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DETAILED LECTURE NOTES

Unit - III

PAGE NO.

Recurrence Relation (Difference equation)

A Relation which involves an independent variable x , a dependent variable y and one or more than one differences $\Delta y, \Delta^2 y, \Delta^3 y, \dots$ is called a Recurrence relation.

$$F(x, y, \Delta y, \Delta^2 y, \dots) = 0 \quad \text{--- (1)}$$

Or.

$$y_{n+2} - 5y_{n+1} + 6y_n = 0$$

$$y_{n+1} - 3y_n = 3^n$$

Order and degree of Recurrence Relation

The order of a Recurrence Relation is equal to the difference between highest and lowest subscripts (arguments) divided by unit difference interval.

$$n+2-n=2, \quad n+1-n=1 \text{ is 1 order.} \\ \text{ie 2 order}$$

The degree of a difference Δ^n is equal to the power of highest term. ie highest term.

Solution of a Recurrence Relation

A Relation between the independent variable and dependent Variable is said to be a solution of a recurrence relation if this Relation satisfies the recurrence relation.

General Solution:-

A general solution of a recurrence relation of order n is a solution which involves n arbitrary constants. A general solution is also called complete primitive.

Particular Solution:- A particular solution is a solution obtained from the general solution by assigning particular values to one or more arbitrary constants.

Linear Recurrence Relation

A recurrence relation of degree n is called linear recurrence relation.

Linear Recurrence Relation with Constant Coefficients

form $c_0 a_r + c_1 a_{r-1} + c_2 a_{r-2} + \dots + c_k a_{r-k} = f(r)$

is a linear recurrence relation of order K where $c_0, c_1, c_2, \dots, c_k$ are constants. ————— (1)

Total Solution

$$a_r = H.S + P.S$$

$$[a_r = a_r^H + a_r^P]$$

Homogeneous Solution

$$a_r + b_1 a_{r-1} + b_2 a_{r-2} = 0 \quad \text{--- (1)}$$

is a linear recurrence relation of order 2

$$\text{let } a_r = m^r \quad \therefore m^r + b_1 m^{r-1} + b_2 m^{r-2} = 0$$

$$\text{Put } r=2 \quad m^2 + b_1 m + b_2 = 0 \quad \text{--- (2)}$$

is called Auxiliary equation.

Find the Roots suppose the Roots are $m = m_1, m_2$

Case-I If the Roots are real and distinct.

$$a_r = C_1 (m_1)^r + C_2 (m_2)^r$$

Case-II If the Roots are real and equal. $m = m_1, m_1$

$$a_r^H = (C_1 + C_2 r) (m_1)^r$$



POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

Case III If the Roots are Imaginary

$$m = \alpha \pm i\beta$$

$$a_r^u = (c_1 \cos \theta + c_2 \sin \theta) (R)^r \quad \text{where } R = \sqrt{\alpha^2 + \beta^2} \\ \theta = \tan^{-1}(\frac{\beta}{\alpha})$$

Particular Solution

$$\boxed{a_r^{(P)} = \frac{R \cdot h \cdot s}{f(m)} = \frac{f(r)}{f(m)}}$$

Q. Solve the following recurrence Relations

- i) $a_r - 6a_{r-1} + 8a_{r-2} = 0$ given $a_0 = 3$, and $a_1 = 2$
- ii) $2a_r - 5a_{r-1} + 2a_{r-2} = 0$ given $a_0 = 0$, $a_1 = 1$

Sol: - given $a_r - 6a_{r-1} + 8a_{r-2} = 0$ — (1)

order is 2

$$\therefore R \cdot h \cdot s = 0$$

$$\text{Put } [a_r = m^r] \quad m^r - 6m^{r-1} + 8m^{r-2} = 0 \quad P \cdot S = 0$$

$$\text{A.O.E.} \quad \text{Put } [r=2] \quad m^2 - 6m + 8 = 0 \quad — (2)$$

$m = 2, 4$ (Real and distinct)

$$a_r^{(u)} = C_1 (2)^r + C_2 (4)^r$$

$$a_r^{(P)} = 0 \quad \therefore R \cdot h \cdot s = 0$$

$$\text{Complete Sol. } a_r = a_r^u + a_r^P = C_1 (2)^r + C_2 (4)^r — (3)$$

$$r=0, \quad a_0 = 3$$

$$3 = C_1 + C_2 — (4)$$

$$r=1, \quad a_1 = 2$$

$$2 = 2C_1 + 4C_2 — (5)$$

$$1 = C_1 + 2C_2 \Rightarrow 1 = 3 - C_2 + 2C_2 \therefore C_2 = -2, C_1 = 5$$

True

$$a_2 = (5)(2)^r - 2(4)^r \quad A$$

Q. Solve the following recurrence Relations

i) $a_r - 6a_{r-1} + 9a_{r-2} = 0$ given $a_0 = 0, a_1 = 1$

ii) $a_r + a_{r-2} = 0$ given $a_0 = 0, a_1 = 1$

iii) $a_r + 2a_{r-1} + 2a_{r-2} = 0$ given $a_0 = 0, a_1 = -1$

iv) $a_r + 6a_{r-1} + 25a_{r-2} = 0$

v) Solve $a_r - 7a_{r-2} - 6a_{r-3} = 0$ with initial condition

Q. $a_0 = 9, a_1 = 10, a_2 = 32,$

vi) $a_r + 3a_{r-1} + 3a_{r-2} + a_{r-3} = 0$ with $a_0 = 1, a_1 = -2$ and $a_2 = -1$

Particular Solutions (a_p^P)

i) Method of Inspection (or Method of undetermined coefficients)

S.No	Terms in $f(r)$	Trial Sol a_p^P
i)	b^r	$A \cdot b^r$
ii)	a Polynomial of degree K ,	$A_0 + A_1 r + A_2 r^2 + \dots + A_K r^K$.
iii)	$b^r \cdot (\text{a Polynomial of degree } m)$	$b^r \cdot (A_0 + A_1 r + A_2 r^2 + \dots + A_K r^K)$
iv)	$\sin br$ or $\cos br$	$A \sin br + B \cos br$
(v)	$a^r \sin br$	$a^r (A \sin br + B \cos br)$
vi)	$a^r \cos br$	$a^r (A \sin br + B \cos br)$



POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

Q. Solve $a_2 s^2 + 5a_1 s + 6a_0 = s^r$ — (1)

Sol:- A.E. becomes

$$m^2 - sm + 6 = 0 \quad (2)$$

Find the Roots

$m = 2, 3$ (Real and distinct)

H.S will be.

$$a_r^{(n)} = C_1 (2)^r + C_2 (3)^r \quad (3)$$

P.S. $a_r^P = \frac{R \cdot h \cdot S}{\cdot} = A (5)^r$ by Trial Method
Assume.

Putting in (1)

$$A (5)^r - 5 \cdot A (5)^{r-1} + 6 \times A (5)^{r-2} = (5)^r$$

$$A (5)^r \left[1 - 5 + \frac{6}{5^2} \right] = (5)^r$$

$$A \left[\frac{6}{25} \right] = 1 \quad \therefore A = \frac{25}{6}$$

Thus P.S will be $a_r^{(P)} = A (5)^r = \frac{25}{6} (5)^r \quad (4)$

Total Sol $a_r = n \cdot S + P.S$

$$\boxed{a_r = C_1 (2)^r + C_2 (3)^r + \frac{25}{6} (5)^r}$$

A

Q. Solve $a_r - 2a_{r-1} = f(r)$ where $f(r) = 7r$
Determine the Particular Sol.

Q. Solve the Recurrence Relation

$$a_r - 5a_{r-1} + 6a_{r-2} = 2^{r+2}, r \geq 2$$

with boundary conditions $a_0 = 1$ and $a_1 = 1$

Sol: - A.E becomes $m^2 - 5m + 6 = 0$

$$a_r^{(h)} = C_1(2)^r + C_2(3)^r \quad m=2,2 \text{ (real and diff)} \quad \textcircled{1}$$

Trial Sol: -

The Particular Sol (Trial Sol) corresponding to the Term 2^{r+2} occurring is $A_0 + A_1 r$

$$a_r^{(P)} = A_0 + A_1 r \quad \textcircled{2}$$

Substituting $\textcircled{2}$ in given eqn

$$(A_0 + A_1 r) - 5(A_0 + A_1(r-1)) + 6(A_0 + A_1(r-2)) = 2^{r+2}$$

$$(A_0 - 5A_0 + 6A_0) + A_1(r - 5(r-1) + 6(r-2)) = 2^{r+2}$$

$$2A_0 + A_1(2r - 5r + 5 + 6r - 12) = 2^{r+2}$$

$$2A_0 + A_1(2r - 7) = 2^{r+2}$$

$$(2A_0 - 7A_1) + 2A_1 r = 2^{r+2}$$

$$2A_0 - 7A_1 = 2 \quad 2A_1 = 1 \quad \therefore A_1 = \frac{1}{2}$$

$$2A_0 - \frac{7}{2} = 2 \quad \therefore 2A_0 = 2 + \frac{7}{2} = \frac{11}{2} \quad \therefore A_0 = \frac{11}{4}$$

$$\therefore P.S \quad a_r^{(P)} = \frac{11}{4} + \frac{1}{2}r$$

$$\text{Total Sol} \quad a_r = C_1(2)^r + C_2(3)^r + \frac{11}{4} + \frac{1}{2}r$$

Ans.



