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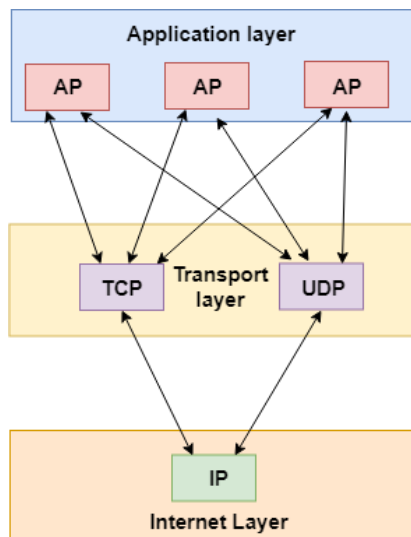
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# Transport Layer

## Functions and its working

### Description

In the open system interconnection (OSI) model, the transport layer is the fourth layer, and is responsible for end-to-end communication over a network. It offers logical communication within a layered architecture of protocols and other network components between application processes running on different hosts. In the open system interconnection (OSI) model, the transport layer is the fourth layer, and is responsible for end-to-end communication over a network.



The transport layer receives message segments from applications in a nutshell and transmits them to the network (Layer 3). This is where the fragments are reassembled and moved to Layer 7 into fully fledged documents. The transport layer is also responsible for the management of error-correction, ensuring continuity and reliability for the end user. This layer allows the host to send and receive data, packets or messages corrected by error over a network and is the part of the network that allows multiplexing.<sup>[1]</sup>

The data is called Segments in the Transport Layer. The transport layer is controlled by the Operating System(OS). It is a part of the OS and by making system calls, communicates with the Application Layer. The Transport Layer is also known as the OSI Heart model.<sup>[6]</sup>

**Keywords :**

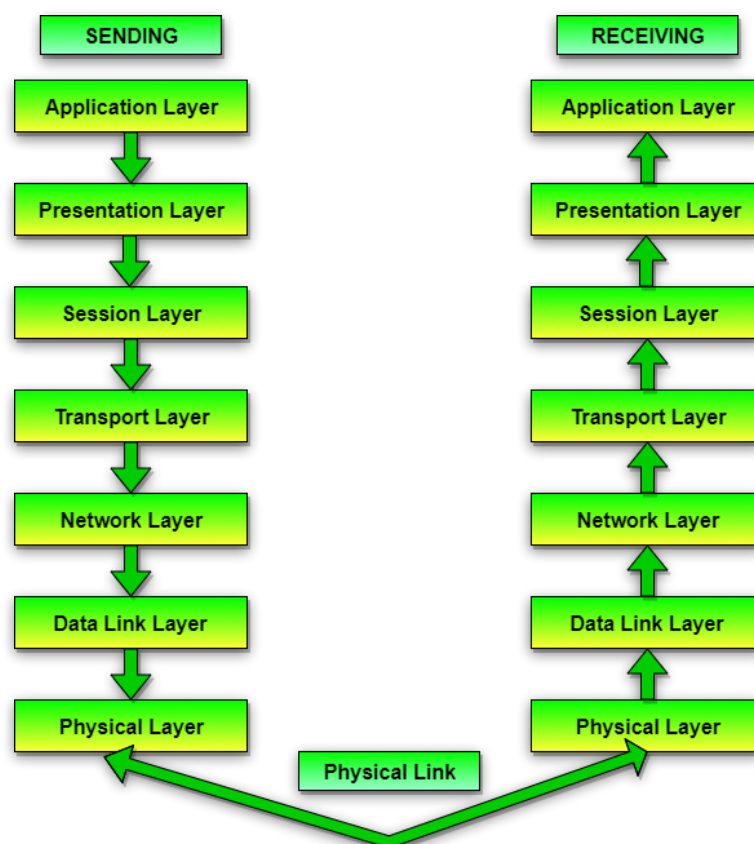
OSI Model, Error correction, Reliability, Connection-Oriented Communication, Same Order Delivery, Data Integrity, Flow Control, Traffic Control, Multiplexing, Byte orientation, end-to-end.

