## Institute Of Technology, Nirma University B.Tech. (CE) Sem VI 2CS601 Theory of Computation

## **Tutorial 6**

Q-1 Design CFG for the following languages:

a. 
$$\{a^i b^j c^k \mid i = j + k\}$$

b. 
$$\{a^i b^j c^k \mid j = i + k\}$$

c. 
$$\{a^i b^j c^k \mid j = i \text{ or } j = k\}$$

d. 
$$\{a^i b^j c^k \mid i = j \text{ or } i = k\}$$

e. 
$$\{a^{i}b^{j}c^{k} | i < j \text{ or } i > k\}$$

f. 
$$\{a^i b^j \mid i \le 2j\}$$

g. 
$$\{a^i b^j \mid i < 2j\}$$

$$h. \{a^i b^j \mid i \le j \le 2i\}$$

Q.2 In each each case, what languages are generated for the following CFGs:

- 1.  $S \rightarrow aSa \mid bSb \mid \land$
- 2...  $S \rightarrow aSa \mid bSb \mid a \mid b$
- 3..  $S \rightarrow aSb \mid bSa \mid \land$
- $4. S \rightarrow aSa \mid bSb \mid aAb \mid bAa \mid$

$$A \rightarrow aAa \mid bAb \mid a \mid b \mid \land$$

- 5. S  $\rightarrow$  aS | bS | a
- 6.  $S \rightarrow SS \mid bS \mid a$
- 7. S  $\rightarrow$  SaS | b
- 8. S  $\rightarrow$  aT | bT |  $\land$ 
  - $T \rightarrow aS|bS$

Q.3 Consider the CFG with productions  $S \rightarrow aSbScS \mid aScSbS \mid bSaScS \mid bScSaS \mid cSaSbS \mid cSbSaS \mid \wedge$ . Does this generate the language  $\{x \in \{a,b,c\}^* \mid n_a(x) = n_b(x) = n_c(x)\}$ ? Prove your answer.

Q:4 Convert the following CFG to Chomsky Normal Form:

1. 
$$S \rightarrow aAbB$$

$$A \rightarrow Ab/b$$

$$B \rightarrow Ba/a$$

- 2.  $S \rightarrow aA \mid bB$ 
  - $A \rightarrow bAA \mid a$
  - $B \rightarrow BBa \mid b$
- 3.  $S \rightarrow aAC$ 
  - A →aB|bAB

do not consider this:

S -> null A -> S

- 4. S→0X1Y
  - $X \rightarrow 0X \mid 0$
  - $Y \rightarrow 1Y/1$
- 5.  $S \rightarrow abSab \mid a \mid aAAb$ 
  - $A \rightarrow bS \mid aAAb \mid c$
- Q:5 Explain the term ambiguity and prove that the following grammar is ambiguous grammar.
  - $S \rightarrow S+S \mid S-S \mid S*S \mid S-S \mid a$
- Q.6 Remove unit productions from the following grammar and generate equivalent grammar:
  - 1.  $S \rightarrow ABC \mid 0$ 
    - $A \rightarrow 1$
    - $B \rightarrow C \mid 0$
    - $C \rightarrow D$
    - $D \rightarrow E$
    - $E \rightarrow 2$
  - 2.  $S \rightarrow ABCD|0$ 
    - $A \rightarrow BC|1$
    - в→с
    - $C\rightarrow D$
    - D→d