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Fast Algorithm for the 3-D DCT-II

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Abstract—Recently, many applications for three-dimensional (3-D) image and video compression have been proposed using 3-D discrete cosine transforms (3-D DCTs). Among different types of DCTs, the type-II DCT (DCT-II) is the most used. In order to use the 3-D DCTs in practical applications, fast 3-D algorithms are essential. Therefore, in this paper, the 3-D vector-radix decimation-in-frequency (3-D VR DIF) algorithm that calculates the 3-D DCT-II directly is introduced. The mathematical analysis and the implementation of the developed algorithm are presented, showing that this algorithm possesses a regular structure, can be implemented in-place for efficient use of memory, and is faster than the conventional row-column-frame (RCF) approach. Furthermore, an application of 3-D video compression-based 3-D DCT-II is implemented using the 3-D new algorithm. This has led to a substantial speed improvement for 3-D DCT-II-based compression systems and proved the validity of the developed algorithm.

Index Terms—Fast multidimensional algorithms, fast 3-D transforms, 3-D DCT, 3-D vector-radix algorithm.

I. INTRODUCTION

SINCE the introduction of one-dimensional discrete cosine transforms (1-D DCTs) in 1974 [1], it has been applied in a wide range of applications, because the DCT performs very close to the statistically optimum Karhunen–Loeve transform [2] in terms of compression performance. These applications are mainly in data, image/video, and multimedia applications [3]–[9]. Additionally, it has been adopted as part of several compression standards [10]–[12].

This has led to the development of a large number of fast algorithms to calculate the 1-D and two-dimensional (2-D) DCTs. These algorithms can be classified into direct and indirect algorithms. The direct algorithms generally have a regular computational structure, which reduces the implementation complexity [13]–[16]. On the other hand, indirect algorithms exploit the relationship between the DCTs and other transforms. These algorithms include the calculation of the DCT through the fast Fourier [17], Hartley [18], and polynomial transforms [19]. These algorithms generally have irregular structures and complex indexing schemes.

Although many algorithms have been developed for fast calculation of the 1-D and 2-D DCTs [13]–[24], algorithm development for the multidimensional discrete cosine transform (m-D DCT) in three and more dimensions is more challenging and has

not been given similar attention. Hence, the three-dimensional (3-D) DCT is usually calculated using the row-column-frame (RCF) approach or through mapping it to 1-D and using other transforms [25]–[29]. However, proper and direct m-D algorithms have a better computational structure, can be more efficient than the RCF approach, and need to be developed.

Owing to the rapid growth in the 3-D applications based on the 3-D DCT [30]–[36], there is a greater need now to develop fast algorithms for the 3-D DCT for such applications. Consequently, this paper concentrates on developing a direct, fast, and efficient algorithm for fast calculation of the type-II 3-D DCT (3-D DCT-II). The 3-D DCT-II was first suggested for 3-D compression applications in 1977 [30], [31], but it did not gain much attention because of the high computation time involved. This is reflected by the huge amount of data associated with its calculation and the lack of fast algorithms in 3-D. Due to the enhancement in software, hardware, and the introduction of fast algorithms, many new applications have been proposed based on the 3-D DCTs [32]–[36]. These applications include hyper-spectral coding systems [32], variable temporal length 3-D DCT coding [33], video coding algorithms [34], adaptive video coding [35], and 3-D compression [36], etc.

Since the 3-D DCT is separable, it is usually calculated by successively applying 1-D fast algorithms over the rows, the columns, and then the frames in what is known as an RCF approach. Other fast algorithms have been developed to reduce the computational cost of the 3-D DCT [24]–[29]. Common to all these algorithms is that they exploit the relationship between the DCT and other transforms such as Fourier, Hartley, and polynomial transforms to calculate the 3-D DCT involving complex indexing and overheads. Direct and true multidimensional algorithms that have regular structure and indexing schemes are preferred [37]–[43].

In this paper, new 3-D vector-radix decimation-in-frequency (VR DIF) algorithm is developed for fast calculation of the 3-D DCT-II. The development shows all the stages of calculation, the arithmetic complexity, and the computer run-times. From the number of arithmetic operations, the 3-D DCT-II VR DIF algorithm is found to require substantially fewer multiplications as compared with the familiar RCF approach. Moreover, based on the computer run-times, it is found to be faster (around 70% for $8 \times 8 \times 8$ 3-D DCT) than the RCF approach. In addition, it has a regular butterfly structure, making it suitable for software and hardware implementations.

The organization of this paper starts with the definition of the 3-D DCT-II in Section II. Section III presents the mathematical derivation and development, as well as the arithmetic complexity of the 3-D DCT-II VR DIF algorithm. Comparisons between the 3-D vector-radix algorithm and other related algorithms are carried out in Section IV. Finally, in Section V, the

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developed 3-D DCT-II VR algorithm is used for the implementation of a 3-D video compression algorithm. The “C” code for the butterfly stage of the computation of the 3-D DCT-II VR DIF algorithm is included in the Appendix. The full “C” code for the 3-DCT-II and 3-D IDCT-II is available from the authors.

II. DEFINITION OF THE FORWARD AND INVERSE 3-D DCT-II

The 3-D discrete cosine transform type-II $X(k_1, k_2, k_3)$, of size $N \times N \times N$, is defined as

$$X(k_1, k_2, k_3) = \frac{8}{N^3} \varepsilon_{k_1} \varepsilon_{k_2} \varepsilon_{k_3} \sum_{n_1=0}^{N-1} \sum_{n_2=0}^{N-1} \sum_{n_3=0}^{N-1} x(n_1, n_2, n_3) \\ \times \cos\left(\frac{\pi}{2N}(2n_1+1)k_1\right) \\ \times \cos\left(\frac{\pi}{2N}(2n_2+1)k_2\right) \\ \times \cos\left(\frac{\pi}{2N}(2n_3+1)k_3\right) \\ k_1, k_2, k_3 = 0, 1, \dots, N-1. \quad (1)$$

