

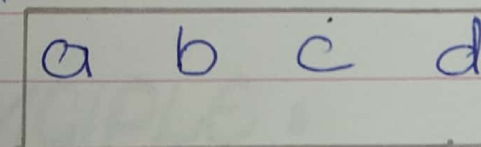
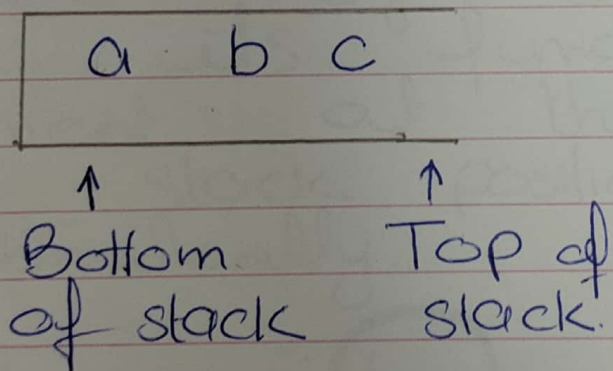
# STACKS:

**What is stack:** A stack is an ordered list with the restrictions that elements are added or deleted from only one end of the list termed as "top of the stack." The other end of the list which lies within it is "inactive" and called "bottom of the stack."

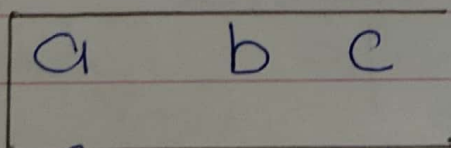
**EXPLANATION:-** IF S is a stack with three elements in it a, b, c where c occupies the top of the stack position, and if d were to be added, the resultant stack contents will

a, b, c, d, where d occupies top of the stack position and for deleting operation, it would throw out the element at the top of the stack automatically. As the Fig illustration given below shows that.

Ordered list of element



Addition of element d into the stack



Removal / deletion of element d from stack.



The point should be kept in mind while dealing with this phenomenon is:

During the insertion of element into stack, it is essential that their identities are specified, whereas in the case of deletion this is not necessary since by the virtue of its functionality, the element at the top of the stack position is deleted automatically.

## LIFO PRINCIPLE:

The stack data structure obeys the principle of LIFO which is the abbreviation of Last in first out. You can say element inserted into the stack

join last and that  
joined last are the  
first to be removed

## STACK OPERATION:

- i PUSH: Insertion of elements ✓
- ii POP: Deletion of elements. ✓

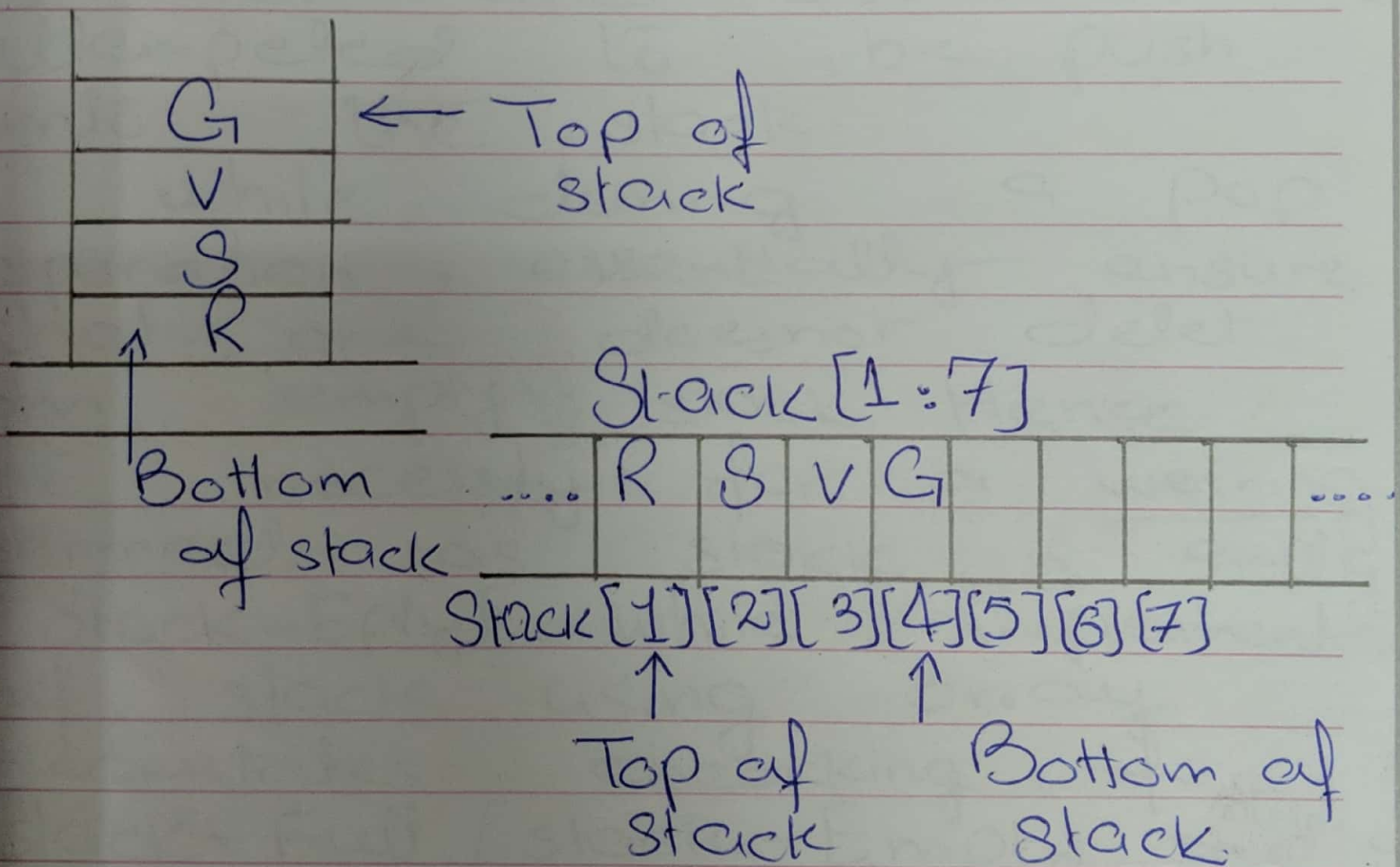
## IMPLEMENTATION:

