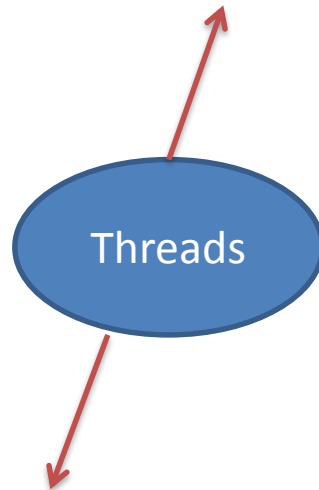
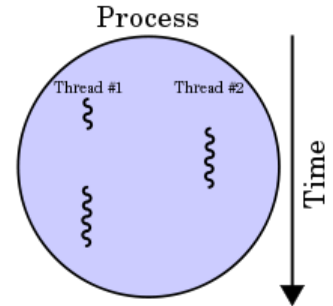


# Thread

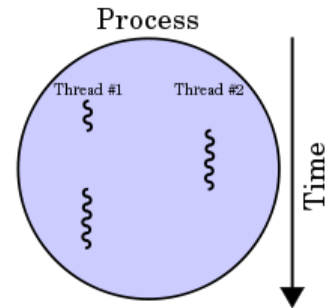
- Thread is **light weight process** created by a process.

**Processes** are used to **execute large**, 'heavyweight' jobs such as working in word, while **threads** are used to **carry out smaller** or 'lightweight' jobs such as auto saving a word document.



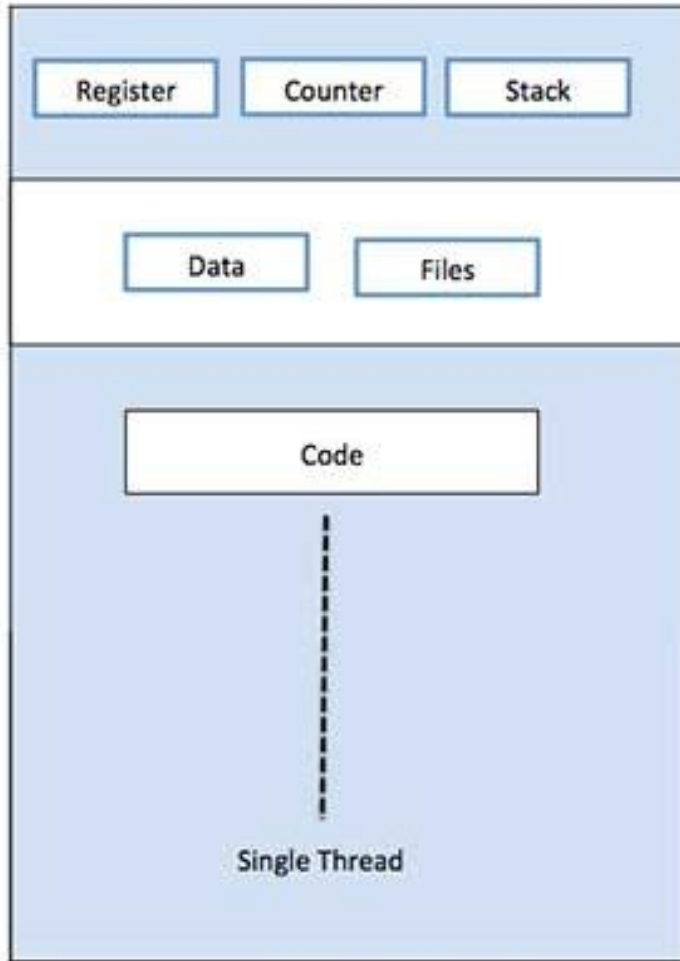
# Thread

- Thread is **light weight process** created by a process.
- Thread is a **single sequence stream** within a process.
- Thread has its own