The inverse 3-D DCT-II (also called 3-D DCT-III) is defined as

$$x(n_1, n_2, n_3) = \sum_{k_1=0}^{N-1} \sum_{k_2=0}^{N-1} \sum_{k_3=0}^{N-1} \varepsilon_{k_1} \varepsilon_{k_2} \varepsilon_{k_3} X(k_1, k_2, k_3) \\ \times \cos\left(\frac{\pi}{2N}(2n_1+1)k_1\right) \\ \times \cos\left(\frac{\pi}{2N}(2n_2+1)k_2\right) \\ \times \cos\left(\frac{\pi}{2N}(2n_3+1)k_3\right) \\ n_1, n_2, n_3 = 0, 1, \dots, N-1 \quad (2)$$

where

$$\varepsilon_{ki} = \begin{cases} \frac{1}{\sqrt{2}}, & \text{for } k_i = 0 \\ 1, & \text{otherwise} \end{cases} \quad i = 1, 2, 3.$$

III. DEVELOPMENT OF THE 3-D DCT-II VR DIF ALGORITHM

The vector-radix algorithms have been applied for fast calculation of numerous multidimensional transforms such as discrete Fourier [42] and Hartley transforms [43]. They have been also used for fast calculation of the 2-D DCTs [41], [44] and the 3-D DCT-III [45]. These algorithms deal with the data as a multidimensional array and calculate the DCT directly, whereas the familiar RCF approach uses algorithms developed for the calculation of 1-D transforms to calculate multidimensional transforms. Direct m-D algorithms are found to be more efficient and faster than the RCF approach [38], [40]–[45].

In this section, for the purpose of fast 3-D video compression applications, the 3-D VR DIF algorithm is developed and analyzed for the 3-D DCT-II.

A. Algorithm Development of the 3-D VR DIF Algorithm for the 3-D DCT-II

For simplicity, the multiplication by the normalization factor $(8/N^3)$ and the factor $(\varepsilon_{k_1} \varepsilon_{k_2} \varepsilon_{k_3})$ can be neglected or delayed

to the last stage. The transform size $N \times N \times N$ is assumed to be power of $2 \times 2 \times 2$ in order to apply the vector-radix algorithm. First, the input data $x(n_1, n_2, n_3)$ needs to be rearranged according to the index mapping as follows [13], [45]:

$$\begin{bmatrix} \tilde{x}(n_1, n_2, n_3) \\ \tilde{x}(n_1, n_2, N - n_3 - 1) \\ \tilde{x}(n_1, N - n_2 - 1, n_3) \\ \tilde{x}(n_1, N - n_2 - 1, N - n_3 - 1) \\ \tilde{x}(N - n_1 - 1, n_2, n_3) \\ \tilde{x}(N - n_1 - 1, n_2, N - n_3 - 1) \\ \tilde{x}(N - n_1 - 1, N - n_2 - 1, n_3) \\ \tilde{x}(N - n_1 - 1, N - n_2 - 1, N - n_3 - 1) \end{bmatrix} \\ = \begin{bmatrix} x(2n_1, 2n_2, 2n_3) \\ x(2n_1, 2n_2, 2n_3 + 1) \\ x(2n_1, 2n_2 + 1, 2n_3) \\ x(2n_1, 2n_2 + 1, 2n_3 + 1) \\ x(2n_1 + 1, 2n_2, 2n_3) \\ x(2n_1 + 1, 2n_2, 2n_3 + 1) \\ x(2n_1 + 1, 2n_2 + 1, 2n_3) \\ x(2n_1 + 1, 2n_2 + 1, 2n_3 + 1) \end{bmatrix} \\ 0 \leq n_1, n_2, n_3 \leq \frac{N}{2} - 1. \quad (3)$$

Replacing $x(n_1, n_2, n_3)$ in (1) by $\tilde{x}(n_1, n_2, n_3)$ in (3), $X(k_1, k_2, k_3)$ can be written as

$$X(k_1, k_2, k_3) = \sum_{n_1=0}^{N-1} \sum_{n_2=0}^{N-1} \sum_{n_3=0}^{N-1} \tilde{x}(n_1, n_2, n_3) \\ \times \cos(\phi_1 k_1) \cos(\phi_2 k_2) \cos(\phi_3 k_3) \quad (4)$$

where $\phi_i = (\pi/2N)(4n_i + 1)$, and $i = 1, 2, 3$

If the even and the odd parts of k_1, k_2 , and k_3 are considered, the general formula for the calculation of the 3-D DCT-II can be expressed as

$$X(2k_1 + i, 2k_2 + j, 2k_3 + l) \\ = \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} \tilde{x}_{ijl}(n_1, n_2, n_3) \\ \times \cos(\phi_1(2k_1 + i)) \cos(\phi_2(2k_2 + j)) \\ \times \cos(\phi_3(2k_3 + l)) \\ ijl = \{000, 001, 010, 011, 100, 101, 110, 111\} \quad (5)$$

where

$$\tilde{x}_{ijl}(n_1, n_2, n_3) = \tilde{x}(n_1, n_2, n_3) + (-1)^l \tilde{x}\left(n_1, n_2, n_3 + \frac{n}{2}\right) \\ + (-1)^j \tilde{x}\left(n_1, n_2 + \frac{n}{2}, n_3\right) \\ + (-1)^{j+l} \tilde{x}\left(n_1, n_2 + \frac{n}{2}, n_3 + \frac{n}{2}\right) \\ + (-1)^i \tilde{x}\left(n_1 + \frac{n}{2}, n_2, n_3\right) \\ + (-1)^{i+l} \tilde{x}\left(n_1 + \frac{n}{2}, n_2, n_3 + \frac{n}{2}\right) \\ + (-1)^{i+j} \tilde{x}\left(n_1 + \frac{n}{2}, n_2 + \frac{n}{2}, n_3\right) \\ + (-1)^{i+j+l} \tilde{x}\left(n_1 + \frac{n}{2}, n_2 + \frac{n}{2}, n_3 + \frac{n}{2}\right) \\ i, j, l = 0 \text{ or } 1. \quad (6)$$

For k_1, k_2 , and k_3 -even, (5) can be written as

$$X(2k_1, 2k_2, 2k_3) = \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} \tilde{x}_{000}(n_1, n_2, n_3) \times \cos(\phi_1 2k_1) \cos(\phi_2 2k_2) \cos(\phi_3 2k_3). \quad (7)$$

For k_1 and k_2 -even and k_3 -odd, (5) can be written as

$$X(2k_1, 2k_2, 2k_3 + 1) = \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} \tilde{x}_{001}(n_1, n_2, n_3) \times \cos(\phi_1 2k_1) \cos(\phi_2 2k_2) \cos(\phi_3 (2k_3 + 1)). \quad (8)$$