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COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO. _____

Q. Solve $a_r - 2a_{r-1} + a_{r-2} = 7$

Sol:- A.E becomes $m^2 - 2m + 1 = 0$, $m=1, 1$ Real and equal
H.S will be. $a_r^{(p)} = (c_1 + c_2 r) (1)^r \quad \text{--- } 1$

By Method of undetermined coefficient

its trial solution will be $a_r = A r^2 \quad \text{--- } 2$

The question arises that why $A r^2$ has been chosen corresponding to 7 and not $A r$ or A , the reason is simple. Since 1 is a double root of characteristic eq^m, or we may consider like this since c_1 and $c_2 r$ appears in $a_r^{(p)}$

Now $7A r^2 - 2A \times 1 \times (r-1)^2 + A(r-2)^2 = 7$

$7A = 7 \quad \therefore A = \frac{7}{2}$

$\therefore P\text{-s will be } a_r^{(P)} = \frac{7}{2} r^2 \quad \text{--- } 3$

Hence the total solution of given recurrence relation is

$$\boxed{a_r = (c_1 + c_2 r) + \frac{7}{2} r^2}$$

I Solve the recurrence relation.

$$a_r - 5a_{r-1} + 6a_{r-2} = 2^r + r, \quad r \geq 2$$

with boundary conditions

Q. Solve $a_r - 4a_{r-1} + 4a_{r-2} = (r+1)^2 \quad r \geq 2$

Generating function

The Method of generating functions is a Powerful method to solve the difference equations (Recurrence Relations)

Definition: — If y_0, y_1, y_2, \dots is a sequence of real numbers. The function y defined for some interval of real numbers containing zero, whose value at t is given by the series

$$Y(t) = \sum_{n=0}^{\infty} y_n t^n = y_0 + y_1 t + y_2 t^2 + \dots + y_n t^n + \dots$$

is called generating function of the sequence $\{y_n\}$.

To discuss the generating function method for solving a linear difference equation: —

Consider a linear recurrence relation of first order

$$y_{n+1} = A y_n + C, \quad n = 0, 1, 2, 3, \dots \quad (1)$$

where A and C are constant.

Consider the generating function

$$Y(t) = \sum_{n=0}^{\infty} y_n t^n = y_0 + y_1 t + y_2 t^2 + \dots \quad (2)$$

Multiplying (1) by t^n and summing from $n=0$ to $n=\infty$ we get

$$\sum_{n=0}^{\infty} y_{n+1} t^n = \sum_{n=0}^{\infty} A y_n t^n + \sum_{n=0}^{\infty} C t^n$$

$$(y_1 + y_2 t + y_3 t^2 + \dots) = A(y_0 + y_1 t + y_2 t^2 + y_3 t^3 + \dots) + C(t + t^2 + t^3 + \dots)$$



POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

$$\frac{1}{t} (y_1 t + y_2 t^2 + y_3 t^3 + \dots) = A(y(t)) + C(1 + t + t^2 + \dots)$$

$$\frac{1}{t} [y(t) - y_0] = A y(t) + C(1-t)^{-1}$$

$$y(t) - y_0 = A t y(t) + \frac{C t}{1-t}$$

$$y(t) - A t y(t) = y_0 + \frac{C t}{1-t}$$

$$y(t)(1-At) = y_0 + \frac{Ct}{1-t} \Rightarrow y(t) = \frac{y_0}{1-At} + \frac{Ct}{(1-t)(1-At)} \quad \text{--- (1)}$$

$$y(t) = y_0(1-At)^{-1} + \frac{Ct}{(1-t)(1-At)}$$

by Partial
Fraction

$$\frac{Ct}{(1-t)(1-At)} = \frac{B}{1-t} + \frac{D}{(1-At)}$$

$$Ct = B(1-At) + D(1-t) \quad \text{Put } t=1$$

$$B = \frac{C}{1-A}$$

$$\text{Put } t = \frac{1}{A}$$

$$\frac{C}{A} = D\left(1-\frac{1}{A}\right) \quad \therefore D = \frac{C}{A\left(1-\frac{1}{A}\right)} = \frac{C}{(A-1)} = -\frac{C}{1-A}$$

$$y(t) = y_0(1-At)^{-1} + \left(\frac{C}{1-A}\right) \left[\frac{1}{1-t} - \frac{1}{1-At}\right]$$

$$= y_0(1+At+A^2t^2+\dots) + \left(\frac{C}{1-A}\right) [(1-t)^{-1} - (1-At)^{-1}]$$

$$\sum_{n=0}^{\infty} y_n t^n = y_0 \sum_{n=0}^{\infty} A^n t^n + \left(\frac{C}{1-A}\right) \left[\sum_{n=0}^{\infty} t^n - \sum_{n=0}^{\infty} A^n t^n \right]$$

$$\sum_{n=0}^{\infty} y_n t^n = y_0 \sum_{n=0}^{\infty} \left[A^n + \left(\frac{C}{1-A}\right) [1-A^n] \right] t^n$$

Equating the coefficients of t^n on both sides

$$y_n = y_0 n! + \frac{c}{1-A} (1-A)^n \quad \text{if } A \neq 1$$

$$\text{If } A=1 \quad y(t) = y_0 + \frac{ct}{(1-t)^2}$$

$$\begin{aligned} \sum_{n=0}^{\infty} y_n t^n &= y_0 (1-t)^{-1} + ct(1-t)^{-2} \\ &= y_0 (1+t+t^2+\dots) + ct(1+2t+3t^2+\dots) \\ &= y_0 \sum_{n=0}^{\infty} t^n + c \sum_{n=0}^{\infty} nt^n \end{aligned}$$

Equating two coefficient of t^n on both sides we have.

$$\boxed{y_n = y_0 + ct^n}$$

$$y_n = \begin{cases} y_0 n! + \frac{c}{1-A} (1-A)^n & \text{if } A \neq 1 \\ y_0 + ct^n & \text{if } A = 1 \end{cases}$$

Q Apply Generating Function Technique to solve
the initial value problem $y_{n+1} - 2y_n = 0$ with $y_0 = 1$

Sol:- given $y_{n+1} - 2y_n = 0 \quad \text{--- (1)}$

Consider the generating function $Y(t)$ is given

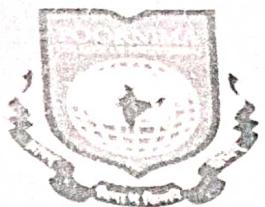
$$Y(t) = \sum_{n=0}^{\infty} y_n t^n = y_0 + y_1 t + y_2 t^2 + \dots$$

Multiplying by t^n of given eqⁿ (1) and summing from
 $n=0$ to $n=\infty$ we get

$$\sum_{n=0}^{\infty} y_{n+1} t^n - 2 \sum_{n=0}^{\infty} y_n t^n = 0$$

$$(y_1 + y_2 t + y_3 t^2 + \dots) - 2(y_0 + y_1 t + y_2 t^2 + \dots) = 0$$

$$\frac{1}{t} [y_1 t + y_2 t^2 + y_3 t^3 + \dots] - 2 Y(t) = 0$$



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COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

$$\frac{1}{t} [Y(t) - y_0] - 2Y(t) = 0 \quad \text{given } y_0 = 1$$

$$\frac{1}{t} [Y(t) - 1] = 2Y(t)$$

$$Y(t) - 1 = 2tY(t) + 1 \quad \therefore Y(t) - 2tY(t) = 1$$

$$(1-2t)Y(t) = 1$$

$$\therefore Y(t) = \frac{1}{1-2t} = (1-2t)^{-1} = 1 + 2t + (2t)^2 + (2t)^3 + \dots$$

$$\sum_{n=0}^{\infty} y_n t^n = \sum_{n=0}^{\infty} (2t)^n$$

equating the coefficients of the both sides

$$y_n = (2)^n \quad \text{which is the required sol.}$$

Q. (2) Solve $y_{n+2} - 7y_{n+1} + 10y_n = 0$ with $y_0 = 0, y_1 = 3$ by the Method of generating function. —①

Sol: — Consider generating function

$$Y(t) = \sum_{n=0}^{\infty} y_n t^n = (y_0 + y_1 t + y_2 t^2 + \dots) \quad \text{—①}$$

