

## Data structure

Data structure are a specific way of organizing data in a specialized format of a computer so that information can be organized processed, retrieve quickly and effectively.

### • Primitive data type

Primitive data type are the most basic data type that are used for representing simple values such as integer, float, character etc.

### • Non-primitive data type

Non-primitive data type is a data type that can store the data of more than one type.  
Ex- stack, queue, array, Linklist.

### (#) Difference Between Primitive and Non-primitive data type

#### → Primitive Data type

- Data types are pre-defined.
- PDT will have certain value.
- Size depends on the type of data structure
- It can start with a lowercase
- Ex- One Numbers and strings.

→ Non-primitive

- It is created by programmers.
- Non-primitive data types can be Null.
- Size is not fixed.
- It can start with uppercase.
- Ex - Array and Linked list.

## ④ Linear and Non-Linear Data Structure

→ Linear Data Structure

- The data elements connect to each other sequentially. A user can traverse each element through single run.
- Linear data structure are easier to implement.
- You can traverse linear data structure in a single run.
- Memory is not utilized in an efficient way.
- In Linear data structure Single level is involved.
- Ex - Array, stack, queue.

→ Non-linear data structure

- Data structure where data elements are not arranged sequentially or linearly are called Non-linear data structure.
- Non-linear data structure are difficult to implement.
- You can traverse non-linear data structure in multiple runs.
- Memory is utilized in an efficient way.
- In Non-linear data structure Multiple level is involved.
- Ex - Tree and graph.

## # Arrays

- An array is a linear data structure that collects elements of the same data type and stores them in contiguous and adjacent memory location.
- Array works on an index system starting from 0 to  $(n-1)$  where n is the size of the array.

Ex :-

Address →	I1	I2	I3	I4
array →	A	R	R	A
Index →	01	02	03	04

↓                      ↓  
lower bound          upper bound

## Types of Array

### # One-dimensional array

One dimensional array contain a single row of element. These arrays are usually indexed from 0 to  $(n-1)$  where n is the size of array.

Element →					
Index →	0	1	2	3	4

### \* Multi-dimensional array

Multi-dimensional array are a powerful data structure used to store and manage data organizationally. This type of array consist of multiple arrays that are arranged hierarchically. They can have any number of dimension.

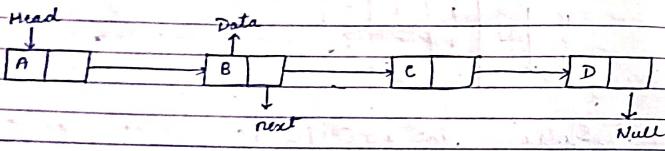
Column →	0	1	2
R	0	1	2
o	1	4	5
w	2	7	8

Syntax :- int arr[m][n]; , int arr[m][n][o];

## ④ Linked List

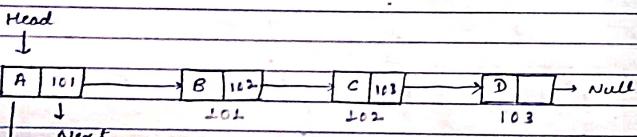
A linked list is a fundamental data structure in computer science. It consists of nodes where each node contains some data and link to next node in the sequence.

- They allow for dynamic memory allocation and efficient insertion and deletion operation.



## \* Singly - Linked List

A singly-linked list is a linear data structure in which the elements are not stored in contiguous memory location and each node is connected to its next element using a pointer.

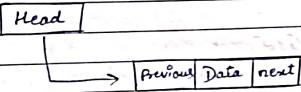


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## \* Doubly - Linked List

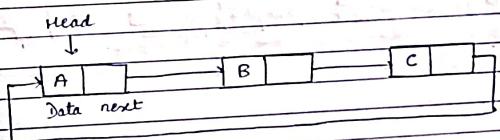
- Doubly-linked list is a complex type of linked list in which a node contains a pointer to the previous as well as the next node in the sequence.
- Therefore a doubly-linked list a node consist of three parts:
  - ↳ node data.
  - ↳ Pointer to the next node in sequence.
  - ↳ Pointer to the previous node.



