

Graphs 2

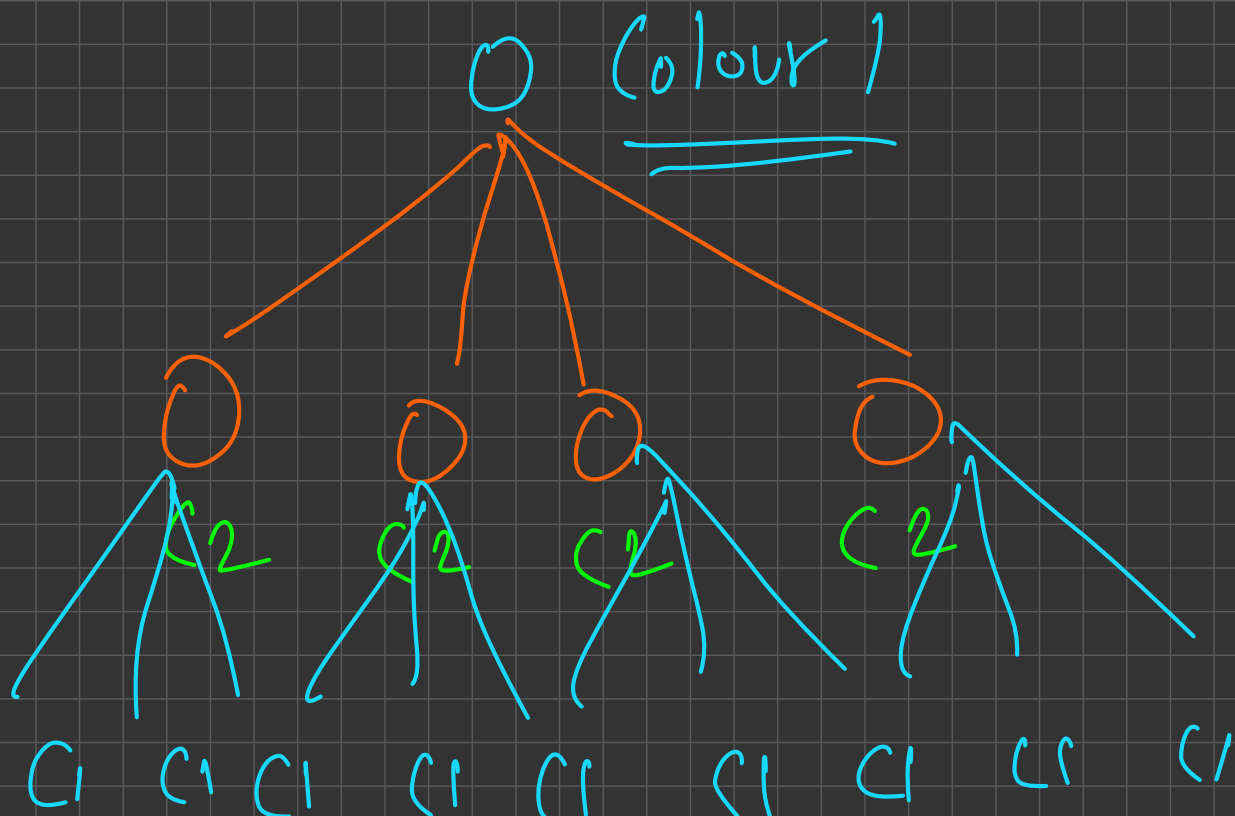
Bipartite Graphs, Dijkstra, Bellman Ford

