(1) What is time complexity of below code & how? bold for ( int my 2 int j=1; int i=0; while (i(n) of 1=1 i=1+2 m-level 1=2 1=1+2+3 1=3 for (i) m(m+1) < mby summation method 2 1 = 1+1+ ..... + In times |T(n) = [m

(99) With recurred relation for function that prints Fibonacci series Solve it to get the time completely what will be open complicity & inly? tor filonoxii drleg f(0) = 0 f(n) = f(n-1) + f(n-3)f(1)=1 by forming tree f(n-2) m-levels f(n-3) f(n-3) f(n. w) " At every function call we get 2 function calls . for an deads We have = 2x2 --- n time T(m) = 2m Maximum Space-Conding Remoin No of calls marinen: n for each cell we have space complinity O(1) (T(n) = O(n) littat considering Recursive stack: lact cell we have time complexity O(1) T(n) = 0(1)

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(93) Write a program which have complexity
            n (loga), no log (loga)
   () nlogn
             void internain 1) of
                     inti,j; int n;
                    for (i=0; i(n, i++)
                     for (j=1; j < m; j*=2)
                        printf ("*");
2.) n3
           baid
                main () {
```

int u, j, K; int n; for (i=0, i<n, i++) for (j=0; j <n; j++) for (KO; KCn; K++) print f ( '+");

3) log (logn) interior void main () of int count: 0; inti; for (1:2, i(m, i=ixi) of country;

Qu.) Solve the following recurrence relation

$$T(n) = T(n/4) + T(n/2) + cn^2$$

$$T(n/4) = T(n/4) + T(n/3) - 1$$

$$T(n/8) = T(n/6) + T(n/4) + T(n/8) - 2$$
At level  $\frac{1}{2}$ 

$$\frac{1}{4} + \frac{n^2}{2^2} + \frac{n^2}{4^2} + \frac{n^2}{8^2} = \left(\frac{5}{16}\right)^2 n^2 c$$

$$\lim_{n \to \infty} \frac{1}{8^3} + \frac{n^2}{16^2} + \frac{n^2}{4^2} + \frac{n^2}{8^2} = \left(\frac{5}{16}\right)^2 n^2 c$$

$$\lim_{n \to \infty} \frac{1}{8^3} + \frac{n^2}{16^3} + \frac{n^2}{16^3} + \frac{1}{16^3} + \frac{1}{16^3}$$

$$T(n) = c^{2}n^{2} \left[ 1 + \frac{5}{16} + \left( \frac{5}{16} \right)^{2} + \cdots + \left( \frac{5}{16} \right)^{2} \right]$$

$$= c^{2}n^{2} \times 1 \times \left( \frac{1 - \left( \frac{5}{16} \right)^{2} + \cdots + \left( \frac{5}{16} \right)^{2} \right)^{2}}{1 - \frac{5}{16}} \right)$$

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Q5) What so time complexity of following func ()? int fun (int n) ( for ( int i=1 ; i <= n ; i++) { for (int j:1; j Kon j j+=i) & 11 Some O(1) task for i j= (n-1)/c times 1+3+5 1+5+9 5 (m-1) (n-1) = (n-1) + (n-1) + (n-1) + (n-1)T(n)= n[1+ \frac{1}{2} + \frac{1}{3} + --- + \frac{1}{2} - 1[1+ \frac{1}{2} + \frac{1}{3} + --- + \frac{1}{n}] = n logn - logn T(n) = O(nlog n)Q6) what should be time complimity of for (int i=2; i = n; i = pow (i,K))

of 11 Some O(1) 3 Where K is constant. Jor i  $g^{k}$   $g^{k^{2}}$   $g^{k^{3}}$   $g^{k^{3}}$ 

Q7) Write a recurrence relation when quick port repeatedly divides arrows into 2 parts of 99%. & 1%. Derive love complexity in this case. Show the recurrence tree while deriving time complexity & find difference in heights of both extrems parts. What do you understand by this analysis?

Given algorithm divide away in 97% & 11 part.

T(n) = T(n-1) + O(1)

m-levila

m-2 m-2

Lowest height = 2 highest height = n o's difference = n-2

The given algorithm produces linear result

- Q8) Arrange following in increasing order of rate of growth

  a) n, n!, log n, log log n, Tn, log n!, nlog n, log 2, 2, 22, 100 < log logn < logn < (logn)<sup>2</sup> < Tn < n < nlogn < logn! (n<sup>2</sup> < 2<sup>n</sup> < 4<sup>n</sup> < 2<sup>2</sup>
  - b) 2(2<sup>n</sup>), 4n, 2n, 1, logn, log (logn), Tegn, log 2n, 2logn, 2, log(n!), n!, n2, nlog n  $1 < \log \log n < \log n < \log n < \log n < 2\log n < n < n < 2n < 1$   $4n < \log(n!) < n^2 < n! < 2^2$
  - c)  $9^{2n}$ ,  $\log_{2}n$ ,  $n\log_{2}n$ ,  $n\log_{2}n$ ,  $\log_{2}n$ ,  $\log_{3}n$ ,