

Selecting a Set

Peter Stuckey

Choosing from a set of objects

- ▶ Many problems require us to select a **subset** from a set of objects that
 - Meets some criteria; and
 - Optimizes some objective function
- ▶ Example 0-1 knapsack (at most one copy of each object)
 - Limit choices of **x** variable
 - Make the **x** an array of Booleans
 - Use a set variable

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0-1Knapsack Model

```
int: n; % number of objects
set of int: OBJ = 1..n;
int: capacity;
array[OBJ] of int: profit;
array[OBJ] of int: size;

array[OBJ] of var 0..1: x;

constraint forall(i in OBJ) (x[i] >= 0);
constraint sum(i in OBJ) (size[i] * x[i])
    <= capacity;
solve maximize sum(i in OBJ) (profit[i] * x[i]);

output ["x = ", show(x), "\n"];
knapsack01.mzn
```

0-1Knapsack Model

```
int: n; % number of objects
set of int: OBJ = 1..n;
int: capacity;
array[OBJ] of int: profit;
array[OBJ] of int: size;

array[OBJ] of var bool: x;

constraint sum(i in OBJ) (size[i] *
    bool2int(x[i])) <= capacity;
solve maximize sum(i in OBJ)
    (profit[i] * bool2int(x[i]));

output ["x = ", show(x), "\n"];
knapsack01bool.mzn
```

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0-1Knapsack Model

```
int: n; % number of objects
set of int: OBJ = 1..n;
int: capacity;
array[OBJ] of int: profit;
array[OBJ] of int: size;

array[OBJ] of var bool: x;

constraint sum(i in OBJ)(size[i] *
    bool2int(x[i])) <= capacity;
solve maximize sum(i in OBJ)
    (profit[i] * bool2int(x[i]));

output ["x = ", show(x), "\n"];
knapsack01bool.mzn
```

0-1Knapsack Model

```
int: n; % number of objects
set of int: OBJ = 1..n;
int: capacity;
array[OBJ] of int: profit;
array[OBJ] of int: size;

var set of OBJ: x;

constraint sum(i in OBJ)(size[i] *
    bool2int(i in x)) <= capacity;
solve maximize sum(i in OBJ)(profit[i] *
    bool2int(i in x));

output ["x = ", show(x), "\n"];
knapsack01set.mzn
```

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0-1Knapsack Model

```
int: n; % number of objects
set of int: OBJ = 1..n;
int: capacity;
array[OBJ] of int: profit;
array[OBJ] of int: size;

var set of OBJ: x;

constraint sum(i in OBJ)(size[i] *
    bool2int(i in x)) <= capacity;
solve maximize sum(i in OBJ)(profit[i] *
    bool2int(i in x));

output ["x = ", show(x), "\n"];
knapsack01set.mzn
```

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0-1Knapsack Model

```
int: n; % number of objects
set of int: OBJ = 1..n;
int: capacity;
array[OBJ] of int: profit;
array[OBJ] of int: size;

var set of OBJ: x;

constraint sum(i in x)(size[i]) <= capacity;
solve maximize sum(i in x)(profit[i]);

output ["x = ", show(x), "\n"];
knapsack01set_concise.mzn
```

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bool2int

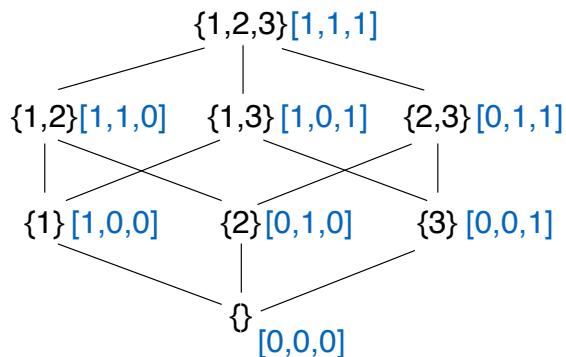
- ▶ Coerce a Boolean to an integer
 - `bool2int(false) = 0`
 - `bool2int(true) = 1`

- ▶ When you use a Boolean where MiniZinc expects an integer, MiniZinc will automatically “add” the `bool2int` coercion
- ▶ Many solvers will use the same internal representation for `x` and `b` where
 - `x = bool2int(b)`

