MCQ on DBMS

	MCQ on DBMS			
1	Acknowledge that	Tables(Y)		
		Fiolds(NI)		
		Fields(N)		
		Records(N)		
		` ,		
		Keys(N)		
2	Determine the phr	Attribute(N)		
<u> </u>	Determine the pin	710010000		
		Tuple(Y)		
		Field(N)		
		rieiu(iv)		
		Instance(N)		
3	Determine the ER	Entity Row(N)		
		,		
		Entity Relatio	nship(Y)	
		Fuelita - Demand	- /NI\	
		Entity Renam	e(N)	
		Entity Relatio	n(N)	
 			(.•/	
4	Which of the follow	Name(N)		
	1	· · /		
	+			
		Street(N)		
	+	15(1)		
	ļ	ID(Y)		

				1	
		Department(N)		
5	Which of the follow	Object of Pol	ation/NI)		
)	Which of the follow	Object of Kei	ation(N)		
		Model of Rela	ation(N)		
		Present Work	ring Model	(NI)	
		Treserie Work	ang woder		
		Thing in Real	World(Y)		
6	Determine the lang	Query(Y)			
		-			
		Assembly(N)			
		Structural(N)			
		, ,			
		Compiler/NI			
		Compiler(N)			
7	Which of the follow	Phone Numb	er(Y)		
		Name(N)			
		runc(ru)			
		Date- of- Birt	h(N)		
		All of these(N)		
_	NAME OF THE STREET	212(21)			
8	Which of the follov	NA(N)			
		Zero(N)			

	1	
		NULL(Y)
		INOLL(1)
		Blank Space(N)
9	Determine that the	Domain(Y)
		2 1 11 (21)
		Relation(N)
		Set(N)
		JCC(14)
		Schema(N)
10	List the descriptive	Entity(N)
		A
		Attribute(Y)
		Relation(N)
		The latter (14)
		Model(N)
11	Determine the rela	Natural Join(N)
		Assignment(N)
		Assignment(N)
L		
		Set Intersection(N)
		'
L		None of these(Y)

12	Determine which jo	Loft Outor lo	in/NI)	
12	Determine which jo	Left Outer Jo	11(IN)	
		Dight Outer L	oin/NI)	
		Right Outer J	oin(iv)	
		1.2.2.2.1.2.1.2.1.2.1.2.1.2.1.2.1.2.1.2		
		Inner Join(Y)		
		Not called a		
		Natural Join(I	N)	
4.2	Data and a little	0.00		
13	Determine which o	&(Y)		
		0//NI)		
		%(N)		
		11(0.1)		
		(N)		
		_		
		None of these	e(N)	
14	Select the option t	Function(N)		
		View(Y)		
		Procedure(N)		
<u> </u>				
<u> </u>		None of these	e(N)	
<u> </u>				
15	Choose the option	Union(N)		
		Set Differenc	e(Y)	
		Projection(N)		

			1	1	T
		Intersection(N)		
		microconon			1
16	Which statement of	Select * from	emp wher	e emp-id=	102;(N)
					<u> </u>
		Select emp-io	d from emp	where em	np-id=106;(N
		Select emp-io	d from emp	;(N)	
		Select emp-io	d where em	p-id=109 a	and firstnam
17	NA/leatiatlea assissa	11: (81)			1
17	What is the equiva	Union(N)			
					1
		Intersection(N.)		
		microconon			
		Set Differenc	:e(N)		1
]		
		Cartesian Pro	oduct(Y)		
18	What is the main o	Data Storage	(N)		
		Data Retrieva	al(Y)		
		D . A4	(21)		
		Data Modelli	ng(N) I		
					1
		Data Security	/NI)		1
		Data Security	/(IN) 		
19	What is an example	Join(N)			1
	Time is an example				1
		Project(Y)			1
	<u>I</u>	1: / 5/556(1/	<u> </u>	<u> </u>	1

		<u> </u>	I			Ì
		Union(N)				
		Onion(N)				
		Intersection(N)			
20	Which statement a	It combines a	all possible	combinatio	ns of tuple	s(N)
		It includes or	ly matchin	g tuples fro	m both rela	ations(N)
			<u> </u>	Ĭ		` ´
		It excludes m	l Satching tur	les from h	th relation	l vs(NI)
		it excludes in	latering tup			
-						
		It is based on	the equali	ty of values	in specifie	d columns(Y)
21	What is an example	Project(N)				
		Select(N)				
		Union(N)				
		Join/V)				
		Join(Y)				
22	In which normal fo	First(Y)				
		Second(N)				
		Third(N)				
		· · ·				
		Fourth(N)				
		i ourth(N)				
<u> </u>						
		<u> </u>	<u> </u>	<u> </u>	<u> </u>	

