Assignment-2 CS-314:Operating Systems Laboratory

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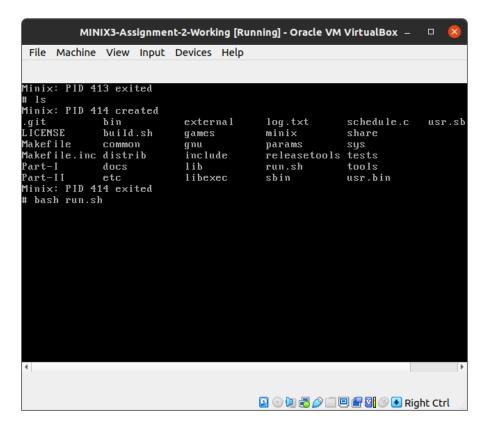
Part-1

In part 1, we have updated the *Minix-3* source code so that anytime a user-level process is brought in by the scheduler, the string "*PID <pid> swapped in*" is printed.

To accomplish this, we added the following lines of C code to the *schedule.c* file, which can be found in the */minix/minix/servers/sched* directory.

```
if(rmp->priority >= USER_Q){
    printf("(200010039, 200010041)Minix: PID %d swapped in\n",
    _ENDPOINT_P(rmp->endpoint));
}
```

Also to make the appropriate changes to the Minix machine we have made *run.sh* which will copy the modified code to appropriate location in the Minix machine and rebuild the Minix OS. The changes are reflected after a successful build and reboot.



Copied files and command to run the run.sh file

```
File Machine View Input Devices Help

Minix: PID 5165 exited
Minix: PID 97 swapped in
Minix: PID 5166 created
Minix: PID 5166 exited
Minix: PID 5166 exited
Minix: PID 5167 created
Minix: PID 5167 created
Minix: PID 5167 exited
Minix: PID 5078 exited
Minix: PID 5078 exited
Minix: PID 5078 exited
Minix: PID 5091 exited
Minix: PID 5080 exited
Minix: PID 5168 created
Minix: PID 5168 created
Minix: PID 5168 exited
Minix: PID 5168 exited
Minix: PID 5168 exited
Minix: PID 5169 exited
Minix: PID 5169 exited
Minix: PID 5169 exited
Minix: PID 5169 exited
Minix: PID 5170 created
Minix: PID 5170 created
Minix: PID 5170 exited
```

Successful build

Changes after reboot

Part-2

In this section, we have used the **UnixBench** benchmark suite to test various work-load combinations in order to understand and analyse the minix scheduler and scheduling orders. The sections that follow discuss the various workloads that are used.

a. workload_mix1.sh

The code for the workload_mix1.sh is as follows:

```
#!/bin/sh
./arithoh.sh &
./fstime.sh &
wait
```

In this workload we have used two workloads, namely *arithoh.sh* and *fstime.sh*.

After inspecting their source code and scheduling order, we discovered that *arithoh.sh* is *CPU bound* and does certain CPU bound arithmetic calculations, whereas *fstime.sh* is *I/O bound* and reads and writes some buffer data into memory.

```
File Machine View Input Devices Help

Minix: PID 57 swapped in

Minix: PID 20395 exited

4:15.68 real 18.10 user 0.25 sys

Minix: PID 20391 exited

Minix: PID 20390 exited

# Minix: PID 20390 exited

# Minix: PID 27 swapped in

Minix: PID 27 swapped in
```

arithoh.sh completed

```
File Machine View Input Devices Help

Minix: PID 118 swapped in

Minix: PID 116 swapped in

Minix: PID 116 swapped in

Minix: PID 118 swapped in

Copy done: 1000004 in 99.8167, score 2504

COUNTI2504:01KBps

TIME199.8

Minix: PID 371 exited

4:12.81 real 3.93 user 28.95 sys

Minix: PID 357 exited

fstime completed

——

Minix: PID 347 exited

Minix: PID 357 swapped in

Minix: PID 116 swapped in

Minix: PID 118 swapped in

Minix: PID 110 swapped in

Minix: PID 112 swapped in

Minix: PID 122 swapped in

Minix: PID 120 swapped in
```

fstime.sh completed

From the above figure we can see that arithoh.sh, with PID 57, being swapped in while fstime.sh, with PID 118 is waiting for I/O response. When fstime.sh becomes ready after I/O response, it is scheduled since it has relatively smaller burst time and it is completed as seen in the above figure After this the arithoh.sh is scheduled until completion as seen from figure.