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Set Representations

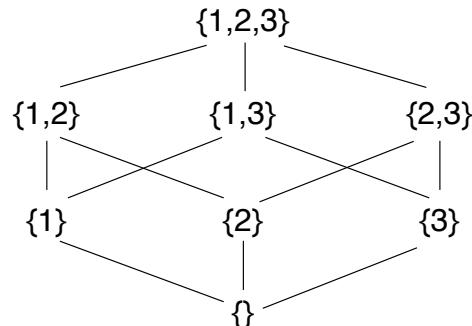
- ▶ Other set representations encode or model the possible values of the set
- ▶ `array[1..3] of var 0..1: x;`



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Set Variables

- ▶ Set variables in MiniZinc choose a set from a given fixed superset
- ▶ `var set of {1, 2, 3}: x;`



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Set Constraints

- ▶ MiniZinc provides the (infix) set operations
 - `in`, (membership e.g. `x in s`)
 - `subset`, `superset`
 - `intersect` (intersection)
 - `union`
 - `card` (cardinality)
 - `diff` (set difference e.g `x diff y = x \ y`)
 - `symdiff` (symmetric difference)
 - e.g. $\{1, 2, 5, 6\} \text{ symdiff } \{2, 3, 4, 5\} = \{1, 3, 4, 6\}$

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Which model is best?

- ▶ Most solvers will treat each model the same
 - CP solvers may treat the last model better since they can combine cardinality reasoning with other set reasoning
- ▶ Model whichever makes it easier to express the constraints
 - for knapsack01 the first version
- ▶ Model using the highest level model
 - the last version

Overview

- ▶ Modeling with sets is common for combinatorial problems
- ▶ There are at least three representations
 - Indicator variables: 0-1 or bool variables
 - Native set variables

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EOF

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Choosing a Set Representation: Example

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SetSelect Question

- ▶ Write a MiniZinc model that given an array of k subsets of numbers $1 \dots n$, chooses a subset of $1 \dots n$ which includes at most one from each subset and maximizes the sum of the chosen set.

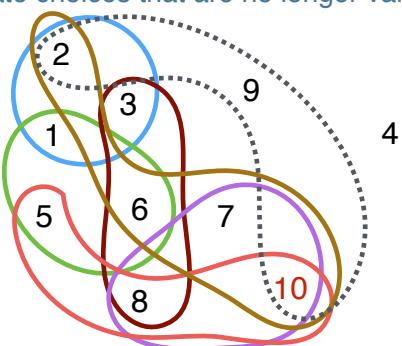
```
n = 10;
k = 7;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselect.dzn
```

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Simple Set Select Algorithm

- ▶ Greedy algorithm

- choose the largest available element
- eliminate choices that are no longer valid



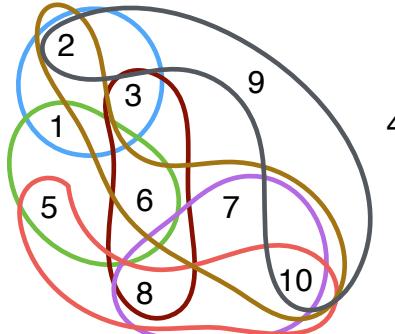
```
n = 10;
k = 7;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselect.dzn
```

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Simple Set Select Algorithm

- ▶ Greedy algorithm

- choose the largest available element
- eliminate choices that are no longer valid



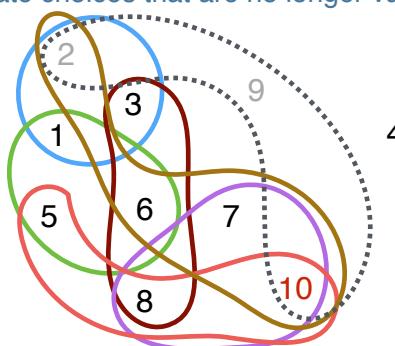
```
n = 10;
k = 7;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselect.dzn
```

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Simple Set Select Algorithm

- ▶ Greedy algorithm

- choose the largest available element
- eliminate choices that are no longer valid



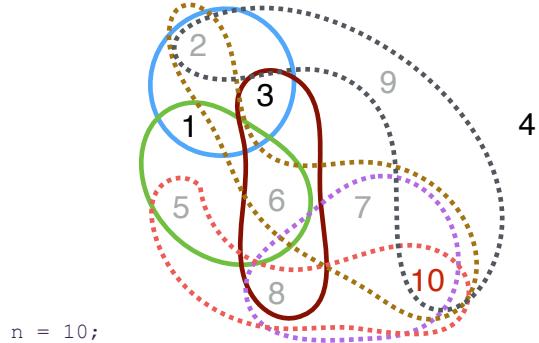
```
n = 10;
k = 7;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselect.dzn
```