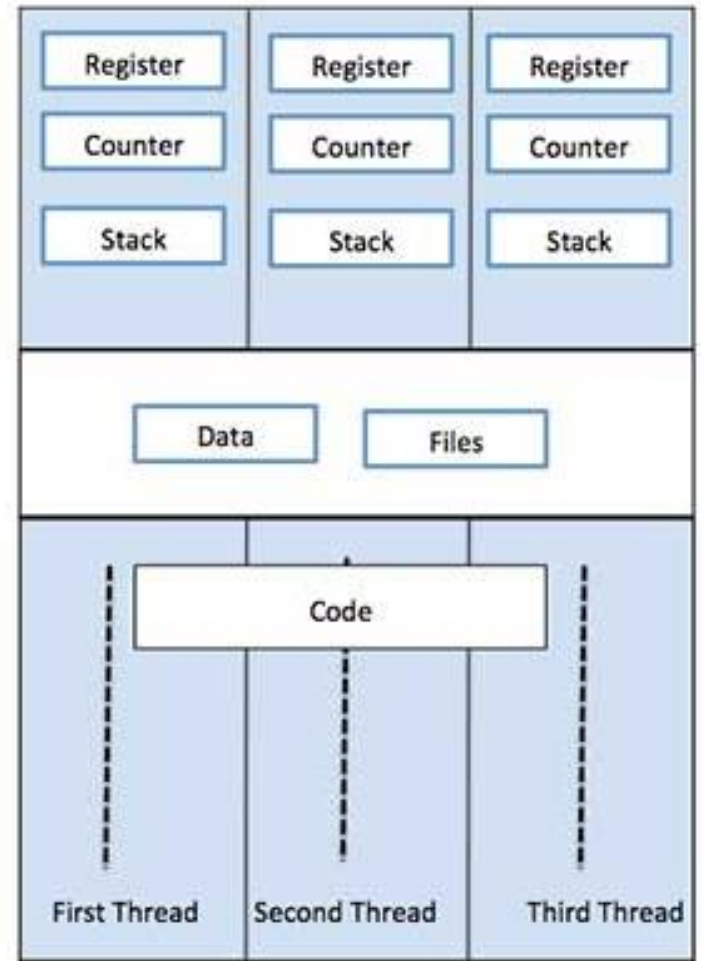


1. **program counter** that **keeps track** of **which instruction to execute next**.
2. **system registers** which **hold its current working variables**.
3. **stack** which **contains the execution history**.

# Single Thread VS Multiple Thread



Single Process P with single thread



Single Process P with three threads

# Similarities between Process & Thread

---

- Like processes threads **share CPU** and **only one thread is running at a time**.
- Like processes threads **within a process execute sequentially**.
- Like processes thread **can create childrens**.
- Like a traditional process, a thread **can be in any one of several states**: running, blocked, ready or terminated.
- Like process threads **have Program Counter, Stack, Registers and State**.

# Dissimilarities between Process & Thread

---

- Unlike processes threads are **not independent** of one another.
- Threads within the same process **share an address space**.
- Unlike processes all threads can **access every address** in the task.
- Unlike processes threads are **design to assist one other**. Note that processes might or might not assist one another because processes may be originated from different users.