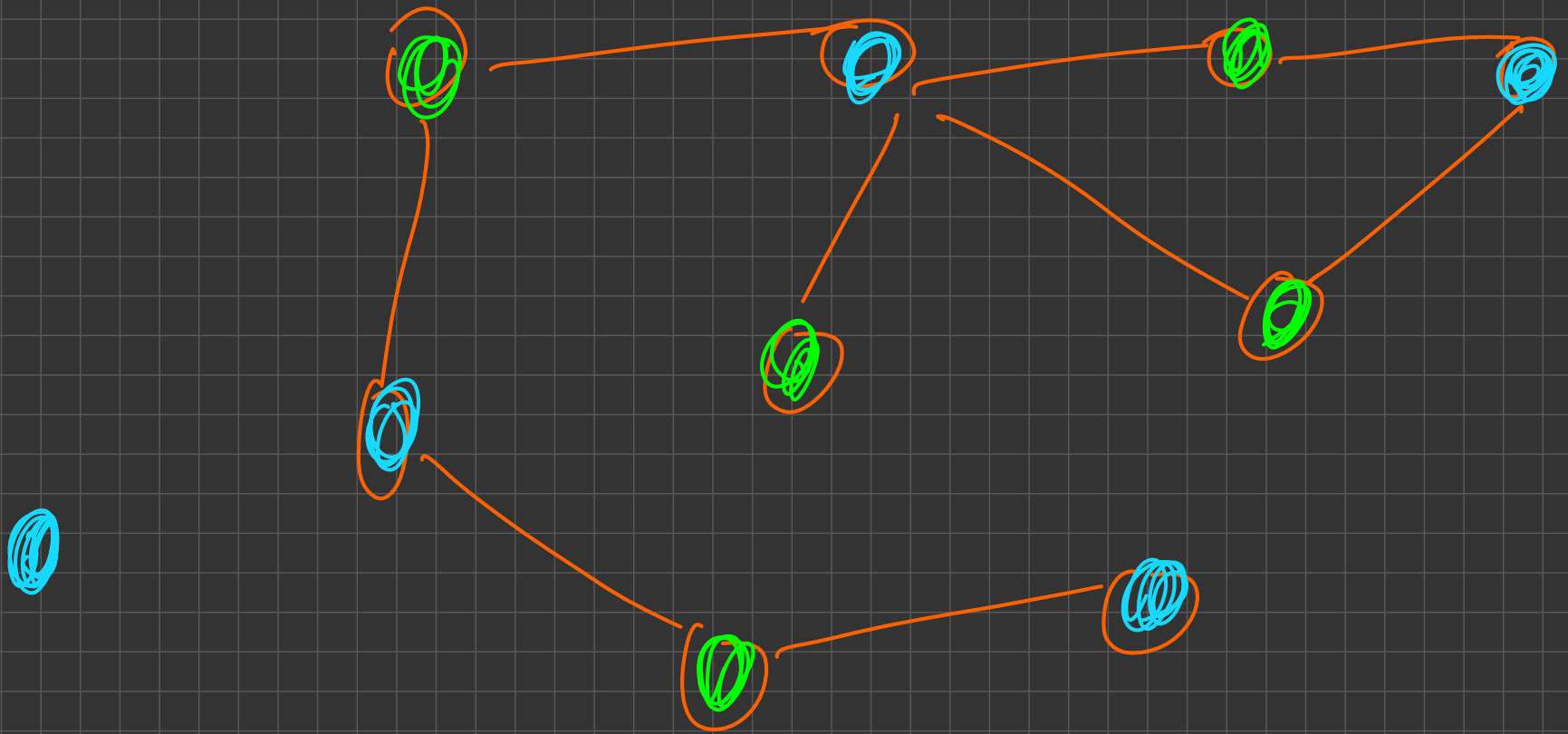
-Priyansh Agarwal

Bi-partite Graphs

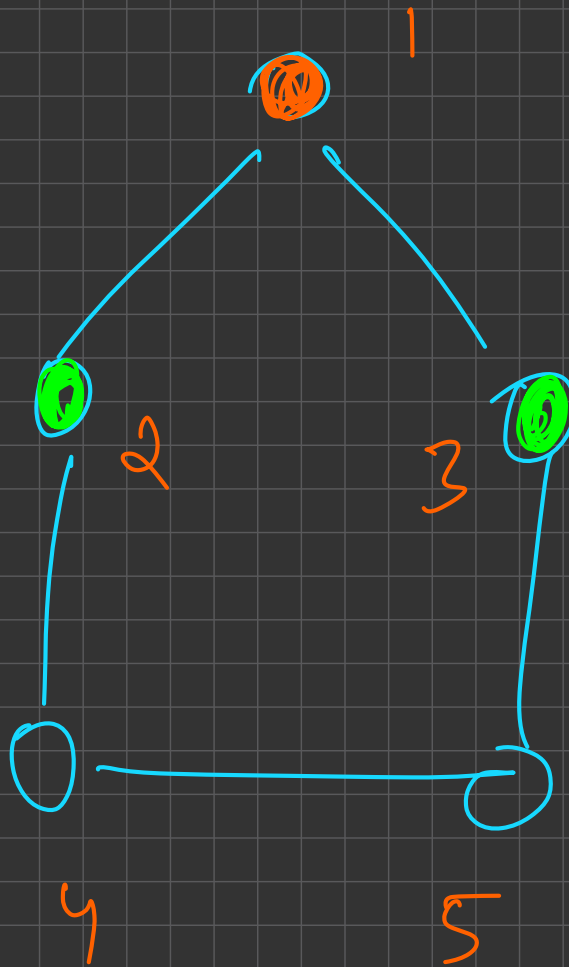
- Definition  2 neighbours don't have same colour
- Algorithm
- Odd Length Cycle
- Tree Property 
- Problem: [Link](#)

A graph can be coloured in a bipartite manner if and only if there is no odd length cycle in the graph