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Simple Set Select Algorithm

► Greedy algorithm

- choose the largest available element
- eliminate choices that are no longer valid

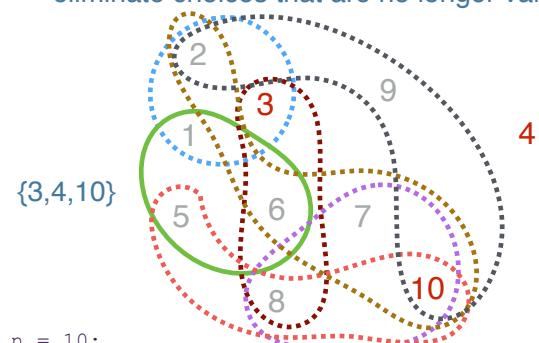


```
6   n = 10;  
    k = 7;  
    s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},  
          {2,9,10}, {5,8,10}, {7,8,10}];
```

Simple Set Select Algorithm

► Greedy algorithm

- choose the largest available element
- eliminate choices that are no longer valid

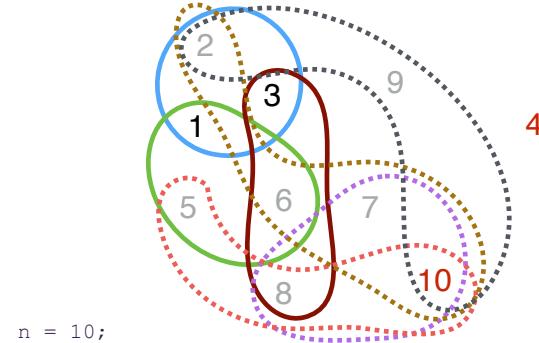


```
8   n = 10;  
    k = 7;  
    s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},  
          {2,9,10}, {5,8,10}, {7,8,10}];
```

Simple Set Select Algorithm

► Greedy algorithm

- choose the largest available element
- eliminate choices that are no longer valid



```
7   n = 10;  
    k = 7;  
    s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},  
          {2,9,10}, {5,8,10}, {7,8,10}];
```

SetSelect Data + Decisions

► Data

```
int; n;  
set of int: OBJ = 1..n;  
int: k;  
set of int: SET = 1..k;  
array[SET] of set of OBJ: s;
```

► Decisions

```
var set of OBJ: x;
```

SetSelect Constraints + Objective

- ▶ At most one intersection

```
forall(i in SET)
    (card(x intersect s[i]) <= 1);
```

- ▶ Objective

```
solve maximize sum(i in x)(i);
```

Solving the model

- ▶ Executing the model

```
$ minizinc setselect.mzn setselect.dzn
```

```
x = {3,4,5,7,9};
-----
=====
```

- ▶ Much better than the greedy algorithm

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SetSelect Revised Question

- ▶ Write a MiniZinc model that given an array of k subsets of numbers $1..n$, chooses a subset of $1..n$ of size u which includes at most one from each subset and maximizes the sum of the chosen set.

```
n = 10;
k = 7;
u = 3;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselectr.dzn
```

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SetSelect Revised Addition

- ▶ Additional constraint: `card(x) = u;`

- ▶ Executing the model

```
$ minizinc setselectr.mzn setselectr.dzn
```

```
x = {4,8,9};
-----
=====
% failures = 237
```

- ▶ But we can model a set of known cardinality differently!

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Choosing a Fixed Cardinality Set

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SetSelect Revised Question

- ▶ Write a MiniZinc model that given an array of k subsets of numbers $1 \dots n$, chooses a subset of $1 \dots n$ of size u which includes at most one from each subset and maximizes the sum of the chosen set.

```
n = 10;
k = 7;
u = 3;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselectrev.dzn
```

Deciding a set of fixed cardinality

- ▶ Instead of a set variable `var set of 1..n: x` with
 - cardinality constraint `card(x) = u`
- ▶ An array of u values
 - `array[1..u] of var 1..n: x`
 - and some other constraints ...
- ▶ Why: suppose $n = 1000$, $u = 4$
- ▶ First representation
 - 1000 Boolean variables
- ▶ Second representation
 - 4 integer variables