Multiplying t^n in given eqⁿ both sides and summing $n=0$ to $n=\infty$

$$\text{thus } \sum_{n=0}^{\infty} y_{n+2} t^n - 7 \sum_{n=0}^{\infty} y_{n+1} t^n + 10 \sum_{n=0}^{\infty} y_n t^n = 0$$

$$(y_2 + y_3 t + y_4 t^2 + \dots) - 7(y_1 + y_2 t + y_3 t^2 + \dots) + 10(y_0 + y_1 t + y_2 t^2 + \dots) = 0$$

$$\frac{1}{t^2} (y_2 + y_3 t + y_4 t^2 + \dots) - \frac{7}{t} (y_1 + y_2 t + y_3 t^2 + \dots) + 10Y(t) = 0$$

$$\frac{1}{t^2} [Y(t) - y_0 - y_1 t] - \frac{7}{t} [Y(t) - y_0] + 10Y(t) = 0$$

$$\frac{1}{t^2} [Y(t) - 3t] - \frac{7}{t} [Y(t)] + 10Y(t) = 0$$

$$Y(t) - 3t - 7tY(t) + 10t^2 Y(t) = 0$$

$$\begin{aligned} (10t^2 - 7t + 1)Y(t) &= 3t \\ Y(t) &= \frac{3t}{10t^2 - 7t + 1} \end{aligned}$$

$$Y(t) = \frac{3t}{(1-2t)(1-5t)} \quad \text{by Partial Fractions}$$

$$Y(t) = \frac{1}{1-5t} - \frac{1}{1-2t}$$

$$Y(t) = (1-5t)^{-1} - (1-2t)^{-1} = (1+5t+(5t)^2+\dots) - (1+2t+(2t)^2+\dots)$$

$$\sum_{n=0}^{\infty} y_n t^n = \sum_{n=0}^{\infty} (5t)^n - \sum_{n=0}^{\infty} (2t)^n$$

equating the coefficient
of the both sides

$$y_n = 5^n - 2^n$$

A
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- Q. Solve $y_{n+2} - y_{n+1} - y_n = 0$ with $y_0 = 0$, $y_1 = 1$ by the method of generating functions.

[Ans.]

$$y_n = -\frac{\sqrt{5}}{5} (-1)^n \left(\frac{\sqrt{5}-1}{2}\right)^n + \frac{\sqrt{5}}{5} \left(\frac{\sqrt{5}+1}{2}\right)^n$$

- Q. Apply the generating function method to solve the following linear difference eqn

$$y_{n+2} - 2y_{n+1} + y_n = 2^n \quad \text{with } y_0 = 2 \text{ and } y_1 = 1$$

[Ans.] $y_n = 3 - 2 \cdot (n+1) + 2^n$



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DETAILED LECTURE NOTES

Binomial Theorem

PAGE NO.

It states that for any positive integer n we have

$$(x+a)^n = n_0 x^n + n_1 x^{n-1} a + n_2 x^{n-2} a^2 + \dots + n_n a^n \\ = \sum_{r=0}^n n_r x^{n-r} \cdot a^r$$

$$(x+a)^n = \sum_{r=0}^n n_r x^{n-r} a^r$$

Binomial coefficient

n_r also written as $C(n, r)$ or $\binom{n}{r}$ is called
Binomial coefficients where

$$n_r = C(n, r) = \binom{n}{r} = \frac{l^m}{r! (l-r)!}$$

Remarks :- when n is a positive integer, we have

$$\text{(i)} \quad (1+x)^n = \sum_{r=0}^n n_r x^r = \sum_{r=0}^n C(n, r) x^r \\ = C(n, 0) + C(n, 1)x + C(n, 2)x^2 + \dots + C(n, r)x^r \\ + C(n, n)x^n$$

here $C(n, r)$ is the coefficient of x^r in the expansion of $(1+x)^n$

$$\text{(ii)} \quad (1-x)^n = \sum_{r=0}^n n_r (-1)^r x^r$$

Q. If m is a fixed positive integer, then prove the following

- i) $C(n,0) + C(n,1) + C(n,2) + \dots + C(n,r) + \dots + C(n,m)$
 $= 2^m$ or $\sum_{k=0}^m C(n,k) = 2^m$
- ii) $C(m,0) + C(m,2) + C(m,4) + \dots = C(n,1) + C(n,3) + \dots = 2^{m-1}$
- iii) $C(m,1) + 2C(m,2) + 3C(m,3)$
 $\dots + iC(m,i) + \dots + mC(m,m) = m2^{m-1}$
- iv) $C(m,0) + 2C(m,1) + 2^2C(m,2) + \dots + 2^mC(m,m) = 3^m$
- v) $C^2(r,0) + C^2(r,1) + C^2(r,2) + \dots + C^2(r,i) + \dots + C^2(r,r) = C(2r,r)$

Sol:- we know

$$(1+x)^m = \sum_{r=0}^m {}^m C_r x^r = \sum_{r=0}^m C(m,r) x^r$$

ii) $(1+x)^m = C(m,0) + C(m,1)x + C(m,2)x^2 + \dots \dots \dots \quad \text{--- (1)}$

Put $x=1$ in eqn ①

$$2^m = C(m,0) + C(m,1) + C(m,2) + \dots \dots \quad \text{--- (ii)} \quad \text{By eqn 1}$$

iii) we know

$$(1-x)^m = \sum_{r=0}^m {}^m C_r (-1)^r x^r = \sum_{r=0}^m C(m,r) (-1)^r x^r$$

$$\Rightarrow C(m,0)x + C(m,1)x + C(m,2)x^2 - C(m,3)x^3 + C(m,4)x^4 - \dots$$

Put $x=+1$

$$+ C(m,0) + C(m,1) + C(m,2) - C(m,3) + C(m,4) - \dots = 0$$

$$C(m,0) + C(m,2) + C(m,4) + \dots = C(m,1) + C(m,3) + C(m,5) + \dots \quad \text{--- (iii)}$$

$$[C(m,0) + C(m,2) + C(m,4) + \dots] = 2^m$$

$$C(m,0) + C(m,2) + C(m,4) + \dots = \frac{2^m}{2} = 2^{m-1} \quad \text{--- (iv)}$$

From ③ $C(m,0) + C(m,2) + C(m,4) + \dots = C(m,1) + C(m,3) + \dots = 2^{m-1}$ A



POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

(iii)

$$C(m,1) + 2C(m,2) + 3C(m,3) + \dots = n c_1 + 2 n c_2 + 3 n c_3 + \dots$$

$$n + 2 \cancel{n(n-1)} + 3 \cancel{\frac{n(n-1)(n-2)}{2}} + \dots$$

$$n + n(n-1) + \cancel{2n(n-1)(n-2)} + \dots$$

$$n \left[1 + n-1 + \frac{(n-1)(n-2)}{2} + \dots \right]$$

$$= n \left[{}^{n-1}c_0 + {}^{n-1}c_1 + {}^{n-1}c_2 + \dots \right] = n[(1+1)^{n-1}]$$

$\stackrel{=} {n 2^{n-1}}$

(iv) we have

$$(1+x)^m = C(m,0) + C(m,1)x + C(m,2)x^2 + \dots$$

Put $x=2$

$$(1+2)^m = C(m,0) + 2C(m,1) + 4C(m,2) + \dots$$

$$3^m = C(m,0) + 2C(m,1) + 4C(m,2) + \dots$$

B-aus

A

✓ we have

$$r c_0 + r c_1 x + r c_2 x^2 + \dots + r c_{r-1} x^{r-1} + r c_r x^r = (1+x)^r \quad \textcircled{1}$$

$$\text{also } r c_0 x^r + r c_1 x^{r-1} + r c_2 x^{r-2} + \dots + r c_{r-1} x^2 + r c_r = (x+1)^r \quad \textcircled{2}$$

Multiplying $\textcircled{1}$ and $\textcircled{2}$

$$r c_0^2 + r c_1^2 + r c_2^2 + \dots + r c_r^2 = 2^r c_r$$

A

Multinomial coefficient

An expression of the form $x_1 + x_2 + \dots + x_r$, where x_1, x_2, \dots, x_r is called a Multinomial. We know that binomial coefficients occur when powers of a binomial (i.e. powers of $x_1 + x_2$) are expanded. In the similar way, when powers of a multinomial are expanded then multinomial coefficients occur.

Multinomial Theorem

It states that for any positive n , we have

$$(x_1 + x_2 + \dots + x_r)^n = \sum_{m_1 + m_2 + \dots + m_r = n} \binom{n}{m_1, m_2, \dots, m_r} x_1^{m_1} x_2^{m_2} \dots x_r^{m_r}$$

where $\binom{n}{m_1, m_2, \dots, m_r} = \frac{L^n}{L_{m_1} L_{m_2} \dots L_{m_r}}$

The quantity $\binom{n}{m_1, m_2, \dots, m_r}$, where $m_1 + m_2 + \dots + m_r = n$ is called Multinomial Coefficient.