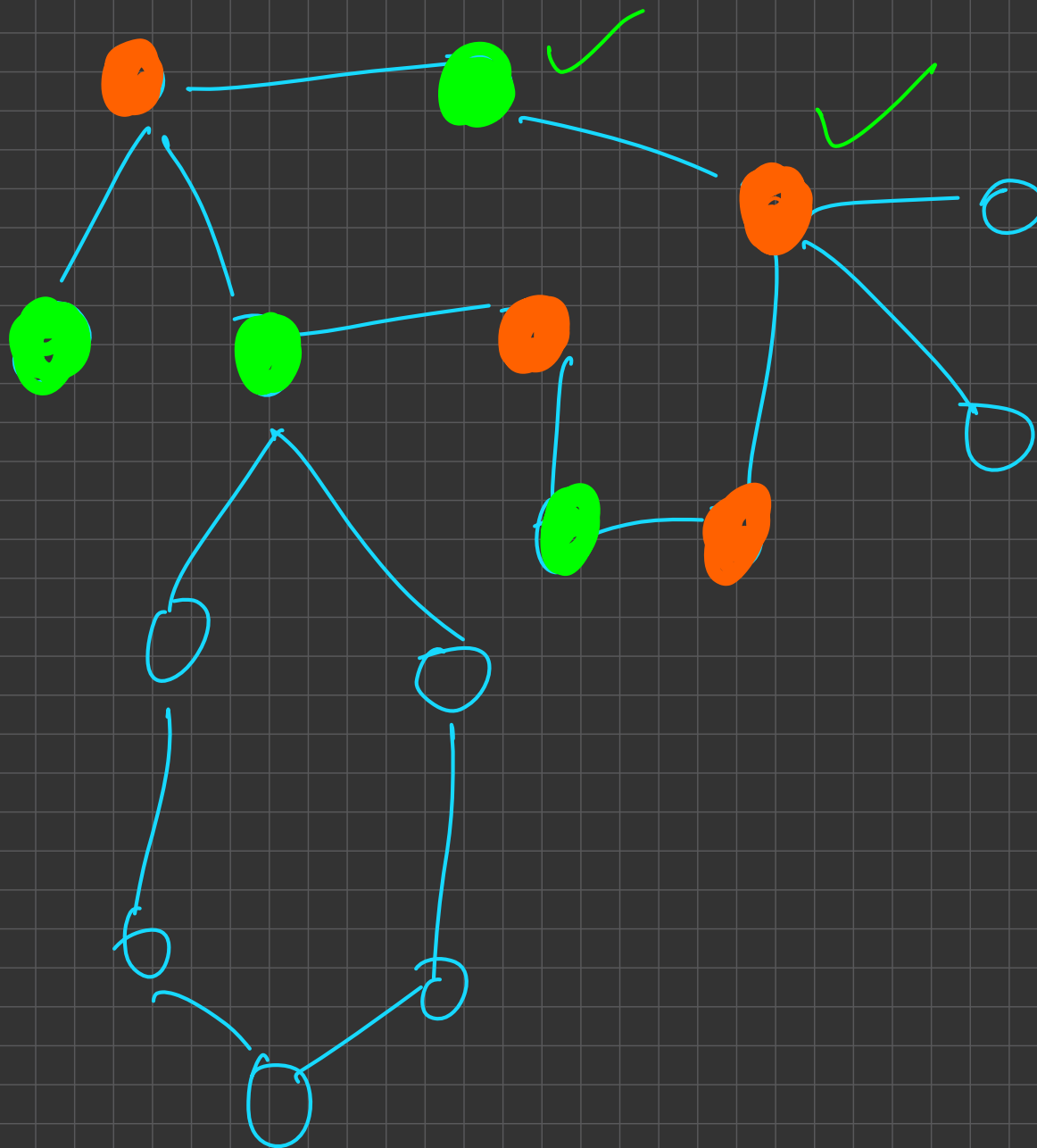




2 neighbours shouldn't have the
same colour

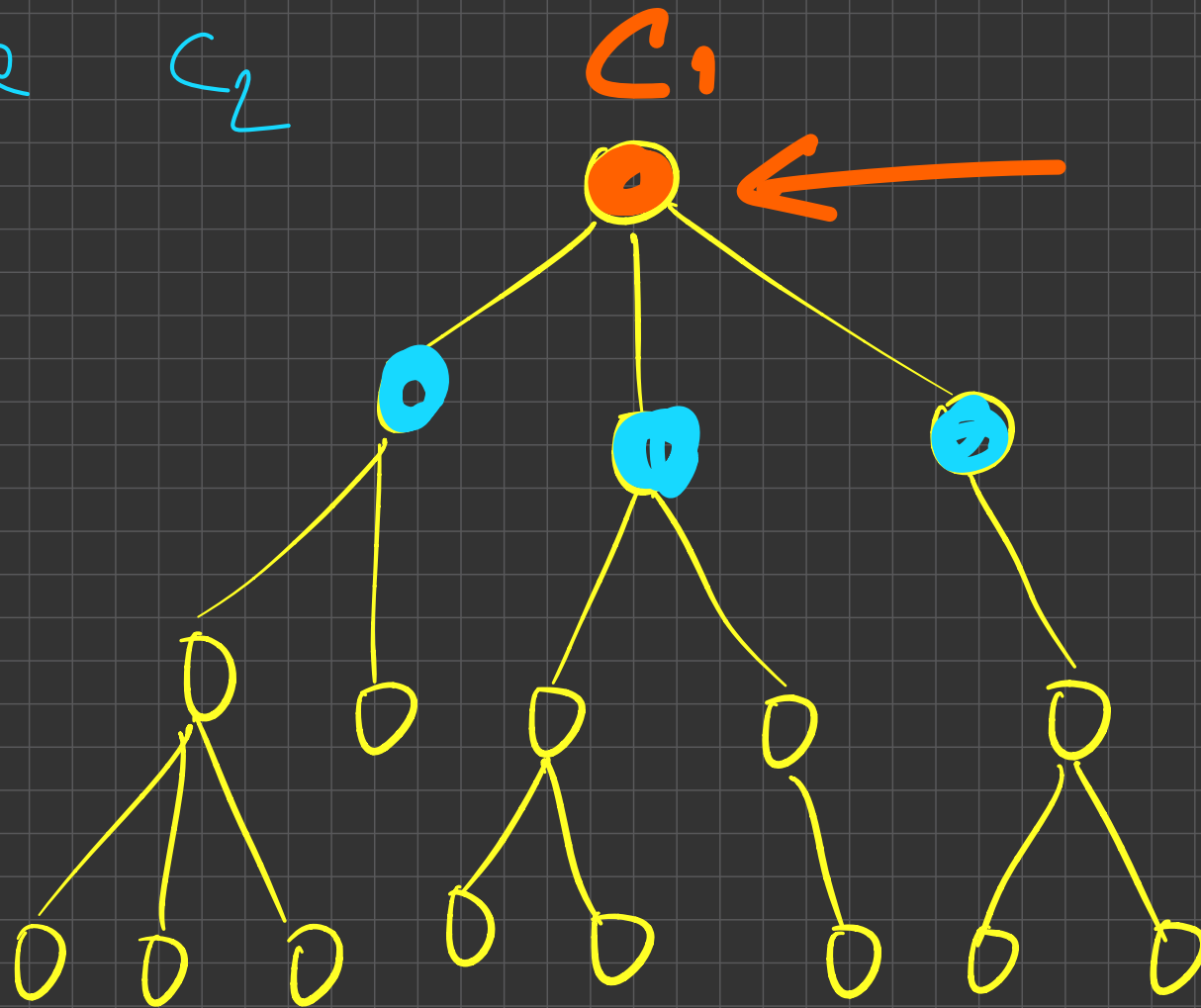


α Bipartite



if level of node = odd $\rightarrow C_1$

else C_2



bool dfs (int curr, edges, visited, color)

if (color[curr] == 0) \rightarrow color[curr] = 1

for (int child : edges[curr])

if (color[child] == 0)

color[child] = color[curr] ^ 1

dfs (child)

else if (color[child] == color[curr])

return false

Dijkstra (most important)



- Single Source Shortest Path Algorithm - Idea + Visualization

- Non-negative edge weights

- Proof/Intuition:

- On every iteration the marked vertex is the one that can never have a better distance later on.