## \* Circular - Linked List

The Circular - linked list is a linked list where all nodes are connected to form a circle. In a circular linked list the first node and the last node are connected to each other which form a circle.

There is no null at the End.

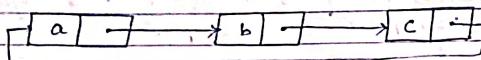


### → Singly Circular linked list

In a singly circular linked list the last node of the list contains a pointer to the first node of the list.

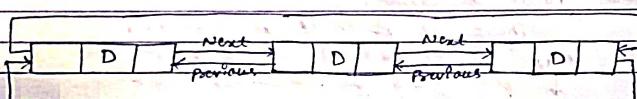
We traverse the singly circular linked list until we reach the same node where we started.

The singly circular linked list has no beginning or end.



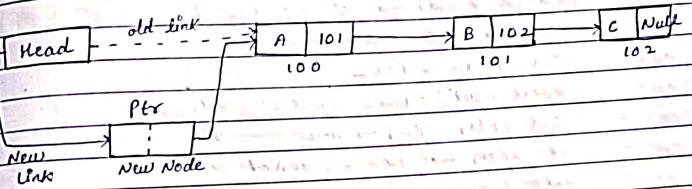
### → Doubly Circular linked list

Doubly circular linked list has all the properties of both circular linked list and doubly linked list in which two consecutive elements are linked or connected by the previous or next pointers and last node points to the first node by pointer and also first node points to the last node by previous pointer.



### \* Insertion (Single Linked List)

#### → At beginning



Step 1 :- If  $\text{ptr} = \text{Null}$  (write overflow)

Step 2 :- set  $\text{New Node} = \text{ptr}$

Step 3 :- Set  $\text{ptr} = \text{ptr} \rightarrow \text{Next}$

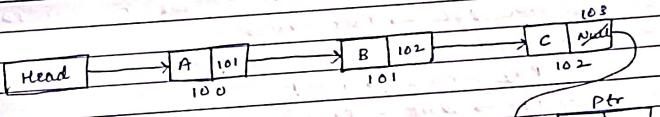
Step 4 :- set  $\text{New Node} \rightarrow \text{Data} = \text{value}$

Step 5 :- set  $\text{New Node} \rightarrow \text{Next} = \text{Head}$

Step 6 :- set  $\text{New Node} = \text{Head}$

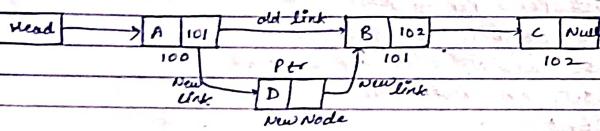
Step 7 :- Exit

#### → At End



Step 1 - If  $\text{Ptr} = \text{Null}$  (write overflow)  
Step 2 - Set  $\text{NewNode} = \text{Ptr}$   
Step 3 - Set  $\text{Ptr} = \text{Ptr} \rightarrow \text{next}$   
Step 4 - Set  $\text{NewNode} \rightarrow \text{Data} = \text{value}$   
Step 5 - Set  $\text{NewNode} \rightarrow \text{Next}! = \text{NULL}$   
Step 6 - Set  $\text{Head} = \text{Ptr} \perp$   
Step 7 - Repeat until  $\text{Ptr} \perp \rightarrow \text{next}! = \text{NULL}$   
Step 8 - Set  $\text{Ptr} \perp = \text{Ptr} \rightarrow \text{next}$   
Step 9 - Set  $\text{Ptr} \rightarrow \text{next} = \text{NewNode}$   
Step 10 - Exit