Using the trigonometric identity

$$\cos(\phi_i (2k_i + 1)) = 2 \cos \phi_i \cos(\phi_i 2k_i) - \cos(\phi_i (2k_i - 1)) \quad i = 1, 2, 3 \quad (9)$$

$X(2k_1, 2k_2, 2k_3 + 1)$ can be decomposed into $(N/2) \times (N/2) \times (N/2)$ -points 3-D DCTs plus some multiplications by twiddle factors and additions. Therefore, $X(2k_1, 2k_2, 2k_3 + 1)$ can be decomposed as

$$X(2k_1, 2k_2, 2k_3 + 1) = \left\{ \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} [2\tilde{x}_{001}(n_1, n_2, n_3) \cos \phi_3] \cos(\phi_1 2k_1) \times \cos(\phi_2 2k_2) \cos(\phi_3 2k_3) \right\} - X(2k_1, 2k_2, 2k_3 - 1). \quad (10)$$

Following the same procedure, $X(2k_1, 2k_2 + 1, 2k_3)$ can be decomposed as

$$X(2k_1, 2k_2 + 1, 2k_3) = \left\{ \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} [2\tilde{x}_{010}(n_1, n_2, n_3) \cos \phi_2] \cos(\phi_1 2k_1) \times \cos(\phi_2 2k_2) \cos(\phi_3 2k_3) \right\} - X(2k_1, 2k_2 - 1, 2k_3) \quad (11)$$

and $X(2k_1 + 1, 2k_2, 2k_3)$ can be calculated as

$$X(2k_1 + 1, 2k_2, 2k_3) = \left\{ \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} [2\tilde{x}_{100}(n_1, n_2, n_3) \cos \phi_1] \times \cos(\phi_1 2k_1) \cos(\phi_2 2k_2) \cos(\phi_3 2k_3) \right\} - X(2k_1 - 1, 2k_2, 2k_3). \quad (12)$$

For k_1 -even and k_2 and k_3 -odd, $X(2k_1, 2k_2 + 1, 2k_3 + 1)$ can be written as

$$X(2k_1, 2k_2 + 1, 2k_3 + 1) = \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} \tilde{x}_{011}(n_1, n_2, n_3) \times \cos(\phi_1 2k_1) \cos(\phi_2 (2k_2 + 1)) \cos(\phi_3 (2k_3 + 1)). \quad (13)$$

Using the trigonometric identity

$$\begin{aligned} & \cos(\phi_i (2k_i + 1)) \cos(\phi_j (2k_j + 1)) \\ &= 4 \cos \phi_i \cos \phi_j \cos(\phi_i 2k_i) \cos(\phi_j 2k_j) \\ & \quad - \cos(\phi_i (2k_i - 1)) \cos(\phi_j (2k_j + 1)) \\ & \quad - \cos(\phi_i (2k_i + 1)) \cos(\phi_j (2k_j - 1)) \\ & \quad - \cos(\phi_i (2k_i - 1)) \cos(\phi_j (2k_j - 1)) \end{aligned} \quad (14)$$

$X(2k_1, 2k_2 + 1, 2k_3 + 1)$ can be converted to $(N/2) \times (N/2) \times (N/2)$ -point 3-D DCT-II as

$$\begin{aligned} X(2k_1, 2k_2 + 1, 2k_3 + 1) &= \left\{ \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} [4\tilde{x}_{011}(n_1, n_2, n_3) \cos \phi_2 \cos \phi_3] \right. \\ & \quad \times \cos(\phi_1 2k_1) \cos(\phi_2 2k_2) \cos(\phi_3 2k_3) \left. \right\} \\ & \quad - X(2k_1, 2k_2 - 1, 2k_3 + 1) - X(2k_1, 2k_2 + 1, 2k_3 - 1) \\ & \quad - X(2k_1, 2k_2 - 1, 2k_3 - 1). \end{aligned} \quad (15)$$

Similarly, $X(2k_1 + 1, 2k_2, 2k_3 + 1)$ and $X(2k_1 + 1, 2k_2 + 1, 2k_3)$ can be decomposed as follows:

$$\begin{aligned} X(2k_1 + 1, 2k_2, 2k_3 + 1) &= \left\{ \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} [4\tilde{x}_{101}(n_1, n_2, n_3) \cos \phi_1 \cos \phi_3] \right. \\ & \quad \times \cos(\phi_1 2k_1) \cos(\phi_2 2k_2) \cos(\phi_3 2k_3) \left. \right\} \\ & \quad - X(2k_1 - 1, 2k_2, 2k_3 + 1) - X(2k_1 + 1, 2k_2, 2k_3 - 1) \\ & \quad - X(2k_1 - 1, 2k_2, 2k_3 - 1) \end{aligned} \quad (16)$$

$$\begin{aligned} X(2k_1 + 1, 2k_2 + 1, 2k_3) &= \left\{ \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} [4\tilde{x}_{110}(n_1, n_2, n_3) \cos \phi_1 \cos \phi_2] \right. \\ & \quad \times \cos(\phi_1 2k_1) \cos(\phi_2 2k_2) \cos(\phi_3 2k_3) \left. \right\} \\ & \quad - X(2k_1 - 1, 2k_2 + 1, 2k_3) - X(2k_1 + 1, 2k_2 - 1, 2k_3) \\ & \quad - X(2k_1 - 1, 2k_2 - 1, 2k_3). \end{aligned} \quad (17)$$

For k_1, k_2 , and k_3 -odd, $X(2k_1 + 1, 2k_2 + 1, 2k_3 + 1)$ can be written as

$$X(2k_1 + 1, 2k_2 + 1, 2k_3 + 1) = \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} \tilde{x}_{111}(n_1, n_2, n_3) \cos(\phi_1 (2k_1 + 1)) \times \cos(\phi_2 (2k_2 + 1)) \cos(\phi_3 (2k_3 + 1)). \quad (18)$$