23	Mhat is a databass	Functional Do	n and an au	NI)	
23	What is a database	Functional De	ерепаенсу	IN)	
		Normalization	2/1/1		
		Normalization	1(Y)		
_		Databasa Ma	مامالنه مرا۱۱		
		Database Mo	aeiiing(N)		
		D	- (81)		
		Decomposition	on(N)		
24	L. Litabara and Co	Et al Nia and I	/4 NJE\/	(8.1)	
24	In which normal fo	First Normai	-orm(INF)(N)	
			15 (2)	15)(51)	
_		Second Norm	al Form(2N	IF)(N)	
			- (2)	(2.1)	
		Third Normal	Form(3NF))(N)	
		Boyce codd N	lormal Forr	n(BCNF)(Y)	
25	What are the true	BCNF is strict	er than 3NI	F(N)	
		3NF removes	transitive I	Dependency	/(N)
		2NF removes	partial dep	endency(N)	
		All of these(Y)		
26	What role does Ca	Define data t	ypes(N)		
		Specify Data	constraints	(N)	
		Describe Rela	tionship Co	onstraints(N)
			-		

						1
		The number	of occuran	cos(V)		
		The number	l occurant	les(1)		
27	Which normal forn	1NF(N)				
		2NF(N)				
		,				
		3NF(Y)				
		3141 (1)				
		4NF(N)				
28	What is a drawbac	Difficult to im	nplement(N	١)		
		Limited supp	ort for com	nplex relation	nship(Y)	
		High redunda	l ncv(NI)			
		Tilgit reduita				
		D			- (01)	
		Poor perform	nane in iarg	ge database	S(N)	
29	What is the purpos	It is used to e	stablish re	lationships	between ta	bles
		It is a primary	key of and	other table(N)	
		It is used for	indexing(N)		
		16.15 4564 16.1	110000000000000000000000000000000000000	,		
		IA ::	: -l +:£: £		/NI)	
		It is a unique	identifier f	or a record	(N)	
30	Among the options	Hierarchical r	model(N)			

Logical model(Y) Physical model(N) Intersection(N) Set difference(N) Cross product(N) Cross product(N) The projection of tuples t based on predicate P(Y) The Cartesian product of tuples t and predicate The Cartesian product of tuples t and predicate
Physical model(N) Physical model(N) Intersection(N) Set difference(N) Cross product(N) In Relational Algeb The set of all tuples t that satisfy predicate P(Y) The projection of tuples t based on predicate P(X) The Cartesian product of tuples t and predicate
Physical model(N) Intersection(N) Set difference(N) Cross product(N) In Relational Algeb The set of all tuples t that satisfy predicate P(Y) The projection of tuples t based on predicate P The Cartesian product of tuples t and predicate
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The projection of tuples t based on predicate P The Cartesian product of tuples t and predicate
The Cartesian product of tuples t and predicate
The Cartesian product of tuples t and predicate
The Cartesian product of tuples t and predicate
The union of tuples t and predicate P(N)
3 What does the nat All combinations of tuples from both relations(
yvinat ages the hadrin combinations of tuples from both relations
What does the hat Air combinations of tuples from both relations
Tuples with matching attribute values(Y)
Tuples with matching attribute values(Y)
Tuples with matching attribute values(Y)
Tuples with matching attribute values(Y)
Tuples with matching attribute values(Y) Tuples with non-matching attribute values(N)

What role do unive	To specify co	nditions fo	some tupl	es(N)
	T:£		!! + ! /	
	To specify co	nditions foi	r all tuples(Y)
	To exclude ce	ertain tuple	s(N)	
	To rename at	tributes(N)		
What is a true state	Δ relationshir	hetween:	 tahles(N)	
Wilde is a crue state	ATCIACIONSIN	between	tabics(iv)	
	The reliability	of one pie	ce of data	on another(
	Data stored in	n multiple l	ocations(N)
	The size of th	 e datahase	[
	1110 3120 01 111			
What is a true state	A key with m	ultiple attri	butes(Y)	
	A Primary Ke	y(N) I		
	A foreign key	(N)		
	r roreign key	(,		
	A key withou	t depender	ncies(N)	
Which one among	Reflexivity Ru	ıle(N)		
	Trancitivity R	ule(NI)		
	ι ι αποιτίνιτη Κ	uic(N)		
	Pseudo-trans	itivity Rule	(Y)	
	What is a true state What is a true state	To specify co To exclude ce To rename at The reliability Data stored i The size of th What is a true state A key with m A Primary Key A foreign key A key withou Which one among Reflexivity Ru	To specify conditions for To exclude certain tuple To rename attributes(N) To rename attributes(N) The reliability of one pies Data stored in multiple I The size of the database The size of the database A Primary Key(N) A foreign key(N) A foreign key(N) A key without depender Which one among Reflexivity Rule(N) Transitivity Rule(N)	A foreign key(N) A key without dependencies(N) Which one among Reflexivity Rule(N)