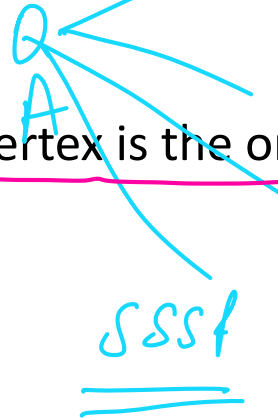
- Code

- Retrieving the shortest path?

- Problems

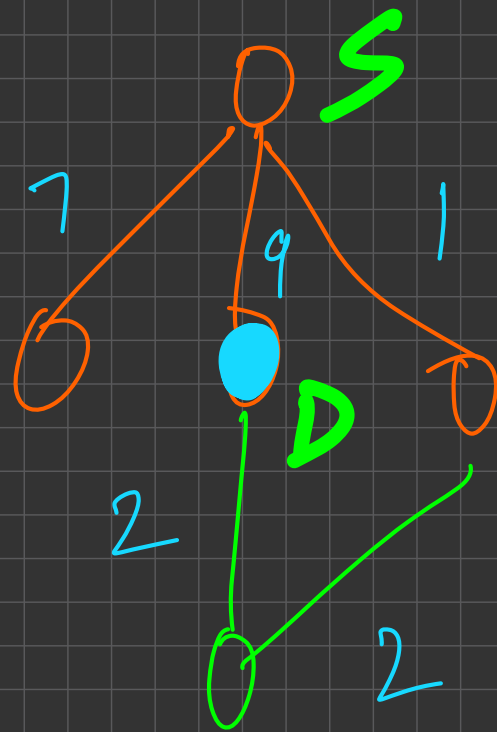
- Google Interview Problem

undirected +
directed
weighted



Dijkstra Tree

✓ = 2



