

# Spear of Longinus

## Shanghai Jiao Tong University

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## 数论

## 扩展欧几里德算法

```

1 LL exgcd(LL a, LL b, LL &x, LL &y){
2     if(!b){
3         x=1; y=0; return a;
4     }else{
5         LL d=exgcd(b, a%b, x, y);
6         LL t=x; x=y; y=t-a/b*y;
7         return d;
8     }
9 }

```

## 中国剩余定理

```

1 LL china(int n, int *a, int *m){
2     LL M=1, d, x=0, y;
3     for(int i=0; i<n; i++){
4         M*=m[i];
5         for(int i=0; i<n; i++){
6             LL w=M/m[i];
7             d=exgcd(m[i], w, d, y);
8             y=(y*M+a[i])%M;
9             x=(x+y*w*M*a[i])%M;
10        }
11        while(x<0) x+=M;
12        return x;
13 }

```

## 中国剩余定理 2

```

1 //merge Ax=B and ax=b to A'x=B'
2 void merge(LL &A, LL &B, LL a, LL b){
3     LL x, y;
4     sol(A, -a, b-B, x, y);
5     A=lcm(A, a);
6     B=(a*y+b)%A;
7     B=(B+A)%A;
8 }

```

## 扩展小步大步

```

1 LL solve2(LL a, LL b, LL p){
2     //a^x=b (mod p)
3     b%=p;
4     LL e=1%p;
5     for(int i=0; i<100; i++){
6         if(e==b) return i;
7         e=e*a%p;
8     }
9     int r=0;
10    while(gcd(a, p)!=1){
11        LL d=gcd(a, p);
12        if(b%d) return -1;
13        p/=d; b/=d; b=b*inv(a/d, p);
14        r++;
15    } LL res=BSGS(a, b, p);
16    if(res==-1) return -1;
17    return res+r;
18 }

```

## 卢卡斯定理

```

1 LL Lucas(LL n, LL m, LL p){
2     LL ans=1;
3     while(n&& m){

```

```

4         LL a=n%p, b=m%p;
5         if(a<b) return 0;
6         ans=(ans*C(a, b, p))%p;
7         n/=p; m/=p;
8     } return ans%p;
9 }

```

## Miller Rabin 素数测试

```

1 const int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
2 bool check(long long n, int base) {
3     long long n2=n-1, res;
4     int s=0;
5     while(n2%2==0) n2>>=1, s++;
6     res=pw(base, n2, n);
7     if((res==1) || (res==n-1)) return 1;
8     while(s--){
9         res=mul(res, res, n);
10        if(res==n-1) return 1;
11    }
12    return 0; // n is not a strong pseudo prime
13 }
14 bool isprime(const long long &n) {
15     if(n==2) return true;
16     if(n<2 || n%2==0) return false;
17     for(int i=0; i<12&&BASE[i]<n; i++){
18         if(!check(n, BASE[i])) return false;
19     }
20     return true;
21 }
22 }
23 }
24 }

```

## Pollard Rho 大数分解

时间复杂度： $\mathcal{O}(n^{1/4})$

```

1 LL prho(LL n, LL c){
2     LL i=1, k=2, x=rand()%(n-1)+1, y=x;
3     while(1){
4         i++; x=(x*x%n+c)%n;
5         LL d=__gcd((y-x+n)%n, n);
6         if(d>1&&d<n) return d;
7         if(y==x) return n;
8         if(i==k) y=x, k<<=1;
9     }
10 }
11 void factor(LL n, vector<LL>&fat){
12     if(n==1) return;
13     if(isprime(n)){
14         fat.push_back(n);
15         return;
16     } LL p=n;
17     while(p>=n) p=prho(p, rand()%(n-1)+1);
18     factor(p, fat);
19     factor(n/p, fat);
20 }

```

## 快速数论变换 (zky)

返回结果：

$$c_i = \sum_{0 \leq j \leq i} a_j \cdot b_{i-j} \pmod{p} \quad (0 \leq i < n)$$

使用说明： $magic$  是  $mod$  的原根

时间复杂度： $\mathcal{O}(n \log n)$

```

1 /*
2 {(mod,G)}={(81788929,7),(101711873,3),(167772161,3)
3             ,(377487361,7),(998244353,3),(1224736769,3)
4             ,(1300234241,3),(1484783617,5)}
5 */
6 int mo=998244353,G=3;
7 void NTT(int a[],int n,int f){
8     for(register int i=0;i<n;i++){
9         if(i<rev[i])
10            swap(a[i],a[rev[i]]);
11     for (register int i=2;i<=n;i<=1){
12         static int exp[maxn];
13         exp[0]=1;exp[1]=pw(G,(mo-1)/i);
14         if(f==-1)exp[1]=pw(exp[1],mo-2);
15         for(register int k=2;k<(i>>1);k++){
16             exp[k]=1LL*exp[k-1]*exp[1]%mo;
17         for(register int j=0;j<n;j+=i){
18             for(register int k=0;k<(i>>1);k++){
19                 register int &pA=a[j+k],&pB=a[j+k+(i>>1)];
20                 register int A=pA,B=1LL*pB*exp[k]%mo;
21                 pA=(A+B)%mo;
22                 pB=(A-B+mo)%mo;
23             }
24         }
25     }
26     if(f==1){
27         int rv=pw(n,mo-2)%mo;
28         for(int i=0;i<n;i++)
29             a[i]=1LL*a[i]*rv%mo;
30     }
31 }
32 void mul(int m,int a[],int b[],int c[]){
33     int n=1,len=0;
34     while(n<m)n<=1,len++;
35     for (int i=1;i<n;i++){
36         rev[i]=(rev[i>>1]>>1)|((i&1)<<(len-1));
37     NTT(a,n,1);
38     NTT(b,n,1);
39     for(int i=0;i<n;i++){
40         c[i]=1LL*a[i]*b[i]%mo;
41     NTT(c,n,-1);
42 }

```

### 原根

```

1 vector<LL>fct;
2 bool check(LL x,LL g){
3     for(int i=0;i<fct.size();i++){
4         if(pw(g,(x-1)/fct[i],x)==1)
5             return 0;
6     return 1;
7 }
8 LL findrt(LL x){
9     LL tmp=x-1;
10    for(int i=2;i*i<=tmp;i++){
11        if(tmp%i==0){
12            fct.push_back(i);
13            while(tmp%i==0)tmp/=i;
14        }
15    }if(tmp>1)fct.push_back(tmp);
16    // x is 1,2,4,p^n,2p^n
17    // x has phi(phi(x)) primitive roots
18    for(int i=2;i<int(1e9);i++){if(check(x,i))
19        return i;
20    return -1;
21 }
22 const int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};

```

```

23 bool check(long long n,int base) {
24     long long n2=n-1,res;
25     int s=0;
26     while(n2%2==0) n2>>=1,s++;
27     res=pw(base,n2,n);
28     if((res==1)|| (res==n-1)) return 1;
29     while(s--){
30         res=mul(res,res,n);
31         if(res==n-1) return 1;
32     }
33     return 0; // n is not a strong pseudo prime
34 }
35 bool isprime(const long long &n) {
36     if(n==2)
37         return true;
38     if(n<2 || n%2==0)
39         return false;
40     for(int i=0;i<12&&BASE[i]<n;i++){
41         if(!check(n,BASE[i]))
42             return false;
43     }
44     return true;
45 }

```

### 线性递推

```

1 //已知  $a_0, a_1, \dots, a_{m-1} \setminus \setminus$ 
2  $a_n = c_0 * a_{n-m} + \dots + c_{m-1} * a_{n-1} \setminus \setminus$ 
3 求  $a_n = v_0 * a_0 + v_1 * a_1 + \dots + v_{m-1} * a_{m-1} \setminus \setminus$ 
4
5 void linear_recurrence(long long n, int m, int a[], int c[], int p) {
6     long long v[M] = {1 % p}, u[M << 1], msk = !n;
7     for(long long i(n); i > 1; i >= 1) {
8         msk <<= 1;
9     }
10    for(long long x(0); msk; msk >>= 1, x <= 1) {
11        fill_n(u, m << 1, 0);
12        int b(!(n & msk));
13        x |= b;
14        if(x < m) {
15            u[x] = 1 % p;
16        }else {
17            for(int i(0); i < m; i++) {
18                for(int j(0), t(i + b); j < m; j++, t++) {
19                    u[t] = (u[t] + v[i] * v[j]) % p;
20                }
21            }
22            for(int i((m << 1) - 1); i >= m; i--) {
23                for(int j(0), t(i - m); j < m; j++, t++) {
24                    u[t] = (u[t] + c[j] * u[i]) % p;
25                }
26            }
27        }
28        copy(u, u + m, v);
29    }
30    //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
31    for(int i(m); i < 2 * m; i++) {
32        a[i] = 0;
33        for(int j(0); j < m; j++) {
34            a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
35        }
36    }
37    for(int j(0); j < m; j++) {
38        b[j] = 0;
39        for(int i(0); i < m; i++) {
40            b[j] = (b[j] + v[i] * a[i + j]) % p;
41        }

```

```

42 }
43 for(int j(0); j < m; j++) {
44     a[j] = b[j];
45 }
46 }

```

直线下整点个数

返回结果：

$$\sum_{0 \leq i < n} \lfloor \frac{a+b \cdot i}{m} \rfloor$$

```

1 //calc \sum_{i=0}^{n-1} [(a+bi)/m]
2 // n,a,b,m >0
3 LL solve(LL n,LL a,LL b,LL m){
4     if(b==0)
5         return n*(a/m);
6     if(a>=m || b>=m)
7         return n*(a/m)+(n-1)*n/2*(b/m)+solve(n,a%m,b%m,m);
8     return solve((a+b*n)/m,(a+b*n)%m,m,b);
9 }

```

数值

高斯消元

```

1 void Gauss(){
2     int r,k;
3     for(int i=0;i<n;i++){
4         r=i;
5         for(int j=i+1;j<n;j++){
6             if(fabs(A[j][i])>fabs(A[r][i]))r=j;
7             if(r!=i)for(int j=0;j<n;j++)swap(A[i][j],A[r][j]);
8             for(int k=i+1;k<n;k++){
9                 double f=A[k][i]/A[i][i];
10                for(int j=i;j<n;j++)A[k][j]-=f*A[i][j];
11            }
12        }
13        for(int i=n-1;i>=0;i--){
14            for(int j=i+1;j<n;j++){
15                A[i][n]-=A[j][n]*A[i][j];
16            }
17            A[i][n]/=A[i][i];
18        }
19        cout<<fixed<<setprecision(3)<<A[i][n]<<" ";
20        cout<<fixed<<setprecision(3)<<A[n-1][n];
21    }
22    bool Gauss(){
23        for(int i=1;i<=n;i++){
24            int r=0;
25            for(int j=i;j<=m;j++){
26                if(a[j][i]){r=j;break;}
27            }
28            if(!r)return 0;
29            ans=max(ans,r);
30            swap(a[i],a[r]);
31            for(int j=i+1;j<=m;j++){
32                if(a[j][i])a[j]^=a[i];
33            }
34            for(int i=n;i>=1;i--){
35                for(int j=i+1;j<=n;j++)if(a[i][j])
36                    a[i][n+1]=a[i][n+1]^a[j][n+1];
37            }
38            return 1;
39        }
40    }
41    LL Gauss(){
42        for(int i=0;i<n;i++)for(int j=0;j<n;j++)A[i][j]%m;
43        for(int i=0;i<n;i++)for(int j=0;j<n;j++)A[i][j]=(A[i][j]+m)%m;
44        LL ans=n%2?-1:1;
45        for(int i=0;i<n;i++){

```

```

42     for(int j=i+1;j<n;j++){
43         while(A[j][i]){
44             LL t=A[i][i]/A[j][i];
45             for(int k=0;k<n;k++)
46                 A[i][k]=(A[i][k]-A[j][k]*t%m+m)%m;
47             swap(A[i],A[j]);
48             ans=-ans;
49         }
50     }ans=ans*A[i][i]%m;
51     return (ans%m+m)%m;
52 }
53 int Gauss(){//求秩
54     int r,now=-1;
55     int ans=0;
56     for(int i = 0; i < n; i++){
57         r = now + 1;
58         for(int j = now + 1; j < m; j++)
59             if(fabs(A[j][i]) > fabs(A[r][i]))
60                 r = j;
61         if (!sgn(A[r][i])) continue;
62         ans++;
63         now++;
64         if(r != now)
65             for(int j = 0; j < n; j++)
66                 swap(A[r][j], A[now][j]);
67
68         for(int k = now + 1; k < m; k++){
69             double t = A[k][i] / A[now][i];
70             for(int j = 0; j < n; j++){
71                 A[k][j] -= t * A[now][j];
72             }
73         }
74     }
75     return ans;
76 }

```

快速傅立叶变换

返回结果：

$$c_i = \sum_{0 \leq j \leq i} a_j \cdot b_{i-j} \quad (0 \leq i < n)$$

时间复杂度： $\mathcal{O}(n \log n)$ 

```

1 typedef complex<double> cp;
2 const double pi = acos(-1);
3 void FFT(vector<cp>&num,int len,int ty){
4     for(int i=1,j=0;i<len-1;i++){
5         for(int k=len;j^=k>=1,~j&k;){
6             if(i<j)
7                 swap(num[i],num[j]);
8         }
9     }
10    for(int h=0;(1<<h)<len;h++){
11        int step=1<<h,step2=step<<1;
12        cp w0(cos(2.0*pi/step2),ty*sin(2.0*pi/step2));
13        for(int i=0;i<len;i+=step2){
14            cp w(1,0);
15            for(int j=0;j<step;j++){
16                cp &x=num[i+j+step];
17                cp &y=num[i+j];
18                cp d=w*x;
19                x=y-d;
20                y=y+d;
21                w=w*w0;
22            }
23        }

