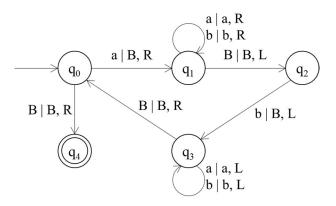
## **Turing Machines**

Design a TM for each of the following languages

1. 
$$L = \{a^nb^n \mid n \ge 0\}$$
  
Let  $w = aaabbb$ 

We define a TM which converts the given string to a smaller one such that it accepts the smaller string only if the initial string is valid.

It can be done by converting w = aaabbb to w = aabb by removing the first 'a' and the last 'b'. We continue this process and the string is accepted if  $w = \varepsilon$  at the end.

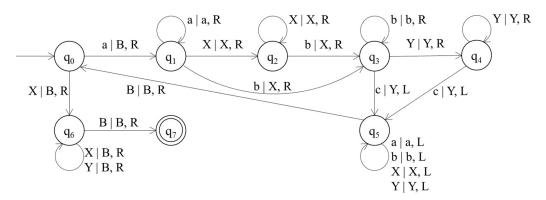


$$\begin{split} M &= (Q, \Sigma, \pmb{\tau}, \delta, \, q_0, \, B, \, F \} \\ Q &= \{q_0, \, q_1, \, q_2, \, q_3, \, q_4\}, \, \Sigma = \{a, \, b\}, \, \pmb{\tau} = \{a, \, b, \, B\}, \, F = \{q_4\} \\ q_0 \text{ is start state, } B \text{ is Blank tape symbol} \end{split}$$

| δ                    | a                       | b                       | В             |
|----------------------|-------------------------|-------------------------|---------------|
| $\rightarrow$ q $_0$ | $(q_1, B, R)$           | -                       | (q4, B, R)    |
| qı                   | $(q_1, a, R)$           | $(q_1, b, R)$           | $(q_2, B, L)$ |
| q <sub>2</sub>       | -                       | (q <sub>3</sub> , B, L) | -             |
| q3                   | (q <sub>3</sub> , a, L) | $(q_3, b, L)$           | $(q_0, B, R)$ |
| *q4                  | -                       | -                       | -             |

ID for w = aaabbb

## 2. $L = \{a^n b^n c^n \mid n \ge 1\}$



$$\begin{split} M &= (Q, \Sigma, \tau, \delta, \, q_0, \, B, \, F \} \\ Q &= \{q_0, \, q_1, \, q_2, \, q_3, \, q_4, \, q_5, \, q_6, \, q_7 \}, \, \Sigma = \{a, \, b\}, \, \tau = \{a, \, b, \, B, \, X, \, Y \}, \, F = \{q_7\} \\ q_0 \text{ is start state, } B \text{ is Blank tape symbol} \end{split}$$

| δ                    | a                       | b             | c             | X             | Y             | В             |
|----------------------|-------------------------|---------------|---------------|---------------|---------------|---------------|
| $\rightarrow$ q $_0$ | $(q_1, B, R)$           | -             | -             | $(q_6, B, R)$ | -             | -             |
| qı                   | $(q_1, a, R)$           | $(q_3, X, R)$ | -             | $(q_2, X, R)$ | -             | -             |
| q <sub>2</sub>       | -                       | $(q_3, X, R)$ | -             | $(q_2, X, R)$ | -             | -             |
| q <sub>3</sub>       | -                       | $(q_3, b, R)$ | $(q_5, Y, L)$ | ı             | $(q_4, Y, R)$ | -             |
| q4                   | -                       | -             | $(q_5, Y, L)$ | -             | $(q_4, Y, R)$ | -             |
| q <sub>5</sub>       | (q <sub>5</sub> , a, L) | $(q_5, b, L)$ | -             | $(q_5, X, L)$ | $(q_5, Y, L)$ | $(q_0, B, R)$ |
| q <sub>6</sub>       | -                       | -             | -             | $(q_6, B, R)$ | $(q_6, B, R)$ | $(q_7, B, R)$ |
| * <b>q</b> 7         | -                       | -             | -             | -             | -             | -             |