Augmentation Rule(N) 38 Which normal form 1NF(Y) 2NF(N) 3NF(N) BCNF(N) BCNF(N) Left Join(Y) Full Outer Join(N) Right Join(N) Right Join(N) If they have same execution time(N) If they have same number of tables joined(N) If they have same number of rows in the WHERE clause(N) If they have same number of rows in the WHERE clause(N) 41 What is the purpos Query Optimization(Y)
38 Which normal forn 1NF(Y) 2NF(N) 3NF(N) BCNF(N) BCNF(N) Left Join(Y) Full Outer Join(N) Right Join(N) Right Join(N) If they have same result set(Y) If they have same number of tables joined(N) If they have same number of rows in the WHERE clause(N) If they have same number of rows in the WHERE clause(N) If they have same number of rows in the WHERE clause(N) 41 What is the purpos Query Optimization(Y)
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39 What kind of join c Inner Join(N) Left Join(Y) Full Outer Join(N) Right Join(N) When are two SQL If they produce same result set(Y) If they have same execution time(N) If they have same number of tables joined(N) If they have same number of rows in the WHERE clause(N) If they have same number of rows in the WHERE clause(N) What is the purpos Query Optimization(Y)
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41 What is the purpos Query Optimization(Y)
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41 What is the purpos Query Optimization(Y)
Normalization(N)

	<u>r</u>			r	_
		Dependency	Inference(I	V)	
		2 0 0 11 0 11 0 1		·,	
		Indexing(N)			
42	What concept is im	No data loss (during und:	ates(N)	
-	Windt concept is iii	140 data 1033			
		No data loss	during joins	s(Y)	
		No data loss	during dala	tions(NI)	
		NO data 1033 (uuring uele	Lions(iv)	
		No data loss	during inse	rtions(N)	
42	Dunida an arranal	CELECT * EDG	N 1 + - - - 1 / N		
43	Provide an example	SELECT * FRC	ivi tableT(i	N)	
		SELECT colum	nn1 FROM 1	table1(N)	
		CELECT	2 50014	 - -2/NI\	
		SELECT colum	INZ FROM I	table2(N)	
		All of these(Y)		
		,	<u>, </u>		
44	Which normal forn	ZNF(N)			
		3NF(N)			
		BCNF(Y)			
		4NF(N)			
	+	INI (IN <i>)</i>			
	-				

45	What type of depe	Functional De	ependency(N)	
		Multivalued [L Dependenc	<u>. </u>	
		Transitive De	pendency(I	N)	
			, , ,		
		Join depende	ncv(N)		
		Join depende	ncy(N)		
46	Identify the main r	Improvo data	cocurity/M	١	
40	luentiny the main r	improve data	security(N)	
		- 1			
		Enhance data	integrity(N	N)	
		Accelerate da	ita retrieva	l(Y)	
		Reduce stora	ge space(N)	
47	Specify the commo	Stacks(N)			
		Arrays(N)			
		Queues(N)			
		B-Trees(Y)			
		-()			
48	Select the option a	Sort data(NI)			
	Select the option a	Sort data(iv)			
		Dotring day	autald (AA)		
		Retrieve data	quickly(Y)		
		Create a back	cup of data	(N)	
Щ_	1	<u> </u>		L	L

						1
			[-
		Encrypt data	(IN) I			
49	Choose the accura	A database e	rror(N)			
		A situation w	here two d	ifferent key	s hash to t	he same location(Y)
]
		A data incons	sistency(N)			1
						1
		A security Bro	each(N)			
		rescourie, Ere				1
						1
50	Determine the pur	To opforce de	l ata intogriti	//NI)		-
50	Determine the pur	ro enforce da	ata integrit I	y(N)		-
			<u> </u>		<u> </u>	
		To provide a	unique idei	ntifier for e	ach record	(N) 1
		To speed up	the retrieva	of record	s(Y)	
		To create a b	ackup of th	e database	(N)	
]
51	Choose the correct	It is created a	utomatical	lly by the D	BMS.(N)	1
				<u> </u>		1
						1
		It is always a	clustered ii	ndex (N)		
		it is divays a		Ιαεχ.(11)		
						1
		It is used to s	nood un th	o rotrioval	of records I	J based on a non-primary key
		it is used to s	реец ир пі І	e retrievari	l	Dased on a non-primary key
						-
		14				 - - (NI)
		It cannot be	created afte	er the creat	ion of the i	table.(N) 1
			<u> </u>			
52	Choose the correct	To sort data(N)			
	1	To encrypt da	. / \	i	i	1

	1		1		
		To generate a	l a unique id	L entifier for	L each recor
		To generate t	a dinique id		
		To map keys	to hash cod	les for stor	age and ret
53	Choose the correct	Collisions are	frequent.(N)	
			<u> </u>	<u> </u>	. (21)
		It always pro	duces uniqi T	ue hash cod	des(N)
		It is sensitive	to changes	in the innu	l ıt data (Y)
		it is scrisitive	lochanges	in the inpe	
		It is complex	and difficul	t to unders	tand.(N)
54	Identify the disadv	Inefficient fo	r range que	ries(N)	
			<u> </u>		
		High insertion	n and delet T	ion costs(Y) I
		Limited supp	ort for ad-h	oc gueries	(NI)
		сиписа зарр		loc queries	
		Requires a se	parate stor	age space	for keys(N)
55	Specify the commo	SQL(Y)			
		(· ()			
		PL/SQL(N)			
		T_SOL(NI)			
		T-SQL(N)			
	 				