```

```

24     if(ty==1)
25         for(int i=0;i<len;i++)
26             num[i]=cp(num[i].real()/(double)len,num[i].imag());
27 }
28 vector<cp> mul(vector<cp>a,vector<cp>b){
29     int len=a.size()+b.size();
30     while((len&-len)!=len)len++;
31     while(a.size()<len)a.push_back(cp(0,0));
32     while(b.size()<len)b.push_back(cp(0,0));
33     FFT(a,len,1);
34     FFT(b,len,1);
35     vector<cp>ans(len);
36     for(int i=0;i<len;i++){
37         ans[i]=a[i]*b[i];
38     }
39     FFT(ans,len,-1);
40     return ans;
}

```

### 单纯形法求解线性规划

返回结果：

$$\max\{c_1 \times m + x_{m+1} \mid x_{m+1} \geq 0, a_n \times m + x_{m+1} \leq b_{n+1}\}$$

```

1 namespace LP{
2     const int maxn=233;
3     double a[maxn][maxn];
4     int Ans[maxn],pt[maxn];
5     int n,m;
6     void pivot(int l,int i){
7         double t;
8         swap(Ans[l+n],Ans[i]);
9         t=-a[l][i];
10        a[l][i]=-1;
11        for(int j=0;j<=n;j++)a[l][j]/=t;
12        for(int j=0;j<=m;j++){
13            if(a[j][i]&&j!=l){
14                t=a[j][i];
15                a[j][i]=0;
16                for(int k=0;k<=n;k++)a[j][k]+=t*a[l][k];
17            }
18        }
19    }
20    vector<double> solve(vector<vector<double>>
21        ↪>A,vector<double>B,vector<double>C){
22        n=C.size();
23        m=B.size();
24        for(int i=0;i<C.size();i++)
25            a[0][i+i]=C[i];
26        for(int i=0;i<B.size();i++)
27            a[i+1][0]=B[i];
28
29        for(int i=0;i<m;i++){
30            for(int j=0;j<n;j++){
31                a[i+1][j+1]=-A[i][j];
32            }
33        }
34        for(int i=1;i<=n;i++)Ans[i]=i;
35
36        double t;
37        for(;;){
38            int l=0;t=-eps;
39            for(int j=1;j<=m;j++)if(a[j][0]<t)t=a[j][0];
40            if(!l)break;
41            int i=0;
42            for(int j=1;j<=n;j++)if(a[l][j]>eps){i=j;break;}
43            if(!i){
44                puts("Infeasible");
45                return vector<double>();
46            }
47        }
48    }
49 }

```

```

45     pivot(l,i);
46 }
47 for(;;){
48     int i=0;t=eps;
49     for(int j=1;j<=n;j++)if(a[0][j]>t)t=a[0][j];
50     if(!i)break;
51     int l=0;
52     t=1e30;
53     for(int j=1;j<=m;j++)if(a[j][i]<=-eps){
54         double tmp;
55         tmp=-a[j][0]/a[j][i];
56         if(t>tmp)t=tmp,l=j;
57     }
58     if(!l){
59         puts("Unbounded");
60         return vector<double>();
61     }
62     pivot(l,i);
63 }
64 vector<double>x;
65 for(int i=n+1;i<=n+m;i++)pt[Ans[i]]=i-n;
66 for(int i=1;i<=n;i++)x.push_back(pt[i]?a[pt[i]][0]:0);
67 return x;
68 }
69 }

```

### 自适应辛普森

```

1 double area(const double &left, const double &right) {
2     double mid = (left + right) / 2;
3     return (right - left) * (calc(left) + 4 * calc(mid) + calc(right)) / 6;
4 }
5
6 double simpson(const double &left, const double &right,
7     const double &eps, const double &area_sum) {
8     double mid = (left + right) / 2;
9     double area_left = area(left, mid);
10    double area_right = area(mid, right);
11    double area_total = area_left + area_right;
12    if (std::abs(area_total - area_sum) < 15 * eps) {
13        return area_total + (area_total - area_sum) / 15;
14    }
15    return simpson(left, mid, eps / 2, area_left)
16        + simpson(mid, right, eps / 2, area_right);
17 }
18
19 double simpson(const double &left, const double &right, const double &eps) {
20     return simpson(left, right, eps, area(left, right));
21 }

```

### 多项式求根

```

1 const double eps=1e-12;
2 double a[10][10];
3 typedef vector<double> vd;
4 int sgn(double x) { return x < -eps ? -1 : x > eps; }
5 double mypow(double x,int num){
6     double ans=1.0;
7     for(int i=1;i<=num;++i)ans*=x;
8     return ans;
9 }
10 double f(int n,double x){
11     double ans=0;
12     for(int i=n;i>=0;--i)ans+=a[n][i]*mypow(x,i);
13     return ans;
14 }
15 double getRoot(int n,double l,double r){

```

```

16     if(sgn(f(n,l))==0)return l;
17     if(sgn(f(n,r))==0)return r;
18     double temp;
19     if(sgn(f(n,l))>0)temp=-1;else temp=1;
20     double m;
21     for(int i=1;i<=10000;++i){
22         m=(l+r)/2;
23         double mid=f(n,m);
24         if(sgn(mid)==0){
25             return m;
26         }
27         if(mid*temp<0)l=m;else r=m;
28     }
29     return (l+r)/2;
30 }
31 vd did(int n){
32     vd ret;
33     if(n==1){
34         ret.push_back(-1e10);
35         ret.push_back(-a[n][0]/a[n][1]);
36         ret.push_back(1e10);
37         return ret;
38     }
39     vd mid=did(n-1);
40     ret.push_back(-1e10);
41     for(int i=0;i+1<mid.size();++i){
42         int t1=sgn(f(n,mid[i])),t2=sgn(f(n,mid[i+1]));
43         if(t1*t2>0)continue;
44         ret.push_back(getRoot(n,mid[i],mid[i+1]));
45     }
46     ret.push_back(1e10);
47     return ret;
48 }
49 int main(){
50     int n; scanf("%d",&n);
51     for(int i=n;i>=0;--i){
52         scanf("%lf",&a[n][i]);
53     }
54     for(int i=n-1;i>=0;--i)
55         for(int j=0;j<=i;++j)a[i][j]=a[i+1][j+1]*(j+1);
56     vd ans=did(n);
57     sort(ans.begin(),ans.end());
58     for(int i=1;i+1<ans.size();++i)printf("%.10f\n",ans[i]);
59     return 0;
60 }

```

## 数据结构

### 平衡的二叉查找树

#### Treap

```

1 #include<bits/stdc++.h>
2 using namespace std;
3 const int maxn=1e5+5;
4 #define sz(x) (x?x->siz:0)
5 struct Treap{
6     struct node{
7         int key,val;
8         int siz,s;
9         node *c[2];
10        node(int v=0){
11            val=v;
12            key=rand();
13            siz=1,s=1;
14            c[0]=c[1]=0;
15        }
16        void rz(){siz=s;if(c[0])siz+=c[0]->siz;if(c[1])siz+=c[1]->siz;}
17    }pool[maxn],*cur,*root;
18    Treap(){cur=pool;}

```

```

19     node* newnode(int val){return *cur=node(val),cur++;}
20     void rot(node *&t,int d){
21         if(!t->c[d])t=t->c[!d];
22         else{
23             node *p=t->c[d];t->c[d]=p->c[!d];
24             p->c[!d]=t;t->rz();p->rz();t=p;
25         }
26     }
27     void insert(node *&t,int x){
28         if(!t){t=newnode(x);return;}
29         if(t->val==x){t->s++;t->siz++;return;}
30         insert(t->c[x>t->val],x);
31         if(t->key<t->c[x>t->val]->key)
32             rot(t,x>t->val);
33         else t->rz();
34     }
35     void del(node *&t,int x){
36         if(!t)return;
37         if(t->val==x){
38             if(t->s>1){t->s--;t->siz--;return;}
39             if(!t->c[0]||!t->c[1]){
40                 if(!t->c[0])t=t->c[1];
41                 else t=t->c[0];
42                 return;
43             }
44             int d=t->c[0]->key<t->c[1]->key;
45             rot(t,d);
46             del(t,x);
47             return;
48         }
49         del(t->c[x>t->val],x);
50         t->rz();
51     }
52     int pre(node *t,int x){
53         if(!t)return INT_MIN;
54         int ans=pre(t->c[x>t->val],x);
55         if(t->val<x)ans=max(ans,t->val);
56         return ans;
57     }
58     int nxt(node *t,int x){
59         if(!t)return INT_MAX;
60         int ans=nxt(t->c[x>t->val],x);
61         if(t->val>x)ans=min(ans,t->val);
62         return ans;
63     }
64     int rank(node *t,int x){
65         if(!t)return 0;
66         if(t->val==x)return sz(t->c[0]);
67         if(t->val<x)return sz(t->c[0])+t->s+rank(t->c[1],x);
68         if(t->val>x)return rank(t->c[0],x);
69     }
70     int kth(node *t,int x){
71         if(sz(t->c[0])>=x)return kth(t->c[0],x);
72         if(sz(t->c[0])+t->s>=x)return t->val;
73         return kth(t->c[1],x-t->s-sz(t->c[0]));
74     }
75 }T;

```

## 坚固的数据结构

### 坚固的平衡树

```

1 #define sz(x) (x?x->siz:0)
2 struct node{
3     int siz,key;
4     LL val,sum;
5     LL mu,a,d;
6     node *c[2],*f;

```

```

7 void split(int ned,node *&p,node *&q);
8 node* rz(){
9     sum=val; siz=1;
10    if(c[0]) sum+=c[0]->sum, siz+=c[0]->siz;
11    if(c[1]) sum+=c[1]->sum, siz+=c[1]->siz;
12    return this;
13 }
14 void make(LL _mu, LL _a, LL _d){
15     sum=sum*_mu+_a*siz+_d*siz*(siz-1)/2;
16     val=val*_mu+_a+_d*sz(c[0]);
17     mu*=_mu; a*=a*_mu+_a; d=d*_mu+_d;
18 }
19 void pd(){
20     if(mu==1&&a==0&&d==0) return;
21     if(c[0]) c[0]->make(mu, a, d);
22     if(c[1]) c[1]->make(mu, a+d+d*sz(c[0]), d);
23     mu=1; a=d=0;
24 }
25 node(){mu=1;}
26 }nd[maxn*2], *root;
27 node *merge(node *p, node *q){
28     if(!p||!q) return p?p->rz():(q?q->rz():0);
29     p->pd(); q->pd();
30     if(p->key<q->key){
31         p->c[1]=merge(p->c[1], q);
32         return p->rz();
33     }else{
34         q->c[0]=merge(p, q->c[0]);
35         return q->rz();
36     }
37 }
38 void node::split(int ned,node *&p,node *&q){
39     if(!ned){p=0;q=this;return;}
40     if(ned==siz){p=this;q=0;return;}
41     pd();
42     if(sz(c[0])>=ned){
43         c[0]->split(ned,p,q); c[0]=0; rz();
44         q=merge(q, this);
45     }else{
46         c[1]->split(ned-sz(c[0])-1,p,q); c[1]=0; rz();
47         p=merge(this, p);
48     }
49 }
50 int main(){
51     for(int i=1;i<=n;i++){
52         nd[i].val=in();
53         nd[i].key=rand();
54         nd[i].rz();
55         root=merge(root, nd+i);
56     }
57 }

```

坚固的字符串  
坚固的左偏树

```

1 int Merge(int x, int y){
2     if (x == 0 || y == 0) return x + y;
3     if (Heap[x].Key < Heap[y].Key) swap(x, y);
4     Heap[x].Ri = Merge(Heap[x].Ri, y);
5     if (Heap[Heap[x].Le].Dis < Heap[Heap[x].Ri].Dis) swap(Heap[x].Le, Heap[x].Ri);
6     if (Heap[x].Ri == 0) Heap[x].Dis = 0;
7     else Heap[x].Dis = Heap[Heap[x].Ri].Dis + 1;
8     return x;
9 }
10
11 for (int i = 0; i <= n; i++){
12     Heap[i].Le = Heap[i].Ri = 0;
13     Heap[i].Dis = 0;

```

```

14     Heap[i].Key = Cost[i];
15 }
16 Heap[0].Dis = -1;

```

树上的魔术师

Link Cut Tree(zky)

```

1 struct LCT{
2     struct node{
3         bool rev;
4         int mx, val;
5         node *f, *c[2];
6         bool d(){return this==f->c[1];}
7         bool rt(){return !f||f->c[0]!=this&&f->c[1]!=this;}
8         void sets(node *x, int d){pd(); if(x)x->f=this; c[d]=x; rz();}
9         void makerv(){rev^=1; swap(c[0], c[1]);}
10        void pd(){
11            if(rev){
12                if(c[0]) c[0]->makerv();
13                if(c[1]) c[1]->makerv();
14                rev=0;
15            }
16        }
17        void rz(){
18            mx=val;
19            if(c[0]) mx=max(mx, c[0]->mx);
20            if(c[1]) mx=max(mx, c[1]->mx);
21        }
22    }nd[int(1e4)+1];
23    void rot(node *x){
24        node *y=x->f; if(!y->rt()) y->f->pd();
25        y->pd(); x->pd(); bool d=x->d();
26        y->sets(x->c[!d], d);
27        if(y->rt()) x->f=y->f;
28        else y->f->sets(x, y->d());
29        x->sets(y, !d);
30    }
31    void splay(node *x){
32        while(!x->rt())
33            if(x->f->rt()) rot(x);
34            else if(x->d()==x->f->d()) rot(x->f), rot(x);
35            else rot(x), rot(x);
36    }
37    node* access(node *x){
38        node *y=0;
39        for(; x=x->f){
40            splay(x);
41            x->sets(y, 1); y=x;
42        }
43        return y;
44    }
45    void makert(node *x){
46        access(x)->makerv();
47        splay(x);
48    }
49    void link(node *x, node *y){
50        makert(x);
51        x->f=y;
52        access(x);
53    }
54    void cut(node *x, node *y){
55        makert(x); access(y); splay(y);
56        y->c[0]=x->f=0;
57        y->rz();
58    }
59    void link(int x, int y){link(nd+x, nd+y);}

