

I) The Assignment problem

- The Assignment problem is a fundamental "combinatorial optimization" problem.
- It is a special type of problem in which the objective function is to find the optimum allocation of a number of jobs(tasks) to an equal number of agents(persons).
- It is assumed that each person can perform each job(task), but with varying efficiency.

* General-form of an Assignment problem :

The assignment problem can be stated in the form of $n \times n$ matrix ; where each entry $[c_{ij}]$ represents the cost of effectiveness matrix i.e. the cost of assigning i^{th} person to j^{th} job.

Jobs

	1	2	3	...	n
1	c_{11}	c_{12}	c_{13}		c_{1n}
2	c_{21}	c_{22}	c_{23}		c_{2n}
3					
...					
...					
...					
n	c_{n1}	c_{n2}	c_{n3}		c_{nn}

Persons

A person can be assigned to n jobs in $n!$ possible ways. One method can be to find all possible $n!$ assignments and evaluate total cost in all cases. Then the assignment with the minimum cost will give the optimal assignment. But this method is extremely infeasible (laborious).

Mathematical Formulation

$$\text{Minimum } Z = \sum_{i=1}^n \sum_{j=1}^n c_{ij} x_{ij}$$

$$\text{where } x_{ij} = \begin{cases} 1; & \text{if } i^{\text{th}} \text{ person is assigned to job "j"} \\ 0; & \text{if } i^{\text{th}} \text{ person is not assigned to job "j"} \end{cases}$$

subject to the conditions:

$$\sum_{i=1}^n x_{ij} = 1, \quad j = 1, 2, 3, \dots, n$$

(It means that only one job is done by the i^{th} person
($i = 1, 2, \dots, n$)

Hungarian Assignment Method (Reduced Matrix method)

- 1) Row reduction (Identify min. element in each row and subtract it from each element of that row).
- 2) Column reduction (Identify min. element in each column and subtract it from each element of that column).

3) Perform row-scan and column-scan respectively, and allocate the job "i" to the person "j" according to the placement of "0".

Q-1:- Find optimal allocation of jobs to persons for the following :-

	1	2	3	4
J ₁	5	9	3	6
J ₂	8	7	8	2
J ₃	6	10	12	7
J ₄	3	10	8	6

Solⁿ:- 1) Row reduction

	1	2	3	4
J ₁	2	6	0	3
J ₂	6	5	6	0
J ₃	0	4	6	1
J ₄	0	7	5	3

2) Column reduction

	1	2	3	4
J ₁	2	2	0	3
J ₂	6	1	6	0
J ₃	0	0	6	1
J ₄	0	3	5	3

3) Row-scan and Column-scan

	1	2	3	4
J1	2	2	0	3
J2	6	1	6	0
J3	X	0	6	1
J4	0	3	5	3

Allocation:-

J1 \rightarrow 3

J2 \rightarrow 4

J3 \rightarrow 2

J4 \rightarrow 1

cost

3

2

10

3

18 Rs.

Q-2

	1	2	3	4	5
J1	11	7	10	17	10
J2	13	21	7	11	13
J3	13	13	15	13	14
J4	18	10	13	16	14
J5	12	8	16	19	10

Solⁿ:-

1) Row reduction

	1	2	3	4	5
J1	4	0	3	10	3
J2	6	14	0	4	6
J3	0	0	2	0	1
J4	8	0	3	6	4
J5	4	0	8	11	2

2) Column reduction

	1	2	3	4	5
J1	4	0	3	10	2
J2	6	14	0	4	5
J3	0	0	2	0	0
J4	8	0	3	6	3
J5	4	0	8	11	1

3) Row-Scan and Column-Scan

	1	2	3	4	5
J1	4	0	3	10	2
J2	6	14	0	4	5
J3	0	X	2	X	X
J4	8	X	3	6	3
J5	4	X	8	11	1

(only 3 assignments have been made)

Procedure :-

- 1) Tick the unassigned rows. (✓)
- 2) If the marked row contains a zero, then mark the corresponding column.
- 3) If the marked column contains an assignment, then mark the corresponding row.
- 4) Draw lines through unticked rows and ticked columns.

	1	2	3	4	5	
J1	4	0	3	10	2	✓
J2	6	14	0	4	5	—
J3	0	X	2	X	X	—
J4	8	X	3	6	3	✓
J5	4	X	8	11	1	✓

	1	2	3	4	5	
J ₁	3	0	2	9	1	✓
J ₂	6	15	0	4	5	
J ₃	0	1	2	7	8	
J ₄	7	0	2	5	2	✓
J ₅	3	0	7	10	0	

	1	2	3	4	5	
J ₁	2	0	1	8	0	✓
J ₂	6	16	0	4	5	
J ₃	0	2	2	0	0	
J ₄	6	0	1	4	1	✓
J ₅	3	1	7	10	0	✓

	1	2	3	4	5	
J ₁	1	0	0	7	0	✓
J ₂	6	17	0	4	6	✓
J ₃	0	3	2	0	0	✓
J ₄	5	0	0	3	1	✓
J ₅	2	1	6	9	0	✓

	1	2	3	4	5
J ₁	0	0	0	6	0
J ₂	5	17	0	3	6
J ₃	0	4	3	0	1
J ₄	4	0	0	2	1
J ₅	1	1	6	8	0

J₁ → 1

J₂ → 3

J₃ → 4

J₄ → 2

J₅ → 5

Total cost :-

$$11 + 7 + 13 + 10 + 10$$

$$= \underline{\underline{51 \text{ Rs.}}}$$