

$$1) T(n) - 4T(n-1) + 4T(n-2) = 0$$

Acc to the form $a_0 t_n + a_1 t_{n-1} + \dots + a_k t_{n-k} = 0$

$$a_0 = 1, a_1 = -4, a_2 = 4, k=2$$

$$\therefore x^2 - 4x + 4 = (x-2)^2 = 0 \quad x = 2, 2$$

$$\text{roots: } r = 2, m = 2$$

$$t_n = C_1 2^n + C_2 n 2^n$$

$$\text{when } n=0, C_1 = t_0$$

$$\therefore t_n = t_0 2^n + C_2 n 2^n$$

$$2) T(n) - 5T(n-1) + 6T(n-2) = 0$$

$$a_0 = 1, a_1 = -5, a_2 = 6 = a_k, k=2$$

$$\therefore x^2 - 5x + 6 = 0$$

$$(x-2)(x-3) = 0$$

$$\text{roots } r_1 = 2, m = 2 \text{ \& } r_2 = 3, m = 1$$

$$\therefore t_n = C_1 2^n + C_2 3^n$$

$$3) T(n) + 5(T(n-1) + 6T(n-2)) = 3n^2 - 2n + 1$$

$$a_0 = 1, a_1 = 5, a_2 = 6, k=2$$

$$b = 1, d = 2, P(n) = 3n^2 - 2n + 1$$

$$\therefore (x^2 + 5x + 6)(x-1)^2 = 0$$

$$(x+2)(x+3)(x-1)^2 = 0$$

$$\text{roots, } r_1 = -2, m = 1, r_2 = -3, m = 1, r_3 = 1, m = 2$$

$$\therefore t_n = C_1 (-2)^n + C_2 (-3)^n + C_3 + C_4 n + C_5 n^2$$

$$③ 1.) T(n) = 8T(n/2) + 1000n^2$$

$$a = 8, b = 2, f(n) = 1000n^2$$

$$n \log_b a = n \log_2 8 = n^3$$

$$\text{Here, } f(n) = O(n^{\log_b a - \epsilon}), \text{ where } \epsilon = 1$$

$$\therefore T(n) = \Theta(n^3)$$

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$$2) T(n) = 25(n/2) + 10n$$

$$a=2 \quad b=2 \quad f(n)=10n$$

$$n^{\log_b a} = n^{\log_2 2} = n$$

$$\text{Here, } f(n) = \Theta(n)$$

$$\therefore T(n) = \Theta(n \log n)$$

$$3) T(n) = 2T(n/2) + n^2$$

$$a=2 \quad b=2 \quad f(n)=n^2$$

$$n^{\log_b a} = n$$

$$\text{Here, } f(n) = n^{(\log_b a + \epsilon)} \text{ where } \epsilon = 1$$

$$\therefore T(n) = \Theta(n^2)$$

$$4) T(n) = 3T(n/3) + n \log n$$

$$a=3, \quad b=3; \quad f(n)=n \log n$$

$$n^{\log_a b} = n$$

Here, $f(n)$ is asymptotically larger than n but not polynomially larger.

\therefore Master method doesn't apply to the recurrence.

$$5) T(n) = 9T(n/3) + n$$

$$a=9 \quad b=3 \quad f(n)=n$$

$$n^{\log_a b} = n^2$$

$$\text{Here } f(n) = O(n^{\log_a b - \epsilon}) \text{ where } \epsilon = 1$$

$$\therefore T(n) = \Theta(n^2)$$

$$6) T(n) = 27T(n/3) + n^3$$

$$a=27 \quad b=3 \quad f(n)=n^3$$

$$\text{Here } f(n) = \Theta(n^3)$$

$$\therefore T(n) = \Theta(n^3 \log n)$$

$$7.) T(n) = 8T(n/2) + n^3 \log n.$$

$$a=8 \quad b=2 \quad f(n) = n^3 \log n.$$

$$n^{\log_b a} = n^3.$$

Here $\frac{f(n)}{n^3} = \frac{1}{\log n}.$

\therefore Master method doesn't apply.

$$8.) T(n) = 2T(n/2) + n \log n$$

$$a=2 \quad b=2 \quad f(n) = n \log n.$$

$$n^{\log_a b} = n$$

Here $\frac{f(n)}{n} = \frac{1}{\log n}$

Master method doesn't apply.

$$9.) T(n) = 0.5 T(n/2) + 1/n.$$

$$a=1/2 \quad b=2 \quad f(n) = 1/n$$

$$n^{\log_a b} = n^{\log_{1/2} 1/2} = 1$$

Here $f(n) = \Theta(1/n)$

$$T(n) = \Theta(\log n / n).$$

$$10.) T(n) = T(n/2) + n(2 - \cos n)$$

$$a=1 \quad b=2 \quad f(n) = n(1 - \cos n)$$

$$n^{\log_a b} = n^0 = 1$$

Here, $f(n) = \Omega(n^{\log_a b + \epsilon})$, $\epsilon = 1$

$$\therefore f(n/b) \leq c f(n)$$

$$\frac{n}{2} (2 - \cos n/2) \leq c (2 - \cos n)$$

let $n = 360$

$$\frac{1}{2} (2 + 1) \leq (2 - 1)$$

$$\frac{3}{2} \leq 1 \quad ; \text{ not possible}$$

$$6.) T(n) = nT^2(n/2)$$

$$\text{let } T(2^i) = 2^i T^2(2^{i-1})$$

$$t_i = 2^{i+2} - 1$$

$$\text{let } u_i = \log t_i = \log(2^{i+2} - 1)$$

$$u_i = i + 2 \log t_{i-1}$$

$$u_i = i + 2 u_{i-1}$$

$$\therefore u_i - 2u_{i-1} = i$$

$$\therefore a_0 = 1 ; a_2 = a_n = 2 ; k=1 ; b=1 ; d=1$$

$$\therefore (x-2)(x-1)^2 = 0$$

$$\text{Roots, } r_1 = 2 \quad m_1 = 1 \quad r_2 = 1 \quad m_2 = 2$$

$$\text{Hence, } u_i = C_1 2^i + C_2 + i C_3$$

$$\text{Here, } u_i - 2u_{i-1} = i$$

$$(C_1 2^i + C_2 + i C_3) - 2(C_1 2^{i-1} + C_2 + (i-1) C_3) = i$$

$$\therefore C_2 - 2C_2 + i C_3 - 2i C_3 - 2C_3 = i$$

$$\therefore C_3 = -1, \quad C_2 = 2$$

$$u_i = C_1 2^i - 2 - i$$

$$\text{Now, } t_i = 2^{u_i}$$

$$T(n) = t \log t_i = 2^{4^i - \log i - 2}$$

$$T(n) = \frac{2^{4^n}}{4n}$$

Now, $T(1) = 1/3 = \frac{2^{L_1}}{4}$

$$2^{L_1} = \frac{4}{3}$$

$$L_1 = \log_2 4/3$$

$$L_1 = 2 - \log_2 3$$

$$L_1 = 2 - \log_2 3$$

$$\therefore T(n) = \frac{2^{2n - n \log_2 3}}{4n}$$

$$\therefore T(n) = \frac{2^{2n}}{4n \cdot 3^n}$$