```



```

59 void cut(int x,int y){cut(nd+x,nd+y);}
60 }T;

Link Cut Tree(Splay)

1 struct node{
2     bool Rev;
3     int c[2], fa;
4 }T[N];
5 inline void Rev(int x){
6     if (!x) return;
7     swap(T[x].c[0], T[x].c[1]);
8     T[x].Rev ^= 1;
9 }
10 inline void Lazy_Down(int x){
11     if (!x) return;
12     if (T[x].Rev) Rev(T[x].c[0]), Rev(T[x].c[1]), T[x].Rev = 0;
13 }
14 void Rotate(int x, int c){
15     int y = T[x].c[c];
16     int z = T[y].c[1 - c];
17     if (T[x].fa){
18         if (T[T[x].fa].c[0] == x) T[T[x].fa].c[0] = y;
19         else T[T[x].fa].c[1] = y;
20     }
21     T[z].fa = x; T[x].c[c] = z;
22     T[y].fa = T[x].fa; T[x].fa = y; T[y].c[1 - c] = x;
23     //Update(x);
24     //Update(y);
25 }
26 int stack[N], fx[N];
27 void Splay(int x){
28     int top = 0;
29     for (int u = x; u; u = T[u].fa)
30         stack[++top] = u;
31     for (int i = top; i >= 1; i--)
32         Lazy_Down(stack[i]);
33     for (int i = 2; i <= top; i++){
34         if (T[stack[i]].c[0] == stack[i - 1]) fx[i] = 0;
35         else fx[i] = 1;
36     }
37     for (int i = 2; i <= top; i += 2){
38         if (i == top) Rotate(stack[i], fx[i]);
39         else {
40             if (fx[i] == fx[i + 1]){
41                 Rotate(stack[i + 1], fx[i + 1]);
42                 Rotate(stack[i], fx[i]);
43             } else {
44                 Rotate(stack[i], fx[i]);
45                 Rotate(stack[i + 1], fx[i + 1]);
46             }
47         }
48     }
49     if (x != stack[top]) Par[x] = Par[stack[top]], Par[stack[top]] = 0;
50     //if (fa == 0) Root = x;
51 }
52 inline int Access(int u){
53     int NXT = 0;
54     while (u){
55         Splay(u);
56         if (T[u].c[1]){
57             T[T[u].c[1]].fa = 0;
58             Par[T[u].c[1]] = u;
59         }
60         T[u].c[1] = NXT;
61         if (NXT){
62             T[NXT].fa = u;
63             Par[NXT] = 0;
64         }
65     }

```

```

64     //Update(u)
65     NXT = u;
66     u = Par[u];
67 }
68 return NXT;
69 }
70 inline void Root(int u){
71     Access(u);
72     Splay(u);
73     Rev(u);
74 }
75 inline void Link(int u, int v){
76     Root(u);
77     Par[u] = v;
78 }
79 inline void Cut(int u, int v){
80     Access(u);
81     Splay(v);
82     if (Par[v] != u){
83         swap(u, v);
84         Access(u);
85         Splay(v);
86     }
87     Par[v] = 0;
88 }
89 inline int Find_Root(int x){
90     Access(x);
91     Splay(x);
92     int y = x;
93     while (T[y].c[0]){
94         Lazy_Down(y);
95         y = T[y].c[0];
96     }
97     return y;
98 }

```

### 可持久化线段树

```

1 struct node1 {
2     int L, R, Lson, Rson, Sum;
3 } tree[N * 40];
4 int root[N], a[N], b[N];
5 int tot, n, m;
6 int Real[N];
7 int Same(int x) {
8     ++tot;
9     tree[tot] = tree[x];
10    return tot;
11 }
12 int build(int L, int R) {
13     ++tot;
14     tree[tot].L = L;
15     tree[tot].R = R;
16     tree[tot].Lson = tree[tot].Rson = tree[tot].Sum = 0;
17     if (L == R) return tot;
18     int s = tot;
19     int mid = (L + R) >> 1;
20     tree[s].Lson = build(L, mid);
21     tree[s].Rson = build(mid + 1, R);
22     return s;
23 }
24 int Ask(int Lst, int Cur, int L, int R, int k) {
25     if (L == R) return L;
26     int Mid = (L + R) >> 1;
27     int Left = tree[tree[Cur].Lson].Sum - tree[tree[Lst].Lson].Sum;
28     if (Left >= k) return Ask(tree[Lst].Lson, tree[Cur].Lson, L, Mid, k);
29     k -= Left;
30     return Ask(tree[Lst].Rson, tree[Cur].Rson, Mid + 1, R, k);
31 }

```



```

31 }
32 int Add(int Lst, int pos) {
33     int root = Same(Lst);
34     tree[root].Sum++;
35     if (tree[root].L == tree[root].R) return root;
36     int mid = (tree[root].L + tree[root].R) >> 1;
37     if (pos <= mid) tree[root].Lson = Add(tree[root].Lson, pos);
38     else tree[root].Rson = Add(tree[root].Rson, pos);
39     return root;
40 }
41 int main() {
42     scanf("%d%d", &n, &m);
43     int up = 0;
44     for (int i = 1; i <= n; i++){
45         scanf("%d", &a[i]);
46         b[i] = a[i];
47     }
48     sort(b + 1, b + n + 1);
49     up = unique(b + 1, b + n + 1) - b - 1;
50     for (int i = 1; i <= n; i++){
51         int tmp = lower_bound(b + 1, b + up + 1, a[i]) - b;
52         Real[tmp] = a[i];
53         a[i] = tmp;
54     }
55     tot = 0;
56     root[0] = build(1, up);
57     for (int i = 1; i <= n; i++){
58         root[i] = Add(root[i - 1], a[i]);
59     }
60     for (int i = 1; i <= m; i++){
61         int u, v, w;
62         scanf("%d%d%d", &u, &v, &w);
63         printf("%d\n", Real[Ask(root[u - 1], root[v], 1, up, w)]);
64     }
65     return 0;
66 }

```

## k-d 树

```

1 long long norm(const long long &x) {
2     // For manhattan distance
3     return std::abs(x);
4     // For euclid distance
5     return x * x;
6 }
7
8 struct Point {
9     int x, y, id;
10
11     const int& operator [] (int index) const {
12         if (index == 0) {
13             return x;
14         } else {
15             return y;
16         }
17     }
18
19     friend long long dist(const Point &a, const Point &b) {
20         long long result = 0;
21         for (int i = 0; i < 2; ++i) {
22             result += norm(a[i] - b[i]);
23         }
24         return result;
25     }
26 } point[N];
27
28 struct Rectangle {
29     int min[2], max[2];
30
31     Rectangle() {

```

```

32     min[0] = min[1] = INT_MAX;
33     max[0] = max[1] = INT_MIN;
34 }
35
36 void add(const Point &p) {
37     for (int i = 0; i < 2; ++i) {
38         min[i] = std::min(min[i], p[i]);
39         max[i] = std::max(max[i], p[i]);
40     }
41 }
42
43 long long dist(const Point &p) {
44     long long result = 0;
45     for (int i = 0; i < 2; ++i) {
46         // For minimum distance
47         result += norm(std::min(min[i], p[i]), max[i]) - p[i]);
48         // For maximum distance
49         result += std::max(norm(max[i] - p[i]), norm(min[i] - p[i]));
50     }
51     return result;
52 }
53 };
54
55 struct Node {
56     Point separator;
57     Rectangle rectangle;
58     int child[2];
59
60     void reset(const Point &p) {
61         separator = p;
62         rectangle = Rectangle();
63         rectangle.add(p);
64         child[0] = child[1] = 0;
65     }
66 } tree[N << 1];
67
68 int size, pivot;
69
70 bool compare(const Point &a, const Point &b) {
71     if (a[pivot] != b[pivot]) {
72         return a[pivot] < b[pivot];
73     }
74     return a.id < b.id;
75 }
76
77 int build(int l, int r, int type = 1) {
78     pivot = type;
79     if (l >= r) {
80         return 0;
81     }
82     int x = ++size;
83     int mid = l + r >> 1;
84     std::nth_element(point + l, point + mid, point + r, compare);
85     tree[x].reset(point[mid]);
86     for (int i = l; i < r; ++i) {
87         tree[x].rectangle.add(point[i]);
88     }
89     tree[x].child[0] = build(l, mid, type ^ 1);
90     tree[x].child[1] = build(mid + 1, r, type ^ 1);
91     return x;
92 }
93
94 int insert(int x, const Point &p, int type = 1) {
95     pivot = type;
96     if (x == 0) {
97         tree[++size].reset(p);
98         return size;
99     }
100     tree[x].rectangle.add(p);

```

```

101     if (compare(p, tree[x].separator)) {
102         tree[x].child[0] = insert(tree[x].child[0], p, type ^ 1);
103     } else {
104         tree[x].child[1] = insert(tree[x].child[1], p, type ^ 1);
105     }
106     return x;
107 }
108
109 // For minimum distance
110 void query(int x, const Point &p, std::pair<long long, int> &answer, int type =
    ↪ 1) {
111     pivot = type;
112     if (x == 0 || tree[x].rectangle.dist(p) > answer.first) {
113         return;
114     }
115     answer = std::min(answer,
        std::make_pair(dist(tree[x].separator, p), tree[x].separator.id));
116     if (compare(p, tree[x].separator)) {
117         query(tree[x].child[0], p, answer, type ^ 1);
118         query(tree[x].child[1], p, answer, type ^ 1);
119     } else {
120         query(tree[x].child[1], p, answer, type ^ 1);
121         query(tree[x].child[0], p, answer, type ^ 1);
122     }
123 }
124 }
125
126 std::priority_queue<std::pair<long long, int> > answer;
127
128 void query(int x, const Point &p, int k, int type = 1) {
129     pivot = type;
130     if (x == 0 ||
131         (int)answer.size() == k && tree[x].rectangle.dist(p) >
            ↪ answer.top().first) {
132         return;
133     }
134     answer.push(std::make_pair(dist(tree[x].separator, p),
        ↪ tree[x].separator.id));
135     if ((int)answer.size() > k) {
136         answer.pop();
137     }
138     if (compare(p, tree[x].separator)) {
139         query(tree[x].child[0], p, k, type ^ 1);
140         query(tree[x].child[1], p, k, type ^ 1);
141     } else {
142         query(tree[x].child[1], p, k, type ^ 1);
143         query(tree[x].child[0], p, k, type ^ 1);
144     }
145 }

```

### 莫队算法

```

1 struct node{
2     int l, r, id;
3     friend bool operator < (const node &a, const node &b){
4         if (a.l / Block == b.l / Block) return a.r < b.r;
5         return a.l / Block < b.l / Block;
6     }
7 }q[N];
8 Block = int(sqrt(n));
9 for (int i = 1; i <= m; i++){
10     scanf("%d%d", &q[i].l, &q[i].r);
11     q[i].id = i;
12 }
13 sort(q + 1, q + 1 + m);
14 Cur = a[1]; /// Hints: adjust by yourself
15 Le = Ri = 1;
16 for (int i = 1; i <= m; i++){
17     while (q[i].r > Ri) Ri++, ChangeRi(1, Le, Ri);

```

```

18     while (q[i].l > Le) ChangeLe(-1, Le, Ri), Le++;
19     while (q[i].l < Le) Le--, ChangeLe(1, Le, Ri);
20     while (q[i].r < Ri) ChangeRi(-1, Le, Ri), Ri--;
21     Ans[q[i].id] = Cur;
22 }

```

### 树状数组 kth

```

1 int find(int k){
2     int cnt=0,ans=0;
3     for(int i=22;i>=0;i--){
4         ans+=(1<<i);
5         if(ans>n || cnt+d[ans]>=k)ans--(1<<i);
6         else cnt+=d[ans];
7     }
8     return ans+1;
9 }

```

### 虚树

```

1 int a[maxn*2],sta[maxn*2];
2 int top=0,k;
3 void build(){
4     top=0;
5     sort(a,a+k,bydfn);
6     k=unique(a,a+k)-a;
7     sta[top++]=1;_n=k;
8     for(int i=0;i<k;i++){
9         int LCA=lca(a[i],sta[top-1]);
10        while(dep[LCA]<dep[sta[top-1]]){
11            if(dep[LCA]>=dep[sta[top-2]]){
12                add_edge(LCA,sta[--top]);
13                if(sta[top-1]!=LCA)sta[top++]=LCA;
14                break;
15            }add_edge(sta[top-2],sta[top-1]);top--;
16        }if(sta[top-1]!=a[i])sta[top++]=a[i];
17    }
18    while(top>1)
19        add_edge(sta[top-2],sta[top-1]),top--;
20    for(int i=0;i<k;i++)inr[a[i]]=1;
21 }

```

### 点分治 (zky)

```

1 int siz[maxn],f[maxn],dep[maxn],cant[maxn],root,All,d[maxn];
2 void makert(int u,int fa){
3     siz[u]=1;f[u]=0;
4     for(int i=0;i<G[u].size();i++){
5         edge e=G[u][i];
6         if(e.v!=fa&&!cant[e.v]){
7             dep[e.v]=dep[u]+1;
8             makert(e.v,u);
9             siz[u]+=siz[e.v];
10            f[u]=max(f[u],siz[e.v]);
11        }
12    }f[u]=max(f[u],All-f[u]);
13    if(f[root]>f[u])root=u;
14 }
15 void dfs(int u,int fa){
16     //Gain data
17     for(int i=0;i<G[u].size();i++){
18         edge e=G[u][i];
19         if(e.v==fa||cant[e.v])continue;
20         d[e.v]=d[u]+e.w;
21         dfs(e.v,u);

