

$$3) T(n) = 2T\left(\frac{n}{2}\right) + \frac{n}{\log n}$$

⇒ Here, $a=2$, $b=2$, $d=1$, $p=-1$

So, $a = b^d$ ($\because 2 = 2^1$)

So, case (II)(ii) is applicable.

$$\begin{aligned} \text{Thus, } T(n) &= \Theta(n^{\log_b a} \log^2 n) \\ &= \Theta(n^{\log_2 2} \log^2 n) \\ &= \underline{\underline{\Theta(n \log^2 n)}}. \end{aligned}$$

$$4) T(n) = \sqrt{2} T\left(\frac{n}{2}\right) + \log n$$

⇒ Here, $a = \sqrt{2}$, $b=2$, $d=0$, $p=1$.

So, $a > b^d$ ($\sqrt{2} > 2^0$)

So, case (III) is applicable.

$$\begin{aligned} \text{Thus, } T(n) &= \Theta(n^{\log_b a}) \\ &= \Theta(n^{\log_2 \sqrt{2}}) \\ &= \Theta(n^{\log_2 2^{\frac{1}{2}}}) \\ &= \underline{\underline{\Theta(\sqrt{n})}} \end{aligned}$$

NOTE:- The following equations cannot be solved using Master method

$$1) T(n) = 2^n T\left(\frac{n}{2}\right) + n^n$$

(a is not a constant; the no. of subproblems should be fixed)

$$2) T(n) = 0.5 T\left(\frac{n}{2}\right) + n$$

($a < 1$; cannot have less than one subproblem)

$$3) T(n) = 64 T\left(\frac{n}{8}\right) - n^2 \log n$$

($f(n)$ cannot be negative)

$$4) T(n) = T\left(\frac{n}{2}\right) + n(2 - \cos n)$$

(trigonometric functions cannot be part of $f(n)$).

Change of Variable method

$$1) T(n) = 2 T(\sqrt{n}) + \log n$$

$$\Rightarrow \text{Let } n = 2^m \Rightarrow m = \log_2 n$$

$$\text{So, } T(2^m) = 2 T(2^{\frac{m}{2}}) + m \longrightarrow \textcircled{1}$$

$$\text{Let } T(2^m) = S(m) \Rightarrow T(2^{\frac{m}{2}}) = S\left(\frac{m}{2}\right)$$

So, eqⁿ $\textcircled{1}$ becomes ;

$$S(m) = 2 S\left(\frac{m}{2}\right) + m$$

We can apply Master method to the above recurrence.

$$a=2, b=2, d=1, p=0.$$

So, $a = b^d$ (case (II) is applicable)

$$\therefore S(m) = \Theta(m^d \log m)$$

$$S(m) = \Theta(m \log m)$$

$$\Rightarrow T(2^m) = \Theta(m \log m)$$

$$\Rightarrow \boxed{T(n) = \Theta(\log_2 n \log \log_2 n)} \quad (\text{putting } m = \log_2 n)$$

$$2) T(n) = 3T(\sqrt{n}) + \log n$$

$$\Rightarrow \text{Let } n = 2^m \Rightarrow m = \log_2 n$$

$$\text{So, } T(2^m) = 3T(2^{\frac{m}{2}}) + m$$

$$\text{Consider } T(2^m) = S(m) \Rightarrow T(2^{\frac{m}{2}}) = S(\frac{m}{2})$$

$$\Rightarrow S(m) = 3S(\frac{m}{2}) + m$$

So, we can apply master method to the above recurrence,

$$a=3, b=2, d=1, p=0.$$

$$\text{Here, } a > b^d \quad (3 > 2^1)$$

$$\text{So, } S(m) = \Theta(m^{\log_2 3})$$

$$\Rightarrow T(2^m) = \Theta(m^{\log_2 3})$$

$$\Rightarrow \boxed{T(n) = \Theta((\log_2 n)^{\log_2 3})}$$