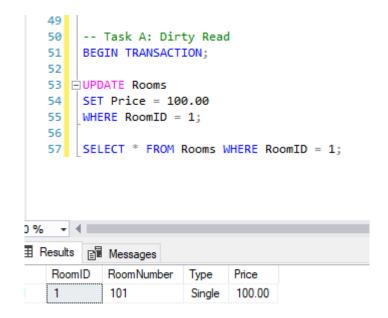
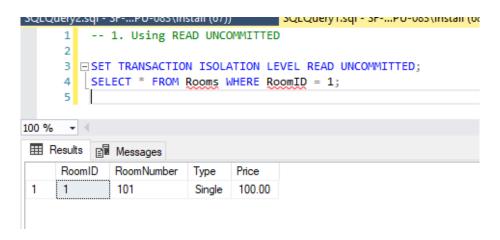
TRANSACTIONS ASSIGNMENT

• Task A: Dirty Read (READ UNCOMMITTED vs READ COMMITTED)

Isolation Level: READ UNCOMMITTED:





Since READ UNCOMMITTED allows dirty reads, **Session 2 sees the updated price** even though it is uncommitted.

Isolation Level: READ COMMITTED:

```
SQLQuery2.sql - SF-...PU-083\install (67))*

SQLQuery1.sql - SF-...I (68)) Executing...*

-- 2 Using READ COMMITED

SET TRANSACTION ISOLATION LEVEL READ COMMITTED;

SELECT * FROM Rooms WHERE ROOMID = 1;

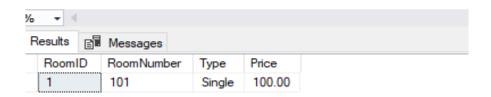
Messages

Messages
```

If **Session 1** has an **open transaction (not committed or rolled back)**, **Session 2 must wait** until the transaction is completed. READ COMMITTED in **Session 2** waits for the exclusive lock to be released before reading.

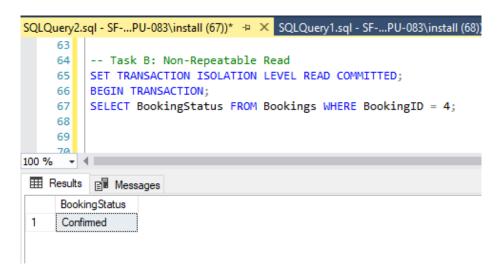
After commit-

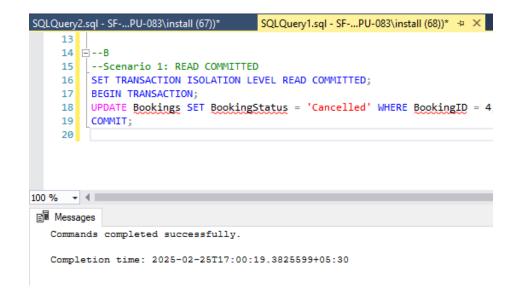
```
2
3 -- 2 Using READ COMMITED
4 □ SET TRANSACTION ISOLATION LEVEL READ COMMITTED;
5 □ SELECT * FROM Rooms WHERE RoomID = 1;
6
```

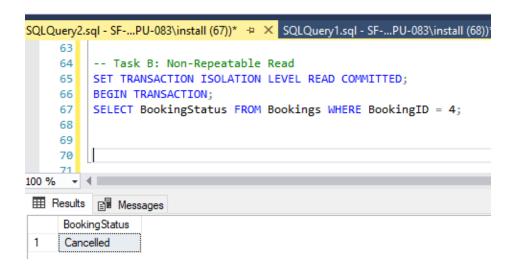


• Task B: Non-Repeatable Read (READ COMMITTED vs REPEATABLE READ)

Isolation Level: READ COMMITTED:







The **READ COMMITTED** isolation level ensures that **Session 1** can read **committed data**. Once **Session 2** commits the change (status update to 'Cancelled'), **Session 1** will see the new status immediately in its next read.

Isolation Level: REPEATABLE READ:

```
| ▶ Execute ■ ✓ XX 🗇 🔒 | X2 X2 🚇 | 🕮 🕮 🖺 | 🧵 🛂 | ಶ 🛫
QLQuery2.sql - SF-...PU-083\install (67))*
                                    SQLQuery1.sql - SF-...l (68)) Executing...* 💠 🗶
   21
   22
        -- Session 2: REPEATABLE READ
        SET TRANSACTION ISOLATION LEVEL REPEATABLE READ;
   23
        BEGIN TRANSACTION;
        UPDATE Bookings SET BookingStatus = 'Cancelled' WHERE RoomID = 4;
   25
   26
   27
   28
   29
00 % → 4 ■
Results Messages
```

The reason **Session 2**'s update statement is still running (blocked) is due to the fact that **REPEATABLE READ** isolation level ensures **consistent reads** by locking the data for the duration of **Session 1**'s transaction. Once **Session 1** commits, **Session 2** will be able to acquire the lock and proceed with its update.

After commit in session 1-

```
QLQuery2.sql - SF-...l (67)) Executing...*
                                     SQLQuery1.sql - SF-...PU-083\install (68))* + X
    21
    22
          -- Session 2: REPEATABLE READ
         SET TRANSACTION ISOLATION LEVEL REPEATABLE READ;
    23
         BEGIN TRANSACTION;
          UPDATE Bookings SET BookingStatus = 'Cancelled' WHERE RoomID = 4;
    25
    26
         COMMIT;
    27
    28
    29
00 %

    Messages

  Commands completed successfully.
  Completion time: 2025-02-25T17:15:22.4013492+05:30
```

• Task C: Phantom Read (REPEATABLE READ vs SERIALIZABLE)

Isolation Level: REPEATABLE READ:

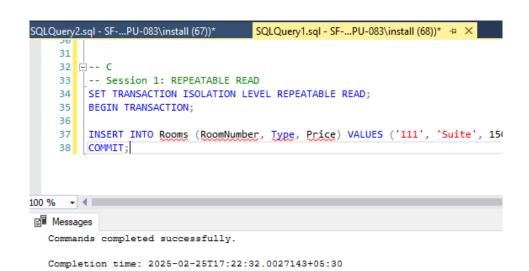
```
SET TRANSACTION ISOLATION LEVEL REPEATABLE READ;
BEGIN TRANSACTION;

SELECT COUNT(*) FROM Rooms WHERE Type = 'Suite';

100 % 
Results Messages

(No column name)

1 3
```



Isolation Level: SERIALIZABLE:

```
SQLQuery1.sql - SF-...PU-083\install (68))* 

SQLQuery1.sql - SF-...
```

Session 2 is blocked from inserting the new row into the 'Suite' category because the SERIALIZABLE lock held by Session 1 prevents any changes to rows that could affect the result of the query in Session 1.