```

```

22     }
23 }
24 void calc(int u){
25     d[u]=0;
26     for(int i=0;i<G[u].size();i++){
27         edge e=G[u][i];
28         if(cant[e.v])continue;
29         d[e.v]=e.w;
30         dfs(e.v,u);
31     }
32 }
33 }
34 void solve(int u){
35     calc(u);cant[u]=1;
36     for(int i=0;i<G[u].size();i++){
37         edge e=G[u][i];
38         if(cant[e.v])continue;
39         All=siz[e.v];
40         f[root=0]=n+1;
41         makert(e.v,0);
42         solve(root);
43     }
44 }
45 All=n;
46 f[root=0]=n+1;
47 makert(1,1);
48 solve(root);

```

### 元芳树

```

1 void tarjan(int u){
2     dfn[u]=low[u]=++tot;
3     for(int i=0;i<G[u].size();i++){
4         edge e=G[u][i];
5         if(dfn[e.v])
6             low[u]=min(low[u],dfn[e.v]);
7         else{
8             S.push(e);
9             tarjan(e.v);
10            if(low[e.v]==dfn[u]){
11
12                if(S.top()==e){
13                    fa[e.v][0]=u;
14                    fw[e.v]=e.w;
15                    S.pop();
16                    continue;
17                }
18
19                Rcnt++;
20                edge ed;
21                do{
22                    ed=S.top();S.pop();
23                    ring[Rcnt].push_back(ed);
24                }while(ed!=e);
25                reverse(ring[Rcnt].begin(),ring[Rcnt].end());
26                int last=ring[Rcnt].back().v;
27                ring[Rcnt].push_back((edge){last,u,Mw[pack(last,u)]});
28            }
29            low[u]=min(low[u],low[e.v]);
30        }
31    }
32 }
33 void up(int u){
34     if(dep[u]||u==1)return ;
35     if(fa[u][0])up(fa[u][0]);
36     dep[u]=dep[fa[u][0]]+1;
37     fw[u]+=fw[fa[u][0]];
38 }

```

```

39 void build(){
40     S.push((edge){0,1,0});
41     tarjan(1);
42
43     for(int i=1;i<=Rcnt;i++){
44         rlen[i]=0;
45         sum[i].resize(ring[i].size());
46         dis[i].resize(ring[i].size());
47         for(int j=0;j<ring[i].size();j++){
48             rlen[i]+=ring[i][j].w;
49             ind[i].push_back(make_pair(ring[i][j].u,j));
50         }
51         sum[i][0]=0;
52         fw[i+n]=0;
53         fa[i+n][0]=ring[i][0].u;
54         for(int j=1;j<ring[i].size();j++){
55             sum[i][j]=sum[i][j-1]+ring[i][j-1].w;
56             dis[i][j]=min(sum[i][j],rlen[i]-sum[i][j]);
57             fw[ring[i][j].u]=dis[i][j];
58             fa[ring[i][j].u][0]=i+n;
59         }
60         sort(ind[i].begin(),ind[i].end());
61     }
62
63     for(int i=1;i<=n+Rcnt;i++)
64         up(i);
65
66     for(int j=1;j<BIT;j++)
67         for(int i=1;i<=n+Rcnt;i++)if(fa[i][j-1])
68             fa[i][j]=fa[fa[i][j-1]][j-1];
69
70 }
71 pair<int,int>second_lca;
72 int lca(int u,int v){
73     if(dep[u]<dep[v])swap(u,v);
74     int d=dep[u]-dep[v];
75     for(int i=0;i<BIT;i++)if(d>>i&1)
76         u=fa[u][i];
77     if(u==v)return u;
78     for(int i=BIT-1;i>=0;i--)if(fa[u][i]!=fa[v][i]){
79         u=fa[u][i];
80         v=fa[v][i];
81     }
82     second_lca=make_pair(u,v);
83     return fa[u][0];
84 }

```

### 图论

#### 点双连通分量 (lyx)

```

1 ///求割点，割点向每个点双连通分量连边
2 void Dfs(int x, int lst){
3     dfn[x] = ++dfc;
4     low[x] = dfn[x];
5     stack[++cnt] = x;
6     int son = 0;
7
8     for (int i = g[x]; i; i = nxt[i]){
9         if (!dfn[adj[i]]){
10             ++son;
11             Dfs(adj[i], i);
12             low[x] = min(low[x], low[adj[i]]);
13             if (low[adj[i]] >= dfn[x]){
14                 int Tmp;
15                 iscut[x] = 1;
16                 ++block;
17                 E[x].push_back(block + n);

```

```

18         do{
19             Tmp = stack[cnt --];
20             belong[Tmp] = block + n;
21
22             E[Tmp].push_back(block + n);
23         }while (Tmp != adj[i]);
24     }
25 }
26 else
27 if ((i ^ lst) != 1) low[x] = min(low[x], dfn[adj[i]]);
28 }
29 if (x == Root && son == 1) iscut[x] = 0, belong[x] = E[x][0];
30 if (x == Root && son == 0){
31     ++block;
32     belong[x] = block + n;
33 }
34 }
35 tot = 1; //!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!1
36 block = 0;
37 cnt = 0;
38 dfc = 0;
39 for (int i = 1; i <= n; i++)
40     if (dfn[i] == 0){
41         Root = i;
42         Dfs(i, 0);
43     }

```

## 2-SAT 问题 (强连通分量)

```

1 int stamp, comps, top;
2 int dfn[N], low[N], comp[N], stack[N];
3
4 void add(int x, int a, int y, int b) {
5     edge[x << 1 | a].push_back(y << 1 | b);
6 }
7
8 void tarjan(int x) {
9     dfn[x] = low[x] = ++stamp;
10    stack[top++] = x;
11    for (int i = 0; i < (int)edge[x].size(); ++i) {
12        int y = edge[x][i];
13        if (!dfn[y]) {
14            tarjan(y);
15            low[x] = std::min(low[x], low[y]);
16        } else if (!comp[y]) {
17            low[x] = std::min(low[x], dfn[y]);
18        }
19    }
20    if (low[x] == dfn[x]) {
21        comps++;
22        do {
23            int y = stack[--top];
24            comp[y] = comps;
25        } while (stack[top] != x);
26    }
27 }
28
29 bool solve() {
30     int counter = n + n + 1;
31     stamp = top = comps = 0;
32     std::fill(dfn, dfn + counter, 0);
33     std::fill(comp, comp + counter, 0);
34     for (int i = 0; i < counter; ++i) {
35         if (!dfn[i]) {
36             tarjan(i);
37         }
38     }
39     for (int i = 0; i < n; ++i) {
40         if (comp[i << 1] == comp[i << 1 | 1]) {

```

```

41         return false;
42     }
43     answer[i] = (comp[i << 1 | 1] < comp[i << 1]);
44 }
45 return true;
46 }

```

## 二分图最大匹配

## Hungary 算法

时间复杂度:  $O(V \cdot E)$ 

```

1 vector<int>G[maxn];
2 int Link[maxn], vis[maxn], T;
3 bool find(int x){
4     for(int i=0; i<G[x].size(); i++){
5         int v=G[x][i];
6         if(vis[v]==T) continue;
7         vis[v]=T;
8         if(!Link[v] || find(Link[v])){
9             Link[v]=x;
10            return 1;
11        }
12    }return 0;
13 }
14 int Hungarian(int n){
15     int ans=0;
16     memset(Link, 0, sizeof Link);
17     for(int i=1; i<=n; i++){
18         T++;
19         ans+=find(i);
20     }return ans;
21 }

```

## Hopcroft Karp 算法

时间复杂度:  $O(\sqrt{V} \cdot E)$ 

```

1 int matchx[N], matchy[N], level[N];
2
3 bool dfs(int x) {
4     for (int i = 0; i < (int)edge[x].size(); ++i) {
5         int y = edge[x][i];
6         int w = matchy[y];
7         if (w == -1 || level[x] + 1 == level[w] && dfs(w)) {
8             matchx[x] = y;
9             matchy[y] = x;
10            return true;
11        }
12    }
13    level[x] = -1;
14    return false;
15 }
16
17 int solve() {
18     std::fill(matchx, matchx + n, -1);
19     std::fill(matchy, matchy + m, -1);
20     for (int answer = 0; ; ) {
21         std::vector<int> queue;
22         for (int i = 0; i < n; ++i) {
23             if (matchx[i] == -1) {
24                 level[i] = 0;
25                 queue.push_back(i);
26             } else {
27                 level[i] = -1;
28             }

```

```

29     }
30     for (int head = 0; head < (int)queue.size(); ++head) {
31         int x = queue[head];
32         for (int i = 0; i < (int)edge[x].size(); ++i) {
33             int y = edge[x][i];
34             int w = matchy[y];
35             if (w != -1 && level[w] < 0) {
36                 level[w] = level[x] + 1;
37                 queue.push_back(w);
38             }
39         }
40     }
41     int delta = 0;
42     for (int i = 0; i < n; ++i) {
43         if (matchx[i] == -1 && dfs(i)) {
44             delta++;
45         }
46     }
47     if (delta == 0) {
48         return answer;
49     } else {
50         answer += delta;
51     }
52 }
53 }

```

## 二分图最大权匹配

时间复杂度： $\mathcal{O}(V^4)$ 

```

1 int labelx[N], labely[N], match[N], slack[N];
2 bool visitx[N], visity[N];
3
4 bool dfs(int x) {
5     visitx[x] = true;
6     for (int y = 0; y < n; ++y) {
7         if (visity[y]) {
8             continue;
9         }
10        int delta = labelx[x] + labely[y] - graph[x][y];
11        if (delta == 0) {
12            visity[y] = true;
13            if (match[y] == -1 || dfs(match[y])) {
14                match[y] = x;
15                return true;
16            }
17        } else {
18            slack[y] = std::min(slack[y], delta);
19        }
20    }
21    return false;
22 }
23
24 int solve() {
25     for (int i = 0; i < n; ++i) {
26         match[i] = -1;
27         labelx[i] = INT_MIN;
28         labely[i] = 0;
29         for (int j = 0; j < n; ++j) {
30             labelx[i] = std::max(labelx[i], graph[i][j]);
31         }
32     }
33     for (int i = 0; i < n; ++i) {
34         while (true) {
35             std::fill(visitx, visitx + n, 0);
36             std::fill(visity, visity + n, 0);
37             for (int j = 0; j < n; ++j) {
38                 slack[j] = INT_MAX;

```

```

39     }
40     if (dfs(i)) {
41         break;
42     }
43     int delta = INT_MAX;
44     for (int j = 0; j < n; ++j) {
45         if (!visity[j]) {
46             delta = std::min(delta, slack[j]);
47         }
48     }
49     for (int j = 0; j < n; ++j) {
50         if (visitx[j]) {
51             labelx[j] -= delta;
52         }
53         if (visity[j]) {
54             labely[j] += delta;
55         } else {
56             slack[j] -= delta;
57         }
58     }
59 }
60 }
61 int answer = 0;
62 for (int i = 0; i < n; ++i) {
63     answer += graph[match[i]][i];
64 }
65 return answer;
66 }

```

## 最大流 (dinic)

时间复杂度： $\mathcal{O}(V^2 \cdot E)$ 

```

1 struct edge{int u,v,cap,flow;};
2 vector<edge>edges;
3 vector<int>G[maxn];
4 int s,t;
5 int cur[maxn],d[maxn];
6 void add(int u,int v,int cap){
7     edges.push_back((edge){u,v,cap,0});
8     G[u].push_back(edges.size()-1);
9     edges.push_back((edge){v,u,0,0});
10    G[v].push_back(edges.size()-1);
11 }
12 bool bfs(){
13     static int vis[maxn];
14     memset(vis,0,sizeof vis);vis[s]=1;
15     queue<int>q;q.push(s);d[s]=0;
16     while(!q.empty()){
17         int u=q.front();q.pop();
18         for(int i=0;i<G[u].size();i++){
19             edge e=edges[G[u][i]];if(vis[e.v]||e.cap==e.flow)continue;
20             d[e.v]=d[u]+1;vis[e.v]=1;q.push(e.v);
21         }
22     }return vis[t];
23 }
24 int dfs(int u,int a){
25     if(u==t||!a)return a;
26     int flow=0,f;
27     for(int &i=cur[u];i<G[u].size();i++){
28         edge e=edges[G[u][i]];
29         if(d[e.v]==d[u]+1&&(f=dfs(e.v,min(a,e.cap-e.flow)))>0){
30             edges[G[u][i]].flow+=f;
31             edges[G[u][i]^1].flow-=f;
32             flow+=f;a-=f;if(!a)break;
33         }

```

```

34     }return flow;
35 }
36 int dinic(){
37     int flow=0,x;
38     while(bfs()){
39         memset(cur,0,sizeof cur);
40         while(x=dfs(s,INT_MAX)){
41             flow+=x;
42             memset(cur,0,sizeof cur);
43         }
44     }return flow;
45 }