Q. Find the coefficient of $x^5 y^3 z^2$ in $(x+y+z)^{10}$

Sol:- The coefficient of $x^5 y^3 z^2$ in $(x+y+z)^{10}$ is

$$= \binom{10}{5, 3, 2} = \frac{L^{10}}{L_5 L_3 L_2} = 2520$$

A

Q. Find the coefficient of $x y^2 z^2 u^2$ in $(x+y+z+u)^7$

Sol:- Required coefficient = $\frac{L^7}{L_1 L_2 L_2 L_2} = 630$

Q. Find the coefficient of $x^5 y^2 z^2$ in $(x+y+z)^9$

756



POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

Po sets

PAGE NO.

Partially ordered sets (Poset)

A Set P together with a Partial order relation R in it is called Partially ordered set or simply an ordered set (or poset) and its denoted by (P, R)

The Relation R is usually denoted by the symbol \leq , thus Partially ordered set may be denoted by (P, \leq)

Note:- A Relation R on a set P is called Partial order relation if

- R is Reflexive i.e. $aRa \forall a \in P$
- R is Anti-Symmetric: i.e. aRb and $bRa \Rightarrow a=b$ for same $a, b \in P$
- R is Transitive: i.e. aRb and $bRc \Rightarrow aRc$ for $a, b, c \in P$

C. Prove that the Relation "a divides b" If there exists an integer c such that $ac=b$ and is denoted by $a|b$ on the set of all positive integers N is a Partial order Relation.

Sol:- i) Reflexive :- $a \in N$ s.t aRa

$$a|a \Rightarrow a \cdot 1 = a \quad \exists \text{ An integer } 1.$$

∴ Hence the Relation " $|$ " is reflexive,

ii) Anti-Symmetric :- aRb , and $bRa \Rightarrow a=b$ for same $a, b \in N$

$$\frac{a}{b} = c_1 \quad \frac{b}{a} = c_2 \quad \therefore c_1 c_2 = 1$$

$$\Rightarrow c_1 = c_2 = 1 \Rightarrow a = b$$

Hence the Relation ' $|$ ' is anti-symmetric

A

Transitive:- let $a, b, c \in N$

$$\frac{a}{b} < \frac{b}{c} \text{ then } \frac{a}{b} = d_1, \quad \frac{b}{c} = d_2$$

$$ad_1 = b \quad bd_2 = c \Rightarrow ad_1d_2 = c$$

$\Rightarrow \exists$ an integer d such that $ad = c$ where $d = d_1d_2$
 $\Rightarrow a \mid c$

Hence the Relation ' \mid ' is Transitive.

Thus the relation ' \mid ' is a Partial order Relation on N .

**

Q. Let N be the set of Positive integers. Prove that the Relation \leq , where \leq has its usual meaning, is a Partial order Relation on N .

d. Let X be the set of all 2×2 real Matrices let

$$x, y \in X \quad x = \begin{bmatrix} x_1 & x_2 \\ x_3 & x_4 \end{bmatrix} \quad y = \begin{bmatrix} y_1 & y_2 \\ y_3 & y_4 \end{bmatrix}$$

and the Relation $x \leq y$ has the meaning

$$x_1 + x_2 + x_3 + x_4 \leq y_1 + y_2 + y_3 + y_4$$

Prove that \leq is not a Partial order relation on the set X . (i)

Sol:- the given Relation \leq is not Anti-symmetric on the set X
 because

$$x \leq y, y \leq x \Rightarrow x_1 + x_2 + x_3 + x_4 \leq y_1 + y_2 + y_3 + y_4$$

$$y_1 + y_2 + y_3 + y_4 \leq x_1 + x_2 + x_3 + x_4$$

$$\Rightarrow x_1 + x_2 + x_3 + x_4 = y_1 + y_2 + y_3 + y_4$$

$$\Rightarrow x \neq y$$

$$\text{For ex:- } x = \begin{bmatrix} 2 & 3 \\ 4 & -5 \end{bmatrix} \quad y = \begin{bmatrix} -2 & -3 \\ 8 & 1 \end{bmatrix} \in X$$

$$\text{where } x_1 + x_2 + x_3 + x_4 = 2 + 3 + 4 - 5 = 4$$

$$y_1 + y_2 + y_3 + y_4 = -2 - 3 + 8 + 1 = 4 ; \text{ Never } \begin{bmatrix} 2 & 3 \\ 4 & -5 \end{bmatrix} \neq \begin{bmatrix} -2 & -3 \\ 8 & 1 \end{bmatrix} = x \neq y$$



Comparability:-

Let (P, \leq) be a Poset. Two elements x and y of the Poset (P, \leq) are said to be Comparable if either $x \leq y$ or $y \leq x$, the elements x and y are called non-comparable.

If x and y are not related i.e. neither $x \leq y$ nor $y \leq x$ and is written as $x \nparallel y$.

Ex In the set of Natural Numbers N with the relation of divisibility. The elements 6 and 12 are comparable since 6 divides 12; but 2 and 5 are non-comparable since neither 2 divides 5 nor 5 divides 2.

Chain:-

Let (P, \leq) be Partially ordered set, A subset of P is called a chain if every two elements in the subset are related. In other words If for each pair of elements x, y in the subset either $x \leq y$ or $y \leq x$, then the subset of P is called a chain. The chain is also known as Totally ordered set or linearly ordered set.

Ex Let $\{x_1, x_2, \dots, x_m\}$ be a chain consisting of m elements (i.e. a finite number of elements). Because of antisymmetry and transitivity in the

Hasse Diagram

A Partial order relation \leq on a finite set P

can be represented by a diagram called the Hasse Diagram

let (P, \leq) be a Partially ordered set (i.e Poset)

let $a, b \in P$, then we write $a < b$ iff and only iff $a \leq b$ and $a \neq b$.

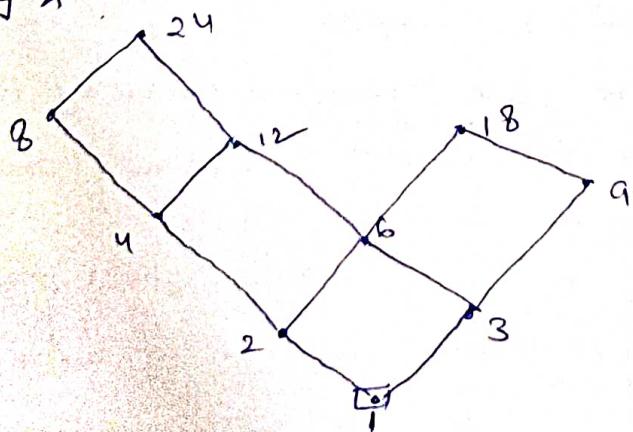
if $a < b$ then we say that b covers a iff $a \leq c \leq b \Rightarrow c = a$ or $c = b$.

In other words $b \in P$ is said to cover $a \in P$. If $a < b$, and no elements of P lies between a and b . i.e there does not exist $c \in P$, such that $a < c < b$. In this case b is also called Immediate Successor of a .

If b is an immediate successor of a , then b is taken at a higher level than a and a straight line is drawn to join a and b in the Hasse diagram

Ex let $A = \{1, 2, 3, 4, 6, 8, 9, 12, 18, 24\}$

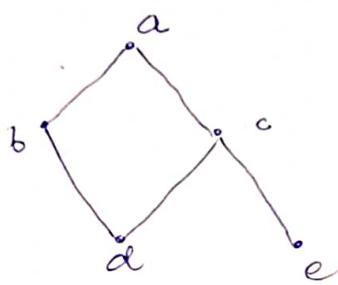
be ordered by the relation "a divides b" the Hasse diagram of A :





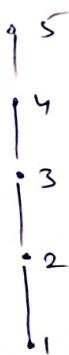
PAGE NO.

- Q. Let $B = \{a, b, c, d, e\}$ the Hasse diagram in the set B ,
and a Partial order on B s.t $d \leq b$ $d \leq a$, $e \leq c$ and so on.



- Q. the Hasse diagram of a finite linearly ordered set

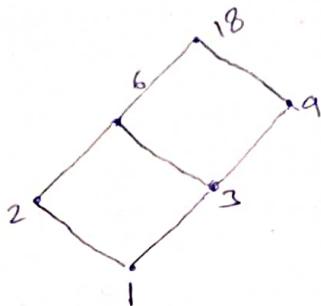
$P_1 = \{1, 2, 3, 4, 5\}$ with the usual relation \leq



- Q. The set D_{18} of Positive integral divisor of 18 with the
relation " | " divides . Find its Hasse diagram.