## Introduction

The most important reference modes in communication networks are :

- OSI reference model
- TCP/IP reference model

There are many users that are based all over the world using computer networks. This model has been developed by ISO (ISO stands for International Organization of Standardization) to ensure national and global data communication. This is called an open system interconnection (OSI) model and is generally referred to as an OSI model. Seven layers consist of the OSI model architecture. In a complete communication structure, it recognises seven layers or levels.<sup>[2]</sup>



## Working Principle

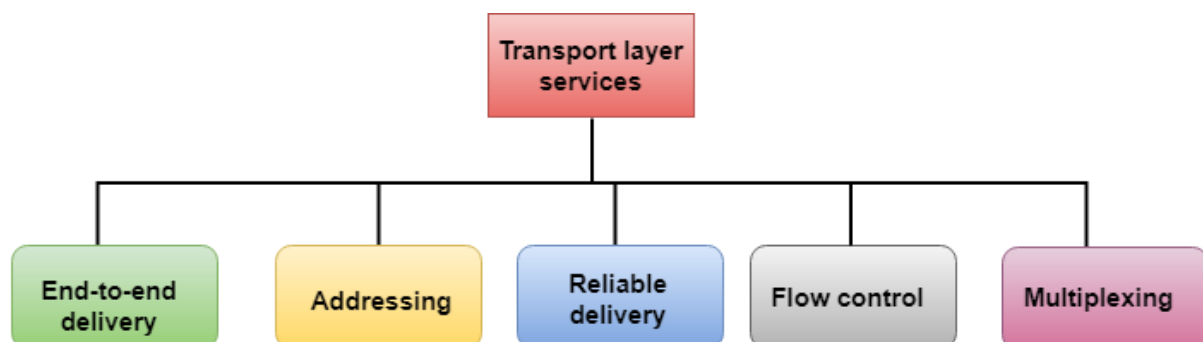
The basic function of the transport layer is to accept the data from the layer above to break it into smaller units, to transfer the data units to the network layer, and to ensure that all the parts arrive at the other end correctly. In addition, all this must be achieved quickly and in such a way that the upper layers are removed from the inevitable changes in hardware technology.

The Transport Layer also defines the type of service to be provided to the network users and, ultimately, to the Session Layer. An error-free point-to-point channel that delivers messages or bytes in the order in which they have been sent is the most common form of transport link. The transport layer, all the way from the source to the destination, is a real end-to-end layer. In other words, a programme on the source machine uses message headers and control messages for a communication with a similar programme on the destination machine.

In short, in order to deliver and receive data without errors, transport layers which is Layer 4, operate transparently within the layers above. Request messages are split into segments which are also called packets on the send side and transferred to the network layer which is Layer 3. Then segments are reassembled into messages by the receiving side and transferred to the application layer which is Layer 7.<sup>[2]</sup>

## Services of Transport Layer

The facilities that the transport layer provides are close to those of the data link layer. Within a single network, the data link layer provides the services, while the transport layer provides the services through a network of multiple networks. The data link layer governs the physical layer, while all the lower layers are managed by the transport layer.<sup>[3]</sup>



- **Connection-Oriented Communication**

Connection-oriented defines a means of transmitting data in telecommunications where end-point devices use a preliminary protocol to create an end-to-end connection before sending any data. Connection-oriented protocol service is often referred to as a "reliable" network service, as it ensures that the correct sequence of data will arrive. A connection-oriented protocol is the Transmission Control Protocol (TCP).<sup>[4]</sup>

A handshake protocol such as TCP is developed by devices at the end points of a network communication to ensure a link is robust before data is exchanged. The downside of this approach is that there is a need for an acknowledgment for each sent message, adding significant network load compared to self-error-correcting packets. When faulty byte streams or datagrams are sent, repetitive requests cause a major slowdown in network speed.<sup>[1]</sup>

- **Data Integrity**

The consistency of data across all distribution layers can be guaranteed using checksums. Those checksums ensure that the transmitted data is the same as the received data and is not corrupt. By demanding retransmission from other layers, lost or corrupted data can be resent.