Using the trigonometric identity

$$\begin{aligned}
& \cos(\phi_i(2k_i + 1)) \cos(\phi_j(2k_j + 1)) \cos(\phi_l(2k_l + 1)) \\
&= 8 \cos \phi_i \cos \phi_j \cos \phi_l \cos(\phi_i 2k_i) \cos(\phi_j 2k_j) \cos(\phi_l 2k_l) \\
&\quad - \cos(\phi_i(2k_i + 1)) \cos(\phi_j(2k_j + 1)) \cos(\phi_l(2k_l - 1)) \\
&\quad - \cos(\phi_i(2k_i + 1)) \cos(\phi_j(2k_j - 1)) \cos(\phi_l(2k_l + 1)) \\
&\quad - \cos(\phi_i(2k_i + 1)) \cos(\phi_j(2k_j - 1)) \cos(\phi_l(2k_l - 1)) \\
&\quad - \cos(\phi_i(2k_i - 1)) \cos(\phi_j(2k_j + 1)) \cos(\phi_l(2k_l + 1)) \\
&\quad - \cos(\phi_i(2k_i - 1)) \cos(\phi_j(2k_j + 1)) \cos(\phi_l(2k_l - 1)) \\
&\quad - \cos(\phi_i(2k_i - 1)) \cos(\phi_j(2k_j - 1)) \cos(\phi_l(2k_l + 1)) \\
&\quad - \cos(\phi_i(2k_i - 1)) \cos(\phi_j(2k_j - 1)) \cos(\phi_l(2k_l - 1))
\end{aligned} \tag{19}$$

$X(2k_1 + 1, 2k_2 + 1, 2k_3 + 1)$ can be converted to $(N/2) \times (N/2) \times (N/2)$ -point 3-D DCT-II as

$$\begin{aligned}
& X(2k_1 + 1, 2k_2 + 1, 2k_3 + 1) \\
&= \left\{ \sum_{n_1=0}^{N/2-1} \sum_{n_2=0}^{N/2-1} \sum_{n_3=0}^{N/2-1} [8\tilde{x}_{111}(n_1, n_2, n_3) \cos \phi_1 \cos \phi_2 \cos \phi_3] \right. \\
&\quad \times \cos(\phi_1 2k_1) \cos(\phi_2 2k_2) \cos(\phi_3 2k_3) \left. \right\} \\
&\quad - X(2k_1 + 1, 2k_2 + 1, 2k_3 - 1) \\
&\quad - X(2k_1 + 1, 2k_2 - 1, 2k_3 + 1) \\
&\quad - X(2k_1 + 1, 2k_2 - 1, 2k_3 - 1) \\
&\quad - X(2k_1 - 1, 2k_2 + 1, 2k_3 + 1) \\
&\quad - X(2k_1 - 1, 2k_2 + 1, 2k_3 - 1) \\
&\quad - X(2k_1 - 1, 2k_2 - 1, 2k_3 + 1) \\
&\quad - X(2k_1 - 1, 2k_2 - 1, 2k_3 - 1).
\end{aligned} \tag{20}$$

B. Arithmetic Complexity of the 3-D DCT-II VR DIF Algorithm

The calculation of the 3-D DCT-II using the 3-D VR DIF algorithm consists of four stages, as shown in Fig. 1. The first stage is the 3-D reordering using the index mapping illustrated by (3). The second stage is the butterfly calculation. Each butterfly calculates eight points together, as illustrated by (7), (10)–(12), (15), (16), and (20), and shown in Fig. 2. In these equations, the parts that are between the parentheses $\{ \}$ are calculated at the butterfly stage, whereas the other parameters are calculated later in the post-additions stage. The whole 3-D DCT calculation needs $\lceil \log_2 N \rceil$ stages, and each stage involves $N^3/8$ butterflies. The whole 3-D DCT requires $[(N^3/8) \log_2 N]$ butterflies to be completed. Each butterfly requires seven real multiplications (including trivial multiplications) and 24 real additions (including trivial additions). Therefore, the total number of real multiplications needed for this stage is $[(7/8)N^3 \log_2 N]$, and the total number of real additions is $[(24/8)N^3 \log_2 N]$. The post-additions (recursive additions) can be calculated directly after the butterfly stage or after the bit-reverse stage, as shown in Fig. 1. This stage

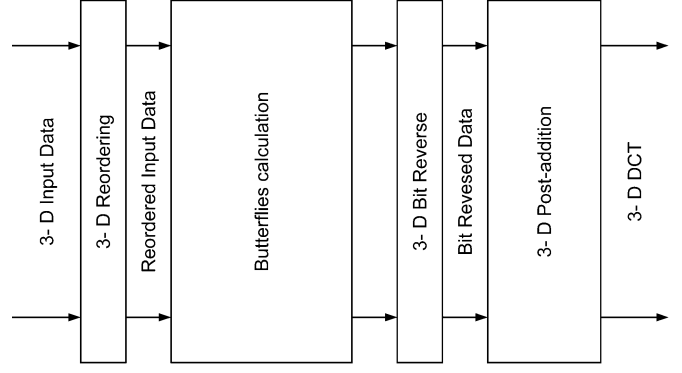


Fig. 1. Stages of the 3-D DCT-II VR DIF algorithm.

involves real additions only. The total number of real additions required for this stage is $[(3/2)N^3 \log_2 N - 3N^3 + 3N^2]$.

Consequently, the total numbers of real multiplications (including trivial multiplications) and additions (including trivial additions) required to calculate the 3-D DCT-II using the 3-D VR DIF algorithm are

$$\text{Multiplications} = \frac{7}{8}N^3 \log_2 N \tag{21}$$

$$\begin{aligned}
\text{Additions} &= \left[\frac{24}{8}N^3 \log_2 N \right] \\
&\quad + \left[\frac{3}{2}N^3 \log_2 N - 3N^3 + 3N^2 \right] \\
&= \frac{9}{2}N^3 \log_2 N - 3N^3 + 3N^2.
\end{aligned} \tag{22}$$

The total number of arithmetic operations (including trivial operations) for the 3-D DCT-II VR DIF algorithm is listed in Table I for different transform sizes.

IV. COMPARISONS OF THE 3-D DCT-II VR ALGORITHM WITH OTHER RELATED ALGORITHMS

In this section, the calculations of the 3-D DCT-II using the 3-D VR algorithm and the RCF approach are compared based on arithmetic operations and computer run times. The developed algorithm is also indirectly compared with existing fast 3-D algorithms for the 3-D DCTs.