		NoSQL(N)			
		·		·	

56	Identify the primar	SOL(N)			
	lacinary and primar	<u> </u>			
		PL/SQL(Y)			
		T-SQL(N)			
		Manage DD 0:		(NI)	
		MongoDB Qเ	iery Langua	lge(N)	
57	Name an open-sou	Oracle Datab	L ase(N)		
	Traine an open see	0.46.6 24.44			
		Microsoft SQ	L Serve(N)		
		PostgreSQL(Y	<u>') </u>		
		IDAA DI 2/NI)			
		IBM Db2(N)			
58	In query processing	Parsing(Y)			
	по чисту рассессия				
		Optimization	(N)		
		Execution(N)			
		Campilation/	N1)		
		Compilation(N)		
59	Specify the compo	Query Optim	izer(Y)		
	, , , , , , , , , , , , , , , , ,	, , - p	. ,		
		Query Parser	(N)		
		Query Execut	or(N)		

	1				
		Query Planne	er(N)		
60	In query optimizati	Minimize the	allery resn	onse timel	V)
-	in query optimizati	IVIIIIIIIZC CIIC	query resp	onse time(',
		Maximize the	storage sp	ace(N)	
		Minimize the	database s	size(N)	
		Maximize the	numher o	f concurrer	t users(N)
		.viaxiiiize tile	a.iibci U	. concurrer	43013(14)
	<u> </u>				
61	is a common	Indexing(N)			
		Normalizatio	n(N)		
		Denormalizat	ion(N)		
		Demormanizat			
		All of the other	- /\/		
		All of the abo	ve(Y)		
62	Identify the comm	Oracle Datab	ase(N)		
		Microsoft SQ	L Server(Y)		
<u> </u>		NAVCOL (NI)			
		MySQL(N)			
		SQLite(N)			
L					
63	Name an open-sou	MySQL(Y)			
	,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	, ,			
		PostaroSOL/N	1\		
		PostgreSQL(N	1)		

		1			1	
	+					
	+	MangaDD(NI)				
		MongoDB(N)				
						-
		Cassandra(N)				
64	join is gen	Nested Loop	Join(N)			
		Hash Join(Y)				
		Merge Join(N	 I)			
		Cross Join(N)				
		,				
65	Define the concept	Pushing dow	n ioin opera	itions in th	ı e guerv pla	ı n(N)
						c
		Pushing dow	n filter con	l ditions in th	l ne query nla	 an(V)
		T darling down			le query pie]
	+	Duching down	 	functions	in the auer	
	+	Pushing dow	ii aggregate	Turictions	Trune quer	y pian(iv)
		D 1: 1			l	(2.1)
		Pushing dow	n sort oper I	ations in th	e query pia	n(N) I
66	Specify the primar	To transform	a query int	to an optim	ized execut	tion plan(N) I
		To validate th	ne syntax o	f the query	(N)	
		To execute th	ne query an	d retrieve	the results(Y)
		To analyze th	e query for	potential	optimizatio	ns(N)
						,

67	Explain the concep	Selecting the	best index	for a given	query to in	nprove performance(Y)
		- 1]
		Selecting the	most recer	nt index cre	eated in the	database(N)
		Colocting all s	 	dayas far a	anon/M)	
		Selecting all a		dexes for a	query(N)	
		Selecting the	index with	the lowest	. cardinality	J _' (N)
		ourouning uno]
68	algorithm is c	Nested Loop	Join(N)			
		Bubble Sort J	oin(N)			
		Merge Join(Y)			
		Quick Sort Jo	in(N)			
69	Dofine the terms "e	Ontinoising		<u> </u>] (\(\)
69	Define the term "c	Optimizing qu	leries base	d on estim	l esour	ce usage(+)
		Optimizing qu	L Jeries hase	d on user r	l references	J (N)
		Ориниги в ч	001103 5030	011 4361 p		
		Optimizing qu	ueries with	out conside	ering resoui	rce usage(N)
		1 01				
		Optimizing qu	ueries base	d on histor	ical data(N)	
70	DBMS use during o	Query Parsing	g(N)			
			<u> </u>			
		Query Rewrit	ing(N)			
		0	in Die C			
		Query Execut	ion Plan Ge	eneration(Y I	<u>)</u>	
	1	1	I	I	I	I