Dijkstra

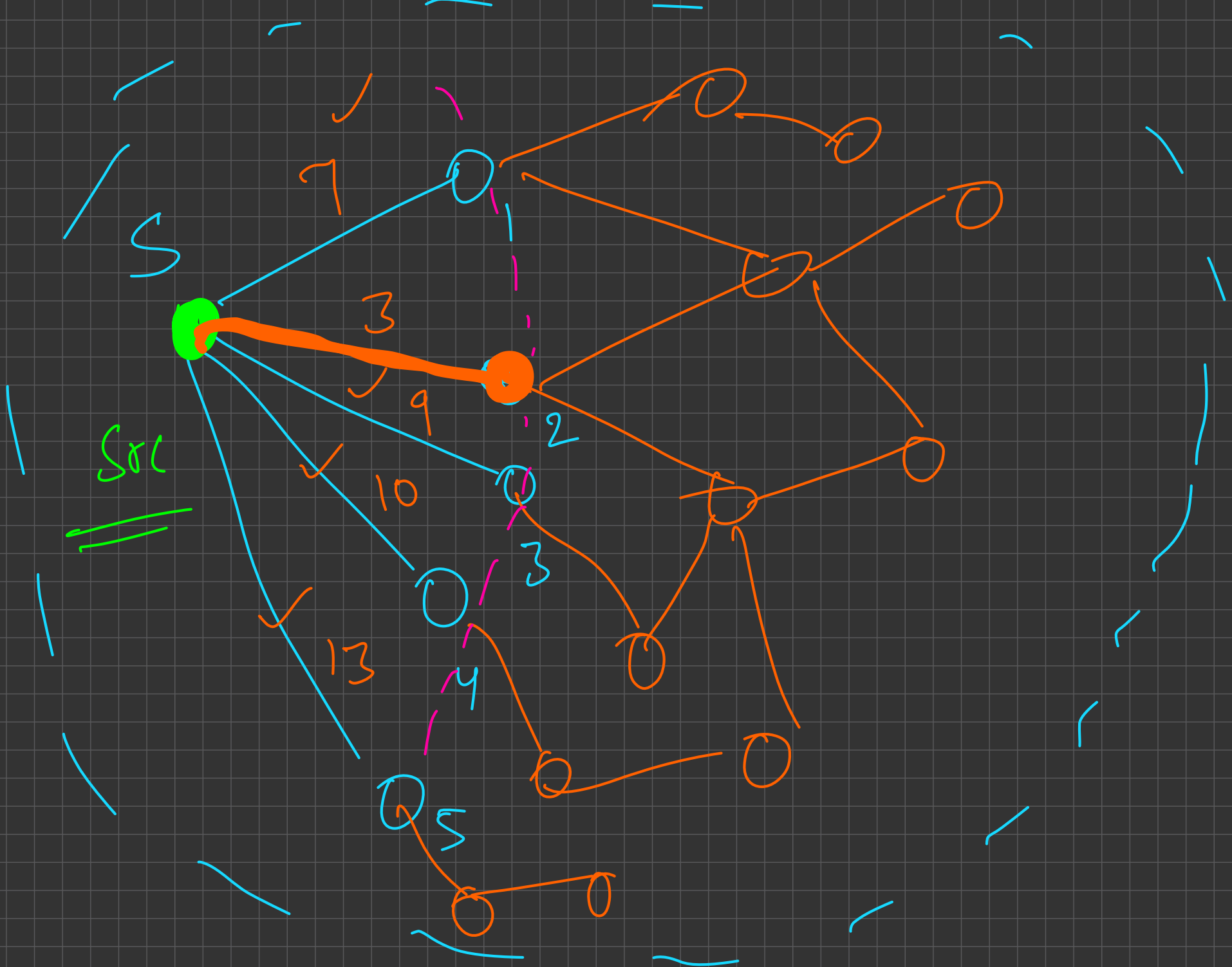
undirected directed

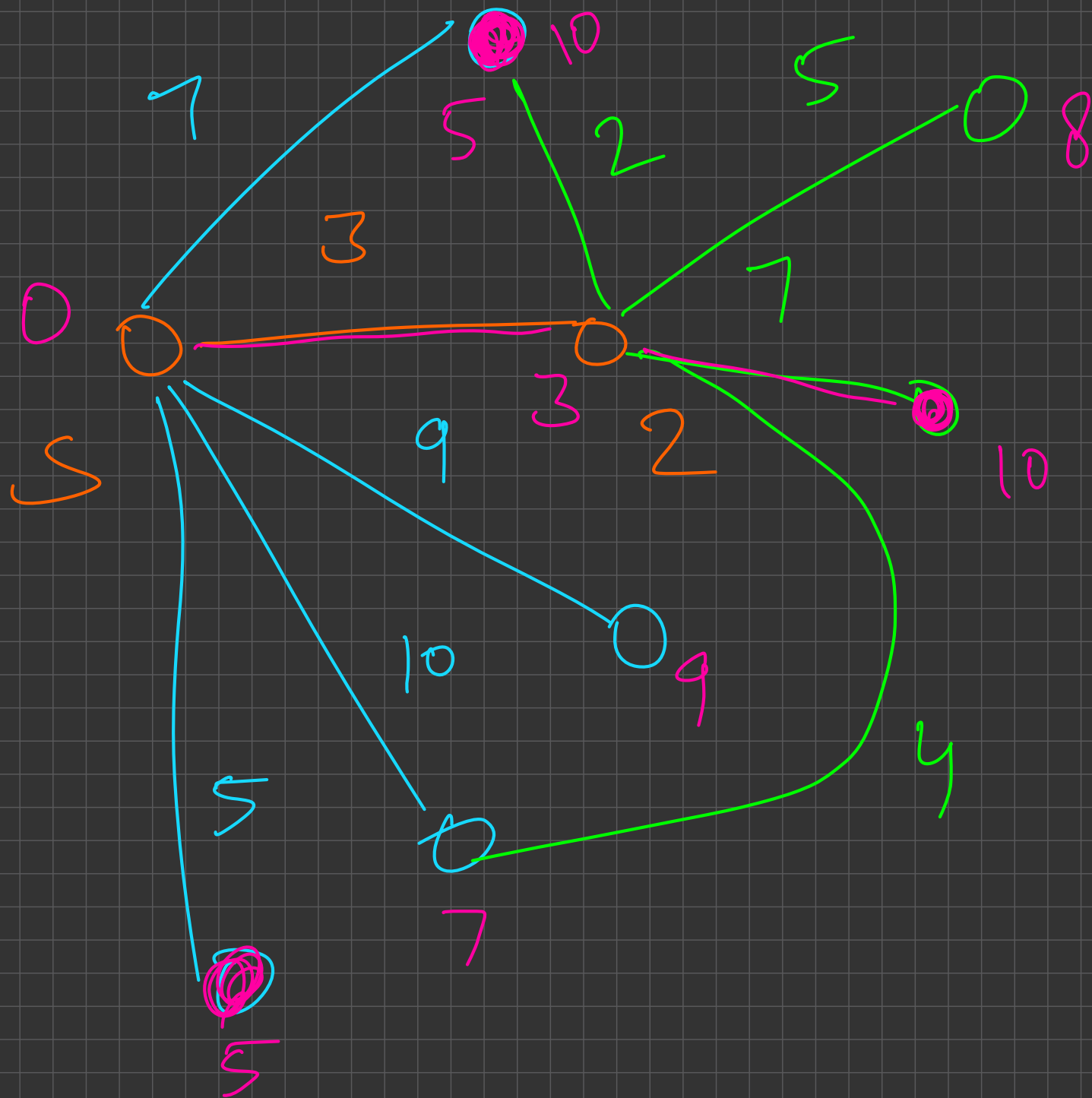
$O(n+m)$

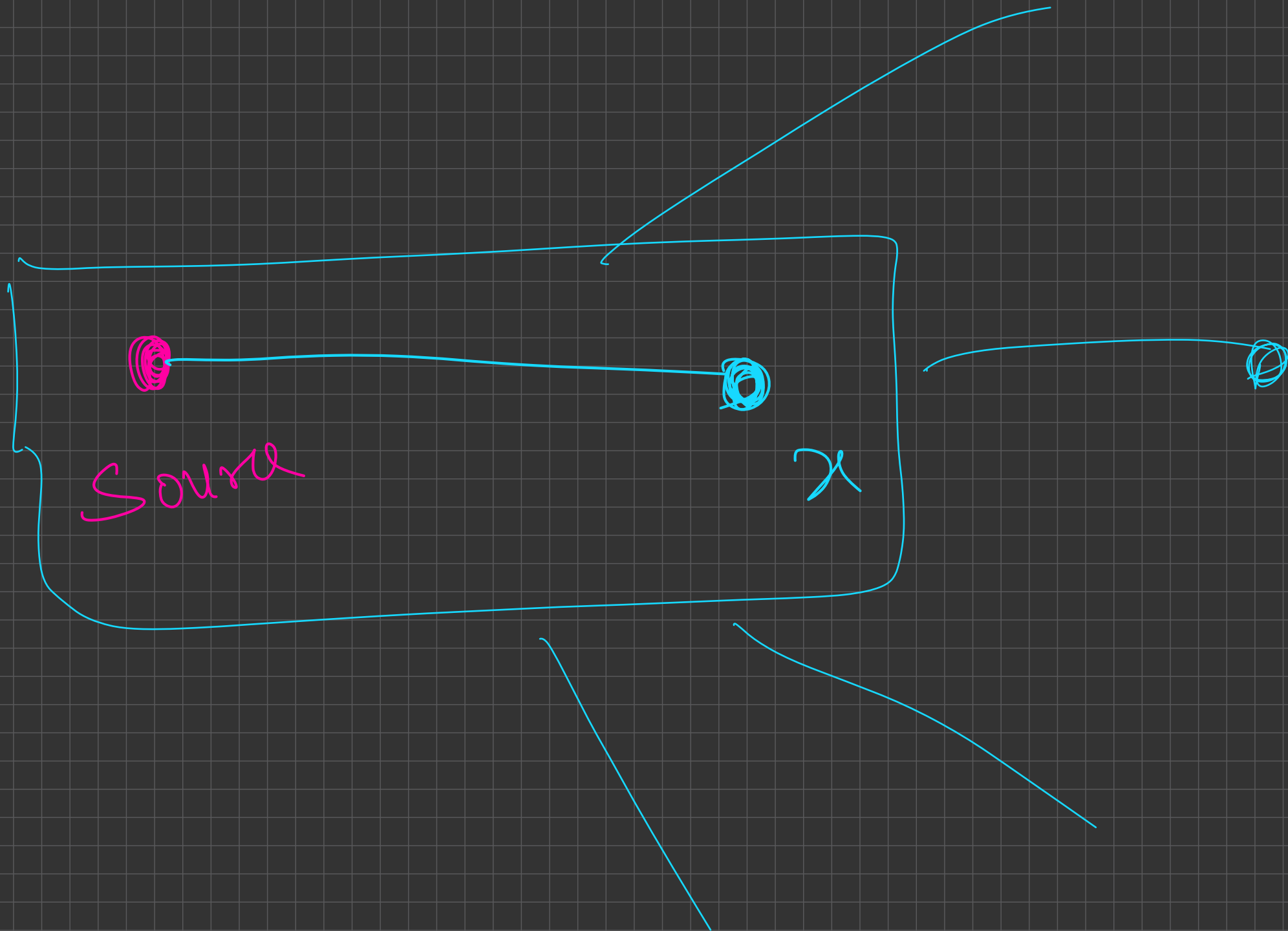
$O((n+m)(\log n))$

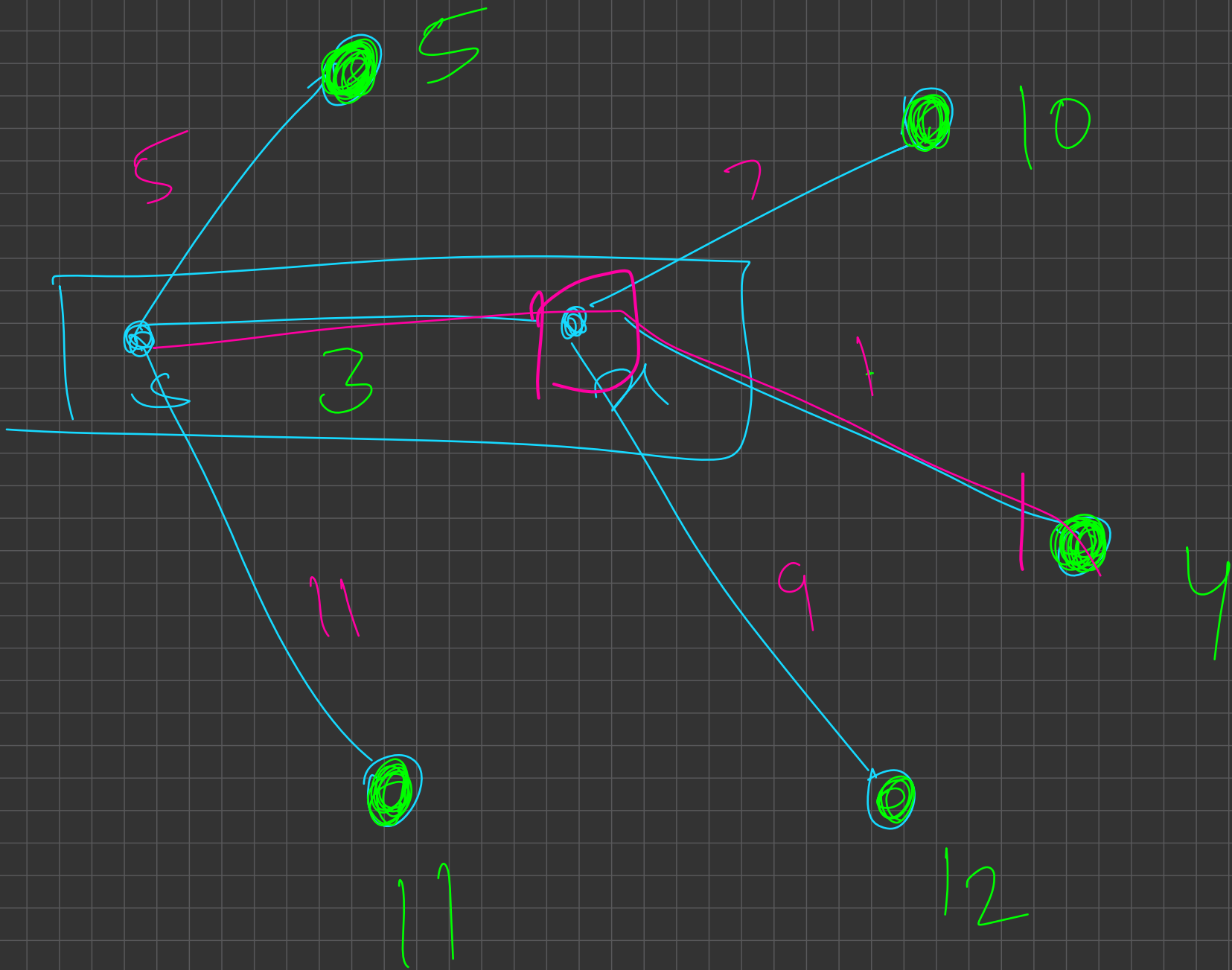
non
