COMP353 Databases

More on SQL Queries

SQL Queries: Review

SQL query has a form

```
SELECT ....
FROM ....
WHERE ....;
```

- The SELECT clause says which Attribute(s) we are interested in
- The **FROM** clause says which **Relation**(s) we refer to
- The WHERE clause says which Tuple(s) we refer to

Case Insensitivity

- SQL is case insensitive
- So, keyword FROM maybe written as:
 - FROM or
 - From or
 - FrOm
- Only in strings, SQL distinguishes between uppercase and the lowercase letters
 - So, the following are different strings:
 - 'FROM'
 - 'From'
 - 'FrOm'

Select Clause

- In place of * in the SELECT clause, we can put any attribute "we wish" to project on
- In the SELECT clause, we can also do renaming

SELECT title **AS** name, length **AS** duration **FROM** Movie;

→ the structure of the query output "appears as": (name, duration)

Select Clause

We can also use a formula in place of an attribute

SELECT title **AS** name, length/60 **AS** lengthInHours **FROM** Movie;

→ the structure of the output: (name, lengthInHours)

Select Clause

 SQL even allows using a constant as an item in the SELECT clause, as shown below.

SELECT title **AS** name, length/60 **AS** length, 'hrs.' **AS** inHours **FROM** Movie;

```
→ The structure of the output: (name, Gone with the wind 1.98 hrs.
King Kong 0.75 hrs.
```

Why? To put some "useful" words into the output that SQL displays

- Two strings are equal if they have the same sequence of characters/symbols
- Strings are compared alphabetically
 - 'fodder' < 'foo'</p>
 - 'bar' < 'bargain'</p>
- WHERE R.A = T.B AND s LIKE p
 - **s** is an attribute of type *string* and **p** is a pattern; e.g. WHERE title LIKE 'Gone%'
 - "Ordinary" characters in p matches ordinary characters in s
 - What about "Special" characters in p: %, _
 - "%" in p matches any sequence of zero or more characters in s
 - "_" in p matches any one character in s

 Suppose we remember a movie "Star something", and we do remember that "the something" has four letters

```
FROM Movie
WHERE title LIKE 'Star _ _ _ ';

SELECT title
FROM Movie
WHERE title LIKE 'Star%';
```

What if the pattern p includes ', %, or _?

■ Find all movies with a possessive ('s) in their title LIKE '%'s%'

SELECT title

FROM Movie

WHERE title LIKE '%"s%';

The convention is that two apostrophes "in a string represent one single apostrophe ("), and not the end of string

- What if **p** involves the *special characters* % or _?
 - We should "escape" their special meaning using "some" escape character
 - **SQL** allows using *any* character as escape character
- s LIKE 'x%%x%' ESCAPE 'x';
 - Here, X is the escape character → x% means the character %, and not its usual meaning (the special character)
 - The pattern 'x%%x%' matches strings: %whatever%

Ordering the Ouput

- We may wish the output of a query to be displayed in some order. This could be done using the SQL clause:
 - ORDER BY < list of attributes >
- E.g., List Disney movies in 1990 by their length, shortest first, and then by the alphabetical order of the titles:

SELECT*

FROM Movie

WHERE studioName = 'Disney' AND year = 1990

ORDER BY length, title;

- Default ordering is ASCending, unless we use the **DESC** keyword
- Ties are broken by the "next" attribute in the ORDER BY list.

Products and Joins

- SQL has a simple way to couple relations in a query
 - How? Simply list each relation in the FROM clause
- All the relations in the FROM clause are coupled through Cartesian product
- Then we can put conditions in the WHERE clause in order to get a desired kind of join

Join (Example, Recall)

Relation schemas:

```
Movie (<u>title</u>, <u>year</u>, <u>length</u>, filmType)
Owns (<u>title</u>, <u>year</u>, <u>studioName</u>)
```

Query:

Find titles and lengths of all movies produced by Disney

Query in SQL:

SELECT Movie.title, Movie.length

FROM Movie, Owns

WHERE Movie.title = Owns.title **AND** Movie.year = Owns.year **AND** Owns.studioName = 'Disney';

- We can apply the common set operations of union, intersection, and difference to relations R and S, if they are compatible.
- When the output of two or more SQL queries are compatible, we may "combine" the queries using:
 - UNION
 - INTERSECT
 - EXCEPT (or MINUS in Oracle)

Relation schemas:

```
Movie (<u>title</u>, <u>year</u>, length, filmType)
StarsIn (<u>title</u>, <u>year</u>, <u>starName</u>)
```

Query:

Find titles and years of movies that appeared in either **Movie** or **StarsIn** relations

Query in SQL:

SELECT title, year

FROM Movie

UNION

SELECT title, year

FROM StarsIn;

Relation schemas:

Star(name, address, gender, birthdate) **Exec**(name, address, cert#, netWorth)

Query:

Find names and addresses of all female movie stars who are also movie executives with a net worth of over \$10,000,000

Query in SQL:

SELECT name, address

FROM Star

WHERE gender = 'F'

INTERSECT

SELECT name, address

FROM Exec

WHERE netWorth > 10000000;

Relation schemas:

Star (name, address, gender, birthdate) **Exec** (name, address, cert#, netWorth)

Query:

Find names and addresses of movie stars who are **not** movie executives

Query in SQL:

SELECT name, address

FROM Star

EXCEPT

//or MINUS in Oracle//

SELECT name, address

FROM Exec;

Duplicate Elimination

Note that in SQL:

- The union, intersection, and difference operations normally eliminate duplicates (the set semantics)
- To retain duplicates, hence preventing duplicate elimination, we must use the keyword ALL after the operator UNION, INTERSECT, and EXCEPT
 - R UNION ALL S (the only bag operation supported in Oracle)
 - R INTERSECT ALL S
 - R EXCEPT ALL S

Retaining Duplicates

R UNION ALL S

The bag of elements that are in \mathbb{R} , \mathbb{S} , or in both. If \mathbb{R} is a bag in which tuple t appears n times, and \mathbb{S} is a bag in which t appears m times, then the number of occurrences of tuple t in bag $\mathbb{R} \cup \mathbb{S}$ is n + m

R INTERSECT ALL S

■ The bag of elements that are in both \mathbf{R} and \mathbf{S} . If \mathbf{R} is a bag in which tuple t appears \mathbf{n} times, and \mathbf{S} is a bag in which t appears \mathbf{m} times, then the number of occurrences of t in bag $\mathbf{R} \cap \mathbf{S}$ is min(n,m)

R EXCEPT ALL S

■ The bag of elements that are in R but not in S. If R is a bag in which tuple t appears n times, and S is a bag in which t appears m times, then the number of occurrences of t in bag R - S is max(0, n - m)

Retaining Duplicates in Union

Relation schemas:

```
Movie (<u>title</u>, <u>year</u>, length, filmType)
StarsIn (<u>title</u>, <u>year</u>, <u>starName</u>)
```

Query:

List the title and year of every movie that appears in **Movie** or **StarsIn**

Query in SQL:

SELECT title, year

FROM Movie

UNION ALL

SELECT title, year

FROM StarsIn;