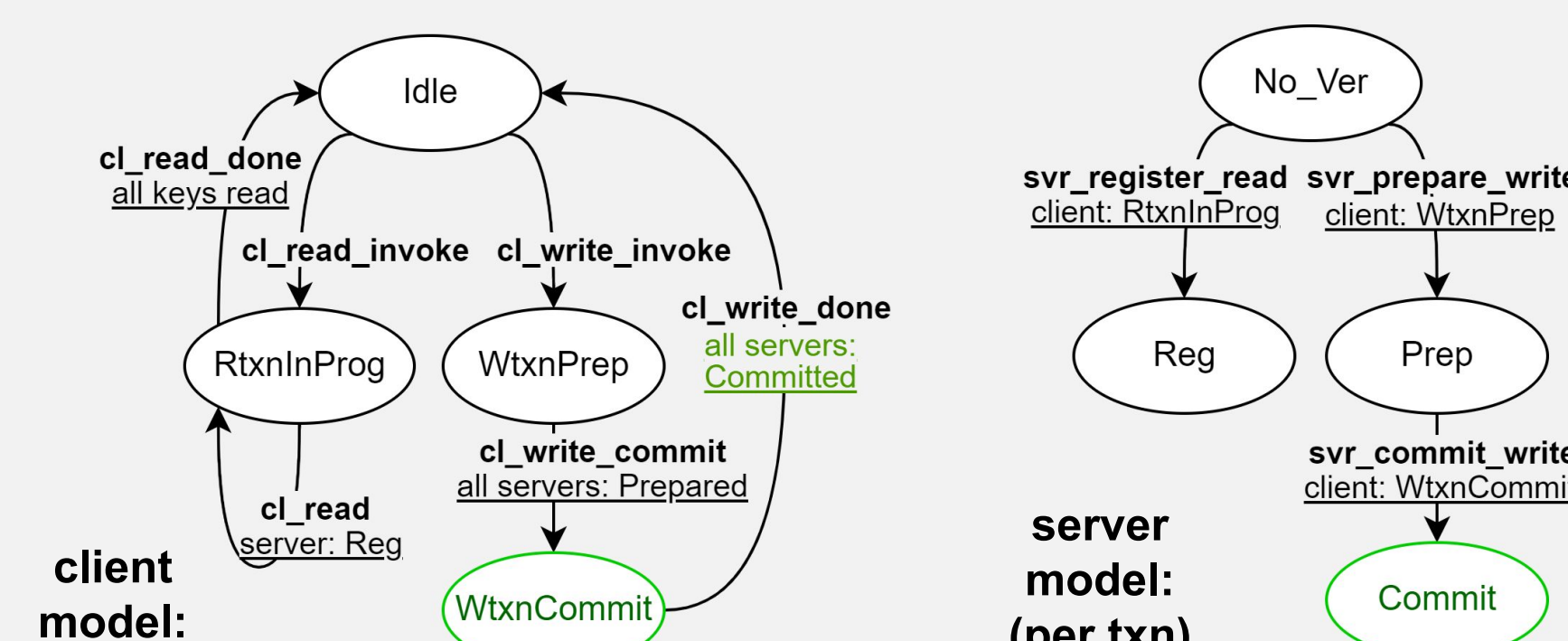


3 Concrete Model: Eiger-PORT+

Eiger-PORT+ protocol:
Timestamp-based (uses Lamport clocks)



- read-only transactions
→ own write above gst
→ or last write below gst
- write-only transactions: two-phase commit