negative

edge weights

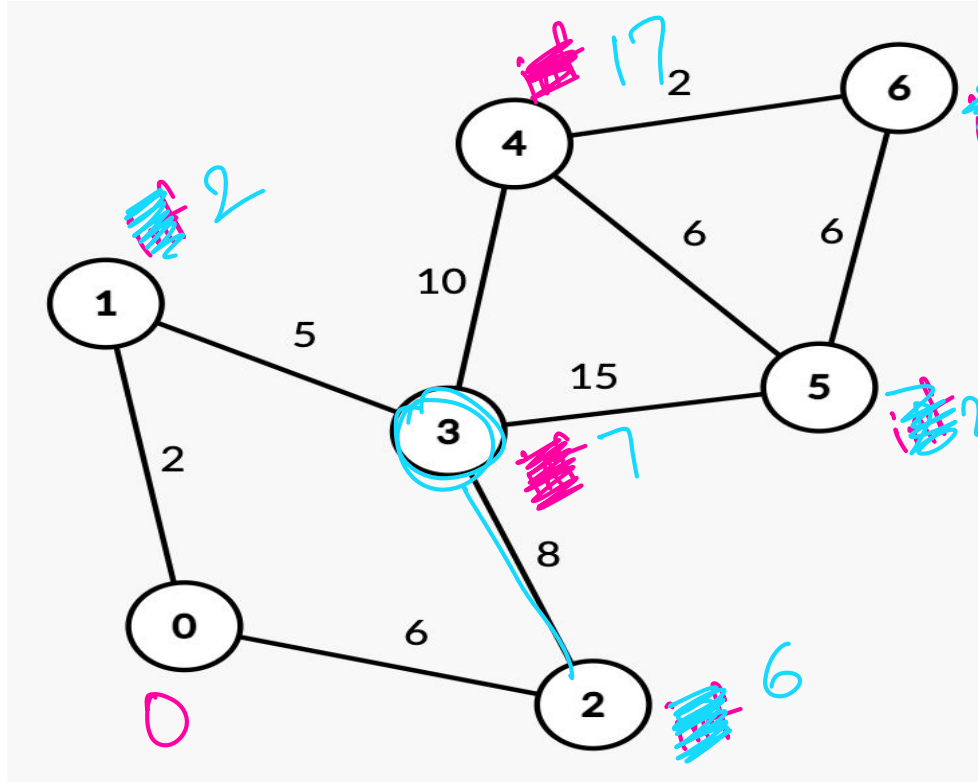








Dijkstra Visualization



✓ 0	0	Sat
✓ 1	2	Sat
2	6	Sat
3	7	Sat
4	17	Sat
5	20	Sat
6	19	Sat

n steps to saturating n nodes

min best answer

$O(\log n)$