Deciding a set of fixed cardinality

- ▶ Consider `array[1..3] of var 1..10: x;`
- ▶ How many possible values
 - 1000
- ▶ And `var set of 1..10: x where card(x) = 3;`
 - $10 * 9 * 8 / 3 * 2 * 1 = 120$
- ▶ **First issue:** some array solutions are **not** sets of cardinality 3
 - e.g. $[1,1,1] = \{1\}$, $[1,2,1] = \{1,2\}$
- ▶ **Solution:** ensure all different
 - `for(i,j in 1..u where i < j)`
 - `(x[i] != x[j]);`

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Critical Issues in Modeling Decisions

- ▶ **Decisions in the problem**
 - are not necessarily the decisions in the model
- ▶ If possible make them identical but
 - what if you are deciding
 - a path in a graph,
 - a tree structure
- ▶ **Critical modeling requirement (1)**
 - decisions in the model (satisfying constraints)
 - are valid decisions in the problem
- ▶ Add constraints to make this so
 - e.g. add constraints forcing array elements \neq

Deciding a set of fixed cardinality

- ▶ Consider `array[1..3] of var 1..10: x;`
 - `for(i,j in 1..u where i < j)`
 - `(x[i] != x[j]);`
- ▶ How many possible values
 - $10 * 9 * 8 = 720$
- ▶ **Second issue:** multiple representations of the same set
 - e.g. $\{1,6,10\} = [1,6,10], [10,1,6], [10,6,1], [1,10,6], [6,10,1], [6,1,10]$
- ▶ **Solution:** ensure ordered
 - `for(i in 1..u-1) (x[i] < x[i+1]);`

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Critical Issues in Modeling Decisions

- ▶ **Multiple decisions in the model**
 - reflect the same decision in the problem
 - e.g. $x = [1,2,7], x = [7,2,1]$ for $x = \{1,2,7\}$
- ▶ **Critical modeling issue (2)**
 - try to have only one set of decisions in the model (satisfying constraints)
 - reflect **valid decisions** in the problem
- ▶ Add constraints to remove all but one set
 - e.g. add constraints forcing array elements $<$

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SetSelect Revised Question

- ▶ Write a MiniZinc model that given an array of k subsets of numbers $1 \dots n$, chooses a subset of $1 \dots n$ of size u which includes at most one from each subset and maximizes the sum of the chosen set.

```
n = 10;
k = 7;
u = 3;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselectrev.dzn
```

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Solving the model

- ▶ Executing the model

```
$ minizinc setselectr2.mzn setselectr.dzn
```

```
x = [5,7,9];
-----
=====
% failures = 22
```

- ▶ This representation makes search easier

SetSelect Revised Model

- ▶ Decisions

```
array[1..u] of var 1..n: x;
for(i in 1..u-1) (x[i] < x[i+1]);
```

- ▶ At most one intersection

```
forall(i in SET)
    (sum(j in 1..u)
        (x[j] in s[i])) % coercion
    <= 1);
```

- ▶ Objective

```
solve maximize sum(x);
```

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Overview

- ▶ There are multiple ways to represent fixed cardinality sets

- var set of OBJ + cardinality constraint
 - good if the solver natively supports sets
 - good when OBJ is not too big
- array[1..u] of var OBJ
 - good when u is small

- ▶ Two critical issues in modelling decisions

- ensure each solution to the model is a solution of the problem
- try to ensure each solution of the problem only has one solution in the model (**symmetry**)

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Choosing a Bounded Cardinality Set

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SetSelect Revised Question

- ▶ Write a MiniZinc model that given an array of k subsets of numbers $1 \dots n$, chooses a subset of $1 \dots n$ of size **at most** u which includes at most one from each subset and maximizes the sum of the chosen set.

```
n = 10;
k = 7;
u = 3;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselectrev.dzn
```

Deciding a set of bounded cardinality

- ▶ Instead of a set variable `var set of OBJ: x` with
 - **cardinality constraint** `card(x) <= u`
- ▶ An array of u values
 - `array[1..u] of var EOBJ: x`
 - **extended OBJ**: `EOBJ = OBJ ∪ { extra-value }`
 - extra value represents: no element
- ▶ For example: `OBJ = 1..n`
 - `EOBJ = 0..n`

