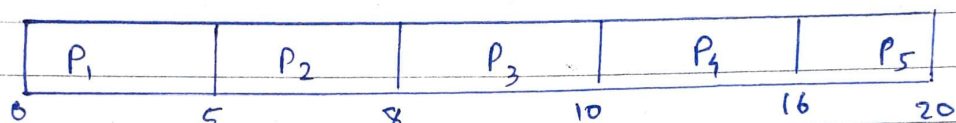


1.) here, given :

Process name:	Arrival time	Service Time
A	0	5
B	1	3
C	2	2
D	3	6
E	4	4

① FCFS

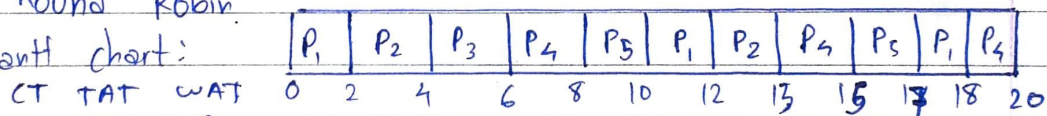
Gantt chart :



	AT	ST	CT	TAT	WAT	
A	0	5	5	5	0	Avg TAT = 9.8 sec
B	1	3	8	7	4	
C	2	2	10	8	6	Avg WT = 5.8 sec
D	3	6	16	13	7	
E	4	4	20	16	12	

② Round Robin

Gantt chart:



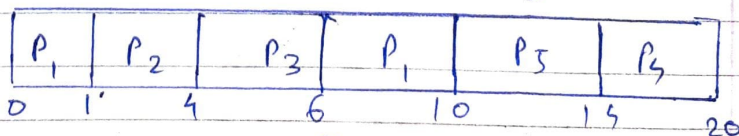
CT	TAT	WAT
14	14	9
13	12	9
6	4	2
20	17	11
18	14	10

$$\text{Avg TAT} = \frac{14+12+4+17+14}{5} = 12.2$$

$$\text{Avg WAT} = \frac{9+9+2+11+10}{5} = 8.2$$

③ SRTN

CT	TAT	WAT
10	10	5
4	3	0
6	4	2
20	17	11
14	10	6

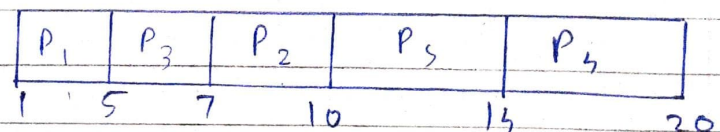


$$\text{Avg TAT} = 8.8$$

$$\text{Avg WT} = 4.8$$

④ SJF

CT	TAT	WAT
5	5	0
10	9	6
7	5	3
20	17	11
14	10	6



$$\text{Avg TAT} = \frac{5+9+5+17+10}{5} = 9.2$$

$$\text{Avg WAT} = \frac{0+6+3+11+6}{5} = 5.2$$

2) here, since  $NEED \text{ matrix} = MAX - ALLOCATION$ ;  
The content of need matrix is as follows,

	MAX	-	ALLOCATION		A	B	C	D
					NEED			
$P_0$	0 0 1 2		0 0 1 2		0	0	0	0
$P_1$	1 7 5 0		1 0 0 0		0	7	5	0
$P_2$	2 3 5 6		1 3 5 4		1	0	0	2
$P_3$	0 6 5 2		0 6 3 2		0	0	2	0
$P_4$	0 6 5 6		0 0 1 4		0	6	4	2

(i) Need matrix obtained:

A	B	C	D
0	0	0	0
0	7	5	0
1	0	0	2
0	0	2	0
0	6	4	2

ii) The system is in safe state.

→ process served in order:  $P_0, P_2, P_1, P_3, P_4$

↳ before serving  $P_0$ , available resources

A	B	C	D
---	---	---	---

1 5 2 0

↳ after serving  $P_0$ ,

1 5 3 2

↳ after serving  $P_2$ ,

2 8 8 6

~~3 8 8 8~~

↳ after serving  $P_1$ ,

3 8 8 6

↳ after serving  $P_3$ ,

3 14 11 8

↳ after serving  $P_4$ ,

3 14 12 12

So the available resources are sufficient for the processes to terminate.

So the system is in safe state.

iii) Yes that request can be granted immediately and still the system will be in safe state.