	_				
		Query Compi	lation(N)		
		, ,	, ,		
71	Determine the min	1(N)			
		2(N)			
		2(11)			
		3(Y)			
		4/81)			
		4(N)			
72	Identify the key ch	Non-leaf nod	es store da	ta entries(N	V)
	, , , , , , , ,				<u> </u>
		Supports dup	licate keys	(N)	
		Height-balan	red tree(Y)		
		ricigite balant			
		Suitable for in	n-memory s	storage(N)	
73	Chassa the starage	D Troo(NI)			
/3	Choose the storage	B-Tree(N)			
		AVL Tree(N)			
		`			
			,		
		Hash Table(N)		
		Linked List(Y)			
		5 2.15 ((1)			
74	Provide the maxim	Non-leaf nod	es store da	ta entries(N	۷)
		n-1(Y)			
		I I I T I			

	·			r	
		2n(N)			
		211(11)			
		2n-1(N)			
		, ,			
	<u> </u>				
75	Identify the proper	Height-balan	ce(Y)		
		Duplicate key	ıc allowed()	NI)	
		Duplicate key		l I	
		Data entries i	n non-leaf	nodes(N)	
				· · ·	
		Leaf nodes at	the same	level(N)	
7.0	I do notification at a second	D. Troo(V)			
76	Identify the storage	B-Tree(Y)			
		AVL Tree(N)			
		, ,			
		Hash Table(N)		
		Lloop/NI)			
		Heap(N)			
L					
77	In a B-Tree, indicat	Node splits in	ito two(Y)		
	,	,	· · ·		
		Node merges	with a sibl	ing node(N)
		Koy is insorts	d in the ne	rant nada/f	\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\
		Key is inserte	а птине ра	i erit 1100e(1	N <i>j</i>
		Key is discard	led(N)		
		, : : : : :	. ,		
	+				
<u> </u>					

78	8 In a B+ Tree, the po Non-leaf nodes(N)					
		Root node(N)				
		Leaf nodes(Y)				
		Lear Hodes(1)	<u>'</u>			
		ntermediate	nodes(N)			
79	Choose the storage	B-Tree(Y)				
		AVL Tree(N)				
		7.1.2 1.1.00(1.1)				
		Dad Dlad To	- / N I \			
		Red-Black Tre	e(N)			
		Heap(N)				
80	Choose the operat	Insertion(Y)				
		,				
		Dolotion (N)				
		Deletion(N)				
		Range search	ing(N)			
		Point query(N	1)			
		, , ,				
81	Select the most off	n(Y)				
-	Select the most off	('/				
		4/2.3				
		n-1(N)				
		2n(N)				

		,			
		2n-1(N)			
82	State the operation	Traversing in	nre-order()	\	
02	State the operation	Traversing in	pre-order(i	N)	
				4>	
		Finding maxir	num eleme	ent(N)	
		Finding media	an(N)		
		Range search	ing(Y)		
02		Daint access	()()		
83	Choose the kinds o	Point queries	(Y)		
		Range querie	s(N)		
		Joins(N)			
		, ,			
		Sorting(N)			
		Joi ting(iv)			
84	Typically, data entr	Non-leaf nod	es(N)		
		Leaf nodes(Y)			
		Root node(N)			
	1				
		Intermediate	nodos(NI)		
	1	mediate	HOUES(IV)		
85	Describe the main	Less disk spac	ce usage(N)		
		Faster inserti	on and dele	etion(N)	

			l	ĺ	Ī	1
		Efficient for r	ange queri	es(Y)		
		Lower height	of the tree	e(N)		
			<u> </u>			
86	Determine the mo	Serializability	'(Y) T			
		Encryption(N	<u> </u> }			
			<u>, </u>			
		Redundancy(N)			
			(2.1)			
		Compression	(N) I			
87	serializability in tra	Ability to exe	L cute transa	l ections in a	l serialized n	 nanner(Y)
		, rusmity to ente				
		Ability to exe	cute transa	ctions con	currently(N)
		A1 111				
		Ability to rec	over from f T	ailures(N)		
		Ability to end	rvpt transa	L ctions(N)		
		7 10 1110				
88	Indicate the purpo	Maximize thr	oughput(N)		
		Minimize late	ency(N) T			
		Ensure serial	l izability wh	l ile maximiz	ing concur	l rencv(Y)
			Lawring Will		S COTICUIT	
		Maximize en	cryption(N)			

89	Which of the follow	Ability to enc	rypt transa	ctions for s	ecure proc	essing(N)
		Ability to pro	cess transa	ctions in a	way that th	e result is the same as if they wer
		Ability to pro	coss transa	ctions with	out any on	J cryption(N)
		Ability to pro	Less transa	T WILLI	Tout any en	
		Ability to pro	cess transa	ctions cond	currently wi	ith no restrictions(N)
90	Select the purpose	To enforce se	rial execut	ion of trans	actions(N)	
		To allow may	imum cone	Lurronov wh	l ilo oncurin	J g data consistency(Y)
		TO allow Illax		Turrency wi	liie ensum	
		To minimize	encryption	overhead(I	<u>v)</u>	
		o prevent tra	nsactions f	rom execut	ing concur	rently(N)
]
01	l., t.,,,,,,,,,	A			()()	
91	In transaction man	A set of instru	actions for	encryption	(Y)	
		A sequence o	f operation	ns from trar	nsactions(N)
		A set of trans	actions exe	ecuting con	currently(N])
		7130001010113				i İ
			<u> </u>	<u> </u>]
		A schedule ha	as no releva	ance in trar	isaction ma	nagement(N)
92	properties	Recoverabilit	y(N)			
		Consistency(I	NI)			
		CONSISTENCY(I	<u>'</u>			
		Atomicity(N)				