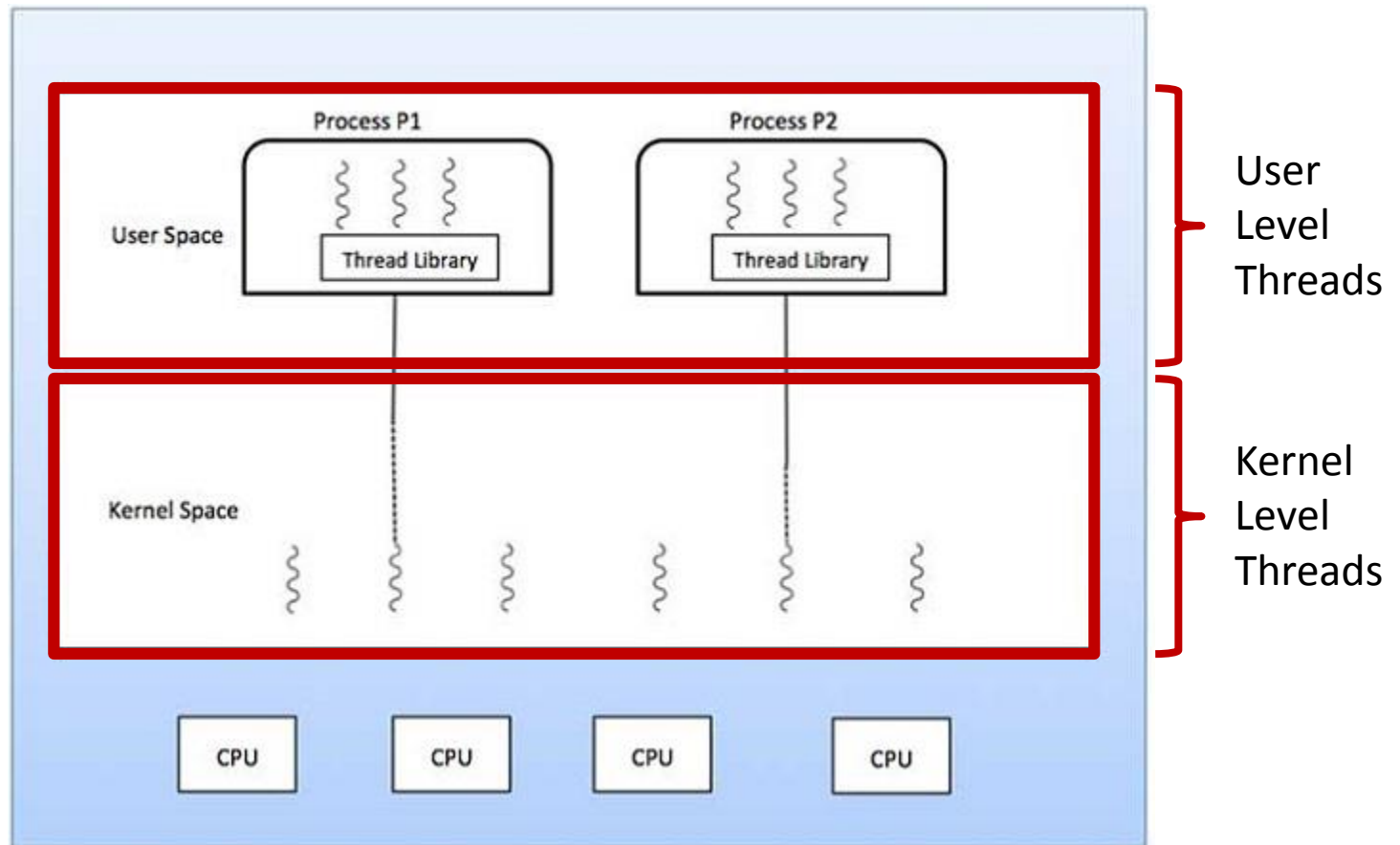
# Advantages of Threads

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- Threads **minimize** the **context switching time**.
- Use of threads **provides concurrency** within a process.
- **Efficient communication**.
- It is more **easy to create** and **context switch** threads.
- Threads can **execute** in **parallel** on multiprocessors.
- With threads, an application can **avoid per-process overheads**
  - Thread creation, deletion, switching easier than processes.
- Threads have **full access** to **address space** (easy sharing).

# Types of Threads

1. Kernel Level Thread
2. User Level Thread



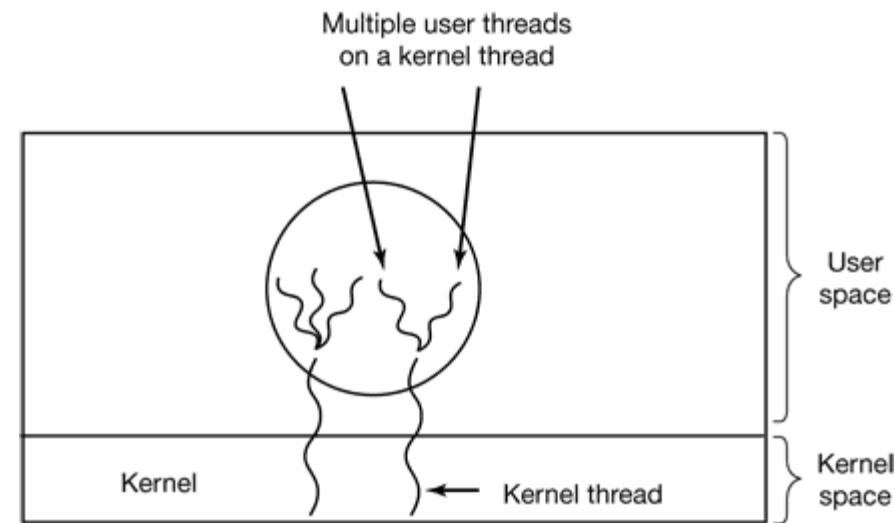
# Types of Threads (Cont...)