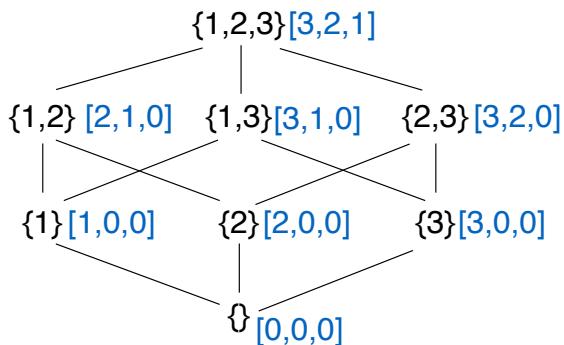
Two critical issues

- ▶ Each solution in the model represents a solution in the problem
 - $[3,0,3]$ ✗ no repeated values
 - $[0,2,0]$ ✓ repeated extra values are OK
- ▶ Each solution in the problem has just one solution representative in the model
 - $[0,2,0], [0,0,2], [2,0,0] = \{2\}$ ✗
 - $[0,1,2], [0,2,1], [1,0,2], [1,2,0], [2,0,1], [2,1,0]$ ✗
- ▶ Add constraints to fix these issues

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Bounded Cardinality

- ▶ Representation of var set of $\{1,2,3\}$: x ;
- ▶ array[1..3] of var 0..3: x ;



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Bounded cardinality constraints

- ▶ Constraints to fix
 - array[1..u] of var 0..n: x
- ▶ Just order the values (decreasing)

```
for(i in 1..u-1) (x[i] > x[i+1]);
```

 - ✗ No representative left for {2} e.g. [2,0,0]

- ▶ No strict ordering

```
for(i in 1..u-1) (x[i] >= x[i+1]);
```

 - ✗ solutions with repeats [3,2,2]

- ▶ Combine the two: repeats of 0 allowed

```
for(i in 1..u-1)
  (x[i] >= bool2int(x[i]!=0)+ x[i+1]);
```

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SetSelect Revised Question

- ▶ Write a MiniZinc model that given an array of k subsets of numbers $1..n$, chooses a subset of $1..n$ of size at most u which includes at most one from each subset and maximizes the sum of the chosen set.

```
n = 10;
k = 7;
u = 3;
s = [{1,5,6}, {2,6,7,10}, {3,6,8}, {1,2,3},
      {2,9,10}, {5,8,10}, {7,8,10}];
setselectrev.dzn
```

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SetSelect Revised Question

► Decisions

```
array[1..u] of var 0..n: x;  
for(i in 1..u-1)  
  (x[i] >= bool2int(x[i]=0)+ x[i+1]);
```

► At most one intersection

```
forall(i in SET)  
  (sum(j in 1..u)  
    (x[j] in s[i]) % 0 is safe  
    <= 1);
```

► Objective

```
solve maximize sum(x); % 0 is safe
```

Overview

- There are multiple ways to represent sets
- var set of OBJ
 - good if the solver natively supports sets
 - good when OBJ is not too big
- array[OBJ] of var bool / 0..1
 - good when OBJ is not too big
- array[1..u] of var OBJ
 - only for fixed cardinality u
 - good when u is small
- array[1..u] of var EOBJ
 - need to represent “null” object

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EOF

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Modeling with MultiSets

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Multisets

- ▶ A **multiset** is a set with duplicates allowed
 - $\text{set } \{3,4,5,4,5,7,4,5,3,3,3\} = \{3,4,5,7\}$, but
 - $\text{multiset } \{\{3,4,5,4,5,7,4,5,3,3,3\}\} \neq \{\{3,4,5,7\}\}$
 - order is not important $= \{\{3,3,3,3,4,4,4,5,5,5,7\}\}$
- ▶ Some problems require choosing a **multiset**
- ▶ Multisets are **not** built in to MiniZinc

Choosing a Multiset

- ▶ Typically modeled analogous to sets

```
array[OBJ] of var int: x;  
forall(i in OBJ) (x[i] >= 0);
```
- ▶ Usually the multiplicity is bounded (by m)

```
array[OBJ] of var 0..m: x;
```
- ▶ If cardinality u is tight we can treat differently

```
array[1..u] of var OBJ0: x;  
forall(i in 1..u-1)  
(x[i] >= x[i+1]);
```

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Choosing a Multiset

- ▶ Knapsack (not 0-1) is an example
- ▶ Production planning is an example
- ▶ Its often not very interesting to think about choosing a multiset
 - **Usually just think about simultaneous choices** of how many of each object

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Overview

- ▶ Multisets also have multiple representations
- ▶ In many discrete optimization problems, the multisite viewpoint is not helpful

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EOF