ID for w = aabbcc

 $Bq_0aabbcc \vdash Bq_1abbcc \vdash aq_1bbcc \vdash aXq_3bcc \vdash aXbq_3cc \vdash aXq_5bYc \vdash aq_5XbYc \vdash Bq_5aXbYc \vdash Bq_5BaXbYc \vdash Bq_0aXbYc \vdash Bq_1XbYc \vdash Xq_2bYc \vdash XXq_3Yc \vdash XXYq_4c \vdash XXq_5YY \vdash Xq_5XYY \vdash Bq_5XXYY \vdash Bq_5BXXY \vdash Bq_0XXYY \vdash Bq_6XYY \vdash Bq_6YY \vdash Bq_6Y \vdash Bq_6B \vdash Bq_7B$ 

3. 
$$L = \{ww^r \mid w \in \{a, b\}\}$$

$$\begin{split} M &= (Q, \Sigma, \pmb{\tau}, \delta, \, q_0, \, B, \, F \} \\ Q &= \{q_0, \, q_1, \, q_2, \, q_3, \, q_4, \, q_5, \, q_6 \}, \, \Sigma = \{a, \, b\}, \, \pmb{\tau} = \{a, \, b, \, B\}, \, F = \{q_6\} \\ q_0 \text{ is start state, } B \text{ is Blank tape symbol} \end{split}$$

| δ                    | a                       | b             | В                       |
|----------------------|-------------------------|---------------|-------------------------|
| $\rightarrow$ q $_0$ | $(q_1, B, R)$           | $(q_3, B, R)$ | $(q_6, B, R)$           |
| q1                   | $(q_1, a, R)$           | $(q_1, b, R)$ | $(q_2, B, L)$           |
| q <sub>2</sub>       | $(q_5, B, L)$           | -             | -                       |
| q3                   | (q <sub>3</sub> , a, R) | $(q_3, b, R)$ | (q <sub>4</sub> , B, L) |
| q <sub>4</sub>       | -                       | (q5, B, L)    | -                       |
| q5                   | (q <sub>5</sub> , a, L) | $(q_5, b, L)$ | $(q_0, B, R)$           |
| * <b>q</b> 6         | -                       | -             | -                       |

Let w = abaaba

4. 
$$L = \{0^n 1^{n+m} 0^m \mid m, n \ge 1\}$$

Let 
$$w = 001110$$

We will now look at 2 ways of solving this problem.

The usual approach:

$$M = (Q, \Sigma, \tau, \delta, q_0, B, F)$$

$$Q = \{q_0,\,q_1,\,q_2,\,q_3,\,q_4,\,q_5,\,q_6,\,q_7\},\, \Sigma = \{0,\,1\},\, \boldsymbol{\tau} = \{0,\,1,\,X,\,Y,\,B\},\, F = \{q_7\}$$

q<sub>0</sub> is start state, B is Blank tape symbol

| δ                    | 0             | 1             | X             | Y             | В             |
|----------------------|---------------|---------------|---------------|---------------|---------------|
| $\rightarrow$ q $_0$ | $(q_1, X, R)$ | -             | -             | $(q_3, Y, R)$ | -             |
| qı                   | $(q_1, 0, R)$ | $(q_2, Y, L)$ | -             | $(q_1, Y, R)$ | -             |
| q <sub>2</sub>       | $(q_2, 0, L)$ | -             | $(q_0, X, R)$ | $(q_2, Y, L)$ | -             |
| q3                   | -             | $(q_4, Y, R)$ | $(q_6, X, R)$ | $(q_3, Y, R)$ | -             |
| q <sub>4</sub>       | (q5, X, L)    | (q4, 1, R)    | (q4, X, R)    | -             | -             |
| q5                   | -             | $(q_5, 1, L)$ | $(q_5, X, L)$ | $(q_3, Y, R)$ | -             |
| <b>q</b> 6           | -             | -             | $(q_6, X, R)$ | -             | $(q_7, B, R)$ |
| *q7                  | -             | -             | -             | -             | -             |

Let w = 001110

 $Bq_0001110 \vdash Xq_101110 \vdash X0q_11110 \vdash Xq_20Y110 \vdash Bq_2X0Y110 \vdash Xq_00Y110$ 

 $\vdash XXq_1Y110 \vdash XXq_1Y110 \vdash XXYq_1110 \vdash XXq_2YY10 \vdash Xq_2XYY10 \vdash$ 

 $XXq_0YY10 \vdash XXYq_3Y10 \vdash XXYYq_310 \vdash XXYYYq_40 \vdash XXYYq_5YX \vdash$ 

 $XXYYYq_3X \vdash XXYYYXq_6B \vdash XXYYYXBq_7B$ 

Optimized approach:

$$001110 \rightarrow 0110 \rightarrow 10 \rightarrow \epsilon$$

If the string begins with 01, replace both of them by blank and move to a different state which accepts strings of the form  $1^{m}0^{m}$ .