```

### 最大流 (sap)

时间复杂度:  $\mathcal{O}(V^2 \cdot E)$

```

1 int g[T], adj[M], nxt[M], f[M];
2 int cnt[T], dist[T], cur[T], fa[T], dat[T];
3 void Ins(int x, int y, int ff, int rf){
4     adj[++tot] = y; nxt[tot] = g[x]; g[x] = tot; f[tot] = ff;
5     adj[++tot] = x; nxt[tot] = g[y]; g[y] = tot; f[tot] = rf;
6 }
7 int sap(int s, int t){
8     int x, sum;
9     for (int i = 1; i <= t; i++){
10         dist[i] = 1;
11         cur[i] = g[i];
12         fa[i] = 0;
13         dat[i] = 0;
14         cnt[i] = 0;
15     }
16     cnt[0] = 1; cnt[1] = t - 1;
17     dist[t] = 0;
18     dat[s] = INF;
19     x = s;
20     sum = 0;
21     while (1){
22         int p;
23         for (p = cur[x]; p; p = nxt[p]){
24             if (f[p] > 0 && dist[adj[p]] == dist[x] - 1) break;
25         }
26         if (p > 0){
27             cur[x] = p;
28             fa[adj[p]] = p;
29             dat[adj[p]] = min(dat[x], f[p]);
30             x = adj[p];
31             if (x == t){
32                 sum += dat[x];
33                 while (x != s){
34                     f[fa[x]] -= dat[t];
35                     f[fa[x] ^ 1] += dat[t];
36                     x = adj[fa[x] ^ 1];
37                 }
38             }
39         } else {
40             cnt[dist[x]]--;
41             if (cnt[dist[x]] == 0) return sum;
42             dist[x] = t + 1;
43             for (int p = g[x]; p; p = nxt[p]){
44                 if (f[p] > 0 && dist[adj[p]] + 1 < dist[x]){
45                     dist[x] = dist[adj[p]] + 1;
46                     cur[x] = p;
47                 }
48             }
49             cnt[dist[x]]++;
50             if (dist[s] > t) return sum;

```

```

51         if (x != s) x = adj[fa[x] ^ 1];
52     }
53 }
54 }
55 /*
56 tot = 1
57 edges' id start from 2
58 remember to clean g
59 t is the number of points
60 */

```

### 上下界网络流

$B(u, v)$  表示边  $(u, v)$  流量的下界,  $C(u, v)$  表示边  $(u, v)$  流量的上界,  $F(u, v)$  表示边  $(u, v)$  的流量。设  $G(u, v) = F(u, v) - B(u, v)$ , 显然有

$$0 \leq G(u, v) \leq C(u, v) - B(u, v)$$

#### 无源汇的上下界可行流

建立超级源点  $S^*$  和超级汇点  $T^*$ , 对于原图每条边  $(u, v)$  在新网络中连如下三条边:  $S^* \rightarrow v$ , 容量为  $B(u, v)$ ;  $u \rightarrow T^*$ , 容量为  $B(u, v)$ ;  $u \rightarrow v$ , 容量为  $C(u, v) - B(u, v)$ 。最后求新网络的最大流, 判断从超级源点  $S^*$  出发的边是否都满流即可, 边  $(u, v)$  的最终解中的实际流量为  $G(u, v) + B(u, v)$ 。

#### 有源汇的上下界可行流

从汇点  $T$  到源点  $S$  连一条上界为  $\infty$ , 下界为 0 的边。按照无源汇的上下界可行流一样做即可, 流量即为  $T \rightarrow S$  边上的流量。

#### 有源汇的上下界最大流

1. 在有源汇的上下界可行流中, 从汇点  $T$  到源点  $S$  的边改为连一条上界为  $\infty$ , 下届为  $x$  的边。 $x$  满足二分性质, 找到最大的  $x$  使得新网络存在无源汇的上下界可行流即为原图的最大流。
2. 从汇点  $T$  到源点  $S$  连一条上界为  $\infty$ , 下界为 0 的边, 变成无源汇的网络。按照无源汇的上下界可行流的方法, 建立超级源点  $S^*$  和超级汇点  $T^*$ , 求一遍  $S^* \rightarrow T^*$  的最大流, 再将从汇点  $T$  到源点  $S$  的这条边拆掉, 求一次  $S \rightarrow T$  的最大流即可。

#### 有源汇的上下界最小流

1. 在有源汇的上下界可行流中, 从汇点  $T$  到源点  $S$  的边改为连一条上界为  $x$ , 下界为 0 的边。 $x$  满足二分性质, 找到最小的  $x$  使得新网络存在无源汇的上下界可行流即为原图的最小流。
2. 按照无源汇的上下界可行流的方法, 建立超级源点  $S^*$  与超级汇点  $T^*$ , 求一遍  $S^* \rightarrow T^*$  的最大流, 但是注意这一次不加上汇点  $T$  到源点  $S$  的这条边, 即不使之改为无源汇的网络去求解。求完后, 再加上那条汇点  $T$  到源点  $S$  上界  $\infty$  的边。因为这条边下界为 0, 所以  $S^*$ ,  $T^*$  无影响, 再直接求一次  $S^* \rightarrow T^*$  的最大流。若超级源点  $S^*$  出发的边全部满流, 则  $T \rightarrow S$  边上的流量即为原图的最小流, 否则无解。

### 最小费用最大流

#### 稀疏图

时间复杂度:  $\mathcal{O}(V \cdot E^2)$

```

1 struct EdgeList {
2     int size;
3     int last[N];
4     int succ[M], other[M], flow[M], cost[M];
5     void clear(int n) {
6         size = 0;
7         std::fill(last, last + n, -1);
8     }
9     void add(int x, int y, int c, int w) {
10         succ[size] = last[x];
11         last[x] = size;
12         other[size] = y;
13         flow[size] = c;
14         cost[size++] = w;
15     }
16 } e;
17
18 int n, source, target;
19 int prev[N];
20
21 void add(int x, int y, int c, int w) {

```



```

22     e.add(x, y, c, w);
23     e.add(y, x, 0, -w);
24 }
25
26 bool augment() {
27     static int dist[N], occur[N];
28     std::vector<int> queue;
29     std::fill(dist, dist + n, INT_MAX);
30     std::fill(occur, occur + n, 0);
31     dist[source] = 0;
32     occur[source] = true;
33     queue.push_back(source);
34     for (int head = 0; head < (int)queue.size(); ++head) {
35         int x = queue[head];
36         for (int i = e.last[x]; ~i; i = e.succ[i]) {
37             int y = e.other[i];
38             if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
39                 dist[y] = dist[x] + e.cost[i];
40                 prev[y] = i;
41                 if (!occur[y]) {
42                     occur[y] = true;
43                     queue.push_back(y);
44                 }
45             }
46         }
47         occur[x] = false;
48     }
49     return dist[target] < INT_MAX;
50 }
51
52 std::pair<int, int> solve() {
53     std::pair<int, int> answer = std::make_pair(0, 0);
54     while (augment()) {
55         int number = INT_MAX;
56         for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
57             number = std::min(number, e.flow[prev[i]]);
58         }
59         answer.first += number;
60         for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
61             e.flow[prev[i]] -= number;
62             e.flow[prev[i] ^ 1] += number;
63             answer.second += number * e.cost[prev[i]];
64         }
65     }
66     return answer;
67 }

```

稠密图

使用条件：费用非负

时间复杂度： $\mathcal{O}(V \cdot E^2)$ 

```

1 struct EdgeList {
2     int size;
3     int last[N];
4     int succ[M], other[M], flow[M], cost[M];
5     void clear(int n) {
6         size = 0;
7         std::fill(last, last + n, -1);
8     }
9     void add(int x, int y, int c, int w) {
10         succ[size] = last[x];
11         last[x] = size;
12         other[size] = y;
13         flow[size] = c;
14         cost[size++] = w;
15     }

```

```

16 } e;
17
18 int n, source, target, flow, cost;
19 int slack[N], dist[N];
20 bool visit[N];
21
22 void add(int x, int y, int c, int w) {
23     e.add(x, y, c, w);
24     e.add(y, x, 0, -w);
25 }
26
27 bool relabel() {
28     int delta = INT_MAX;
29     for (int i = 0; i < n; ++i) {
30         if (!visit[i]) {
31             delta = std::min(delta, slack[i]);
32         }
33         slack[i] = INT_MAX;
34     }
35     if (delta == INT_MAX) {
36         return true;
37     }
38     for (int i = 0; i < n; ++i) {
39         if (visit[i]) {
40             dist[i] += delta;
41         }
42     }
43     return false;
44 }
45
46 int dfs(int x, int answer) {
47     if (x == target) {
48         flow += answer;
49         cost += answer * (dist[source] - dist[target]);
50         return answer;
51     }
52     visit[x] = true;
53     int delta = answer;
54     for (int i = e.last[x]; ~i; i = e.succ[i]) {
55         int y = e.other[i];
56         if (e.flow[i] > 0 && !visit[y]) {
57             if (dist[y] + e.cost[i] == dist[x]) {
58                 int number = dfs(y, std::min(e.flow[i], delta));
59                 e.flow[i] -= number;
60                 e.flow[i ^ 1] += number;
61                 delta -= number;
62                 if (delta == 0) {
63                     dist[x] = INT_MIN;
64                     return answer;
65                 }
66             } else {
67                 slack[y] = std::min(slack[y], dist[y] + e.cost[i] - dist[x]);
68             }
69         }
70     }
71     return answer - delta;
72 }
73
74 std::pair<int, int> solve() {
75     flow = cost = 0;
76     std::fill(dist, dist + n, 0);
77     do {
78         do {
79             fill(visit, visit + n, 0);
80         } while (dfs(source, INT_MAX));
81     } while (!relabel());
82     return std::make_pair(flow, cost);
83 }

```



## 一般图最大匹配

时间复杂度:  $\mathcal{O}(V^3)$ 

```

1 int match[N], belong[N], next[N], mark[N], visit[N];
2 std::vector<int> queue;
3
4 int find(int x) {
5     if (belong[x] != x) {
6         belong[x] = find(belong[x]);
7     }
8     return belong[x];
9 }
10
11 void merge(int x, int y) {
12     x = find(x);
13     y = find(y);
14     if (x != y) {
15         belong[x] = y;
16     }
17 }
18
19 int lca(int x, int y) {
20     static int stamp = 0;
21     stamp++;
22     while (true) {
23         if (x != -1) {
24             x = find(x);
25             if (visit[x] == stamp) {
26                 return x;
27             }
28             visit[x] = stamp;
29             if (match[x] != -1) {
30                 x = next[match[x]];
31             } else {
32                 x = -1;
33             }
34         }
35         std::swap(x, y);
36     }
37 }
38
39 void group(int a, int p) {
40     while (a != p) {
41         int b = match[a], c = next[b];
42         if (find(c) != p) {
43             next[c] = b;
44         }
45         if (mark[b] == 2) {
46             mark[b] = 1;
47             queue.push_back(b);
48         }
49         if (mark[c] == 2) {
50             mark[c] = 1;
51             queue.push_back(c);
52         }
53         merge(a, b);
54         merge(b, c);
55         a = c;
56     }
57 }
58
59 void augment(int source) {
60     queue.clear();
61     for (int i = 0; i < n; ++i) {
62         next[i] = visit[i] = -1;
63         belong[i] = i;
64         mark[i] = 0;
65     }

```

```

66     mark[source] = 1;
67     queue.push_back(source);
68     for (int head = 0; head < (int)queue.size() && match[source] == -1; ++head) {
69         int x = queue[head];
70         for (int i = 0; i < (int)edge[x].size(); ++i) {
71             int y = edge[x][i];
72             if (match[x] == y || find(x) == find(y) || mark[y] == 2) {
73                 continue;
74             }
75             if (mark[y] == 1) {
76                 int r = lca(x, y);
77                 if (find(x) != r) {
78                     next[x] = y;
79                 }
80                 if (find(y) != r) {
81                     next[y] = x;
82                 }
83                 group(x, r);
84                 group(y, r);
85             } else if (match[y] == -1) {
86                 next[y] = x;
87                 for (int u = y; u != -1; ) {
88                     int v = next[u];
89                     int mv = match[v];
90                     match[v] = u;
91                     match[u] = v;
92                     u = mv;
93                 }
94                 break;
95             } else {
96                 next[y] = x;
97                 mark[y] = 2;
98                 mark[match[y]] = 1;
99                 queue.push_back(match[y]);
100             }
101         }
102     }
103 }
104
105 int solve() {
106     std::fill(match, match + n, -1);
107     for (int i = 0; i < n; ++i) {
108         if (match[i] == -1) {
109             augment(i);
110         }
111     }
112     int answer = 0;
113     for (int i = 0; i < n; ++i) {
114         answer += (match[i] != -1);
115     }
116     return answer;
117 }

```

## 无向图全局最小割

时间复杂度:  $\mathcal{O}(V^3)$ 

注意事项: 处理重边时, 应该对边权累加

```

1 int node[N], dist[N];
2 bool visit[N];
3
4 int solve(int n) {
5     int answer = INT_MAX;
6     for (int i = 0; i < n; ++i) {
7         node[i] = i;
8     }
9     while (n > 1) {

```

```

10     int max = 1;
11     for (int i = 0; i < n; ++i) {
12         dist[node[i]] = graph[node[0]][node[i]];
13         if (dist[node[i]] > dist[node[max]]) {
14             max = i;
15         }
16     }
17     int prev = 0;
18     memset(visit, 0, sizeof(visit));
19     visit[node[0]] = true;
20     for (int i = 1; i < n; ++i) {
21         if (i == n - 1) {
22             answer = std::min(answer, dist[node[max]]);
23             for (int k = 0; k < n; ++k) {
24                 graph[node[k]][node[prev]] =
25                     (graph[node[prev]][node[k]] += graph[node[k]][node[max]]);
26             }
27             node[max] = node[--n];
28         }
29         visit[node[max]] = true;
30         prev = max;
31         max = -1;
32         for (int j = 1; j < n; ++j) {
33             if (!visit[node[j]]) {
34                 dist[node[j]] += graph[node[prev]][node[j]];
35                 if (max == -1 || dist[node[max]] < dist[node[j]]) {
36                     max = j;
37                 }
38             }
39         }
40     }
41 }
42 return answer;
43 }