Soln:- D_{18} is the set of divisors of 18

$$D_{18} = \{1, 2, 3, 6, 9, 18\}$$

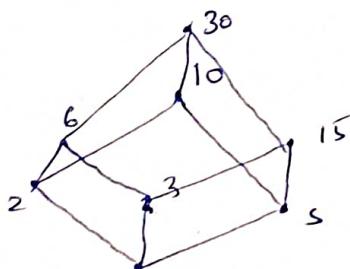


Q. The Hasse diagram of $\{3, 5, 7\}$ with the relation " $|$ " drawn below:-

$$\begin{matrix} & & \\ \bullet & \bullet & \bullet \\ 3 & 5 & 7 \end{matrix}$$

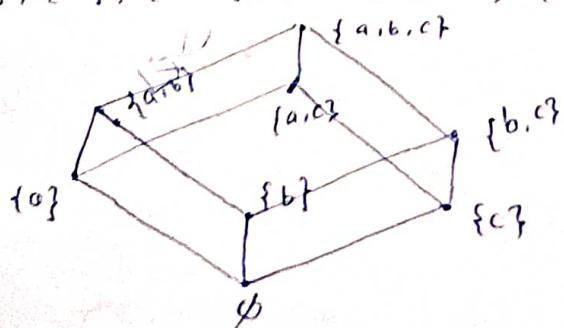
Q. The Hasse diagram of the set D_{30} of positive integral divisors of 30 with the relation " $|$ "

Soln:- $D_{30} = \{1, 2, 3, 5, 6, 10, 15, 30\}$



Q. Draw the Hasse diagram of the subsets of $S = \{a, b, c\}$ with the inclusion relation " \subseteq ".

$$P(S) = \{\{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}, \emptyset\}$$





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DETAILED LECTURE NOTES

PAGE NO.

Maximal element:-

let (P, \leq) be a Partially ordered set (Poset). An element m in P is said to be Maximal element if

$$x \leq m \quad \forall x \Rightarrow m = x \text{ (for any } x \in P)$$

obviously, a Maximal element is that element which is not smaller than any element.

Ex ① let (P, \leq) be a Partially ordered set, where $P = \{1, 2, 3, 4, 5\}$ and \leq is the relation of division which is Partially ordered set P . 3, 4, 5 are Maximal elements of P , because they do not divide the other.

Ex ② If $P = \{1, 2, 3, 4, 5\}$ is Partially ordered by the relation \leq where \leq has its usual meaning less than or equal to, then 5 is the only Maximal element in P because it is the biggest integer.

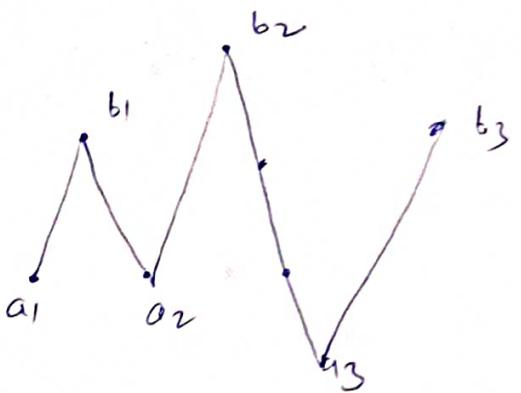
Minimal element:-

let (P, \leq) be a Partially ordered set. An element m in P is said to be a minimal element if for any $x \in P$

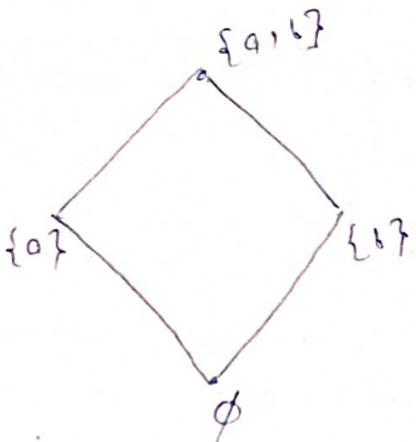
$$\text{s.t. } x : m \leq x \Rightarrow x = m$$

obviously a Minimal element is that element which is not greater than any element.

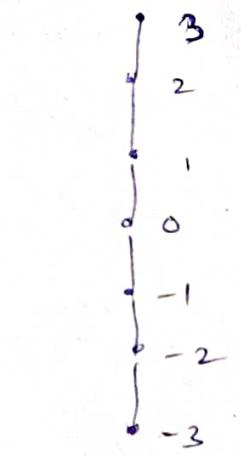
Ex. are Maximal and minimal elements



(a) a_1, a_2, a_3 — Minimal
 b_1, b_2, b_3 — Maximal



(b) ϕ is Minimal
 $\{a, b\}$ is Maximal



(c) ($I \subseteq$)
either
Maximal
or Minimal
element

Least Upper Bound (LUB)

Let (P, \leq) be a partially ordered set (poset) and A CP. If UGP is such an element that

$$a \leq u \quad \forall a \in A$$

Then u is called an upper bound of A .

If UGP is such that an element that

$$\text{i)} \quad a \leq u \quad \forall a \in A \quad \text{ii)} \quad u \leq v \quad \text{for all upper bounds } v \text{ of } A.$$

Then u is called the least upper bound or supremum of A and symbolically it is written as

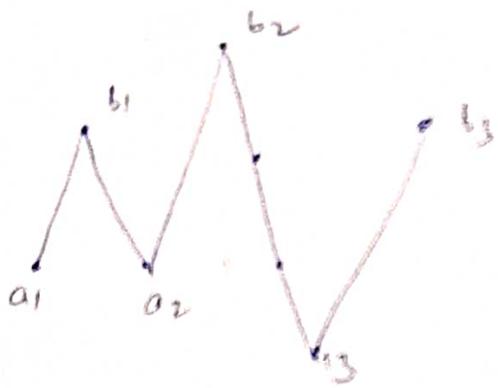
$$u = l.u.b.(A) \text{ or } u = \sup(A)$$

The supremum of a set $\{a, b\}$ of two elements is denoted by

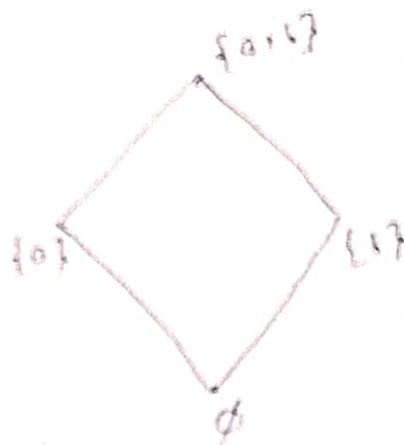
$a \vee b$ i.e

$$a \vee b = d.u.b.\{a, b\}$$

Ex. on Maximal and minimal elements



(a) a_1, a_2, a_3 — minimal
 b_1, b_2, b_3 — maximal



(b) ϕ is minimal
 $\{0,1\}$ is maximal



(c) $(I \leq)$
either
Maximal
or Minimal
element

Least Upper Bound (LUB)

Let (P, \leq) be a partially ordered set (Poset) and $A \subseteq P$. If UGP is such an element that

$$a \leq u \quad \forall a \in A$$

Then u is called an upper bound of A .

If UGP is such that an element that

- i) $a \leq u \quad \forall a \in A$
- ii) $u \leq v$ for all upper bounds v of A .

Then u is called the least upper bound or supremum of A and symbolically it is written as

$$u = l.u.b.(A) \text{ or } u = \sup(A)$$

The supremum of a set $\{a, b\}$ of two elements is denoted by

$a \vee b$ i.e.

$$a \vee b = d.u.b.\{a, b\}$$



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DETAILED LECTURE NOTES

PAGE NO.

Greatest Lower Bound (GLB)

Let (P, \leq) be a partially ordered set (Poset) and let $A \subset P$.

If $\exists p \in P$ is such that

$$l \leq a \quad \forall a \in A$$

Then l is called a lower bound of A .

If $\exists p \in P$ is such that an element that

i) $l \leq a \quad \forall a \in A$

ii) $l \leq b$ for all lower bounds of A then l is called the greatest lower bound or Infimum of A and symbolically it is written as

$$l = \text{glb}(A) \text{ or } l = \text{inf}(A)$$

thus Infimum of a set $\{a, b\}$ of two elements is denoted by $a \wedge b$ or

$$a \wedge b = \text{glb}\{a, b\}$$

Note:- The greatest lower bounds of a set A may be stated as:- If a lower bound of A succeeds every lower bound of A , then it is called the greatest lower bound of A .

LATTICE

let L be a non empty set closed under two binary operations called meet and join denoted respectively by \wedge and \vee .

Then L is called a lattice if the following axioms hold where a, b, c are any elements in L .

[L₁] Commutative law:-

$$(a \wedge b) = b \wedge a$$

$$(a \vee b) = b \vee a$$

[L₂] Associative law:-

$$(i) a \wedge (b \wedge c) = (a \wedge b) \wedge c \quad (ii) a \vee (b \vee c) = (a \vee b) \vee c$$

[L₃] Absorption law:-

$$(i) a \wedge (a \vee b) = a \quad (ii) a \vee (a \wedge b) = a$$

we shall denote the lattice by (L, \wedge, \vee)

Q let X be a collection of sets closed under intersection and union then prove that (X, \cap, \cup) is a lattice.