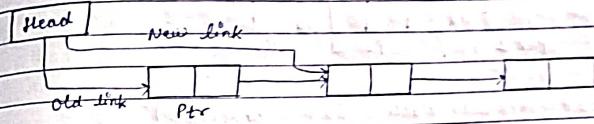
→ At specific place



Step 1 - If  $\text{Ptr} = \text{NULL}$  (write overflow)  
Step 2 - Set  $\text{NewNode} = \text{Ptr}$   
Step 3 - Set  $\text{NewNode} \rightarrow \text{Data} = \text{value}$   
Step 4 - Set  $\text{Head} = \text{Ptr} \perp$   
Step 5 - Set  $\text{Ptr} \perp = \text{Ptr} \perp \rightarrow \text{next}!$   
Step 6 - Set  $\text{Ptr} \rightarrow \text{Next} = \text{Ptr} \perp \rightarrow \text{next}$   
Step 7 - Set  $\text{Ptr} \perp = \text{NewNode} \rightarrow \text{Next} = \text{Ptr}$   
Step 8 - Set  $\text{Ptr} = \text{NewNode}$   
Step 9 - Exit

### \* Deletion

→ At beginning

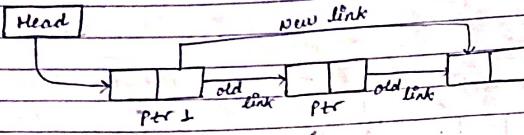


Step 1 - If  $\text{Head} = \text{NULL}$  (write underflow)  
Step 2 - Set  $\text{Ptr} = \text{Head}$   
Step 3 - Set  $\text{Head} = \text{Head} \rightarrow \text{Next}$   
Step 4 - free  $\text{Ptr}$   
Step 5 - Exit

→ At End

Step 1 - If  $\text{Head} = \text{NULL}$  (write underflow)  
Step 2 - Set  $\text{Ptr} = \text{Head}$   
Step 3 - Repeat until  $\text{Ptr} \rightarrow \text{Next}! = \text{NULL}$   
Step 4 - Set  $\text{Ptr} \perp = \text{Ptr} \perp$   
Step 5 - Set  $\text{Ptr} \perp \rightarrow \text{Next} = \text{NULL}$   
Step 6 - free  $\text{Ptr}$   
Step 7 - Exit

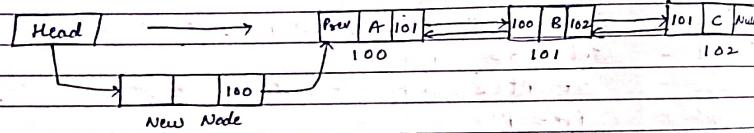
→ At Specific place



- Step1 - If Head = Null (write underflow)
- Step2 - Set Ptr = Head
- Step3 - Set Ptr = Ptr → Next
- Step4 - Set Ptr<sub>l</sub> = Ptr
- Step5 - Set Ptr → next = Ptr<sub>l</sub> → Next
- Step6 - free Ptr
- Step7 - Exit

#### \* Insertion (Doubly linked list)

→ At beginning



- Step1 - If Ptr = Null (write overflow)
- Step2 - Set NewNode = Ptr

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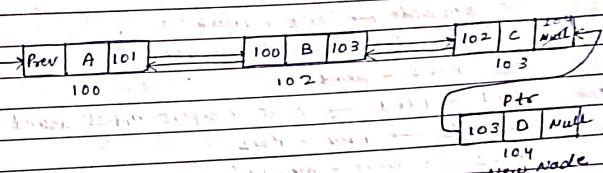
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- Step4 - set NewNode → prev = Null
- Step5 - set NewNode → next = Head
- Step6 - Set Head → prev = NewNode
- Step7 - Set Head = NewNode
- Step8 - Exit

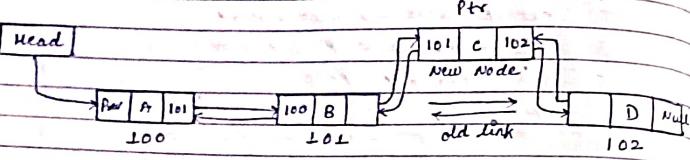
→ At End

Head



- Step1 - If Ptr = Null - (write overflow)
- Step2 - set NewNode = Ptr
- Step3 - set NewNode → Data = value
- Step4 - set NewNode → Next = Null
- Step5 - set Ptr<sub>l</sub> = Head
- Step6 - set Ptr = Ptr<sub>l</sub> → Next != Null
- Step7 - set NewNode → Prev = Ptr<sub>l</sub>
- Step8 - set Ptr<sub>l</sub> → next = NewNode
- Step9 - Exit
- Step10 -