If it begins with 00, replace both of them by blank and change the first 1 to 0. It is equivalent to removing a 0 and a 1. i. e. 00011110 = BB001110.

Repeat the same for the other half.

This method can be tweaked easily to work for cases where m, n = 0.

$$M = (Q, \Sigma, \tau, \delta, q_0, B, F)$$

$$Q = \{q_0,\,q_1,\,q_2,\,q_3,\,q_4,\,q_5,\,q_6,\,q_7,\,q_8,\,q_9\},\, \Sigma = \{0,\,1\},\, \boldsymbol{\tau} = \{0,\,1,\,B\},\, F = \{q_9\}$$

| δ                    | 0                       | 1             | В             |
|----------------------|-------------------------|---------------|---------------|
| $\rightarrow$ q $_0$ | $(q_1, B, R)$           | -             | -             |
| q1                   | $(q_2, B, R)$           | $(q_4, B, R)$ | -             |
| q <sub>2</sub>       | $(q_2, 0, R)$           | $(q_3, 0, L)$ | -             |
| q3                   | $(q_3, 0, L)$           | -             | $(q_0, B, R)$ |
| q <sub>4</sub>       | -                       | (q5, B, R)    | -             |
| q5                   | $(q_8, B, R)$           | $(q_6, B, R)$ | -             |
| q <sub>6</sub>       | (q <sub>7</sub> , 1, L) | $(q_6, 1, R)$ | -             |
| <b>q</b> 7           | -                       | $(q_7, 1, L)$ | $(q_4, B, R)$ |
| q8                   | -                       | -             | (q9, B, R)    |
| * <b>q</b> 9         | -                       | -             | -             |

Let w = 00011111100

 $Bq_000011111100 \vdash BBq_10011111100 \vdash BBBq_2011111100 \vdash BBB0q_211111100 \vdash$ 

BBBq300111100 \( \bar{B}Bq3B00111100 \) \( \bar{B}q\_000111100 \) \( \bar{B}q\_10111100 \) \( \bar{B}q\_10111100 \)

 $Bq_2111100 \vdash Bq_3011100 \vdash Bq_3B011100 \vdash Bq_0011100 \vdash Bq_111100 \vdash Bq_41100$ 

 $\vdash BBq_5100 \vdash BBBq_600 \vdash BBBq_7B10 \vdash Bq_410 \vdash Bq_50 \vdash Bq_8B \vdash Bq_9B$ 

It is clear that this approach is faster since the ID is smaller for a longer string.

5. 
$$L = \{w \mid n_a(w) = n_b(w)\}$$

$$M = (Q, \Sigma, \boldsymbol{\tau}, \delta, \, q_0, \, B, \, F \}$$

$$Q = \{q_0,\,q_1,\,q_2,\,q_3,\,q_4,\,q_5\},\, \Sigma = \{a,\,b\},\, \boldsymbol{\tau} = \{a,\,b,\,X,\,Y,\,B\},\, F = \{q_5\}$$

 $q_0$  is start state, B is Blank tape symbol

| δ                    | a                       | b             | X                       | Y                       | В                       |
|----------------------|-------------------------|---------------|-------------------------|-------------------------|-------------------------|
| $\rightarrow$ q $_0$ | $(q_1, X, R)$           | $(q_0, b, R)$ | $(q_0, X, R)$           | $(q_0, Y, R)$           | (q <sub>4</sub> , B, L) |
| q1                   | (q <sub>1</sub> , a, R) | $(q_1, b, R)$ | $(q_1, X, R)$           | $(q_1, Y, R)$           | $(q_2, B, L)$           |
| q <sub>2</sub>       | (q <sub>2</sub> , a, L) | $(q_3, Y, L)$ | $(q_2, X, L)$           | $(q_2, Y, L)$           | -                       |
| q3                   | $(q_3, a, L)$           | $(q_3, b, L)$ | $(q_3, X, L)$           | $(q_3, Y, L)$           | $(q_0, B, R)$           |
| q4                   | -                       | -             | (q <sub>4</sub> , B, L) | (q <sub>4</sub> , B, L) | $(q_5, B, R)$           |
| *q5                  | -                       | -             | -                       | -                       | -                       |