Sol :- [closure law:-] by hypothesis, X is closed under intersection and union of sets in X i.e. If $A, B \in X$ then

$$A, B \in X \Rightarrow A \cap B, A \cup B \in X \quad \forall A, B \in X$$

[L₁] Commutative law:- let $A, B \in X$, then by algebra of sets

$$(i) A \cap B = B \cap A \quad (ii) A \cup B = B \cup A \quad \forall A, B \in X$$

Hence intersection and union of sets in X follow commutative law.

[L₂] Associative law:- let $A, B, C \in X$

$$(i) (A \cap B) \cap C = A \cap (B \cap C) \quad \forall A, B, C \in X$$

$$(ii) (A \cup B) \cup C = A \cup (B \cup C)$$



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DETAILED LECTURE NOTES

PAGE NO.

Isomorphic (similar) ordered sets

Let (X, \leq) and (Y, \leq) be two partially ordered sets. A one-one onto (ie injective) function $f: X \rightarrow Y$ is said to be a similarity mapping from X onto Y if f preserves the order relation i.e. f satisfies the following conditions

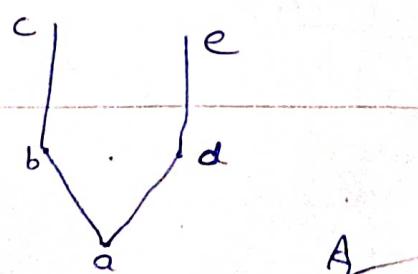
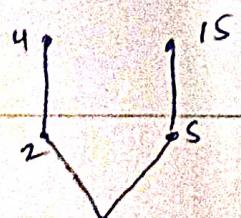
- i) If $a \leq a'$ then $f(a) \leq f(a')$ for every pair $a, a' \in X$
- and ii) If $a \neq a'$ then $f(a) \neq f(a')$ for every pair $a, a' \in X$

i.e. If $a, a' \in X$ are non-comparable then $f(a)$ and $f(a')$ are non-comparable.

If $f: X \rightarrow Y$ is one-one onto function and preserves the order relation then X and Y are called Isomorphic (or similar) ordered sets. and is written as $X \cong Y$.

Ex Let $X = \{1, 2, 4, 5, 15\}$ be an ordered set with the relation divisibility and let $Y = \{a, b, c, d, e\}$ be isomorphic to X .
Say the following function is similarly mapping from X onto Y

$$f = \{(1, a), (2, b), (4, c), (5, d), (15, e)\}$$



Well ordered sets

An ordered set S is called well-ordered if every subset of S has first element.

Ex- 1. The set N of natural numbers with the usual order \leq is well ordered, since every subset of N has a first element. Moreover N itself has a first element 1.

Ex 2. The set Z of all integers with the usual order \leq is not well ordered, since the set of negative integers, which was a subset of Z , has no least (first) element.

Dual statements :-

The dual of any statement in a lattice $(L \wedge V)$ is defined to be the statement that is obtained by interchanging \wedge and \vee for example the dual of $a \wedge (b \vee a) = a \vee a$ is $a \vee (b \wedge a) = a \wedge a$

Ex Theorem ① (Idempotent law)

In a lattice $(L \wedge V)$ for each $a \in L$

$$\text{i)} a \wedge a = a \quad \text{ii)} a \vee a = a$$

Proof:- let a be any arbitrary element of L then

(by Absorption law)

$$\text{i)} a \wedge a = a \wedge (a \vee (a \wedge b)) = a$$

$$a \wedge a = a \quad \forall a \in L$$

$$\text{ii)} a \vee a = a \vee (a \wedge (a \vee b)) \\ = a$$

$$a \vee a = a \quad \forall a \in L$$

A



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DETAILED LECTURE NOTES

PAGE NO.

PDNF and PCNF without Truth Table

Principal Disjunctive Normal form

P, Q → propositional s.

Minterms :- $P \wedge Q$ $\neg P \wedge Q$ $P \wedge \neg Q$

Maxterms :- $P \vee Q$ $P \vee \neg Q$ $\neg P \vee Q$ $\neg P \vee \neg Q$

PDNF :- Sum of minterms

DNF :- Sum of elementary product

PCNF :- Product of Maxterms

CNF :- Product of elementary sum

$$1. P \leftrightarrow Q \equiv (P \wedge Q) \vee (\neg P \wedge \neg Q) \text{ DNF}$$

$$P \leftrightarrow Q \equiv (P \rightarrow Q) \wedge (Q \rightarrow P) \text{ CNF}$$

$$P \rightarrow Q \equiv (\neg P \vee Q)$$

- 2) Introduce missing factors 3) Avoid duplicates.

Q) obtain the PDNF of the following

$$(i) P \vee \neg Q \quad (ii) P \rightarrow Q \quad (iii) Q \vee (P \wedge \neg Q)$$

$$\begin{aligned} \underline{\text{Sol:}} \quad (P \vee \neg Q) &= P \wedge (Q \vee \neg Q) \vee (\neg Q \wedge (P \vee \neg P)) \\ &= (P \wedge Q) \vee (P \wedge \neg Q) \vee (\neg Q \wedge P) \vee (\neg Q \wedge \neg P) \\ &= (P \wedge Q) \vee (P \wedge \neg Q) \vee (\neg P \wedge \neg Q) \quad \underline{\text{S}} \end{aligned}$$

$$P \vee P \in P$$

$$\begin{aligned}
 b) \quad P \rightarrow Q &= \neg P \vee Q \\
 &= [\neg P \wedge (\neg Q \vee Q)] \vee [Q \wedge (P \vee \neg P)] \\
 &= (\neg P \wedge Q) \vee (\neg P \vee \neg Q) \vee (Q \wedge P) \vee (Q \wedge \neg P) \\
 &= (\neg P \wedge Q) \vee (\neg P \vee \neg Q) \vee (P \wedge Q) \quad \underline{\text{A}}
 \end{aligned}$$

(d) $(P \wedge Q) \vee (\neg P \wedge \neg R) \vee (Q \wedge R)$

Sol:-

$$\begin{aligned}
 &[(P \wedge Q) \wedge (\neg R \vee R)] \vee [(\neg P \wedge R) \wedge (\neg Q \vee Q)] \vee [(Q \wedge R) \wedge (\neg P \wedge \neg P)] \\
 &[(P \wedge Q \wedge R) \vee (P \wedge Q \wedge \neg R)] \vee [(\neg P \wedge R \wedge \neg Q) \vee (\neg P \wedge R \wedge Q)] \vee [(Q \wedge R \wedge \neg P) \\
 &\quad \vee (Q \wedge R \wedge P)] \\
 &[(P \wedge Q \wedge R) \vee (P \wedge Q \wedge \neg R) \vee (\neg P \wedge R \wedge \neg Q) \vee (\neg P \wedge R \wedge Q) \vee (Q \wedge R \wedge \neg P) \\
 &\quad \vee (Q \wedge R \wedge P)] \quad \underline{\text{A}}
 \end{aligned}$$

Q. obtain the PCNF of the following

(a) $P \wedge Q$ (b) $P \leftarrow Q$ (c) $P \wedge Q$ (d) $(\neg P \rightarrow R) \wedge (Q \leftarrow P)$

(a) $P \wedge Q \equiv [P \vee (Q \wedge \neg Q)] \wedge [\neg Q \vee (P \wedge \neg P)]$

$$\begin{aligned}
 &= (P \vee Q) \wedge (P \vee \neg Q) \wedge (\neg Q \vee P) \wedge (\neg Q \vee \neg P) \\
 &= (P \vee Q) \wedge (P \vee \neg Q) \wedge (\neg P \vee \neg Q) \quad \underline{\text{A}}
 \end{aligned}$$

(b) $P \leftarrow Q \equiv (P \rightarrow Q) \wedge (Q \rightarrow P)$

$$\begin{aligned}
 &\equiv (\neg P \vee Q) \wedge (\neg Q \vee P) \quad \underline{\text{A}}
 \end{aligned}$$

(c) $P \wedge Q \equiv [P \vee (Q \wedge \neg Q)] \wedge [Q \vee (\neg P \wedge P)]$

$$\begin{aligned}
 &= (P \vee Q) \wedge (P \vee \neg Q) \wedge [(\neg Q \vee \neg P) \wedge (Q \vee P)] \\
 &= (P \vee Q) \wedge (P \vee \neg Q) \wedge (\neg P \vee Q) \quad \underline{\text{A}}
 \end{aligned}$$



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COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

