

**Due 23:59 Nov 6 (Sunday).** There are 100 points in this assignment.

Submit your answers (**must be typed**) in pdf file to CourSys

<https://coursys.sfu.ca/2022fa-cmpt-705-x1/>. Submissions received after 23:59 of Nov 6 will get penalty of reducing points: 20 and 50 points deductions for submissions received at  $[00 : 00, 00 : 10]$  and  $(00 : 10, 00 : 30]$  of Nov 6, respectively; no points will be given to submissions after  $00 : 30$  of Nov 6.

1. (Chapter 7 Problem 11 of the text book) 20 points

The Forward-Edge-Only algorithm computes a flow in a flow networks as follows: it searches for  $s - t$  paths in a graph  $\tilde{G}_f$  consisting only of arcs  $e$  for which  $f(e) < c_e$ , and it terminates when there is no augmenting path consisting entirely of such arcs (the algorithm may choose a forward arc path arbitrarily, provided it terminates only when there are no forward arc paths). A claim for the Forward-Edge-Only algorithm is that there is a constant  $k > 1$  (independent of the particular input flow network), so that on every instance of the Maximum Flow problem, the Forward-Edge-Only algorithm is guaranteed to find a flow of value at least  $1/k$  times the maximum flow value (regardless of how it chooses its forward edge paths). Decide whether you think this claim is true or false, and give a proof of either the claim or its negation.

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**Answer**

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## 2. (Chapter 7 Problem 17 of the text book) 20 points

There is a communication network of  $n$  nodes carrying data from a source node  $s$  to a destination node  $t$ . There are  $k$  link-disjoint paths from  $s$  to  $t$  in the network when all links function formally. Now an attacker disabled a minimum number of links that  $t$  is not reachable from  $s$ . You are sitting at node  $s$  and using **ping** command to find the disabled links (command **ping**( $v$ ) tells whether node  $v$  is reachable or not from  $s$ ). Assume that you have the topology information of the network and you can decide which node to **ping** based on the previous results of the **ping** commands. Give an algorithm which uses  $O(k \log n)$  **ping** commands to find the disabled links.

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**Answer**

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3. (Chapter 7 Problem 21 of the text book) 20 points

A test instance for a WiFi network consists of  $n$  laptops  $\{c_1, \dots, c_n\}$  and  $n$  access points  $\{p_1, \dots, p_n\}$ . A test set  $T$  for the test instance is a set of pairs  $(c_i, p_j)$  with the following properties:

- (i) if  $(c_i, p_j)$  in  $T$  then laptop  $c_i$  is within the access range of access point  $p_j$ ,
  - (ii) each laptop appears in at least one pair of  $T$ , and
  - (iii) each access point appears in at least one pair of  $T$ .
- (a) Give an example of a test instance for which there is no test set of size  $n$ .
- (b) Give a polynomial time algorithm which, given a test instance and integer  $k$ , decides whether there is a test set of size at most  $k$ .

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**Answer**

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## 4. (Chapter 7 Problem 31 of the text book) 20 points

There are  $n$  boxes  $1, \dots, n$ , each box  $i$  has size  $s(i)$ . A box  $i$  can be put inside box  $j$  if  $s(i) < s(j)$ . For any two boxes  $i$  and  $i'$  with  $s(i) < s(j)$  and  $s(i') < s(j)$ , both  $i$  and  $i'$  can not be put inside  $j$  if  $i$  is not put inside  $i'$  or  $i'$  is not put inside  $i$ . But for a sequence of boxes  $i_1, i_2, \dots, i_k$  with  $s(i_1) < s(i_2) < \dots < s(i_k)$ ,  $i_1$  can be put inside  $i_2$ , then  $i_2$  put inside  $i_3, \dots$ , and finally  $i_{k-1}$  put inside  $i_k$ . In this case, all boxes  $i_1, \dots, i_{k-1}$  are inside box  $i_k$  and only  $i_k$  is visible. The nesting arrangement for a set of  $n$  boxes to put one box inside another such that the number of visible boxes is minimized. Give a polynomial time algorithm which, given a set of  $n$  boxes  $1, \dots, n$  and  $s(1), \dots, s(n)$ , solves the nesting arrangement problem.

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**Answer**

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5. 20 points

Figure 1 gives a flow network  $G$  and a function  $f : E(G) \rightarrow R^+$ . The capacity of each arc appears as a label next to the arc, the value assigned to each arc by  $f$  is in the box next to the arc.

- (a) Is  $f$  a flow on  $G$ ? If yes, why? and give the residual graph  $G_f$  w.r.t.  $f$ . If no, why?  
(b) Implement Ford-Fulkerson Algorithm for the maximum flow problem and show the result of your implementation for  $G$ .

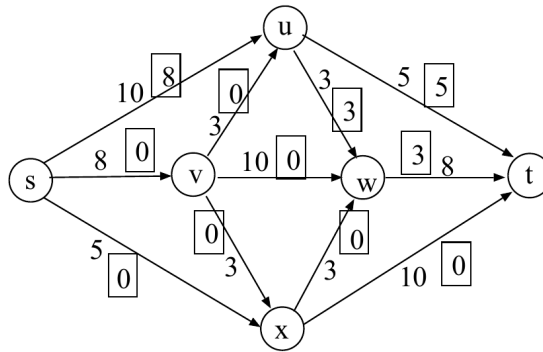


Figure 1: Figure for question 5.

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**Answer**

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