b. workload_mix2.sh

The code for the workload_mix2.sh is as follows:

```
#!/bin/sh
./arithoh.sh &
./syscall.sh &
wait
```

We used two workloads in this workload: *arithoh.sh* and *syscall.sh*. We discovered that arithoh.sh and syscall.sh are both CPU constrained after inspecting their sources and scheduling order. arithoh.sh performs CPU-bound arithmetic operations, whereas syscall.sh runs a loop that calls the system.

```
MINIX3-Assignment-2-Working [Running] - Oracle VM VirtualBox
       Machine View Input Devices Help
59.23 real 1
Minix: PID 20412 exited
                                                                   0.21 sys
rithoh completed
1inix: PID 20410 exited
Minix: PID 77 swapped
                 77 swapped
77 swapped
 inix: PID
                 77 swapped
77 swapped
 inix: PID
linix: PID 77 swapped in
  inix: PID 20415 exited
1:19.83 real 1'
Minix: PID 20413 exited
syscall completed
                                     17.50 user
                                                                 43.78 sys
Minix: PID 20411 exited
Minix: PID 20409 exited
```

arithoh.sh and syscall.sh completed

From the figure we can see that syscall.sh, with PID 77 being swapped one after the other depending upon their priority. Here arithoh.sh is completed before syscall.sh

c. workload_mix3.sh

The code for the workload mix3.sh is as follows:

```
#!/bin/sh
./arithoh.sh &
```

```
./spawn.sh & Wait
```

we utilised two workloads in this workload: *arithoh.sh* and *spawn.sh*. After inspecting their source code and scheduling order, we discovered that arithoh.sh and spawn.sh are both CPU constrained. arithoh.sh performs CPU-bound arithmetic operations, whereas spawn.sh creates (forks) new processes that exit immediately.

spawn.sh completed

```
MINIX3-Assignment-2-Working [Running] - Oracle VM VirtualBox
 File Machine View Input Devices Help
Minix: PID 202 swapped in
                                   i n
Minix: PID 202 swapped
 1inix: PID 202 swapped
Minix: PID 202 swapped
 linix: PID 202 swapped
 1inix: PID 150 swapped
Minix: PID 202 swapped
Minix: PID 202 swapped
Minix: PID 202 swapped in
Minix: PID 202 swapped in
Minix: PID 202 swapped in
 linix: PID 227 exited
48.21 real
Minix: PID 225 exited
                                 18.33 user
                                                           0.16 sys
 rithoh completed
Minix: PID 223 exited
Minix: PID 222 exited
 Minix: PID 27 swapped in
Minix: PID 27 swapped in
```

arithoh.sh completed

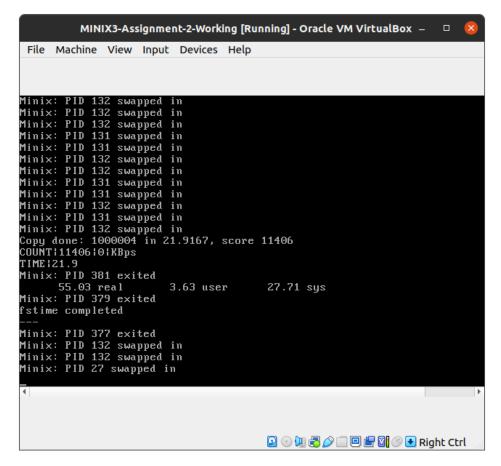
We can see from the above image that spawn.sh is constantly creating new processes and scheduling them to reduce response time. According to the source code, these newly generated processes quit instantly. After a while, spawn.sh and the newly generated process will exit. The arithoh.sh with PID 202 is then scheduled till completion, as seen in the figure.

d. workload mix4.sh

The code for the workload_mix4.sh is as follows:

```
#!/bin/sh
./fstime.sh &
./syscall.sh &
wait
```

In this workload we have used two workloads, namely *fstime.sh* and *syscall.sh*. After looking at their source and the order of scheduling, we came to know that fstime.sh is I/O bound and syscall.sh is CPU bound.