Partial order in a Lattice:-

Let L be a lattice. Then we define a Partial order \leq on L . as follows

$$a \leq b \text{ if } a \wedge b = a$$

Analogously we could define

$$a \leq b \text{ if } a \vee b = b$$

Theorem ②

Let L be a lattice. Then

i) $a \wedge b = a$ iff $a \vee b = b$

ii) The Relation $a \leq b$ (defined by $a \wedge b = a$ or $a \vee b = b$)
is a Partial order on L)

Proof:- Let (L, \wedge, \vee) be a lattice and $a, b \in L$ are arbitrary elements.

i) only if condition: Suppose $a \wedge b = a$ then

$$\begin{aligned} a \vee b &= (a \wedge b) \vee b && \therefore \underline{a \wedge b = a} \\ &= b \vee (a \wedge b) && \text{commutative law} \\ &= b \vee (b \wedge a) && \text{by Absorption law} \end{aligned}$$

$$a \vee b = b$$

Therefore $a \wedge b = a$ ————— (1)

ii) Condition :- Suppose $a \wedge b = b$ then

$$a \wedge b = b \text{ then}$$

$$a \wedge b = a \wedge (a \vee b)$$

$$= a$$

$\therefore a \vee b = b$
by Absorption law

Therefore $a \vee b = b \Rightarrow a \wedge b = a$ — (2)

From (1) and (2) we obtain

$$a \wedge b = a \Leftrightarrow a \vee b = b$$

$$a \wedge b = a \text{ if and only if } a \vee b = b$$

(ii) Suppose the relation \leq is defined as follows

$$a \leq b \text{ if } a \wedge b = a$$

Now we see that

i) Reflexivity:- For each $a \in L$ we have

$$a \wedge a = a \quad (\text{Idempotent law})$$

$$a \leq a \quad \forall a \in L$$

Hence the relation \leq is reflexive in L .

ii) Anti-Symmetry:- If $a, b \in L$ and $a \leq b, b \leq a$ then

$$a \leq b, b \leq a \Rightarrow a \wedge b = a, b \wedge a = b$$

$$\Rightarrow a = b$$

$\therefore [a \wedge b = b \wedge a]$

Hence the relation \leq is anti-symmetric in L .

iii) Transitivity:- If $a, b, c \in L$ are such that $a \leq b, b \leq c$, then

$$a \leq b, b \leq c \Rightarrow a \wedge b = a, b \wedge c = b \quad (\text{Replacing } b \text{ with } b \wedge c)$$

$$\Rightarrow a \wedge (b \wedge c) = a$$

(Associative law: using (2a))

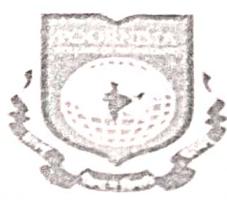
$$\Rightarrow (a \wedge b) \wedge c = a$$

$\therefore [a \wedge b = a]$

$$\Rightarrow a \wedge c = a$$

$$\Rightarrow a \leq c$$

Hence the relation \leq is Transitive in L



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COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

Lattice

A Poset $(A \leq)$ is called Lattice if every pair of elements a and b has both a least upper bound (LUB) and greatest lower bound (glb).

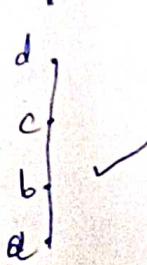
$$glb\{a, b\} = a \wedge b \quad \{ \text{meet of } a \wedge b \}$$

$$lub\{a, b\} = a \vee b \quad \{ \text{join of } a \vee b \}$$

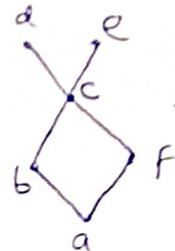
Every chain is a lattice.



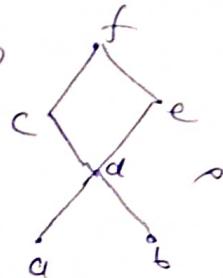
(i)



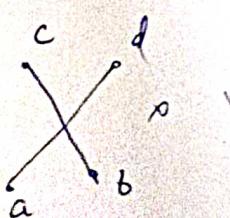
(ii)



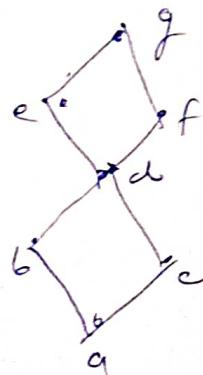
(iii)



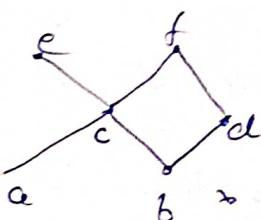
(iv)



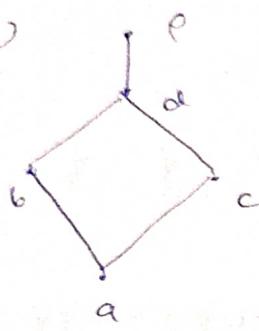
(v)



(vi)



Q. 2)



lub table

v	a	b	c	d	e
a					
b	a	b	c	d	e
c		b	b	d	e
d			c	c	e
e				d	e

glb table

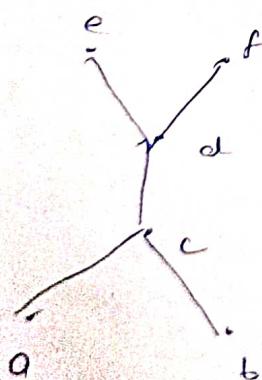
glb:-

v	a	b	c	d	e
a					
b	a	a	a	a	a
c		b	b	b	b
d			c	c	c
e				d	d

This Poset is a lattice.

A

(2)



v	a	b	c	d	e	f
a						
b	a					
c		b				
d			c			
e				d		
f					e	

glb:-

Since all the entries of lub and glb are not 0 or 1
of lub and glb are not complete.
This Poset is not a lattice

A

v	a	b	c	d	e	f
a	-	a	a	a	a	a
b	-	b	b	b	b	b
c	a	b	c	c	c	c
d	a	b	c	d	d	d
e	a	b	c	d	e	e
f	a	b	c	d	f	f



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COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

i) Idempotent :-

$$i) a \vee a = a \quad ii) a \wedge a = a$$

ii) Commutative :-

$$i) a \vee b = b \vee a \quad ii) a \wedge b = b \wedge a$$

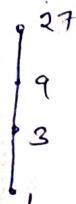
iii) Associative law

$$i) a \vee (b \vee c) = (a \vee b) \vee c \quad ii) a \wedge (b \wedge c) = (a \wedge b) \wedge c$$

iv) Absorption law :-

$$i) a \vee (a \wedge b) = a \quad ii) a \wedge (a \vee b) = a$$

check given Hasse diagram is a lattice or not.



i) Idempotent :-

$$a=3, \quad i) a \vee a = a$$

$$3 \vee 3 = \text{lub}\{3, 3\} = 3$$

$$3 \wedge 3 = \text{glb}\{3, 3\} = 3$$

ii) Commutative :-

$$\text{let } a=1, b=3$$

$$a \vee b = \text{lub}\{1, 3\} = 3$$

$$b \vee a = \text{lub}\{3, 1\} = 3$$

$$a \wedge b = \text{glb}\{1, 3\} = 1$$

$$b \wedge a = \text{glb}\{3, 1\} = 1$$

iii) Associative :- let $a=1, b=3, c=9$

$$(a \vee (b \vee c)) = a \vee (b \vee c) = 1 \vee (3 \vee 9) = 1 \vee 9 = 9$$

$$(a \vee b) \vee c = (1 \vee 3) \vee 9 = 3 \vee 9 = 9$$

(iv) Absorption law :- let $a=1, b=3$

$$a \vee (a \wedge b) = 1 \vee (1 \wedge 3) = 1 \vee 1 = 1 \quad \text{Because}$$

Hence given Hasse diagram is a lattice,

Types of Lattice :-

Sub Lattice

Let (L, \leq) be a Lattice. A subset M of L is said to be a Sub-Lattice of L if M is closed with respect to meet (\wedge) and join (\vee) which means that for each pair of elements $x, y \in M$, $x \wedge y$ and $x \vee y$ are contained in M .

In other words, a non empty subset M of a Lattice L is said to be a Sub-Lattice of L if M itself is a Lattice with respect to the binary operations meet (\wedge) and join (\vee) of L .

Ex let N be the set of positive integers and let there be defined a relation ' $|$ ' on N , where the meaning of $a|b$ is 'a divides b' ie there exists an integer n such that $b=na$, let D_m be the set of all divisors of m . Prove that

- (a) $(N, '|')$ is a lattice
- (b) D_m is a sub-lattice of the lattice N .

(b) let $D_m = \{x \in N : x|m\}$ that is D_m be the set of all divisors of the positive integer m

The binary relations meet (\wedge) and join (\vee) on D_m are defined as follows

$$x \wedge y = \text{H.C.F} \{x, y\} \quad x \vee y = \text{L.C.M} \{x, y\}$$

Then it is clear that for each pair of elements x, y in D_m

$$x, y \in D_m \Rightarrow x \wedge y, x \vee y \in D_m$$

That is D_m is closed with respect to the binary operations \wedge and \vee hence D_m is a sub-lattice of the lattice N .