		Conflict Corio	li=abili+v/\/\		
		Conflict Seria	iizabiiity(1)		
93	In a transaction ma	Operations th	at do not a	effect the d	 atabase(Ni
93	in a cransaction ma	Operations ti	iat do not a		
		Operations th	nat can be e	executed in	ı anv order
					<u> </u>
		Operations th	nat operate	on differe	nt data iter
		Operations th	nat execute	sequentia	lly(N)
0.4	Character than a state of	1 1 . 1 1 .			
94	Choose the metho	Lock-based co	oncurrency	control(Y)	
		Encryption-ba	sed concu	rrency con	trol(N)
		Life yperon be	asca correa	Tremey com	
		Compression	-based con	currency co	ontrol(N)
		Recovery-bas	ed concurr	ency contr	ol(N)
95	Choose the benefit	Reduced enci	ryption ove	rhead(N)	
		Increased thr	oughput/N	\	
		increased till	ougriput(iv)	
		Lower latency	/(N)		
			/ (/		
		Enhanced sec	curity(Y)		
96	Declare that the co	Two-phase lo	cking(N)		
		<u> </u>		<u> </u>	ļ
		Strict two-ph	ase locking	(N)	

]
		Optimistic co	ncurrency	(V)		
		Optimistic co	licuiteticy			
		Contain alder		1 - 1 (81)		
		Serializable s	napshot isc	lation(N)		
97	A transaction in a	Indivisible(Y)				
		Reversible(N)				
		Concurrent(N	1)			
		Synchronous	(N)			
		,	,			
98	Determine the pol	Strict two-nh	L ase locking	[
	Determine the por	i Strict two pii	asc locking			
-		Conservative	two phace	locking(N)		
		Conservative	two-phase	locking(iv)		
-		Diagram tura		-i(NI)		
		Rigorous two	-pnase loci	(ing(in) I		
		Optimistic co	ncurrency	control(Y)		
99	What feature of rig	Locks are acq	uired and ı	released in	two phases	5(N)
		Locks can be	acquired a	nd released	at any tim	e(N)
		Transactions	can acquire	e both read	and write	ı locks simultaneously(N)
]
		Transactions	releace all	lacks anly s	fter they c	J ommit(Y)
		Tansactions	i cicase aii	TOCKS OTHY C	arter they t	
 						
						J

100	A schedule is consi	Transactions are executed in any order(N)				
		All transactio	ns commit	successfull	y(N)	
		After a transa	action com	mits, it is st	ill possible	to recover to a consistent state(Y
		Transactions	do not inte	rfere with	each other	(N)
101	ensures that	Atomicity(Y)				
		Consistency(N)			
		Durability(N)				
		Isolation(N)				
102	A schedule that is	Serializability	without ar	ny conflicts	(N)	
		The absence	of deadlocl	ks(N)		
		Consistency l	out not isol	ation(N)		
		Serializability	equivalent	t to a serial	schedule(Y)
103	If two transactions	Executed in a	ny order(Y)		
		Executed seq	uentially o	nly(N)		
			,	, , , , , , , , , , , , , , , , , , ,		
		Executed in p	arallel only	/(N)		
		<u>'</u>	, 	1		

		Executed witl	n a shared	lock(N)	
104	Other transactions	Shared lock(Y	<u>'</u>		
	other transactions	Sharea lock()	,		
		Francisco Incil	-/81)		
		Exclusive lock	(IV)		
		Read lock(N)			
		Write lock(N)			
105	Write the drawbac	Deadlocks ca	n occur(N)		
100	Time the did wood	Deda Tooks ear			
		It is too strict	and reduct	es concurre	ncy(Y)
		It requires ad	ditional sto	rage(N)	
		It does not er	sure data	consistency	(N)
106	When transactions	Strict schedul	e(N)		
100	TTTTCTT CT GTTGGCCTGTTG	30.100 30.100 41.	C(1.1)		
		Corial cabad	lo(V)		
		Serial schedu	ie(1)		
		Timestamp so	chedule(N)		
		Conservative	schedule(N	1)	
107	concurrency	Two-phase lo	cking(N)		
	concarrency	priase io	-·····················/		
-		Ombigation			
		Optimistic co	ncurrency	control(Y)	