A. Comparison of the 3-D DCT-II VR Algorithm With the RCF Approach Based on the Arithmetic Complexity

The arithmetic complexity of the 3-D VR algorithm for the 3-D DCT-II was obtained in Section III-B. On the other hand, the 3-D DCT-II can be calculated using the RCF approach based on 1-D algorithms. Since the 3-D VR algorithm is based on the idea of the radix-2 in 1-D, the 3-D VR is compared with the 3-D RCF approach based on the 1-D radix-2 algorithm [13], [14], [22]. This algorithm has been used as the basis for comparison purposes [40], [41], and it involves

$$\begin{aligned}
\text{Multiplications} &= \left[\frac{1}{2}N \log_2 N \right] \quad \text{and} \\
\text{Additions} &= \left[\frac{3}{2}N \log_2 N - N + 1 \right].
\end{aligned}$$

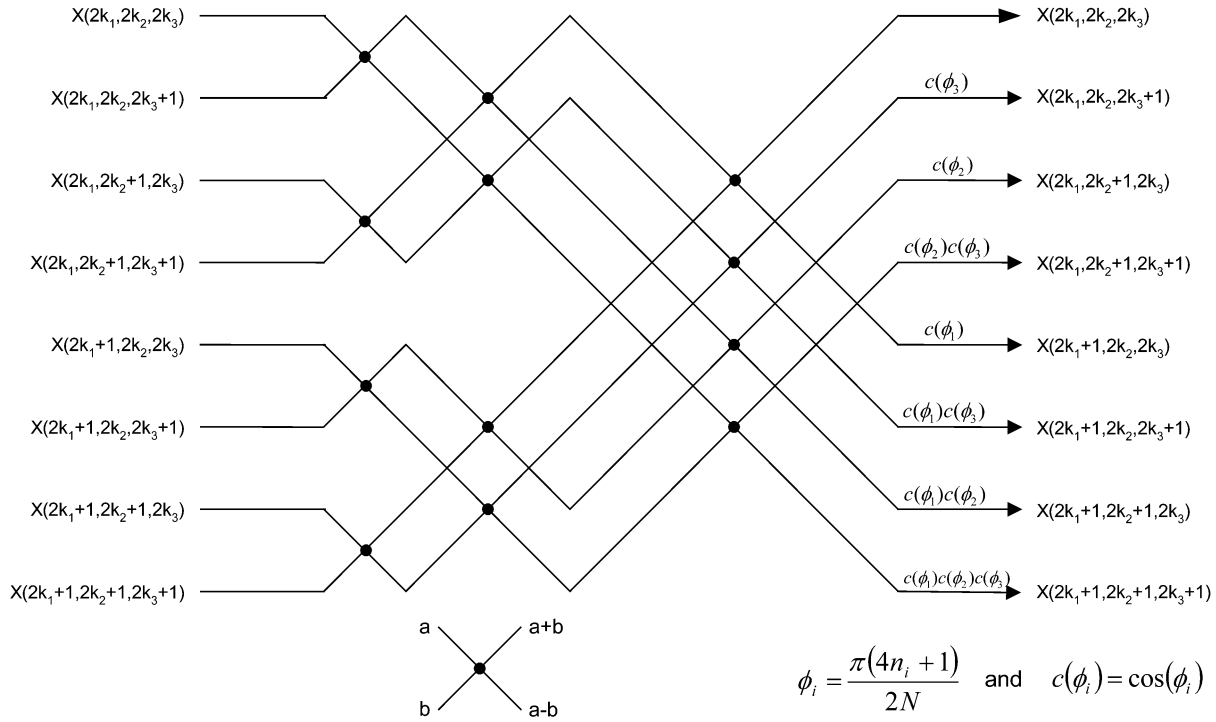


Fig. 2. Single butterfly of the 3-D DCT-II VR DIF algorithm.

The 1-D algorithm needs to be applied, successively, over the rows, columns, and frames, involving

$$\begin{aligned} \text{Multiplications} &= N^2 \left[\frac{1}{2} N \log_2 N \right] + N^2 \left[\frac{1}{2} N \log_2 N \right] \\ &+ N^2 \left[\frac{1}{2} N \log_2 N \right] = \left[\frac{3}{2} N^3 \log_2 N \right] \end{aligned} \quad (23)$$

$$\begin{aligned} \text{Additions} &= N^2 \left[\frac{3}{2} N \log_2 N - N + 1 \right] \\ &+ N^2 \left[\frac{3}{2} N \log_2 N - N + 1 \right] \\ &+ N^2 \left[\frac{3}{2} N \log_2 N - N + 1 \right] \\ &= \left[\frac{9}{2} N^3 \log_2 N - 3N^3 + 3N^2 \right]. \end{aligned} \quad (24)$$

The arithmetic operations are detailed in Table I for different transform sizes.

From Table I and Fig. 3, it can be seen that the total number of multiplications associated with the 3-D DCT VR algorithm is less than that associated with the RCF approach by more than 40%. In addition, the RCF approach involves matrix transpose and more indexing and data swapping than the new algorithm. This makes the 3-D DCT VR algorithm more efficient and better suited for 3-D applications that involve the 3-D DCT-II such as video compression and other 3-D image processing applications [30]–[36].

B. Comparison of the 3-D DCT-II VR Algorithm With the RCF Approach Based on Computer Run-Times

The developed vector-radix algorithm and the RCF approach have been implemented for the 3-D DCT-II using “C” pro-

TABLE I
OPERATIONS COUNT FOR THE CALCULATION OF THE 3-D DCT-II
USING THE 3-D VR AND THE RCF ALGORITHMS

Transform Size	Mults/point		Additions/point		Mults+adds/point	
	3-D VR	RCF	3-D VR	RCF	3-D VR	RCF
8×8×8	2.625	4.5	10.875	10.875	13.5	15.375
16×16×16	3.5	6	15.188	15.188	18.688	21.188
32×32×32	4.375	7.5	19.594	19.594	23.969	27.094
64×64×64	5.25	9	24.047	24.047	29.297	33.047
128×128×128	6.125	10.5	28.523	28.523	34.648	39.023
256×256×256	7	12	33.012	33.012	40.012	45.012
512×512×512	7.875	13.5	37.506	37.506	45.381	51.006
2 ¹⁰ ×2 ¹⁰ ×2 ¹⁰	8.75	15	42.003	42.003	50.753	57.003
2 ¹¹ ×2 ¹¹ ×2 ¹¹	9.625	16.5	46.502	46.502	56.126	63.001
2 ¹² ×2 ¹² ×2 ¹²	10.5	18	51.001	51.001	61.501	69.001

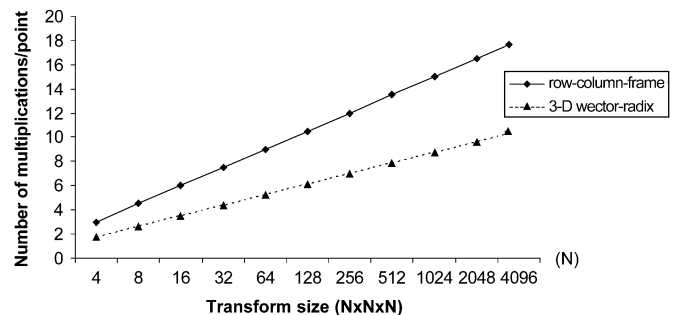


Fig. 3. Comparison between the 3-D vector-radix algorithm and the row-column-frame approach for the 3-D DCT-II (number of multiplications as shown in Table I).