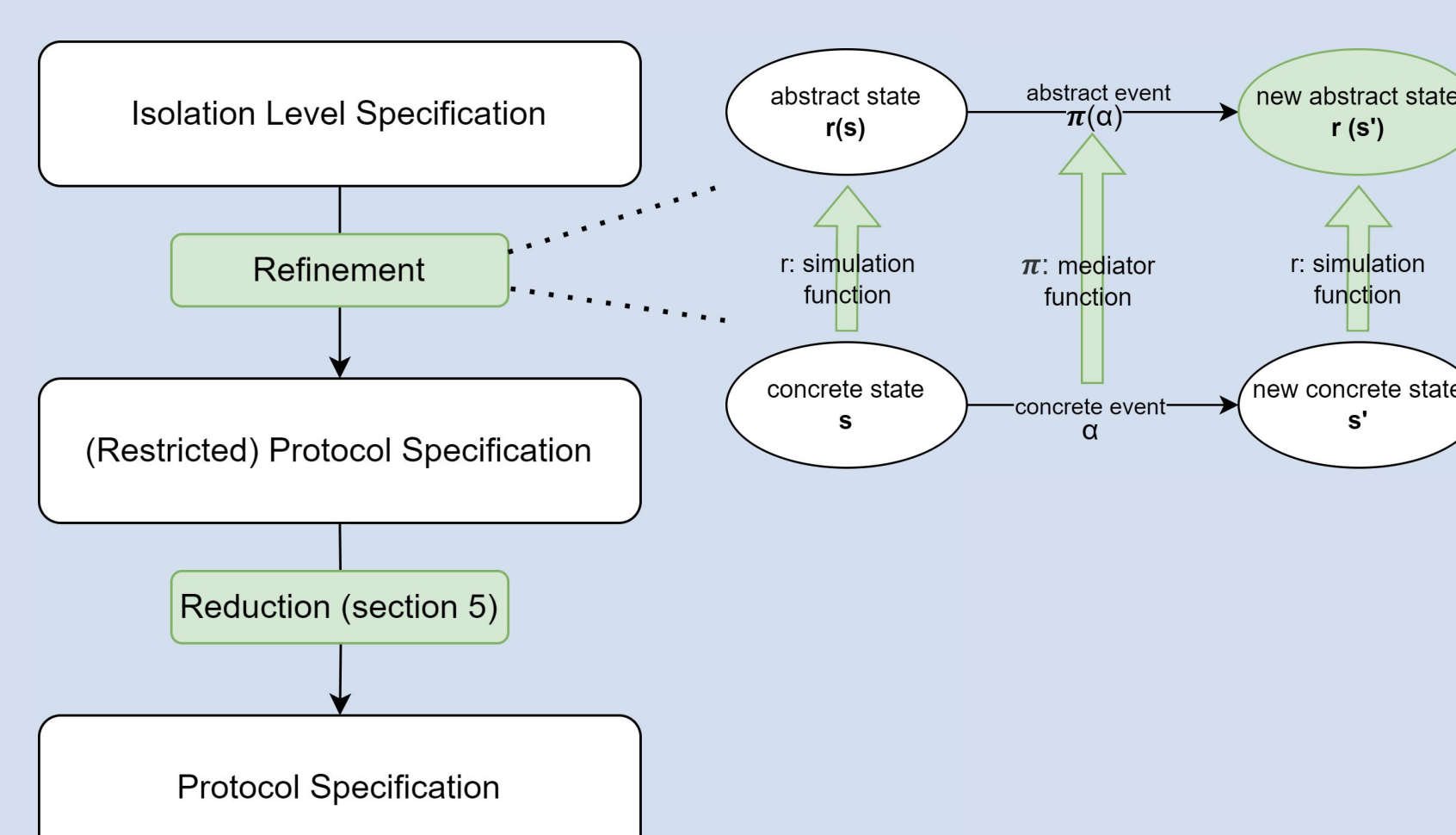


4 Correctness Proof & Invariants

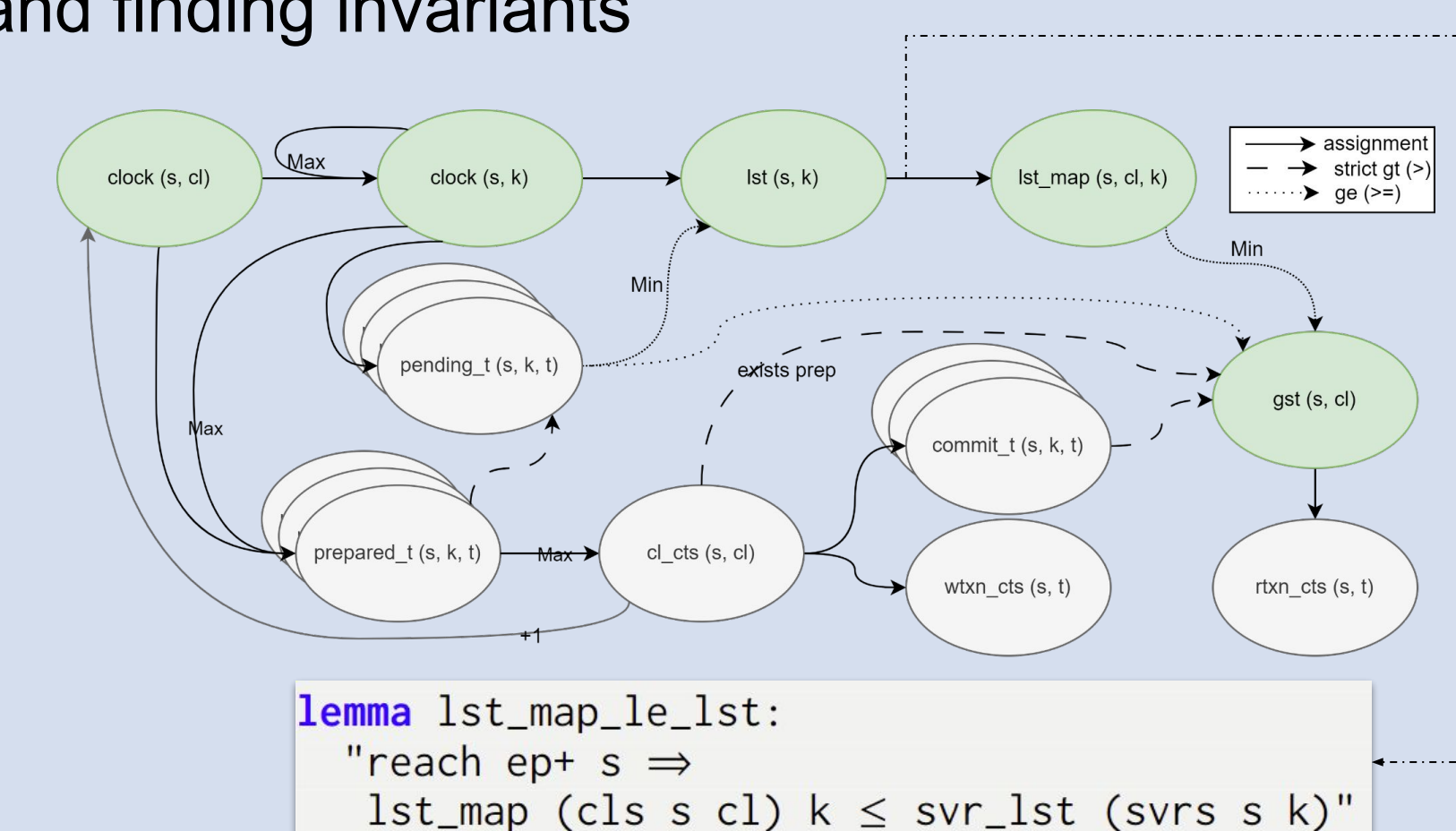
Proof guarantee:

refinement (reach (protocol)) \subseteq reach (IL)

- refinement mapping:
 $r : K$ and U reconstructed as shown above
 π : client_write_commit and client_read_done mapped to Atomic commit
- proof obligations:
 canCommit_{IL} : needs invariants (below)
 LWW : needs reduction

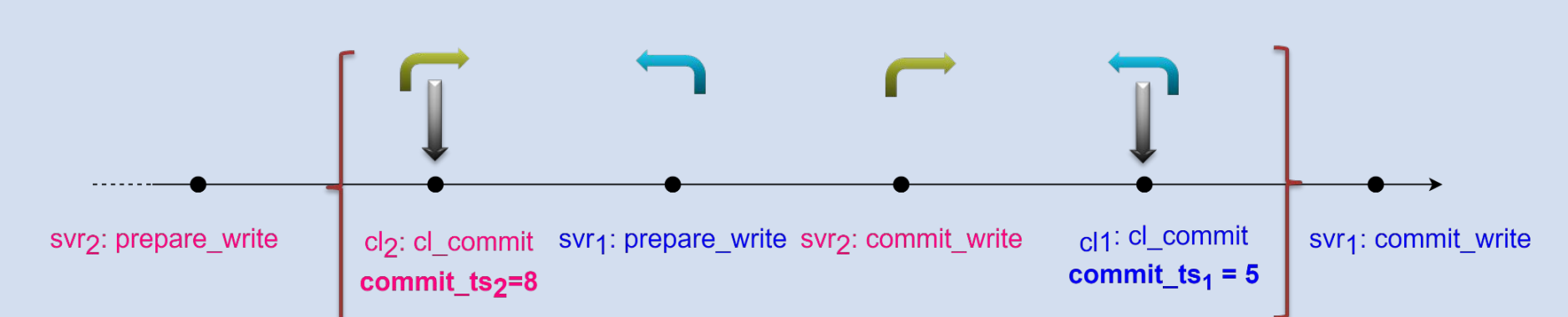


- The relation of different timestamps in the model and finding invariants



5 Inverted Commits & Reduction

Inverted commits: pairs of client commits in protocol executions not ordered by **commit timestamps**.



Can occur for causally independent concurrent transactions.

Problem: Inverted commits would require **inserting** rather than **appending** a transaction's version to the version list.

Can not be simulated by the abstract model.

Hence, refinement alone is not enough for verifying the protocol.