		Strict two-ph	 ase locking	[(N)		
		р				
		Serializable s	napshot iso I	olation(N)		
108	A transaction goes	Growing pha	se and shrir	nking phase	e(Y)	
		Read phase a	l ınd write pl	l nase(N)		
		Locking phase	o and unloc	king phase	(NI)	
		Locking phase		King phase	(14)	
		Begin phase a	and end ph	ase(N)		
109	lidentify the conflic	A situation w	here two tr	ransactions	try to acce	ss the same data item(Y)
		A situation w	here a tran	saction is a	horted(NI)	
		A situation w	nere a tran	380001113 8	borteu(iv)	
		A situation w	here a tran	saction is c	ommitted(I	N)
		A situation w	here a tran	saction is r	olled back(I	N)
110	is a conflict	A schedule th	nat preserve	es the orde	r of transac	tions(N)
110	is a connec	A seriedale ti	lat preserve	es the orde	l or transac	(14)
		A schedule th	nat does no I	t have any	conflicting (operations(N)
		A schedule th	nat is free fr	om deadlo	cks(N)	
		A schedule th	at is equiv	lent to sor	ne serial se	hadula(V)
		A scriedule (I	at is equive	10 301	ile sellal sc	 -

111	A serializable sche	Concurrently	[N)			
		Sequentially(Y)			
			•			
		Randomly(N)				
		,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,				
		Asynchronou	slv(N)			
		7.07.10.11.01.00	., (,			
112	Determine whatev	Consistency()	N)Isolation(ı (N)Durabilit	ı :v(Y)Atomic	J citv(N)
						mance(N)Time-series(Y)Document
			<u> </u>		<u>-</u>	(Y)Optimized for transactional pro
_		<u> </u>	<u> </u>			ew Concurrency Control(N)Multi-
						a redundancy(N)Simplifying data
	,					y(Y)To improve data retrieval perf
	Determine the data			<u>.</u>		• • •
-		· ` ` <i>'</i> · · ·	• • •	· · ·		d, Update, Delete(N)Create, Retri
	Examine the datab					
121			-		-	ty(N)Handling of unstructured dat
-			-			nounts of data(N)To ensure data c
_	Identify which of the					
	,					J ng data in the database(N)Defining
_						sical schema(N)External schema(Y
_	•					ate(N)Inconsistent state(Y)
						MS(N)Concurrency-control compo
			_			
_	Identify what state					transaction and stop transaction(
	· ·	· · · ·	<u> </u>	<u> </u>		• • • • • • • • • • • • • • • • • • • •
	Identify which of the		-	-		i
—			•			ph(N)All of the mentioned((N)
	State that what do	<u> </u>			-	
	,					ximize database performance(N)1
				•	•	ta redundancy(N)Data indexing(N
135						ta confidentiality(Y)Data availabili
		· ·	-	-		indancy(N)Improved data integrit
						(N)Data archiving(N)Data normali
—						availability(N)Monitoring and acco
_	Identify the compo					
		· · · · · · · · · · · · · · · · · · ·	-			undancy(N)Mitigation of vulnerabi
		·				a redundancy(N)Enhances data a
		<u> </u>				edundancy(N)Protection of sensiti
143	Determine the prir	To identify pe	rformance	bottleneck	(s(Y)To revi	ew and evaluate security measure

144	Select the encrypti	Symmetric encryption(Y)Asymmetric encryption(N)Hybrid encryption(N)One-tir
145	Find out what action	Implementing strong authentication mechanisms(N)Applying input validation a
146	Determine the prir	Maximizing data redundancy(N)Ensuring data confidentiality(N)Minimizing dow
147	In the context of d	Providing all users with maximum access rights(N)Providing users with the least
148	State the advantag	Enhanced data redundancy(N)Improved data availability(N)Mitigation of unauth
149	Identify the role do	Enhancing data integrity(N)Improving data availability(N)Protecting sensitive in
150	Select the encrypti	Symmetric encryption(N)Asymmetric encryption(Y)Hybrid encryption(N)One-tir
151	Select the strategy	Implementing data validation checks(Y)Increasing data redundancy(N)Optimizir
152	Establish in a datak	Enhancing data redundancy(N)Monitoring and detecting potential unauthorized

re executed in some sequential order(Y)