TABLE II
COMPUTER RUN TIMES FOR THE 3-D VR ALGORITHM AND THE
RCF APPROACH FOR THE CALCULATION OF THE 3-D DCT-II
ON A P-4 COMPUTER USING THE VC++ COMPILER

Transform Size	3-D VR algorithm (sec)	RCF approach (sec)	Saving %
2×2×2	0.00000185	5.255E-06	64.795
4×4×4	0.000015	0.0000566	73.498
8×8×8	0.000135	0.0006155	78.067
16×16×16	0.0013	0.006055	78.53
32×32×32	0.0145	0.05805	75.022
64×64×64	0.28	0.6055	53.757
128×128×128	2.839	5.858	51.536

gramming language. The code for the butterfly stage of the 3-D VR DIF algorithm is shown in the Appendix. The computer run-time comparison between the developed algorithm and the RCF approach is carried out on two different systems. The first system has a Pentium 4 (P-4) processor with speed of 2000 MHz and 1024 MB RAM, and the second system has a Pentium-III (P-III) processor with speed of 1000 MHz and 256 MB RAM. The run time in the first system has been calculated using Visual C++ (VC++) Version (6), whereas in the second system, VC++ Version (6) and Cygwin have been used. Therefore, Table II shows the run times on the first system using VC++ compiler. Using the second system, the run times for the 3-D VR and the 3-D RCF algorithms are shown in Table III using VC++ and Cygwin compilers, respectively. The times in Tables II and III represent the average values obtained by repeated execution of the algorithm.

The results of the computer run times confirm the results of the comparison of the arithmetic operations. The 3-D VR algorithm is even better in the comparison of the computer run times. This is due to the fact that the 3-D DCT VR algorithm has a better structure than the RCF approach, and in addition, unlike the RCF approach, it does not involve matrix transpose.

C. Comparison of the 3-D VR Algorithm With Other Existing 3-D DCT Algorithms

A limited number of fast algorithms has been developed to reduce the computational cost of the 3-D and higher dimensions [25]–[29]. In this section, these algorithms are discussed and compared with the 3-D DCT-II VR algorithm presented in this paper.

In [27], the authors have introduced a fast algorithm for the calculation of the n -dimensional DCT by mapping the n -dimensional DCT of size $N_1 \times N_2 \times \dots \times N_n$ ($N_1 = N_2 = \dots = N_n = N$) into N sets of $(n - 1)$ -D DCTs. This algorithm is based on the conversion of the n -dimensional DCT into 1-D DCTs with a reduction in the arithmetic operations. Unlike the developed algorithm in this paper, this algorithm cannot be performed in-place, which leads to large memory requirements and

program sizes. In addition, as stated in [26], the indexing and the post-addition stage become very complicated when ($N > 16$).

In [26], the authors have introduced a polynomial transform (PT)-based algorithm for the calculation of the multidimensional DCT. The main idea of this algorithm is to use the PT to convert the multidimensional DCT into a series of 1-D DCTs directly. Other similar algorithms were introduced in [28] and [29]. Although, these algorithms require fewer multiplications than the 3-D VR algorithm, they are difficult to implement [40] and rarely used in video codecs [24].

In general, these algorithms are based on mapping the multidimensional transform into 1-D cosine and polynomial transforms. They involve fewer multiplications but have irregular computational structure involving more complex indexing and overhead [24], [40]. They do not have the butterfly structure and cannot be calculated in place; therefore, their overall complexity is higher than our algorithm. The main consideration in choosing a fast algorithm is computational and structural complexities. As the technology of computers and DSPs advances, the execution time of arithmetic operations (multiplications and additions) has become very fast, and regular computational structure becomes the most important factor [40].

Therefore, although the proposed 3-D VR algorithm does not achieve the theoretical lower bound on the number of multiplications [46], it has the simplest computational structure among all 3-D DCT algorithms. It can be implemented in place using a single butterfly and possesses the properties of the Couley–Tukey FFT in 3-D. Hence, the 3-D VR presents the best choice for reducing arithmetic operations in the calculation of the 3-D DCT-II while keeping the simple structure that characterizes butterfly style Couley–Tukey-based algorithms.

V. IMPLEMENTATION OF VIDEO COMPRESSION USING THE 3-D DCT-II

In recent years, a considerable amount of research has focused on image and video compression. Compression plays a significant role in signal/image processing and communications. Recently, the 3-D DCT has been proposed for the implementation of a new 3-D compression algorithm called the XYZ-video compression algorithm [9]. This algorithm takes advantage of the statistical behavior of video data both in spatial and temporal domains. This eliminates the need for motion estimation, which is required for the implementation of the MPEG standard. As the authors stated, the heart of this algorithm is the 3-D DCT, and the development of fast and efficient algorithms for the 3-D DCT is essential for the practical use of this technique.

The idea of this algorithm is that it takes a full motion digital video stream and divides it into groups of eight frames. Each group is considered to be a 3-D image, where X and Y are the spatial component, and Z is the temporal component. Each frame in the image is divided into 8×8 blocks, forming $8 \times 8 \times 8$ cubes. Each cube is then independently encoded. The XYZ-video encoder consists of three stages: the 3-D DCT-II, the quantizer, and the entropy encoder. In XYZ decoding, the steps from the encoding process are inverted and implemented in the reverse order, as shown in Fig. 4.