Solution:

- We introduce a restricted protocol model that doesn't produce inverted commits.
- We use **reduction** to transform any protocol execution into one of the restricted model such that:
reach (protocol) = reach (restricted protocol).
- This is achieved by commuting independent concurrent events to eliminate inverted commits. (see arrows on the execution above)

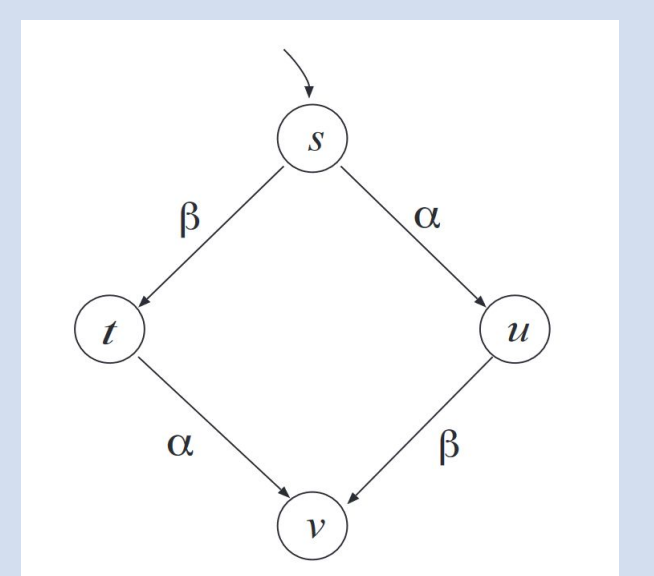
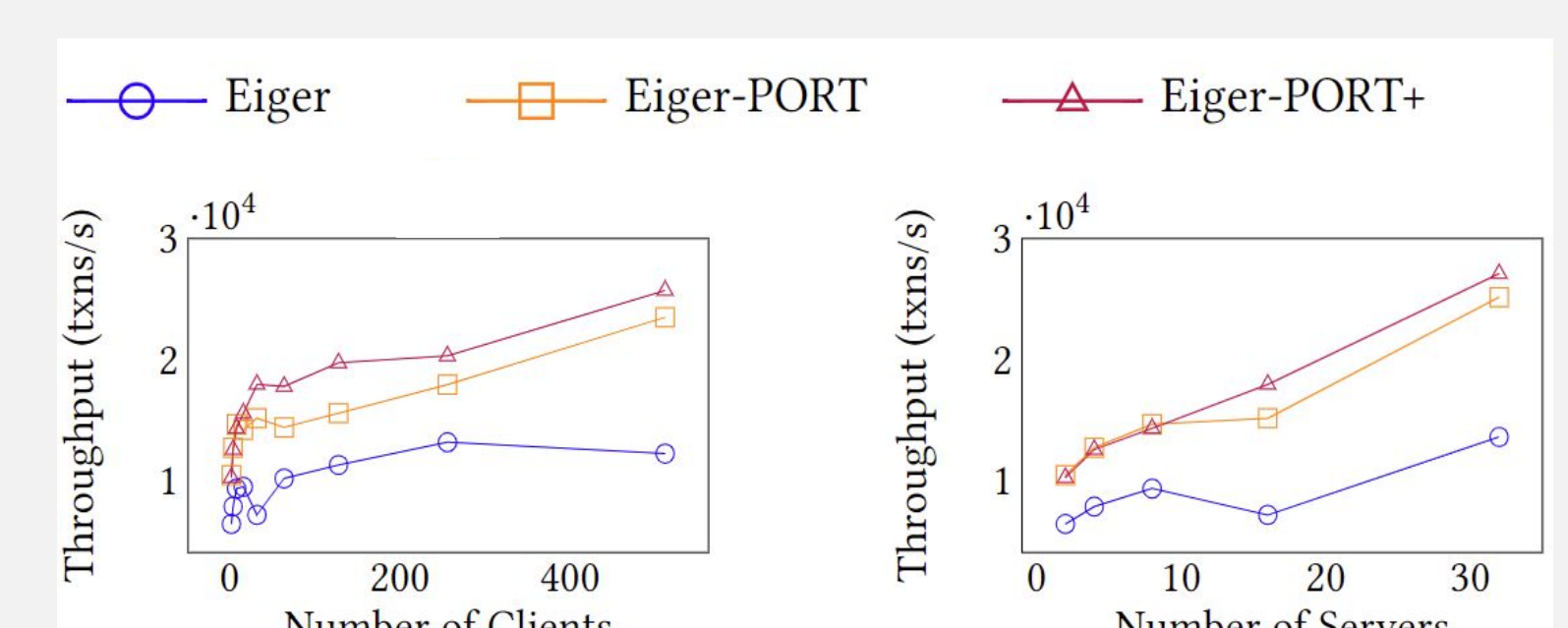


Image source: C. Baier, J. Katoen, 2008, Principles of model checking, p. 595, Figure 8.1.

6 Conclusions and Discussion

- Our Eiger-PORT+ protocol provides TCCv, thus refuting an open conjecture.
- Eiger-PORT+ outperforms state-of-the-art



- Refinement is not always enough
- We deductively verify that Eiger-PORT+ satisfies TCCv, using a combination of refinement and reduction.