```

### 哈密尔顿回路 ( ORE 性质的图 )

ORE 性质 :

$$\forall x, y \in V \wedge (x, y) \notin E \quad s.t. \quad deg_x + deg_y \geq n$$

返回结果 : 从顶点 1 出发的一个哈密尔顿回路

使用条件 :  $n \geq 3$

```

1 int left[N], right[N], next[N], last[N];
2
3 void cover(int x) {
4     left[right[x]] = left[x];
5     right[left[x]] = right[x];
6 }
7
8 int adjacent(int x) {
9     for (int i = right[0]; i <= n; i = right[i]) {
10         if (graph[x][i]) {
11             return i;
12         }
13     }
14     return 0;
15 }
16
17 std::vector<int> solve() {
18     for (int i = 1; i <= n; ++i) {
19         left[i] = i - 1;
20         right[i] = i + 1;
21     }
22     int head, tail;
23     for (int i = 2; i <= n; ++i) {
24         if (graph[i][i]) {
25             head = 1;

```

```

26         tail = i;
27         cover(head);
28         cover(tail);
29         next[head] = tail;
30         break;
31     }
32 }
33 while (true) {
34     int x;
35     while (x = adjacent(head)) {
36         next[x] = head;
37         head = x;
38         cover(head);
39     }
40     while (x = adjacent(tail)) {
41         next[tail] = x;
42         tail = x;
43         cover(tail);
44     }
45     if (!graph[head][tail]) {
46         for (int i = head, j; i != tail; i = next[i]) {
47             if (graph[head][next[i]] && graph[tail][i]) {
48                 for (j = head; j != i; j = next[j]) {
49                     last[next[j]] = j;
50                 }
51                 j = next[head];
52                 next[head] = next[i];
53                 next[tail] = i;
54                 tail = j;
55                 for (j = i; j != head; j = last[j]) {
56                     next[j] = last[j];
57                 }
58                 break;
59             }
60         }
61     }
62     next[tail] = head;
63     if (right[0] > n) {
64         break;
65     }
66     for (int i = head; i != tail; i = next[i]) {
67         if (adjacent(i)) {
68             head = next[i];
69             tail = i;
70             next[tail] = 0;
71             break;
72         }
73     }
74 }
75 std::vector<int> answer;
76 for (int i = head; ; i = next[i]) {
77     if (i == 1) {
78         answer.push_back(i);
79         for (int j = next[i]; j != i; j = next[j]) {
80             answer.push_back(j);
81         }
82         answer.push_back(i);
83         break;
84     }
85     if (i == tail) {
86         break;
87     }
88 }
89 return answer;
90 }

```

必经点树

```

1 vector<int>G[maxn],rG[maxn],dom[maxn];
2 int n,m;
3 int dfn[maxn],rdfs[maxn],dfs_c,semi[maxn],idom[maxn],fa[maxn];
4 struct ufsets{
5     int fa[maxn],best[maxn];
6     int find(int x){
7         if(fa[x]==x)
8             return x;
9         int f=find(fa[x]);
10        if(dfn[semi[best[x]]]>dfn[semi[best[fa[x]]]])
11            best[x]=best[fa[x]];
12        fa[x]=f;
13        return f;
14    }
15    int getbest(int x){
16        find(x);
17        return best[x];
18    }
19    void init(){
20        for(int i=1;i<=n;i++){
21            fa[i]=best[i]=i;
22        }
23    }uf;
24    void init(){
25        uf.init();
26        for(int i=1;i<=n;i++){
27            semi[i]=i;
28            idom[i]=0;
29            fa[i]=0;
30            dfn[i]=rdfs[i]=0;
31        }
32        dfs_c=0;
33    }
34    void dfs(int u){
35        dfn[u]=++dfs_c;
36        rdfs[dfn[u]]=u;
37        for(int i=0;i<G[u].size();i++){
38            int v=G[u][i];
39            if(!dfn[v]){
40                fa[v]=u;
41                dfs(v);
42            }
43        }
44    }
45    void tarjan(){
46        for(int i=n;i>1;i--){
47            int tmp=1e9;
48            int y=rdfs[i];
49            for(int i=0;i<rG[y].size();i++){
50                int x=rG[y][i];
51                tmp=min(tmp,dfn[semi[uf.getbest(x)]]);
52            }
53            semi[y]=rdfs[tmp];
54            int x=fa[y];
55            dom[semi[y]].push_back(y);
56            uf.fa[y]=x;
57            for(int i=0;i<dom[x].size();i++){
58                int z=dom[x][i];
59                if(dfn[semi[uf.getbest(z)]]<dfn[x])
60                    idom[z]=uf.getbest(z);
61                else
62                    idom[z]=semi[z];
63            }
64            dom[x].clear();
65        }
66        semi[rdfs[1]]=1;
67        for(int i=2;i<=n;i++){
68            int x=rdfs[i];

```

```

70         if(idom[x]!=semi[x])
71             idom[x]=idom[idom[x]];
72     }
73     idom[rdfs[1]]=0;
74 }
75 init();
76 dfs(1);
77 tarjan();

```

## 字符串

### 模式匹配

#### KMP 算法

```

1 void build(char *pattern) {
2     int length = (int)strlen(pattern + 1);
3     fail[0] = -1;
4     for (int i = 1, j; i <= length; ++i) {
5         for (j = fail[i - 1]; j != -1 && pattern[i] != pattern[j + 1]; j =
6             ↪ fail[j]);
7         fail[i] = j + 1;
8     }
9 }
10 void solve(char *text, char *pattern) {
11     int length = (int)strlen(text + 1);
12     for (int i = 1, j; i <= length; ++i) {
13         for (j = match[i - 1]; j != -1 && text[i] != pattern[j + 1]; j = fail[j]);
14         match[i] = j + 1;
15     }
16 }
17 ///Hint: 1 - Base

```

#### 扩展 KMP 算法

返回结果：

$$next_i = lcp(text, text_{i...n-1})$$

```

1 void solve(char *text, int length, int *next) {
2     int j = 0, k = 1;
3     for (; j + 1 < length && text[j] == text[j + 1]; j++);
4     next[0] = length - 1;
5     next[1] = j;
6     for (int i = 2; i < length; ++i) {
7         int far = k + next[k] - 1;
8         if (next[i - k] < far - i + 1) {
9             next[i] = next[i - k];
10        } else {
11            j = std::max(far - i + 1, 0);
12            for (; i + j < length && text[j] == text[i + j]; j++);
13            next[i] = j;
14            k = i;
15        }
16    }
17 }
18 /// 0 - Base

```

#### AC 自动机

```

1 struct Node{
2     int Next[30], fail, mark;
3 }Tree[N];
4
5 void Init(){
6     memset(Tree, 0, sizeof Tree);

```

```

7   cnt = 1;
8
9   for (int i = 1; i <= n; i++){
10      char c;
11      int now = 1;
12      scanf("%s", s + 1);
13      int Length = strlen(s + 1);
14      for (int j = 1; j <= Length; j++){
15         c = s[j];
16         if (Tree[now].Next[c - 'a']) now = Tree[now].Next[c - 'a']; else
17            Tree[now].Next[c - 'a'] = ++ cnt, now = cnt;
18      }
19   }
20 }
21
22 void Build_Ac(){
23   int en = 0;
24   Q[0] = 1;
25   for (int fi = 0; fi <= en; fi++){
26      int now = Q[fi];
27      for (int next = 0; next < 26; next++){
28         if (Tree[now].Next[next])
29            {
30               int k = Tree[now].Next[next];
31               if (now == 1) Tree[k].fail = 1; else
32               {
33                  int h = Tree[now].fail;
34                  while (h && !Tree[h].Next[next]) h = Tree[h].fail;
35                  if (!h) Tree[k].fail = 1;
36                  else Tree[k].fail = Tree[h].Next[next];
37               }
38               Q[++ en] = k;
39            }
40      }
41   }
42   /// Hints : when not match , fail = 1
43

```

### 后缀三姐妹 后缀数组

```

1 struct Sa{
2   int heap[N],s[N],sa[N],r[N],tr[N],sec[N],m,cnt;
3   int h[19][N];
4   void Prep(){
5     for (int i=1; i<=m; i++) heap[i]=0;
6     for (int i=1; i<=n; i++) heap[s[i]]++;
7     for (int i=2; i<=m; i++) heap[i]+=heap[i-1];
8     for (int i=n; i>=1; i--) sa[heap[s[i]]--]=i;
9     r[sa[1]]=1; cnt=1;
10    for (int i=2; i<=n; i++){
11       if (s[sa[i]]!=s[sa[i-1]]) cnt++;
12       r[sa[i]]=cnt;
13    }
14    m=cnt;
15  }
16  void Suffix(){
17    int j=1;
18    while (cnt<n){
19       cnt=0;
20       for (int i=n-j+1; i<=n; i++) sec[++cnt]=i;
21       for (int i=1; i<=n; i++) if (sa[i]>j) sec[++cnt]=sa[i]-j;
22       for (int i=1; i<=n; i++) tr[i]=r[sec[i]];
23       for (int i=1; i<=m; i++) heap[i]=0;
24       for (int i=1; i<=n; i++) heap[tr[i]]++;
25       for (int i=2; i<=m; i++) heap[i]+=heap[i-1];
26       for (int i=n; i>=1; i--) sa[heap[tr[i]]--]=sec[i];
27       tr[sa[1]]=1; cnt=1;

```

```

28     for (int i=2; i<=n; i++){
29        if ((r[sa[i]]!=r[sa[i-1]]) || (r[sa[i]+j]!=r[sa[i-1]+j])) cnt++;
30        tr[sa[i]]=cnt;
31     }
32     for (int i=1; i<=n; i++) r[i]=tr[i];
33     m=cnt; j=j+j;
34   }
35 }
36 void Calc(){
37   int k=0;
38   for (int i=1; i<=n; i++){
39      if (r[i]==1) continue;
40      int j=sa[r[i]-1];
41      while ((i+k<=n) && (j+k<=n) && (s[i+k]==s[j+k])) k++;
42      h[0][r[i]]=k;
43      if (k) k--;
44   }
45   for (int i=1; i<19; i++){
46      for (int j=1; j+(1<<i)-1<=n; j++){
47         h[i][j]=min(h[i-1][j],h[i-1][j + (1<<(i-1)) + 1]);
48      }
49   }
50   int Query(int L,int R){
51     L=r[L], R=r[R];
52     if (L>R) swap(L,R);
53     L++;
54     int l0 = Lg[R-L+1];
55     return min(h[l0][L],h[l0][R-(1<<l0)+1]);
56   }
57   void Work(){
58     Prep(); Suffix(); Calc();
59   }
60 }P,S;/// Hints : 1 - Base

```

### 后缀数组 (dc3)

```

1 1//`DC3 待排序的字符串放在 r 数组中, 从 r[0] 到 r[n-1], 长度为 n, 且最大值小于 m.`
2 2//`约定除 r[n-1] 外所有的 r[i] 都大于 0, r[n-1]=0.`
3 3//`函数结束后, 结果放在 sa 数组中, 从 sa[0] 到 sa[n-1]。`
4 4//`r 必须开长度乘 3`
5 5#define maxn 10000
6 6#define F(x) ((x)/3+((x)%3==1?0:tb))
7 7#define G(x) ((x)<tb?(x)*3+1:((x)-tb)*3+2)
8
9 int wa[maxn],wb[maxn],wv[maxn],wss[maxn];
10 int s[maxn*3],sa[maxn*3];
11 int c0(int *r,int a,int b)
12 {
13     return r[a]==r[b]&&r[a+1]==r[b+1]&&r[a+2]==r[b+2];
14 }
15 int c12(int k,int *r,int a,int b)
16 {
17     if(k==2) return r[a]<r[b]||r[a]==r[b]&&c12(1,r,a+1,b+1);
18     else return r[a]<r[b]||r[a]==r[b]&&wv[a+1]<wv[b+1];
19 }
20 void sort(int *r,int *a,int *b,int n,int m)
21 {
22     int i;
23     for(i=0;i<n;i++) wv[i]=r[a[i]];
24     for(i=0;i<m;i++) wss[i]=0;
25     for(i=0;i<n;i++) wss[wv[i]]++;
26     for(i=1;i<m;i++) wss[i]+=wss[i-1];
27     for(i=n-1;i>=0;i--) b[--wss[wv[i]]]=a[i];
28 }
29 void dc3(int *r,int *sa,int n,int m)
30 {
31     int i,j,*rn=r+n,*san=sa+n,ta=0,tb=(n+1)/3,tbc=0,p;
32     r[n]=r[n+1]=0;

```

```

33     for(i=0;i<n;i++)
34         if(i%3!=0) wa[tbc++]=i;
35     sort(r+2,wa,wb,tbc,m);
36     sort(r+1,wb,wa,tbc,m);
37     sort(r,wa,wb,tbc,m);
38     for(p=1,rn[F(wb[0])]=0,i=1;i<tbc;i++)
39         rn[F(wb[i])]=c0(r,wb[i-1],wb[i])?p-1:p++;
40     if (p<tbc) dc3(rn,san,tbc,p);
41     else for (i=0;i<tbc;i++) san[rn[i]]=i;
42     for (i=0;i<tbc;i++)
43         if(san[i]<tb) wb[ta++]=san[i]*3;
44     if(n%3==1) wb[ta++]=n-1;
45     sort(r,wb,wa,ta,m);
46     for(i=0;i<tbc;i++)
47         wv[wb[i]=G(san[i])]=i;
48     for(i=0,j=0,p=0;i<ta && j<tbc;p++)
49         sa[p]=c12(wb[j]%3,r,wa[i],wb[j])?wa[i++]:wb[j++];
50     for(;i<ta;p++) sa[p]=wa[i++];
51     for(;j<tbc;p++) sa[p]=wb[j++];
52 }
53
54 int main(){
55     int n,m=0;
56     scanf("%d",&n);
57     for (int i=0;i<n;i++) scanf("%d",&s[i]),s[i]++,m=max(s[i]+1,m);
58     printf("%d\n",m);
59     s[n]=0;
60     dc3(s,sa,n,m);
61     for (int i=0;i<n;i++) printf("%d ",sa[i]);printf("\n");
62 }