TABLE III
COMPUTER RUN TIMES FOR THE 3-D VR AND THE RCF ALGORITHMS FOR THE CALCULATION
OF THE 3-D DCT-II ON P-III COMPUTER USING CYGWIN AND VC++ COMPILERS

Transform Size	3-D VR algorithm (sec)		RCF (sec)		Saving %	
	Cygwin	VC++	Cygwin	VC++	Cygwin	VC++
2×2×2	22×10 ⁻⁷	275 ×10 ⁻⁸	549 ×10 ⁻⁸	731 ×10 ⁻⁸	59.927	62.38
4×4×4	193 ×10 ⁻⁷	205 ×10 ⁻⁷	604 ×10 ⁻⁷	741 ×10 ⁻⁷	68.046	72.335
8×8×8	192 ×10 ⁻⁶	2 ×10 ⁻⁴	631 ×10 ⁻⁶	746 ×10 ⁻⁶	69.572	73.19
16×16×16	165 ×10 ⁻⁵	19 ×10 ⁻⁴	604 ×10 ⁻⁴	711 ×10 ⁻⁵	72.682	73.277
32×32×32	275 ×10 ⁻⁴	295 ×10 ⁻⁴	633 ×10 ⁻⁴	716 ×10 ⁻⁴	56.556	58.799
64×64×64	0.494	0.5305	0.687	0.741	28.093	28.408
128×128×128	4.697	4.571	6.016	6.6045	21.925	30.79

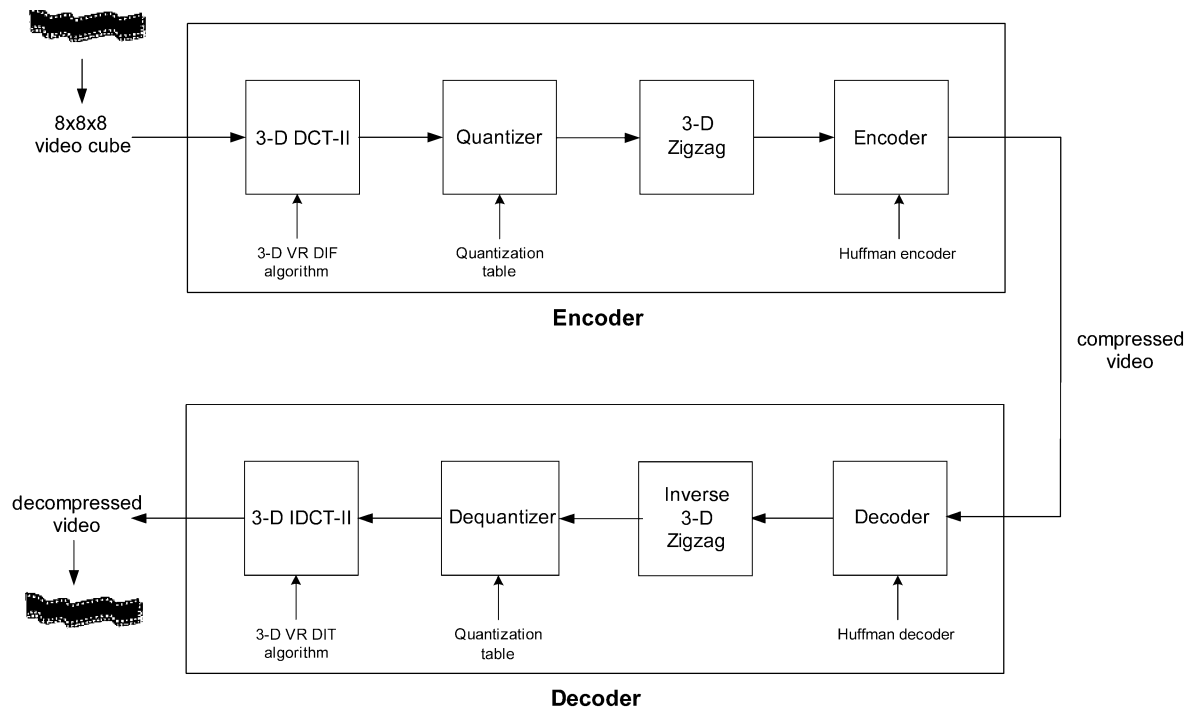


Fig. 4. Encoding and decoding process of the XYZ codec.

The XYZ-video compression algorithm shown in Fig. 4 has been implemented, and the new 3-D DCT VR DIF algorithm has been used to calculate the 3-D DCT-II. The 3-D DCT-II coefficients have been scaled using the quantizer introduced in [36]. After quantization, a 3-D zigzag traversal is performed. The last part of the encoding process is the encoder. In this example, the Huffman coding technique described in [4] has been used.

In the XYZ-decoding, the same steps used for encoding is inverted and then reversed, where first, the Huffman decoder is implemented on the compressed video data, and then, the dequantization is applied. Finally, the inverse 3-D DCT-II or the 3-D DCT-III developed in [45] is used to obtain the decompressed frames, where the 3-D DCT-III VR DIT algorithm has been used to calculate the inverse 3-D DCT-II. Fig. 5(a) shows eight frames, each with size (128 × 128), from a video stream,

whereas Fig. 5(b) shows the same video sequences after compression and decompression using the XYZ algorithm. The related errors between the original frames and the decompressed frames, multiplied by 16 in order to be visible, are shown in Fig. 5(c).

For this example, the compression ratio is 27:1, the normalized root mean square (NRMS) is 0.063, and the peak signal-to-noise ratio (PSNR) is 30.9531 dB. This example has been carried out using VC++ Version (6) running on a P-III computer with speed of 1000 MHz and 256 MB RAM. The timing for the encoding and decoding process of the XYZ-video compression algorithm was 6.76 s. The XYZ-video compression algorithm has also been implemented using the RCF approach for the calculation of the forward and inverse 3-D DCTs-II, where the timing was 11.31 s.