A common basic method of implementation is to make use of the other fundamental data structure "Arrays". while arrays are sequential data structures the other of employing linked data structures of have been successfully attempted and applied.

The fig shows an array base implementation of stack. Considering the fact that stacks are uni-



dimensionally ordered list and so arrays which despite their multi-dimensional structure are inherently associated with a one-dimensional set of memory location



This fig shows a stack of 4 elements R, S, V, G represented by an array stack[1:7]. Usually, if stack is represented

as array stack[1:n] then  $n$  elements can be store in it no more then  $n$  could be. So there is essential to issue a signal termed as STACK-FULL when element whose number is over and above  $n$  are attempted to be push into the stack.

while during a pop operation, essentially ensure that one doesnot delet an empty stack. Hence the necessity for a warning termed as stack is empty "Stack-Empty". while implementation of stack using array necessitates checking for stack-Full / stack-Empty condition during push/pop operations respectively, the implementation of stack with linked data structure dispenses with these testing conditions



# PUSH IMPLEMENTATION:

Let  $\text{stack}[1:N]$  an array implementation of stack and  $\text{top}$  be a variable recording the current top of stack position,  $\text{top}$  is initialize to 0.  $\text{item}$  is element to be added.  $n$  is maximum capacity.

## Algorithm:

Push(STACK,  $n$ ,  $\text{top}$ ,  $\text{item}$ )  
if  $(\text{top} = n)$  then STACK-FULL  
else

$\text{top} = \text{top} + 1$   
     $\text{stack}[\text{top}] = \text{item}$ .

end PUSH

# POP IMPLEMENTATION:

In this case, as we know, no element identity need be specified since by default element is deleted. Here item is used as an output variable which stores a copy of element removed.

## Algorithm:

```
Pop(stack, top, item)
if (top = 0) then STACK EMPTY;
else { item = STACK[TOP];
      top = top - 1;
    }
```

end top



# STACK APPLICATIONS:

1. Recursive Programming
2. Evaluation of Expression

## RECURSIVE PROGRAMING:

If a function call it self more than one time in its own body then this process is called recursion. and this function is called recursive function.

Consider the recursive pseudo-code for factorial computation as shown below,

```
FACTORIAL(n)
if (n=0) then Factorial = 1;
else { x = n - 1;
      y = FACTORIAL(x);
      FACTORIAL = n * y;
    }
end FACTORIAL
```



**TAIL RECURSION:** A special case of recursion where a recursive call to the function turns out to be the last action in calling function. But notice the recursive call needs to be the last executed statement in the function not necessarily the last statement.

## 2: EVALUATING EXPRESSION:

### INFIX, PREFIX and POST FIX

The evaluation of expression is an important feature of computer design. When we write or understand an arithmetic expression —  $(A+B) \uparrow C * D + E$ , we do so by following the scheme of  $(operand) \langle operator \rangle (operand)$  (i.e.) an OPERATOR



is procedured and succeed by an  $\langle \text{operand} \rangle$ . Such an expression is termed INFIX EXPRESSION. It is already known how infix expression is used in programming. It is accorded rule of Hierarchy, Precedence and associativity to ensure that the computer does not misinterpret the expression but compute its value in a unique way.

Actually a computer reworks on infix expression to produce an equivalent expression which follow the scheme of  $\langle \text{operand} \rangle \langle \text{operand} \rangle \langle \text{operator} \rangle$  known as postfix.

A third one category is one which follows the scheme of  $\langle \text{operator} \rangle \langle \text{operand} \rangle \langle \text{operand} \rangle$ . and it is known as Prefix expression.

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```
import java.util.*;

class stack

{
    int stk[] = new int [4];
    int top, item;

    Scanner in = new Scanner(System.in);

    stack()
    {
        top = -1;
    }

    void push ()
    {
        if (top == 3)
            System.out.println("Stack is full");
        else
        {
            top++;
            System.out.print("Enter the number = ");
            item = in.nextInt();
            stk[top] = item;
        }
    }

    void pop ()
    {
        if (top == -1)
            System.out.println("Stack is empty");
        else
```

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Stack

```
{ item = stk[top];
    top--;
    System.out.println("item = " + item);
}}

class apstack
{
    public static void main(String arg [])
    {
        int ch;
        stack s = new stack();
        Scanner in = new Scanner(System.in);
        do
        {
            System.out.print("Enter the 1 to push and 2 to pop = ");
            ch = in.nextInt();
            switch (ch)
            {
                case 1:
                    s.push();
                    break;
                case 2:
                    s.pop();
                    break;
                default:
                    System.out.print("wrong");
            }
        }
        while (ch != 0);
    }
}
```

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```
System.out.print("Enter  
to continue = ");
ch = in.nextInt();
while (ch != 0);
}
```

the 0 to exit & any No