```

### 后缀自动机

多串 LCS 对一个串建后缀自动机。其他串在上面匹配，因为是求所有串的公共子串，所以每个点记录每个串最长匹配长度的最小值。最后找到所有点中最长的一个即可。一个注意事项就是，当走到一个点时，还要更新它的 parent 树上的祖先的匹配长度，数组开两倍啦啦啦！

各长度子串出现次数最大值 给一个字符串 S，令  $F(x)$  表示 S 的所有长度为 x 的子串中，出现次数的最大值。构建字符串的自动机，对于每个节点，right 集合大小就是出现次数，maxs 就是它代表的最长长度，那么我们用  $|right(x)|$  去更新  $f[maxs[x]]$  的值，最后从大到小用  $f[i]$  去更新  $f[i-1]$  的值即可

```

1 struct Node{
2     int len, fail;
3     int To[30];
4 }T[N];
5 int Lst, Root, tot, ans;
6 char s[N];
7 int Len[N], Ans[N], Ord[N];
8 void Add(int x, int l){
9     int Nt = ++tot, p = Lst;
10    T[Nt].len = l;
11    for (;p && !T[p].To[x]; p = T[p].fail) T[p].To[x] = Nt;
12    if (!p) T[Nt].fail = Root; else
13    if (T[p].To[x].len == T[p].len + 1) T[Nt].fail = T[p].To[x];
14    else{
15        int q = ++tot, qt = T[p].To[x];
16        T[q] = T[qt];
17        T[q].len = T[p].len + 1;
18        T[qt].fail = T[Nt].fail = q;
19        for (;p && T[p].To[x] == qt; p = T[p].fail) T[p].To[x] = q;
20    }
21    Lst = Nt;
22 }
23 bool cmp(int a, int b){
24     return T[a].len < T[b].len;
25 }
26 int main(){
27     scanf("%s", s + 1);
28     int n = strlen(s + 1);

```

```

29     ans = n;
30     Root = tot = Lst = 1;
31     for (int i = 1; i <= n; i++)
32         Add(s[i] - 'a' + 1, i);
33     for (int i = 1; i <= tot; i++)
34         Ord[i] = i;
35     sort(Ord + 1, Ord + tot + 1, cmp);
36     for (int i = 1; i <= tot; i++)
37         Ans[i] = T[i].len;
38     bool flag = 0;
39     while (scanf("%s", s + 1) != EOF){
40         flag = 1;
41         int n = strlen(s + 1);
42         int p = Root, len = 0;
43         for (int i = 1; i <= tot; i++) Len[i] = 0;
44         for (int i = 1; i <= n; i++){
45             int x = s[i] - 'a' + 1;
46             if (T[p].To[x]) len++, p = T[p].To[x];
47             else {
48                 while (p && !T[p].To[x]) p = T[p].fail;
49                 if (!p) p = Root, len = 0;
50                 else len = T[p].len + 1, p = T[p].To[x];
51             }
52             Len[p] = max(Len[p], len);
53         }
54         for (int i = tot; i >= 1; i--){
55             int Cur = Ord[i];
56             Ans[Cur] = min(Ans[Cur], Len[Cur]);
57             if (Len[Cur] && T[Cur].fail)
58                 Len[T[Cur].fail] = T[Cur].fail.len;
59         }
60     }
61     if (flag){
62         ans = 0;
63         for (int i = 1; i <= tot; i++){
64             ans = max(ans, Ans[i]);
65         }
66     }
67     printf("%d\n", ans);
68     return 0;
69 }

```

### 回文三兄弟

#### 马拉车

```

1 void Manacher(){
2     R[1] = 1;
3     for (int i = 2, j = 1; i <= length; i++){
4         if (j + R[j] <= i){
5             R[i] = 0;
6         } else {
7             R[i] = min(R[j * 2 - i], j + R[j] - i);
8         }
9         while (i - R[i] >= 1 && i + R[i] <= length
10            && text[i - R[i]] == text[i + R[i]]){
11             R[i]++;
12         }
13         if (i + R[i] > j + R[j]){
14             j = i;
15         }
16     }
17 }
18 length = 0;
19 int n = strlen(s + 1);
20 for (int i = 1; i <= n; i++){
21     text[++length] = '*';
22     text[++length] = s[i];

```

```

23     }
24     text[++length] = '*';
25 // Hints: 1 - Base

```

## 回文自动机 (zky)

```

1 struct PAM{
2     int tot,last,str[maxn],nxt[maxn][26],n;
3     int len[maxn],suf[maxn],cnt[maxn];
4     int newnode(int l){
5         len[tot]=l;
6         return tot++;
7     }
8     void init(){
9         tot=0;
10        newnode(0); // tree0 is node 0
11        newnode(-1); // tree-1 is node 1
12        str[0]=-1;
13        suf[0]=1;
14    }
15    int find(int x){
16        while(str[n-len[x]-1]!=str[n])x=suf[x];
17        return x;
18    }
19    void add(int c){
20        str[++n]=c;
21        int u=find(last);
22        if(!nxt[u][c]){
23            int v=newnode(len[u]+2);
24            suf[v]=nxt[find(suf[u])][c];
25            nxt[u][c]=v;
26        }last=nxt[u][c];
27        cnt[last]++;
28    }
29    void count(){
30        for(int i=tot-1;i>=0;i--)cnt[suf[i]]+=cnt[i];
31    }
32 }P;
33 int main(){
34     P.init();
35     for(int i=0;i<n;i++)
36         P.add(s[i]-'a');
37     P.count();

```

## 循环串最小表示

```

1 string sol(char *s){
2     int n=strlen(s);
3     int i=0,j=1,k=0,p;
4     while(i<n&&j<n&&k<n){
5         int t=s[(i+k)%n]-s[(j+k)%n];
6         if(t==0)k++;
7         else if(t<0)j+=k+1,k=0;
8         else i+=k+1,k=0;
9         if(i==j)j++;
10    }p=min(i,j);
11    string S;
12    for(int i=p;i<p+n;i++)S.push_back(s[i%n]);
13    return S;
14 }

```

## 计算几何

## 二维基础

## 点类

```

1 struct P{
2     double x,y;

```

```

3     P turn90(){return P(-y,x);}
4 };
5 double det(P a,P b,P c){
6     return (b-a)*(c-a);
7 }
8 P intersect(L l1,L l2){
9     double s1=det(l1.a,l1.b,l2.a);
10    double s2=det(l1.a,l1.b,l2.b);
11    return (l2.a*s2-l2.b*s1)/(s2-s1);
12 }
13 P project(P p,L l){
14     return l.a+l.v()*((p-l.a)^l.v())/l.v().len2();
15 }
16 double dis(P p,L l){
17     return fabs((p-l.a)*l.v())/l.v().len();
18 }

```

## 圆类

```

1 struct C{
2     P o;
3     double r;
4     C(){ }
5     C(P _o,double _r):o(_o),r(_r){ }
6 };
7 // 求圆与直线的交点
8 //turn90() P(-y,x)
9 double fix(double x){return sgn(x)?x:0;}
10 bool intersect(C a, L l, P &p1, P &p2) {
11     double x = ((l.a - a.o)^(l.b - l.a)),
12            y = (l.b - l.a).len2(),
13            d = x * x - y * ((l.a - a.o).len2() - a.r * a.r);
14     if (sgn(d) < 0) return false;
15     d = max(d, 0.0);
16     P p = l.a - ((l.b - l.a) * (x / y)), delta = (l.b - l.a) * (sqrt(d) / y);
17     p1 = p + delta, p2 = p - delta;
18     return true;
19 }
20 // 求圆与圆的交点, 注意调用前要先判定重圆
21 bool intersect(C a, C b, P &p1, P &p2) {
22     double s1 = (a.o - b.o).len();
23     if (sgn(s1 - a.r - b.r) > 0 || sgn(s1 - fabs(a.r - b.r)) < 0) return false;
24     double s2 = (a.r * a.r - b.r * b.r) / s1;
25     double aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
26     P o = (b.o - a.o) * (aa / (aa + bb)) + a.o;
27     P delta = (b.o - a.o).norm().turn90() * sqrt(fix(a.r * a.r - aa * aa));
28     p1 = o + delta, p2 = o - delta;
29     return true;
30 }
31 // 求点到圆的切点, 按关于点的顺时针方向返回两个点
32 bool tang(const C &c, const P &p0, P &p1, P &p2) {
33     double x = (p0 - c.o).len2(), d = x - c.r * c.r;
34     if (d < eps) return false; // 点在圆上认为没有切点
35     P p = (p0 - c.o) * (c.r * c.r / x);
36     P delta = ((p0 - c.o) * (-c.r * sqrt(d) / x)).turn90();
37     p1 = c.o + p + delta;
38     p2 = c.o + p - delta;
39     return true;
40 }
41 // 求圆到圆的外公切线, 按关于 c1.o 的顺时针方向返回两条线
42 vector<L> extan(const C &c1, const C &c2) {
43     vector<L> ret;
44     if (sgn(c1.r - c2.r) == 0) {
45         P dir = c2.o - c1.o;
46         dir = (dir * (c1.r / dir.len())).turn90();
47         ret.push_back(L(c1.o + dir, c2.o + dir));

```



```

48     ret.push_back(L(c1.o - dir, c2.o - dir));
49 } else {
50     P p = (c1.o * -c2.r + c2.o * c1.r) / (c1.r - c2.r);
51     P p1, p2, q1, q2;
52     if (tang(c1, p, p1, p2) && tang(c2, p, q1, q2)) {
53         if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);
54         ret.push_back(L(p1, q1));
55         ret.push_back(L(p2, q2));
56     }
57 }
58 return ret;
59 }
60 // 求圆到圆的内共切线, 按关于 c1.o 的顺时针方向返回两条线
61 vector<L> intan(const C &c1, const C &c2) {
62     vector<L> ret;
63     P p = (c1.o * c2.r + c2.o * c1.r) / (c1.r + c2.r);
64     P p1, p2, q1, q2;
65     if (tang(c1, p, p1, p2) && tang(c2, p, q1, q2)) { // 两圆相切认为没有切线
66         ret.push_back(L(p1, q1));
67         ret.push_back(L(p2, q2));
68     }
69     return ret;
70 }

```

## 凸包

```

1 vector<P> convex(vector<P>p){
2     sort(p.begin(),p.end());
3     vector<P>ans,S;
4     for(int i=0;i<p.size();i++){
5         while(S.size()>=2
6             && sgn(det(S[S.size()-2],S.back(),p[i]))<=0)
7             S.pop_back();
8         S.push_back(p[i]);
9     }//dw
10    ans=S;
11    S.clear();
12    for(int i=(int)p.size()-1;i>=0;i--){
13        while(S.size()>=2
14            && sgn(det(S[S.size()-2],S.back(),p[i]))<=0)
15            S.pop_back();
16        S.push_back(p[i]);
17    }//up
18    for(int i=1;i+1<S.size();i++){
19        ans.push_back(S[i]);
20    }
21 }

```

## 半平面交

```

1 struct P{
2     int quad() const { return sgn(y) == 1 || (sgn(y) == 0 && sgn(x) >= 0); }
3 };
4 struct L{
5     bool onLeft(const P &p) const { return sgn((b - a)*(p - a)) > 0; }
6     L push() const { // push out eps
7         const double eps = 1e-10;
8         P delta = (b - a).turn90().norm() * eps;
9         return L(a - delta, b - delta);
10    };
11 };
12 bool sameDir(const L &l0, const L &l1) {
13     return parallel(l0, l1) && sgn((l0.b - l0.a)^(l1.b - l1.a)) == 1;
14 }
15 bool operator < (const P &a, const P &b) {
16     if (a.quad() != b.quad())
17         return a.quad() < b.quad();

```

```

18     else
19         return sgn((a*b)) > 0;
20 }
21 bool operator < (const L &l0, const L &l1) {
22     if (sameDir(l0, l1))
23         return l1.onLeft(l0.a);
24     else
25         return (l0.b - l0.a) < (l1.b - l1.a);
26 }
27 bool check(const L &u, const L &v, const L &w) {
28     return w.onLeft(intersect(u, v));
29 }
30 vector<P> intersection(vector<L> &l) {
31     sort(l.begin(), l.end());
32     deque<L> q;
33     for (int i = 0; i < (int)l.size(); ++i) {
34         if (i && sameDir(l[i], l[i - 1])) {
35             continue;
36         }
37         while (q.size() > 1
38             && !check(q[q.size() - 2], q[q.size() - 1], l[i]))
39             q.pop_back();
40         while (q.size() > 1
41             && !check(q[1], q[0], l[i]))
42             q.pop_front();
43         q.push_back(l[i]);
44     }
45     while (q.size() > 2
46         && !check(q[q.size() - 2], q[q.size() - 1], q[0]))
47         q.pop_back();
48     while (q.size() > 2
49         && !check(q[1], q[0], q[q.size() - 1]))
50         q.pop_front();
51     vector<P> ret;
52     for (int i = 0; i < (int)q.size(); ++i)
53         ret.push_back(intersect(q[i], q[(i + 1) % q.size()]));
54     return ret;
55 }