- **Multiplexing and Demultiplexing**

Transmission through a network with multiple packet streams from unrelated applications or from other sources (multiplexing) involves some very dedicated control mechanisms that are present in the transport layer. This multiplexing enables simultaneous applications to be used over a network, such as when the same device opens different internet browsers. In the OSI model, in the service layer, multiplexing is done.<sup>[1]</sup>

Demultiplexing is needed to obtain the data coming from different processes on the receiver side. Transport gathers the data segments from the network layer and delivers them to the required process running on the computer of the recipient.<sup>[5]</sup>

- **Traffic Control**

Traffic control are also called as Congestion Control. Bandwidth and processing speed limits are placed on digital communications networks, which can mean a large amount of capacity for data congestion on the network. Nearly any aspect of a network can be affected by this network congestion. The transport layer will recognise the signs of overloaded nodes and decreased flow rates and take the necessary measures to correct these problems.

- **Byte orientation**

A communication protocol in which complete bytes are used as control codes is a byte-oriented framing protocol. A character-oriented protocol is also called. UART communication, for instance, is byte-oriented.

Some applications choose to receive byte streams rather than packets; the transport layer enables, if necessary, byte-oriented data streams to be transmitted.

- **Flow Control**

To prevent the sender from overwhelming the recipient, flow control is used. If too much data is overloaded by the receiver, then the receiver discards the packets and demands that the packets be retransmitted. This raises network congestion and thus, reduces the efficiency of the system. For flow regulation, the transport layer is responsible. It uses the sliding window protocol, which allows the transmission of information more effective and manages the data flow so that the receiver is not overloaded. The protocol for sliding windows is byte based rather than frame oriented.

- **Same Order Delivery**

Generally, the network layer does not guarantee that data packets will arrive in the same order that they have been sent, but this is also a beneficial function. Typically, this is achieved by using section numbering, with the receiver passing them in order to the submission. This will provoke obstruction of head-of-line.

Here, Same order delivery Ensure that by assigning them a number, packets are always delivered in strict sequence. Although the network layer is accountable, by reordering them, the transport layer will correct any inconsistencies in sequence caused by packet drops or system disruption.

## Congestion Control

### In computer networks<sup>[7]</sup>

**Congestion** : A condition that arises when the message traffic is so high in the network layer that it slows down the response time of the network.

Congestion's effects :

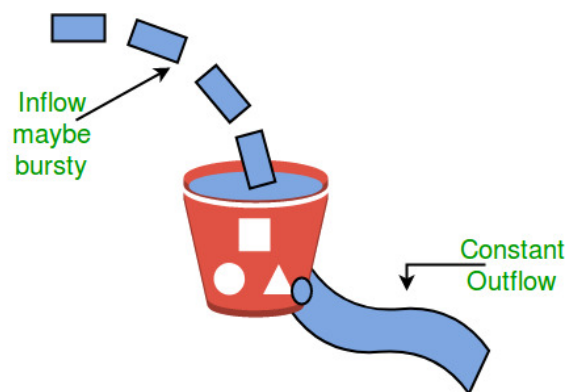
- *Performance decreases as delay increases.*
- *Retransmission happens if the delay increases, making the situation worse.*

### Algorithms for Congestion control :

- Leaky Bucket Algorithm

Let us consider an illustration to understand :

Let us Imagine a bucket which has a small hole in the bottom. The outflow is at a constant rate, no matter what rate of water enters the bucket. When the bucket is full of water, additional water spills over the sides and is lost.



Similarly, a leaky bucket is used in each network interface and the following steps are involved in the leaky bucket algorithm:

1. *The packet will be dropped into the bucket when the host wishes to send the packet.*
2. *At a constant rate, the bucket leaks, meaning packets are transmitted at a constant rate by the network interface.*
3. *The leaky bucket transforms bursty traffic to uniform traffic.*
4. *In reality, the bucket is a finite queue with a finite rate of production.*

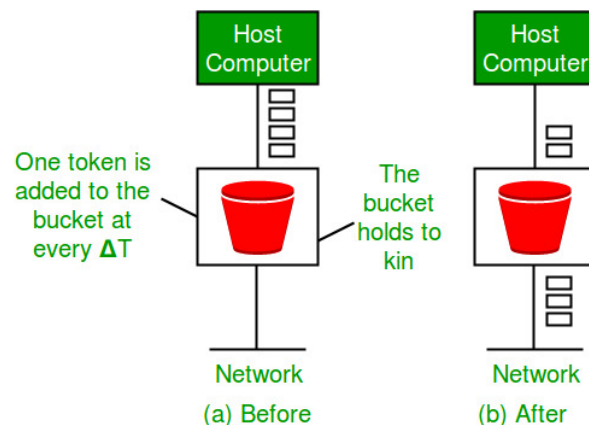
- Token Bucket Algorithm

Need for the Algorithm token bucket:—

The leaky bucket algorithm, no matter how bursty the traffic is, enforces performance patterns at the average rate. But we need a versatile algorithm in order to cope with the bursting traffic so that the data is not lost. One such algorithm is an algorithm for token buckets.