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t-oriented(N)
ocessing(N)Uses a hierarchical data model(N)Supports flexible schemas(N)
Value Concurrency Control(N)Multi-Vendor Concurrency Control(N)
retrieval(N)Enhancing data consistency(N)
formance(N)To enhance data durability(N)To increase data redundancy(N)
ieve, Update, Display(N)Compute, Retrieve, Update, Delete(N)
ta(N)High performance(N)
consistency(N)To handle concurrent transactions(N)
g the structure of the database(Y)Sorting data in the database(N)
1)
onent of the DBMS(Y)Transaction-management component of the DBMS(N)Buffer management component in
[N)Get transaction and post transaction(N)Read transaction and write transaction(N)
To reduce storage costs(N)To increase database size(N)
1)
ity(N)
ty(N)Restriction of unauthorized access(Y)
ization(N)
ountability(Y)Increased data integrity(N)
ilities(Y)Data encryption(N)
vailability(N)Streamlines access management(Y)
ive information(Y)Enhanced database performance(N)
es(N)To optimize data storage(N)To enhance data redundancy(N)
```

me pad encryption(N)

 $nd\ parameterized\ queries (Y) Increasing\ data\ redundancy (N) Encrypting\ all\ database\ contents (N)$

vntime and data loss(Y)Enhancing data availability(N)

t possible access rights necessary for their job functions(Y)Increasing data redundancy(N)Encrypting all database horized access(Y)Increased data encryption(N)

formation by disguising original data(Y)Increasing data redundancy(N)

me pad encryption(N)

ng database performance(N)Data obfuscation(N)

d access or attacks(Y)Improving data availability(N)Increasing data encryption(N)



se contents(N)

Sr. Short/Long questions

- 1 Define Data Integrity and elucidate its importance within a database management system.
- 2 Explain the hierarchical data model and provide an illustrative example depicting its representation of relationships in a real-life scenario.
- 3 Define the relational model and exemplify how this model structures a table.
- 4 Develop a basic Entity-Relationship Diagram (ERD) for a library system, identifying entities, attributes, relationships, and explain its constituent
- 5 Describe the network data model, highlighting its advantages and disadvantages in comparison to the hierarchical and relational models.
- 6 Highlight the fundamental distinctions between SQL and NoSQL databases, emphasizing their respective query languages.
- 7 Explain the purpose and syntax of the SELECT statement in SQL, emphasizing its role in data retrieval.
- 8 Discuss query processing within a DBMS, outlining its primary stages.
- 9 Discuss the significance of query parsing and optimization in enhancing database performance during query processing.
- 10 Explain the importance of query optimization in improving overall database performance.
- 11 Compare and contrast open-source and commercial DBMS, emphasizing differences in licensing, support, and cost.
- 12 Discuss a major advantage and limitation of using open-source DBMS in comparison to commercial alternatives.
- 13 Compare various query languages used in databases.
- 14 Explain key components of query optimization briefly.
- 15 Define relational algebra and relational calculus.
- 16 Differentiate between various types of relational calculus.
- 17 Explain Armstrong's Axioms.
- 18 Describe B-Tree concerning Database Management Systems.
- 19 Explain how a B-Tree manages insertions and deletions while maintaining balance.
- 20 Illustrate how a B-Tree facilitates efficient search operations.
- 21 Explain the storage policy of a B+ Tree concerning duplicate keys.
- 22 Compare the storage efficiency of a B-Tree and a linked list for sorted data.
- 23 Highlight the primary advantage of using B-Trees in database systems.
- 24 Establish the main benefit of employing B-Trees in database systems with diverse data access patterns.
- 25 Explain how a B+ Tree manages duplicate keys within its storage structure.
- 26 Assess the primary advantage of implementing B-Trees in database systems that encounter diverse data access patterns.
- 27 Define the concept of a Transaction within the framework of a database management system.
- 28 Detail the structure of a Schedule in Transaction Processing.
- 29 Elaborate on serializability in the realm of transaction processing.
- 30 Clarify the notion of a conflicting operation within a transaction schedule.
- 31 Define and discuss a serial schedule, elucidating its relevance in transaction processing.
- 32 Describe a serializable schedule and justify its significance in database systems.

- 33 Explain the construction and utilization of a precedence graph for schedule analysis.
- 34 Define the concept of a conflict-serializable schedule.
- 35 Analyze the contribution of the Two-Phase Locking (2PL) protocol to concurrency control.
- 36 Elaborate on the concept of a dirty read in transactions.
- 37 Describe the ACID properties and emphasize their importance in transaction processing.
- 38 Elucidate the purpose of a security policy in database design.
- 39 Explain the principle of least privilege and how it enhances database security.
- 40 Detail how a database firewall reinforces security policy enforcement.
- 41 Describe the role of Database Encryption Key Management in implementing security policies.
- 42 Illustrate Database Clustering and its role in ensuring high availability within security policies.
- 43 Explain NoSQL databases and highlight their differences from traditional relational databases.
- 44 Define blockchain technology and its potential impact on database design.
- 45 Evaluate the purpose of graph databases and identify scenarios where they prove beneficial.
- 46 Explain the significance of Secure Socket Layer (SSL) in enhancing database security.
- 47 Elaborate on the concept of data integrity within a security policy.
- 48 Develop a simple Entity-Relationship Diagram (ERD) for a library system, highlighting entities, attributes, relationships, and their components.
- 49 Discuss the significance of each element in ensuring data integrity and effective database design.
- 50 Explain various aspects of data definition in Database Management Systems, including data types, constraints, attributes, and keys, providing ε
- 51 Explain the concept of data independence within a database system.
- 52 Compare and contrast SQL and NoSQL languages using examples.
- 53 Delve into the realm of query optimization.
- 54 Evaluate the techniques employed for optimizing the execution plan of queries.