fstime.sh completed

```
MINIX3-Assignment-2-Working [Running] - Oracle VM VirtualBox -
 File Machine View Input Devices Help
       PID 132 swapped
Minix: PID 132 swapped in
Minix: PID 132 swapped in
Minix: PID 132 swapped
Minix: PID 132 swapped
Minix: PID 132 swapped
                         in
1inix: PID 132 swapped
Minix: PID 132 swapped
Minix: PID 132 swapped
linix: PID 132
                swapped
Minix: PID 132 swapped
Minix: PID 132 swapped
Minix: PID 132 swapped
                         i n
1inix: PID 132 swapped
Minix: PID 132 swapped
Minix: PID 132 swapped in
Minix: PID 132 swapped in
Minix: PID 382 exited
                        19.63 user
    1:34.21 real
                                          43.21 sys
Minix: PID 380 exited
syscall completed
Minix: PID 378 exited
Minix: PID 376 exited
```

syscall.sh completed

From the above figure, we can see that fstime.sh, with PID 131 is waiting for I/O response, while syscall.sh, with PID 132 is swapped in. So once fstime.sh is swapped in, it runs until completion. After this, syscall.sh is swapped in and runs until completion.

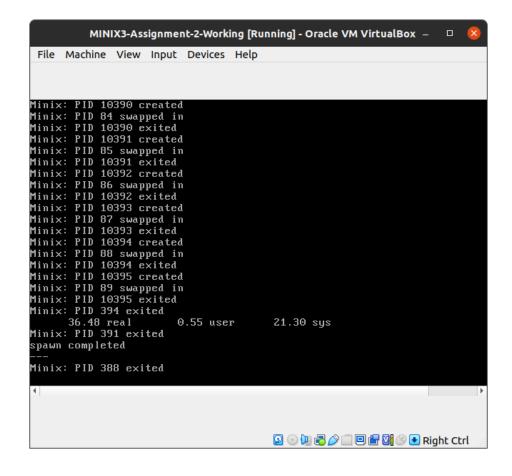
e. workload_mix5.sh

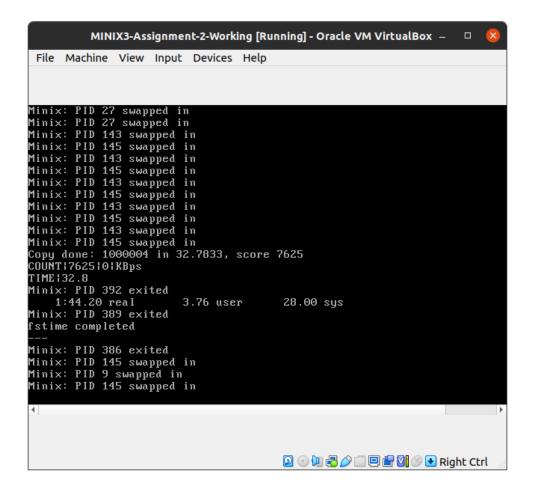
The code for the workload_mix5.sh is as follows:

```
#!/bin/sh
./fstime.sh &
./pipe.sh &
./spawn.sh &
Wait
```

In this workload we have used a three workloads, namely fstime.sh, pipe.sh and spawn.sh. After looking at their source and the order of

scheduling, I came to know that both fstime.sh is I/O bound. The pipe.sh does some I/O bound, spawn.sh is CPU bound.





```
MINIX3-Assignment-2-Working [Running] - Oracle VM VirtualBox -
 File Machine View Input Devices Help
Minix: PID 392 exited
1:44.20 real
Minix: PID 389 exited
                                      3.76 user
                                                              28.00 sys
fstime completed
Minix: PID 386 exited
Minix: PID 145 swapped in
Minix: PID 145 swapped in Minix: PID 9 swapped in Minix: PID 145 swapped in Minix: PID 9 swapped in Minix: PID 145 swapped in Minix: PID 9 swapped in Minix: PID 145 swapped in Minix: PID 393 exited
2:05.71 real
Minix: PID 390 exited
                                     7.43 user
                                                              56.51 sys
pipe completed
Minix: PID 387 exited
Minix: PID 385 exited
# Minix: PID 27 swapped in
Minix: PID 28 swapped in
Minix: PID 27 swapped in
Minix: PID 28 swapped in
```

spawn.sh with PID 89 is swapped with fstime.sh with PID 143. The fstime.sh is swapped with pipe.sh having PID 145.