This algorithm's steps can be defined as follows:

1. *Tokens are tossed into the bucket at regular intervals.*
2. *The bucket has the greatest ability.*
3. *A token is extracted from the bucket if there is a ready packet, and the packet is sent.*
4. *If the bucket does not have a token, the packet cannot be sent.*



Let's use an instance to grasp,

We see a bucket in Figure (A) containing three tokens, with five packets waiting for transmission. It must catch and kill one token for a packet to be transmitted. We see that three of the five packets got through in figure (B), but the other two are stuck waiting for more tokens to be produced.

**Ways in which token bucket is superior to leaky bucket<sup>[7]</sup> :**

The algorithm of the leaky bucket governs the rate at which packets are implemented in the network, but it is rather conservative in nature. In the token bucket algorithm, some versatility is added. Algorithm tokens are created in the token bucket at any tick (up to a certain limit). It must catch a token in order for an incoming packet to be transmitted and the transmission takes place at the same rate. Therefore if tokens are available, some of the busy

packets are transmitted at the same rate and thus introduces some amount of device flexibility.

Formula :  $\mathbf{M} * \mathbf{s} = \mathbf{C} + \mathbf{\rho} * \mathbf{s}$

where,

- $\mathbf{M}$  – Maximum output rate
- $\mathbf{s}$  – is time taken
- $\mathbf{C}$  – Capacity of the token bucket in bytes
- $\mathbf{\rho}$  – Token arrival rate

Let's use an instance to grasp,

• A computer on a 10Mbps network is regulated by a token bucket. The token bucket is filled at a rate of 2Mbps. It is initially filled to capacity with 16Megabits. What is the maximum duration for which the computer can transmit at the full 10Mbps?<sup>[8]</sup>

→ Token arrival rate  $\mathbf{\rho}$  is 2 Mbps, which is given in the question.

Token bucket's capacity  $\mathbf{C}$  is 16 Mbits.

Maximum output rate  $\mathbf{M}$  is 10 Mbps.

As per formula  $\mathbf{M} * \mathbf{s} = \mathbf{C} + \mathbf{\rho} * \mathbf{s}$ ,

We need to find time taken  $\mathbf{s}$  which will be equal to  $\mathbf{C}/(\mathbf{M}-\mathbf{\rho})$

So, the maximum burst time taken  $\mathbf{s} = \mathbf{16}/(\mathbf{10}-\mathbf{2}) = \mathbf{2\ seconds}^{[8]}$

In the formula above to determine the maximum burst time,  $\mathbf{r}$  is subtracted from  $\mathbf{M}$ . The explanation for this subtraction is that at the rate of  $\mathbf{r}$ , new tokens are inserted when transmission continues at the full transmission rate of  $\mathbf{M}$ .

## Acknowledgement

As a transport layer protocol, I propose Secure TCP as one of the basic units for the realisation of secure communication mechanisms on the Internet. For data integrity and confidentiality, it offers security services. In addition, it retains compatibility (interoperability) with the classic TCP and offers a process/process contact security service negotiation service.

A cypher is an algorithm for encrypting and decrypting information in cryptology, the discipline concerned with the study of cryptographic algorithms. Where Data Encryption Standard is the archetypal block cipher : an algorithm that transforms a fixed-length string of



plaintext bits into another ciphertext bitstring of the same length through a series of complex operations.<sup>[10]</sup>

When we add Stable TCP to high-speed network environments, this performance is not enough. It can however be improved by implementing fast cypher technologies such as Data Encryption Standard (DES) chips or other hardware-implemented cyphers, because its efficiency is limited by the overhead of encryption/decryption processes.<sup>[9]</sup>

**References :**

Click on the link to go to the corresponding web-page :

1. [Reference from techopedia.](#)
2. [Reference from studytonight.](#)
3. [Resource from javatpoint.](#)
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