→ because after serving that request, this results in the value of Available being (1, 1, 0, 0)

	Allocation	Max	Need	Available
P <sub>0</sub>	0012	0012	0000	1100
P <sub>1</sub>	1420	1750	0330	
P <sub>2</sub>	1354	2356	1002	
P <sub>3</sub>	0632	0652	0020	
P <sub>4</sub>	0014	0656	0642	

$$\text{work} = \text{available} = 1100$$

$$\Rightarrow \text{work} = \text{work} + \text{allocation}$$

$$\text{work: } P_0 = 1100 + 0012 = 1112$$

$$P_2 = 1112 + 1354 = 2466$$

$$P_3 = 2466 + 0632 = 2 \ 10 \ 9 \ 8$$

$$P_4 = 2 \ 10 \ 9 \ 8 + 0 \ 0 \ 14 = 2 \ 10 \ 10 \ 12$$

$$P_1 = 2 \ 10 \ 10 \ 12 + 1 \ 4 \ 2 \ 0 = 3 \ 14 \ 12 \ 12$$

∴ Safe sequence = P<sub>0</sub> P<sub>2</sub> P<sub>3</sub> P<sub>4</sub> P<sub>1</sub>

3) here given page trace = 4, 3, 2, 1, 4, 3, 5, 4, 3, 2, 1, 5

→ FIFO:

	4	3	2	1	4	3	5	4	3	2	1	5
F <sub>1</sub>	4	4	4	4	4	4	5	5	5	5	1	1
F <sub>2</sub>		3	3	3	3	3	3	4	4	4	4	5
F <sub>3</sub>			2	2	2	2	2	2	3	3	3	3
F <sub>4</sub>				1	1	1	1	1	1	2	2	2
faults	F	F	F	F			F	F	F	F	F	F

total page faults = 10

→ OPTIMAL.

	4	3	2	1	4	3	5	4	3	2	1	5
F <sub>1</sub>	4	4	4	4	4	4	4	4	4	4	1	1
F <sub>2</sub>		3	3	3	3	3	3	3	3	3	3	3
F <sub>3</sub>			2	2	2	2	2	2	2	2	2	2
F <sub>4</sub>				1	1	1	5	5	5	5	5	5
	F	F	F	F			F				F	

total page faults = 6

→ LRU

	4	3	2	1	4	3	5	4	3	2	1	5
F <sub>1</sub>	4	4	4	4	4	4	4	4	4	4	4	5
F <sub>2</sub>		3	3	3	3	3	3	3	3	3	3	3
F <sub>3</sub>			2	2	2	2	5	5	5	5	1	1
F <sub>4</sub>				1	1	1	1	1	1	2	2	2
faults	F	F	F	F			F			F	F	F

total page faults = 8



4) here,

Given memory portions of size: (in order)

100 KB, 500 KB, 200 KB, 300 KB, 600 KB

processes: (in order)

212 KB, 417 KB, 112 KB, 426 KB

→ First fit:

212 KB is put in 500 KB partition

417 KB is put in 600 KB partition

112 KB is put in 288 KB partition

(∵ a new partition =  $500 \text{ KB} - 212 \text{ KB}$ )

426 KB must wait.

→ Best-fit

212 KB is put in 600 KB partition

417 KB is put in 500 KB partition

112 KB is put in 200 KB partition

426 KB is put in 600 KB partition

→ Worst fit

212 KB is put in 600 KB partition

417 KB is put in 500 KB partition

112 KB is put in 388 KB partition

426 KB must wait.

∴ In this problem, best-fit turns out to be the best.

5) here,

Given: total number of cylinders: 300 [0-299]