```

## 最小圆覆盖

```

1 point operator*(line A,line B){
2     point u=B.p-A.p;
3     double t=(B.v*u)/(B.v*A.v);
4     return A.p+A.v*t;
5 }
6 point get(point a,point b){
7     return (a+b)/2;
8 }
9 point get(point a,point b,point c){
10    if(a==b)return get(a,c);
11    if(a==c)return get(a,b);
12    if(b==c)return get(a,b);
13    line AB0=(line){(a+b)/2,Rev(a-b)};
14    line BC0=(line){(c+b)/2,Rev(b-c)};
15    return AB0*BC0;
16 }
17 random_shuffle(p+1,p+1+n);
18 O=p[1];r=0;
19 for(int i=2;i<=n;i++){
20     if(dis(p[i],O)<r+1e-6)continue;
21     O=get(p[1],p[i]);r=dis(O,p[i]);
22     for(int j=1;j<i;j++){
23         if(dis(p[j],O)<r+1e-6)continue;
24         O=get(p[i],p[j]);r=dis(O,p[i]);
25         for(int k=1;k<j;k++){
26             if(dis(p[k],O)<r+1e-6)continue;

```



```

27         0=get(p[i],p[j],p[k]);r=dis(0,p[i]);
28     }
29 }
30 }

```

## 多边形

判断点在多边形内部

```

1 bool InPoly(P p,vector<P>poly){
2     int cnt=0;
3     for(int i=0;i<poly.size();i++){
4         P a=poly[i],b=poly[(i+1)%poly.size()];
5         if(OnLine(p,L(a,b)))
6             return false;
7         int x=sgn(det(a,p,b));
8         int y=sgn(a.y-p.y);
9         int z=sgn(b.y-p.y);
10        cnt+=(x>0&&y<=0&&z>0);
11        cnt-=(x<0&&z<=0&&y>0);
12    }
13    return cnt;
14 }

```

## 其他

斯坦纳树

```

1 priority_queue<pair<int, int> > Q;
2 // m is key point
3 // n is all point
4 for (int s = 0; s < (1 << m); s++){
5     for (int i = 1; i <= n; i++){
6         for (int s0 = (s&(s-1)); s0 ; s0=(s&(s0-1)))
7             f[s][i] = min(f[s][i], f[s0][i] + f[s - s0][i]);
8         for (int i = 1; i <= n; i++) vis[i] = 0;
9         while (!Q.empty()) Q.pop();
10        for (int i = 1; i <= n; i++){
11            Q.push(mp(-f[s][i], i));
12        }
13        while (!Q.empty() && Q.top().first != -f[s][Q.top().second]) Q.pop();
14        if (Q.empty()) break;
15        int Cur = Q.top().second; Q.pop();
16        for (int p = g[Cur]; p; p = nxt[p]){
17            int y = adj[p];
18            if ( f[s][y] > f[s][Cur] + 1){
19                f[s][y] = f[s][Cur] + 1;
20                Q.push(mp(-f[s][y], y));
21            }
22        }
23    }
24 }

```

## 无敌的读入优化

```

1 namespace Reader {
2     const int L = (1 << 20) + 5;
3     char buffer[L], *S, *T;
4     __inline bool getchar(char &ch) {
5         if (S == T) {
6             T = (S = buffer) + fread(buffer, 1, L, stdin);
7             if (S == T) {
8                 ch = EOF;
9                 return false;
10            }
11        }
12        ch = *S ++;

```

```

13         return true;
14     }
15     __inline bool getint(int &x) {
16         char ch;
17         for (; getchar(ch) && (ch < '0' || ch > '9'); );
18         if (ch == EOF) return false;
19         x = ch - '0';
20         for (; getchar(ch), ch >= '0' && ch <= '9'; )
21             x = x * 10 + ch - '0';
22         return true;
23     }
24 }
25 Reader::getint(x);
26 Reader::getint(y);

```

## 最小树形图

```

1 const int maxn=1100;
2
3 int n,m , g[maxn][maxn] , used[maxn] , pass[maxn] , eg[maxn] , more ,
4     queue[maxn];
5 void combine (int id , int &sum ) {
6     int tot = 0 , from , i , j , k ;
7     for ( ; id!=0 && !pass[id] ; id=eg[id] ) {
8         queue[tot++]=id ; pass[id]=1;
9     }
10    for ( from=0; from<tot && queue[from]!=id ; from++);
11    if ( from==tot ) return ;
12    more = 1 ;
13    for ( i=from ; i<tot ; i++) {
14        sum+=g[eg[queue[i]]][queue[i]] ;
15        if ( i!=from ) {
16            used[queue[i]]=1;
17            for ( j = 1 ; j <= n ; j++) if ( !used[j] )
18                if ( g[queue[i]][j]<g[id][j] ) g[id][j]=g[queue[i]][j] ;
19        }
20    }
21    for ( i=1; i<=n ; i++) if ( !used[i] && i!=id ) {
22        for ( j=from ; j<tot ; j++){
23            k=queue[j];
24            if ( g[i][id]>g[i][k]-g[eg[k]][k] ) g[i][id]=g[i][k]-g[eg[k]][k];
25        }
26    }
27 }
28
29 int mdst( int root ) { // return the total length of MDST
30     int i , j , k , sum = 0 ;
31     memset ( used , 0 , sizeof ( used ) ) ;
32     for ( more =1; more ; ) {
33         more = 0 ;
34         memset ( eg,0,sizeof(eg) ) ;
35         for ( i=1 ; i <= n ; i ++ ) if ( !used[i] && i!=root ) {
36             for ( j=1 , k=0 ; j <= n ; j ++ ) if ( !used[j] && i!=j )
37                 if ( k==0 || g[j][i] < g[k][i] ) k=j ;
38             eg[i] = k ;
39         }
40         memset(pass,0,sizeof(pass));
41         for ( i=1; i<=n ; i++) if ( !used[i] && !pass[i] && i!= root ) combine (
42             i , sum ) ;
43     }
44     for ( i =1; i<=n ; i ++ ) if ( !used[i] && i!= root ) sum+=g[eg[i]][i];
45     return sum ;

```

## DLX

```

1 int n,m,K;
2 struct DLX{
3     int L[maxn],R[maxn],U[maxn],D[maxn];
4     int sz,col[maxn],row[maxn],s[maxn],H[maxn];
5     bool vis[233];
6     int ans[maxn],cnt;
7     void init(int m){
8         for(int i=0;i<=m;i++){
9             L[i]=i-1;R[i]=i+1;
10            U[i]=D[i]=i;s[i]=0;
11        }
12        memset(H,-1,sizeof H);
13        L[0]=m;R[m]=0;sz=m+1;
14    }
15    void Link(int r,int c){
16        U[sz]=c;D[sz]=D[c];U[D[c]]=sz;D[c]=sz;
17        if(H[r]<0)H[r]=L[sz]=R[sz]=sz;
18        else{
19            L[sz]=H[r];R[sz]=R[H[r]];
20            L[R[H[r]]]=sz;R[H[r]]=sz;
21        }
22        s[c]++;col[sz]=c;row[sz]=r;sz++;
23    }
24    void remove(int c){
25        for(int i=D[c];i!=c;i=D[i])
26            L[R[i]]=L[i],R[L[i]]=R[i];
27    }
28    void resume(int c){
29        for(int i=U[c];i!=c;i=U[i])
30            L[R[i]]=R[L[i]]=i;
31    }
32    int A(){
33        int res=0;
34        memset(vis,0,sizeof vis);
35        for(int i=R[0];i!=R[i]){if(!vis[i]){
36            vis[i]=1;res++;
37            for(int j=D[i];j!=i;j=D[j])
38                for(int k=R[j];k!=j;k=R[k])
39                    vis[col[k]]=1;
40        }
41        return res;
42    }
43    void dfs(int d,int &ans){
44        if(R[0]==0){ans=min(ans,d);return;}
45        if(d+A())>ans)return;
46        int tmp=23333,c;
47        for(int i=R[0];i!=R[i]){
48            if(tmp>s[i])tmp=s[i],c=i;
49        }
50        for(int i=D[c];i!=c;i=D[i]){
51            remove(i);
52            for(int j=R[i];j!=i;j=R[j])remove(j);
53            dfs(d+1,ans);
54            for(int j=L[i];j!=i;j=L[j])resume(j);
55            resume(i);
56        }
57    }
58    void del(int c){
59        L[R[c]]=L[c];R[L[c]]=R[c];
60        for(int i=D[c];i!=c;i=D[i])
61            for(int j=R[i];j!=i;j=R[j])
62                U[D[j]]=U[j],D[U[j]]=D[j],--s[col[j]];
63    }
64    void add(int c){
65        R[L[c]]=L[R[c]]=c;
66        for(int i=U[c];i!=c;i=U[i])
67            for(int j=L[i];j!=i;j=L[j])
68                ++s[col[U[D[j]]]=D[U[j]]=j];

```

```

68    }
69    bool dfs2(int k){
70        if(!R[0]){
71            cnt=k;return 1;
72        }
73        int c=R[0];
74        for(int i=R[0];i!=R[i])
75            if(s[c]>s[i])c=i;
76        del(c);
77        for(int i=D[c];i!=c;i=D[i]){
78            for(int j=R[i];j!=i;j=R[j])
79                del(col[j]);
80            ans[k]=row[i];if(dfs2(k+1))return true;
81            for(int j=L[i];j!=i;j=L[j])
82                add(col[j]);
83        }
84        add(c);
85        return 0;
86    }
87 }dlx;
88 int main(){
89     dlx.init(n);
90     for(int i=1;i<=m;i++)
91         for(int j=1;j<=n;j++)
92             if(dis(station[i],city[j])<mid-eps)
93                 dlx.Link(i,j);
94     dlx.dfs(0,ans);
95 }

```

## 某年某月某日是星期几

```

1 int solve(int year, int month, int day) {
2     int answer;
3     if (month == 1 || month == 2) {
4         month += 12;
5         year--;
6     }
7     if ((year < 1752) || (year == 1752 && month < 9) ||
8         (year == 1752 && month == 9 && day < 3)) {
9         answer = (day + 2 * month + 3 * (month + 1) / 5 + year + year / 4 + 5) % 7;
10    } else {
11        answer = (day + 2 * month + 3 * (month + 1) / 5 + year + year / 4
12                - year / 100 + year / 400) % 7;
13    }
14    return answer;
15 }

```

枚举大小为  $k$  的子集使用条件： $k > 0$ 

```

1 void solve(int n, int k) {
2     for (int comb = (1 << k) - 1; comb < (1 << n); ) {
3         // ...
4         int x = comb & -comb, y = comb + x;
5         comb = (((comb & ~y) / x) >> 1) | y;
6     }
7 }

```

## 环状最长公共子串

```

1 int n, a[N << 1], b[N << 1];
2
3 bool has(int i, int j) {
4     return a[(i - 1) % n] == b[(j - 1) % n];
5 }

```

```

6
7 const int DELTA[3][2] = {{0, -1}, {-1, -1}, {-1, 0}};
8
9 int from[N][N];
10
11 int solve() {
12     memset(from, 0, sizeof(from));
13     int ret = 0;
14     for (int i = 1; i <= 2 * n; ++i) {
15         from[i][0] = 2;
16         int left = 0, up = 0;
17         for (int j = 1; j <= n; ++j) {
18             int upleft = up + 1 + !!from[i - 1][j];
19             if (!has(i, j)) {
20                 upleft = INT_MIN;
21             }
22             int max = std::max(left, std::max(upleft, up));
23             if (left == max) {
24                 from[i][j] = 0;
25             } else if (upleft == max) {
26                 from[i][j] = 1;
27             } else {
28                 from[i][j] = 2;
29             }
30             left = max;
31         }
32         if (i >= n) {
33             int count = 0;
34             for (int x = i, y = n; y; ) {
35                 int t = from[x][y];
36                 count += t == 1;
37                 x += DELTA[t][0];
38                 y += DELTA[t][1];
39             }
40             ret = std::max(ret, count);
41             int x = i - n + 1;
42             from[x][0] = 0;
43             int y = 0;
44             while (y <= n && from[x][y] == 0) {
45                 y++;
46             }
47             for (; x <= i; ++x) {
48                 from[x][y] = 0;
49                 if (x == i) {
50                     break;
51                 }
52                 for (; y <= n; ++y) {
53                     if (from[x + 1][y] == 2) {
54                         break;
55                     }
56                     if (y + 1 <= n && from[x + 1][y + 1] == 1) {
57                         y++;
58                         break;
59                     }
60                 }

```

```

61     }
62     }
63 }
64 return ret;
65 }

```

## LLMOD

```

1 LL multiplyMod(LL a, LL b, LL P) { // `需要保证 a 和 b 非负`
2     LL t = (a * b - LL((long double)a / P * b + 1e-3) * P) % P;
3     return t < 0 : t + P : t;
4 }

```

## Java

### 基础模板

```

1 public class Main {
2     public static void main(String[] args) {
3         InputReader in = new InputReader(System.in);
4         PrintWriter out = new PrintWriter(System.out);
5     }
6 }
7 public static class cmp implements Comparator<edge>{
8     public int compare(edge a, edge b){
9         if(a.w < b.w) return 1;
10        if(a.w > b.w) return -1;
11        return 0;
12    }
13 }
14 class InputReader {
15     public BufferedReader reader;
16     public StringTokenizer tokenizer;
17     public InputReader(InputStream stream) {
18         reader = new BufferedReader(new InputStreamReader(stream), 32768);
19         tokenizer = null;
20     }
21     public String next() {
22         while (tokenizer == null || !tokenizer.hasMoreTokens()) {
23             try {
24                 tokenizer = new StringTokenizer(reader.readLine());
25             } catch (IOException e) {
26                 throw new RuntimeException(e);
27             }
28         }
29         return tokenizer.nextToken();
30     }
31     public int nextInt() {
32         return Integer.parseInt(next());
33     }
34 }

```