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DETAILED LECTURE NOTES

PAGE NO.

Ex Let $L = \{1, 2, 3, 4, 6, 8, 9, 12, 18, 24\}$ be ordered by the relation ' $|$ ' where $x|y$ means ' x divides y '. Show that D_{24} the set of all divisors of the integer 24 of L is a sublattice of the lattice $(L, |)$.

Sol:- The given lattice is

$$L = \{1, 2, 3, 4, 6, 8, 9, 12, 18, 24\}$$

which is Partially ordered by the Relation ' $|$ '

The set of all divisors of $24 \in L$ is the set D_{24} given

$$\text{by } D_{24} = \{1, 2, 3, 4, 6, 8, 12, 24\}$$

Now we defined two binary operations Meet(\wedge) and Join(\vee) in D as follows.

$$x \wedge y = \text{H.C.F}\{x, y\} \quad x \vee y = \text{L.C.M}\{x, y\}$$

\wedge	1	2	3	4	6	8	12	24
1	1	1	1	1	1	1	1	1
2	1	2	1	2	2	2	2	2
3	1	1	3	1	3	1	3	3
4	1	2	1	4	2	4	4	4
6	1	2	3	2	6	2	6	6
8	1	4	1	4	2	8	4	8
12	1	2	3	4	6	4	12	12
24	1	2	3	4	6	8	12	24

\vee	1	2	3	4	6	8	12	24
1	1	2	3	4	6	8	12	24
2	2	2	6	4	6	8	12	24
3	3	6	3	12	6	24	12	24
4	4	4	12	4	12	8	12	24
6	6	6	6	12	6	24	12	24
8	8	8	24	8	24	8	24	24
12	12	12	12	12	12	24	72	24
24	24	24	24	24	24	24	24	24

Since all 64-64 entries of Composition Tables of \wedge and \vee are elements of D_{24} and so for each pair of elements $x, y \in D_{24}$ we have $x \wedge y \in D_{24} \Rightarrow x \wedge y, x, y \in D_{24}$ Hence D_{24} is closed under \wedge and \vee .

Hence (D_{24}, \leq) is a sublattice of the lattice (L, \leq)

Complete order

A linear order \leq on a set X is called complete if every non-empty subset of X which is bounded above have a supremum in X (which need not belong to that subset).

Ex every well order is complete.

Complete Lattice

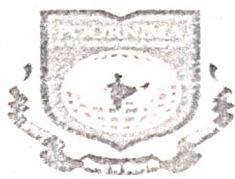
let (L, \leq) be a lattice. Then L is said to be complete if every subset A (finite or infinite) of L , $\wedge A$ and $\vee A$ exist in L . Thus in every complete lattice (L, \leq) there exist a greatest element g and a least element f .

for further discussion we shall use $a \wedge b$ in place of $\inf\{a, b\}$ and $a \vee b$ in place of $\sup\{a, b\}$ for $a, b \in L$

Distributive lattice :-

$$i) a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

$$ii) a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$



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DETAILED LECTURE NOTES

Permutation

PAGE NO.

Suppose an event E can occur in m different ways and associated with each way of occurring of E another event F can occur in n different ways. Then the Total Number of occurrence of the two events in the given order is $m \times n$.

Addititonal Principle If an event E can occur in m ways and another event F can occur in n ways and suppose that both cannot occur together, then E or F can occur in $m+n$ ways.

Permutation :-

A Permutation is an arrangement of objects in a definite order.

Permutation of n different objects. The Number of Permutations of n objects Taken all at a Time denoted by the Symbol $n P_m$ is given by

$$\boxed{n P_m = L^n}$$

$$L^n = n(n-1)(n-2) \dots 3 \cdot 2 \cdot 1$$

The number of Permutations of n objects Taken r at a time w
o or $\leq m$ denoted by ${}^m P_r$ is given by m .

$${}^m P_r = \frac{L^m}{L^{n-r}} \quad \text{we assume that } L^0 = 1$$

when repetition of objects is allowed the number of Permutations of n things Taken all at a time when repetition of objects is allowed is m^n .

Q. ① Find the Number of different Permutations of the set {1, 2}

Sol:- the set {1, 2} has two different elements 1 and 2

Thus there are $2!$ i.e. 2 permutations
These permutations are 12 and 21.

Ex Find the Number of different Permutations of the set {1, 2, 3}

Sol:- Total Number of Permutations = $3P_3 = \frac{L^3}{L^0} = 6$,

Clearly the permutations are

123, 132, 312 321, 213, 231,

A

Permutations with Repetition

Total Number of Permutations when out of m Objects in a set.

P. Objects are exactly alike of one kind (i.e P Objects are repeated)

Q. Objects are exactly alike of second kind, r Objects are alike of third kind, and the remaining $m - (P+Q+r)$ are dissimilar, i.e. $= \frac{L^m}{L^P, L^Q, L^r}$



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DETAILED LECTURE NOTES

PAGE NO.

Q How many numbers can be formed taking any 3 out of the 5 digits 3, 5, 8, 6, 7 ?

Sol:- clearly each arrangement of 3 digits will give us a new number.

$$\therefore \text{Required Number} = 5P_3 = \frac{15}{12} = 5 \times 4 \times 3 = 60$$

Q In how many ways can the letters of the word 'ASPERITY' be arranged. Also find the number of ways so that the vowels may never be separated.

Sol:- The word 'ASPERITY' contains 8 distinct letters

i) Total number of Arrangements = $8P_8 = 18 = 40320$

ii) If the vowels are never separated then let us put vowels within brackets (AEI), then we have only 6 things namely S,P,R,T,Y, (AEI) to arrange.

The Number of arrangement of the letters in which vowels stand together in the order (AEI)

$$= 6P_6 = 16 = 720$$

Q. In how many ways can the letters of the word 'ALLAHABAD' be arranged?

Sol: - The word 'ALLAHABAD' has 4 A's, 2 L's and three different letters namely H, B and D.

Hence the Total Number of Arrangements of the nine letters

$$= \frac{9!}{4! \cdot 2!} = \frac{9 \times 8 \times 7 \times 6 \times 5}{2} = 7560$$

A
S

Circular Permutations:

The permutations discussed above may be called Linear Permutations since the objects are being arranged in a line. If instead, we arrange objects in a circle or in a simple closed curve, then the permutations so obtained are called Circular Permutations.

To obtain number of circular permutations when n objects are placed in a circle, we fix one object and arrange remaining $(n-1)$ objects. Thus the Total Number of circular permutations = $(n-1)!$

Ex In how many ways ~~in which~~ can 9 persons be seated at a round Table?

Sol: - Total Number of ways in which one of the nine persons occupy a fixed position = $8! = 40320$

A

Q. If $2m+1 P_{m-1} : 2n-1 P_m = 3:5$ find m.



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DETAILED LECTURE NOTES

PAGE NO.

Combination

An unordered arrangement or unordered selection of objects from a finite set of objects is called a combination.

From from the definition it follows that the order of objects in a combination does not matter while in Permutation order of objects matters.

The number of Combinations (or Selections) of r ($r \leq n$) Objects taken from a set of n objects is denoted by ${}^n C_r$ or by $C(n, r)$ and is given by

$${}^n C_r = C(n, r) = \frac{L^n}{Lr \cdot L^{n-r}} \quad [\text{Repetition of objects is not allowed}]$$

Theorem:- To Prove that

$${}^n P_r = Lr \cdot {}^n C_r$$

$$\begin{aligned} LHS &= {}^n P_r = \frac{L^n}{L^{n-r}} = \frac{Lr \cdot L^{n-r}}{Lr \cdot L^{n-r}} = Lr \cdot {}^n C_r \\ &= RHS \end{aligned}$$

A

Theorem (2)

To Prove that

$${}^m C_r = {}^m C_{m-r}$$

Proof:-

$${}^m C_r = \frac{L^m}{L^r L^{m-r}} \rightarrow ①$$

$${}^m C_{m-r} = \frac{L^m}{L^{m-r} L^{m-m+r}} = \frac{L^m}{L^r L^{m-r}} = {}^m C_r \rightarrow ②$$

From ① and ②

$$\boxed{{}^m C_{m-r} = {}^m C_r}$$

∴

Theorem (3)

To Prove that

$${}^m C_r = {}^{m-1} C_r + {}^{m-1} C_{r-1}$$

Theorem (4) Let 'n' and 'r' be Positive integers with
 $r \leq m$ to Prove that

$${}^{m-1} C_r = {}^m C_r + {}^m C_{r-1} \quad (\text{Replace } n \rightarrow m-1 \text{ in Theorem 3})$$

degree of
difference
divided

3/2

between two
satisfies two
sides to be a
recurrence