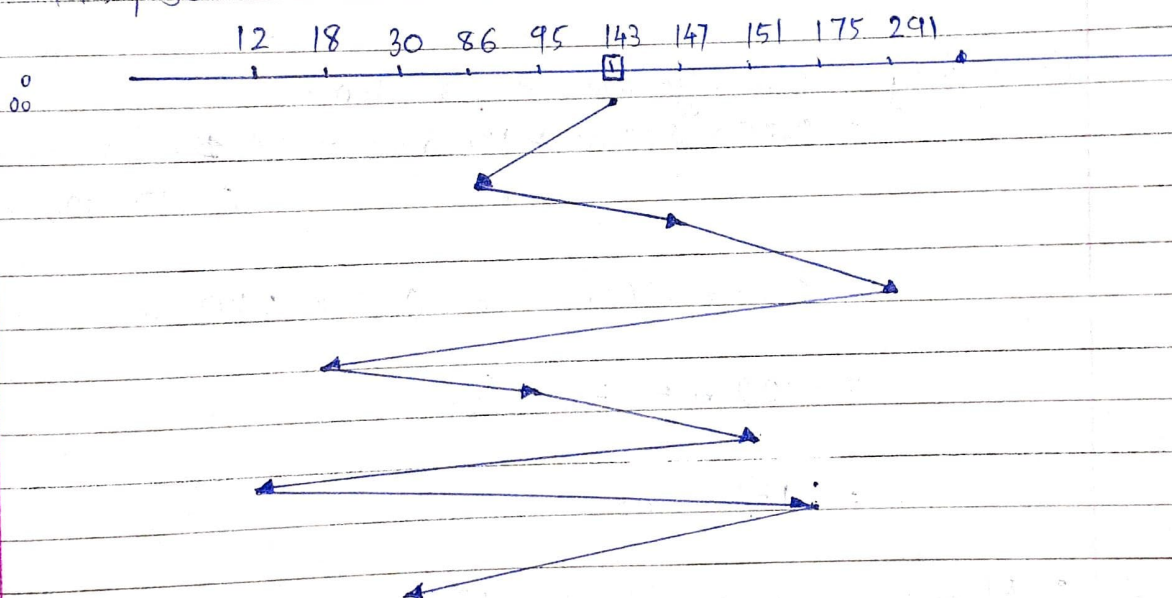
drive currently serving at: 143

previous request served was at: 15

queue of pending request: 86, 147, 291, 18, 95, 151, 12, 175 and 30

(1) FCFS algorithm

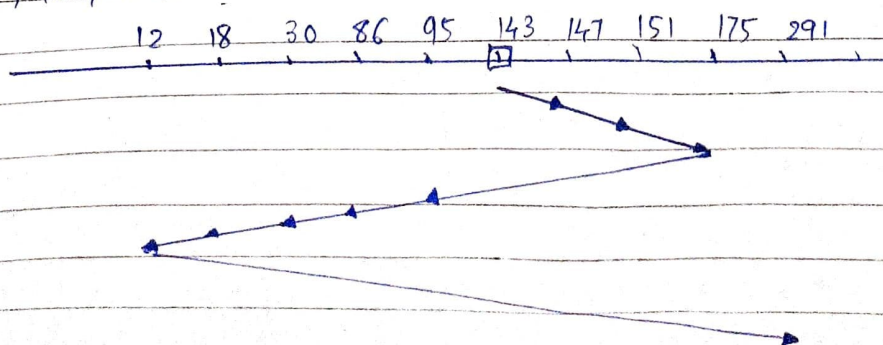
request served queue: 143, 86, 147, 291, 18, 95, 151, 12, 175, 30



$$\therefore \text{Total distance} = 57 + 61 + 144 + 273 + 77 + 56 + 139 + 163 + 145 = 1115$$

(2) SSF algorithm

request served queue: 143, 147, 151, 175, 95, 86, 30, 18, 12, 291





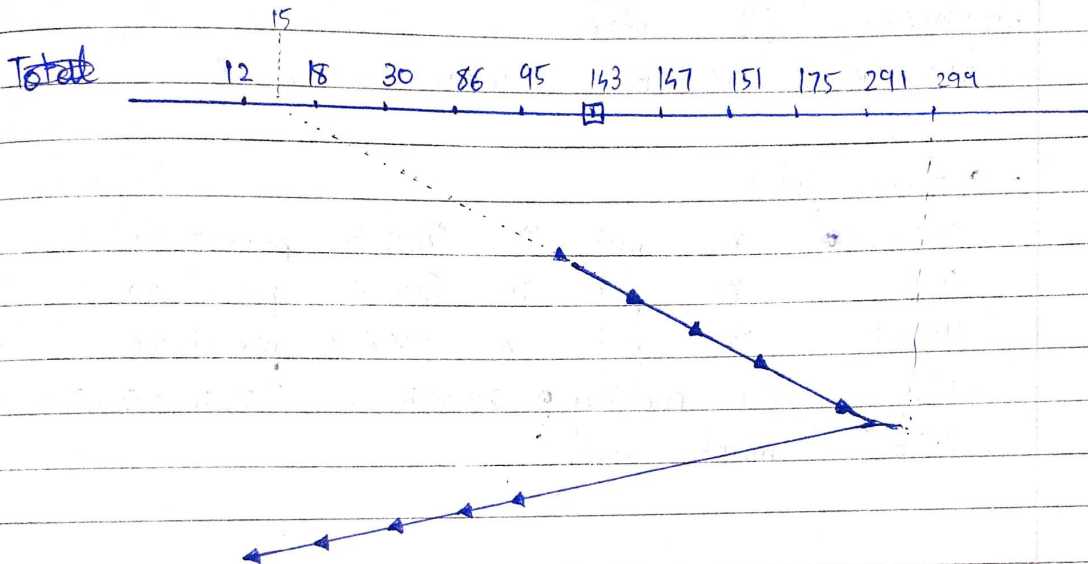
Date \_\_\_\_\_

Page \_\_\_\_\_

Total distance:  $4 + 4 + 24 + 80 + 9 + 56 + 12 + 6 + 279 = 474$

(3) SCAN algorithm

request served queue: 143, 147, 151, 175, 291, 95, 86, 30, 18, 12



Total distance:

$$(299 - 143) + (299 - 12)$$

$$= 443$$

Date \_\_\_\_\_

Page \_\_\_\_\_

6.) here,

Given: Number of level of page table = 1  
 TLB access time = 20 ns  
 Main memory access time = 100 ns  
 TLB hit ratio = 85% = 0.85

Calculating TLB miss ratio:  
 $= 1 - \text{TLB hit ratio}$   
 $= 1 - 0.85$   
 $= 0.15$

Calculating ~~TLB~~ effective access time:  
 $= 0.85 \times \{20 \text{ ns} + 100 \text{ ns}\} + 0.15 \times \{20 \text{ ns} + (1+1) \times 100 \text{ ns}\}$   
 $= 0.85 \times 120 \text{ ns} + 0.15 \times 220 \text{ ns}$   
 $= 102 \text{ ns} + 33 \text{ ns}$   
 $= 135 \text{ ns}$

∴ Effective access time = 135 ns.