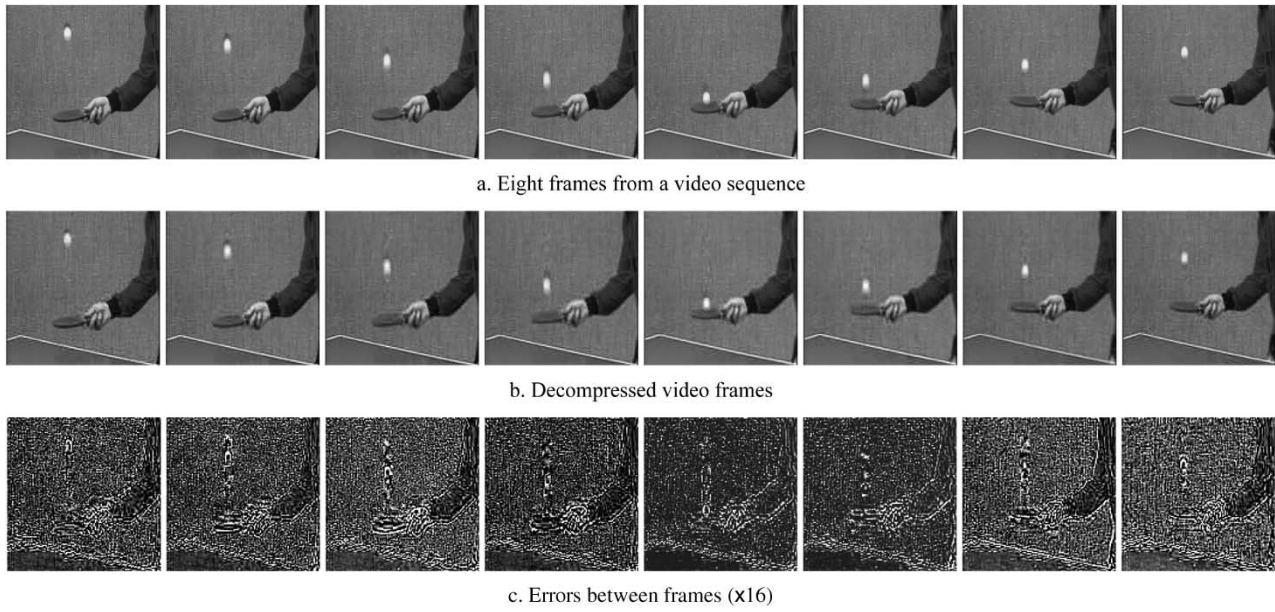


Fig. 5. Coding of video sequence using the XYZ-compression algorithm.

VI. CONCLUSION

In this paper, the 3-D vector-radix DIF algorithm has been introduced for fast calculation of the 3-D DCT-II. The mathematical derivation for this algorithm has been developed, and the total number of multiplications and additions has been calculated, showing that the 3-D DCT-II VR algorithm achieves a significant saving—more than 40%—in the total number of multiplications, as compared with the RCF approach while keeping the number of additions the same. Based on the computer run times, it is also shown that the new 3-D DCT VR algorithm is faster and can save up to 70% in the run time of the $8 \times 8 \times 8$ 3-D DCT-II (a typical size for 3-D image and video compression algorithms). Furthermore, the developed algorithm has a regular structure and can be implemented in-place. Hence, it is suitable for applications where fast computation of the 3-D DCT is required, as in 3-D image and video compression.

For the evaluation of the developed algorithm, the XYZ-video compression algorithm, which uses the 3-D DCT-II, has been implemented, and an example for eight video frames has been presented. The timing results of this application have demonstrated that a significant saving in computer run time can be achieved if the new developed algorithm is used, as compared with the familiar RCF approach.

APPENDIX

Authors: O. Alshibami and S. Boussakta

3-D Vector—radix decimation-in-frequency algorithm

for the 3-D DCT-II.

in[n][n][n] Input array to be transformed

n size of the input array.

m $m = \log_2 n$ (number of stages)

The 3-D DCT is calculated in-place and stored in in[n][n][n]

DCT3D(float in[n][n][n], int n, int m)

```
{
    ns = n << 1;
    ns_2 = n >> 1;
    pi = 3.1415927;
    invr2 = 0.7071067814;
    /* 3-D Reordering Stage */
    /* Butterfly calculation Stage */
    for(stage = 1; stage <= m; stage++)
    {
        ns = ns >> 1;
        ns2 = ns >> 1;
        r2_n = pi/(ns << 1);
        for (i1 = 0; i1 < n; i1 = i1 + ns)
            for (i2 = 0; i2 < n; i2 = i2 + ns)
                for (i3 = 0; i3 < n; i3 = i3 + ns)
                {
                    for (k1 = 0; k1 < ns2; k1++)
                    {
                        s11 = k1 + i1;
                        s13 = s11 + ns2;
                        Cfac = (cos(r2_n * ((k1 << 2) + 1))) << 1;
                        for (k2 = 0; k2 < ns2; k2++)
                        {
                            s21 = k2 + i2;
                            s23 = s21 + ns2;
                            Cfac1 = (cos(r2_n * ((k2 << 2) + 1))) << 1;
                            for (k3 = 0; k3 < ns2; k3++)
                            {
                                s31 = k3 + i3;
                                s33 = s31 + ns2;
                                Cfac2 = (cos(r2_n * ((k3 << 2) + 1))) << 1;
                                p0 = in[s11][s21][s31] + in[s11][s21][s33];
                                p1 = in[s11][s21][s31] - in[s11][s21][s33];
                                p2 = in[s11][s23][s31] + in[s11][s23][s33];
                                p3 = in[s11][s23][s31] - in[s11][s23][s33];
```

```

p4 = in[s13][s21][s31] + in[s13][s21][s33];
p5 = in[s13][s21][s31] - in[s13][s21][s33];
p6 = in[s13][s23][s31] + in[s13][s23][s33];
p7 = in[s13][s23][s31] - in[s13][s23][s33];
p8 = p0 + p2;
p9 = p1 + p3;
p10 = p0 - p2;
p11 = p1 - p3;
p0 = p4 + p6;
p1 = p5 + p7;
p2 = p4 - p6;
p3 = p5 - p7;
Cfacm = Cfac1 * Cfac2;
in[s11][s21][s31] = p8 + p0;
in[s13][s21][s31] = (p8 - p0) * Cfac;
in[s11][s23][s31] = (p10 + p2) * Cfac1;
in[s11][s21][s33] = (p9 + p1) * Cfac2;
in[s13][s23][s31] = (p10 - p2) * Cfac * Cfac1;
in[s13][s21][s33] = (p9 - p1) * Cfac * Cfac2;
in[s11][s23][s33] = (p11 + p3) * Cfacm;
in[s13][s23][s33] = (p11 - p3) * Cfac * Cfacm;
    }
    }
    }
    }
    }
/* 3-D Bit-reverse Stage */
/* Post-additions Stage */
/* Multiplications by transform factor */
in[k1][k2][0] = in[k1][k2][0] * p0;
}

```

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