→ At Specific Place



Step 1 If  $\text{Ptr} = \text{Null}$  (write overflow)

Step 2 Set  $\text{NewNode} = \text{Ptr}$

Step 3 Set  $\text{NewNode} \rightarrow \text{Data} = \text{value}$

Step 4 Set  $\text{NewNode} \rightarrow \text{New Node} \rightarrow \text{Next}$

Step 5 Set  $\text{Ptr} \leftarrow = \text{Head}$

Step 6  $\text{ptr} \leftarrow = \text{ptr} \rightarrow \text{Next}$  (Repeat until desired track is found)

Step 7  $\text{NewNode} \rightarrow \text{Prev} = \text{ptr} \leftarrow$

Step 8  $\text{ptr} \leftarrow \rightarrow \text{Next} = \text{Newnode}$

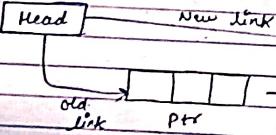
Step 9  $\text{ptr} \leftarrow \rightarrow \text{next} \rightarrow \text{Prev} = \text{New Node}$

Step 10  $\text{New Node} \rightarrow \text{Next} = \text{ptr} \leftarrow \rightarrow \text{next}$

Step 11 End

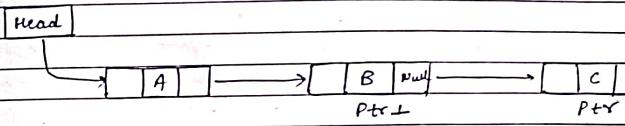
\* Deletion

→ At beginning



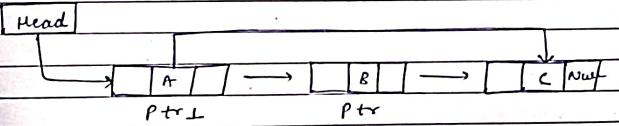
- Page No. \_\_\_\_\_  
Date \_\_\_\_\_
- Step 1 - If  $\text{Head} = \text{Null}$  (write underflow)
  - Step 2 - set  $\text{Ptr} = \text{Head}$
  - Step 3 - set  $\text{Head} = \text{Head} \rightarrow \text{Next}$
  - Step 4 - set  $\text{Head} \rightarrow \text{prev} = \text{Null}$
  - Step 5 - free  $\text{ptr}$
  - Step 6 - Exit

→ At End



- Page No. \_\_\_\_\_  
Date \_\_\_\_\_
- Step 1 - If  $\text{Head} = \text{Null}$
  - Step 2 - set  $\text{Ptr} = \text{Head}$
  - Step 3 -  $\text{ptr} \rightarrow \text{Next} = \text{Null}$  (Repeat Until)
  - Step 4 - set  $\text{ptr} \leftarrow = \text{ptr} \rightarrow \text{Null}$
  - Step 5 - set  $\text{ptr} \leftarrow \rightarrow \text{ptr} \rightarrow \text{prev} = \text{Null}$
  - Step 6 - free  $\text{ptr}$
  - Step 7 - Exit

→ At Specific Point



- Step 1 - If  $\text{head} = \text{null}$
- Step 2 - Set  $\text{head} = \text{ptr}$
- Step 3 - Set  $\text{ptr} = \text{ptr} \rightarrow \text{next}$
- Step 4 - Set  $\text{ptr\_l} = \text{ptr} \rightarrow \text{next}$
- Step 5 - Set  $\text{ptr} \rightarrow \text{next} \rightarrow \text{prev} = \text{ptr\_l}$
- Step 6 - Set  $\text{ptr\_l} \rightarrow \text{next} = \text{ptr} \rightarrow \text{next}$
- Step 7 - free  $\text{ptr}$
- Step 8 - Exit