iterate over the neighbours of this node

update the answer

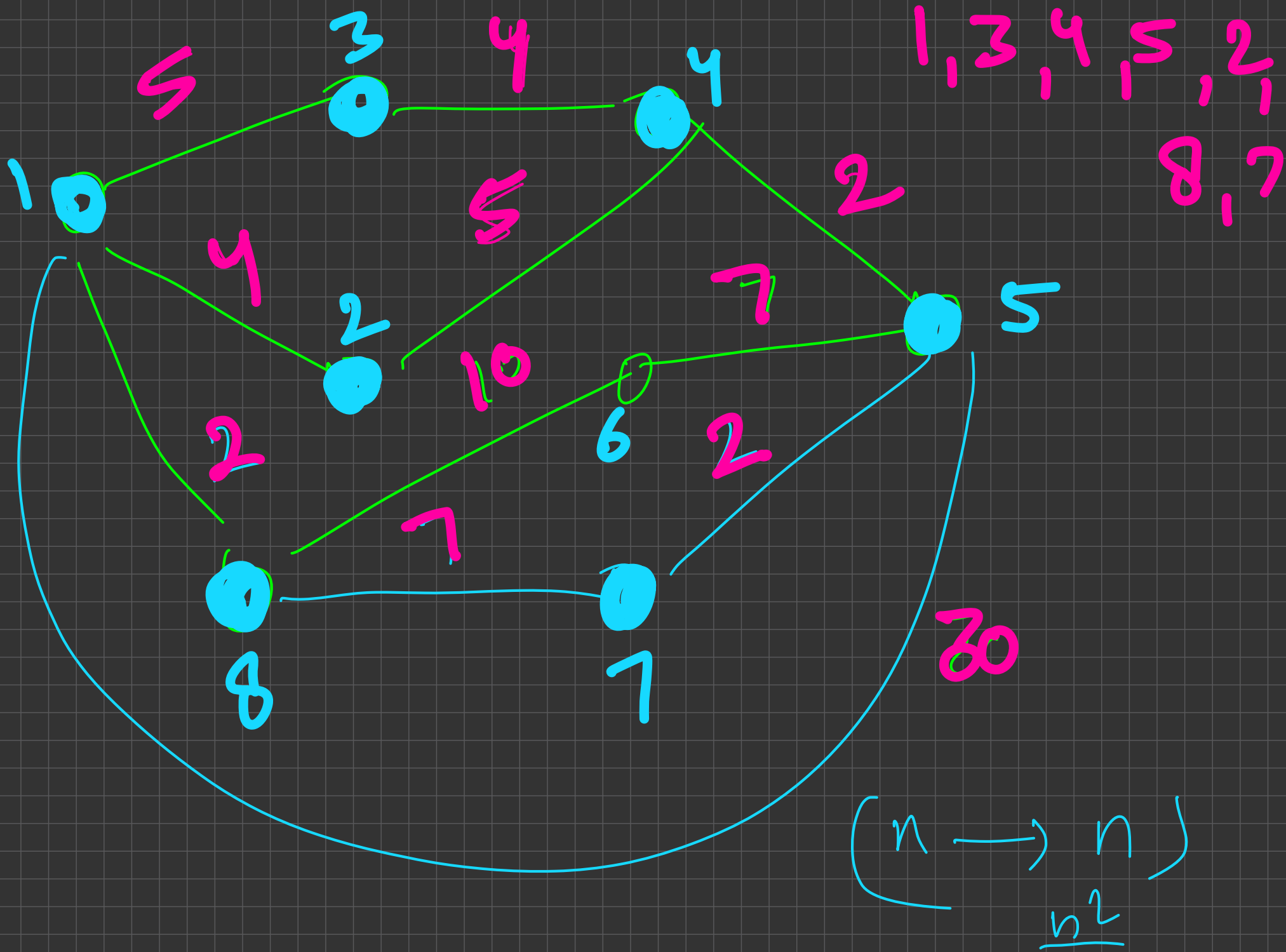
$n \log n$

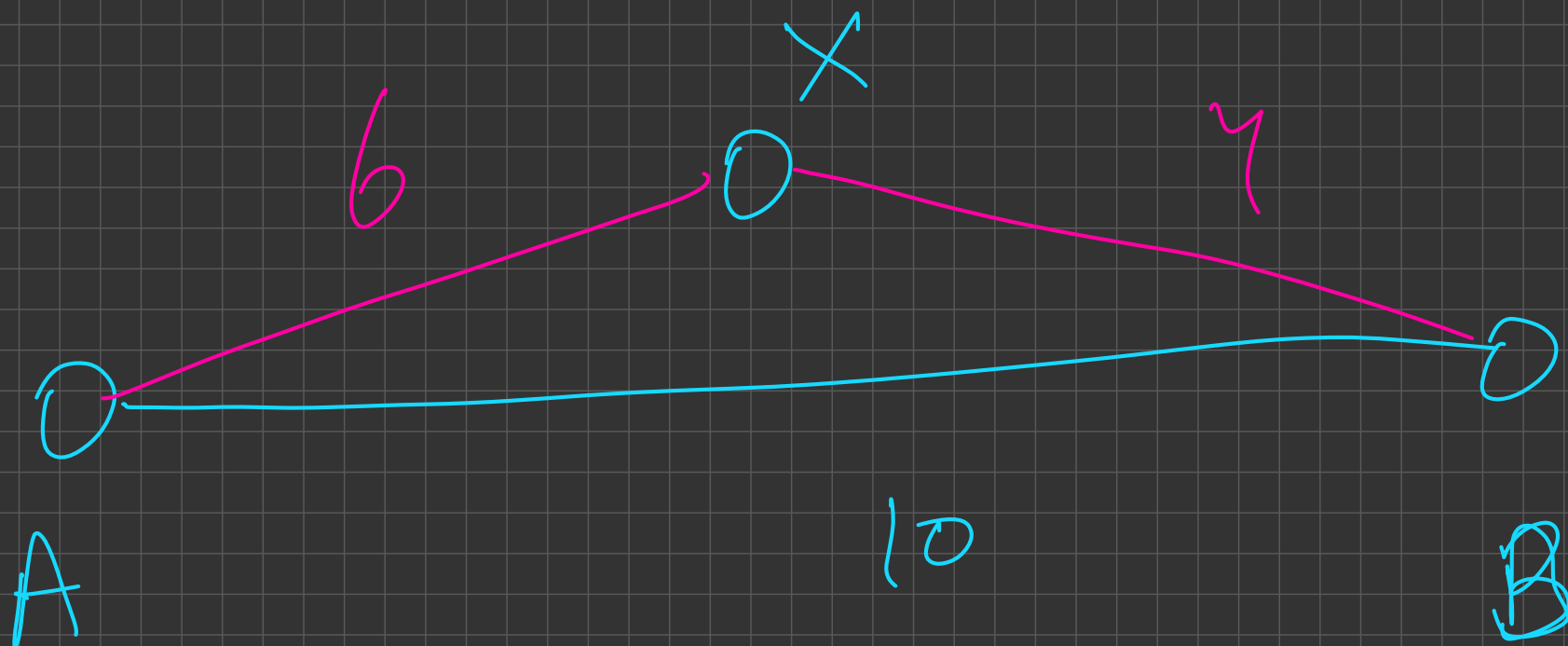
+

$m \cdot \log n$

$O((n+m) \cdot \log n)$

Given a graph with n ($n \leq 10^5$) nodes and m ($m \leq 10^5$) edges and 2 nodes A and B . Find out all the nodes which lie on any one of the shortest paths from A to B .





if $\text{dist}(A, X) + \text{dist}(B, X)$

$$= \text{dist}(A, B)$$