| USER LEVEL THREAD   | KERNEL LEVEL THREAD   |
|---|---|
| User thread are <b>implemented by users</b> .   | Kernel threads are <b>implemented by OS</b> .   |
| <b>OS doesn't recognize</b> user level threads.   | Kernel threads are <b>recognized by OS</b> .  |
| <b>Implementation</b> of User threads is <b>easy</b> .  | <b>Implementation</b> of Kernel thread is <b>complex</b> .  |
| <b>Context switch</b> time is <b>less</b> .   | <b>Context switch</b> time is <b>more</b> .   |
| Context switch <b>requires no hardware support</b> .  | Context switch <b>requires hardware support</b> .   |
| If one user level thread perform blocking operation then <b>entire process</b> will be <b>blocked</b> . | If one kernel thread perform blocking operation then <b>another thread</b> with in same process <b>can continue execution</b> . |
| Example : Java thread, POSIX threads.   | Example : Window Solaris  |



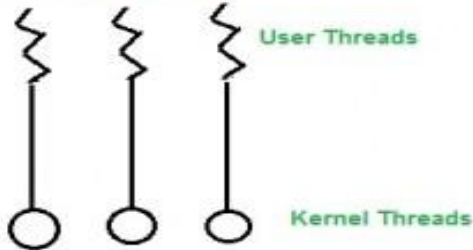
# Hybrid Thread

- Combines the advantages of user level and kernel level thread.
- It **uses kernel level thread** and then **multiplex user level thread on to** some or all of **kernel threads**.
- **Gives flexibility** to programmer that how many kernel level threads to use and how many user level thread to multiplex on each one.
- Kernel is aware of only kernel level threads and schedule it.



# Multi threading models

One to One Model

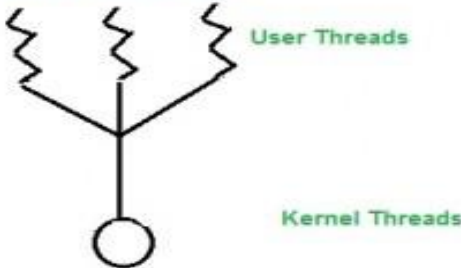


## One to One Model

Each user threads mapped to one kernel thread.

Problem with this model is that creating a user thread requires the corresponding kernel thread.

Many to One Model

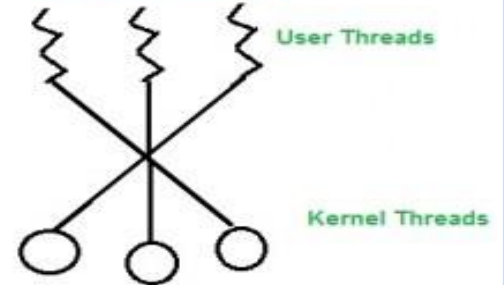


## Many to One Model

Multiple user threads mapped to one kernel thread.

Problem with this model is that a user thread can block entire process because we have only one kernel thread.

Many to Many Model



## Many to Many Model

Multiple user threads multiplex to more than one kernel threads.

Advantage with this model is that a user thread can not block entire process because we have multiple kernel thread.

# Pthread function calls

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1. `Pthread_create`:- Create a new thread
2. `Pthread_exit`:- Terminate the calling thread
3. `Pthread_join`:- Wait for a specific thread to exit
4. `Pthread_yield`:- Release the CPU to let another thread run
5. `Pthread_attr_init`:- Create and initialize a thread's attribute structure
6. `Pthread_destroy`:- Remove a thread's attribute structure