

Tutorial - 4

Q1 $T(n) = 3T(n/2) + n^2$

$$T_n = aT(n/b) + f(n)$$

$$a = 3, b = 2$$

$$(a \geq 1 \text{ and } b > 1)$$

(Condition satisfied)

$$c = \log_b a = \log_2 3 = 1.58$$

$$n^c = n^2$$

case 3 satisfied : $f(n) > n^c$

$$T(n) = O(n^2)$$

2.

$$T(n) = 4T(n/2) + n^2$$

$$T(n) = aT(n/b) + f(n)$$

$$a = 4, b = 2$$

$$(a \geq 1, b > 1)$$

(satisfied)

$$c = \log_b a = \log_2 4 = 2 \log_2 2 = 2$$

$$n^c = n^2$$

case 2 satisfied

$$T(n) = O(n^2 \log n)$$

3. $T(n) = T(n/2) + 2^n$

$$T(n) = aT(n/b) + f(n)$$

$$a = 1, b = 2$$

$$a \geq 1, b > 1$$

(satisfied)

$$c = \log_b a = \log_2 1 = \log_2 2^0 = 0 \log_2 2 = 0$$

$$n^c = n^0 = 1$$

$$f(n) = 2^n$$

Case 3 satisfied $f(n) > n^c$

$$T(n) = \Theta(2^n)$$

4. $T(n) = 2^n T(n/2) + n^n$

$$T(n) = a T(n/b) + f(n)$$

$$a = 2^n$$

$$a \geq 1 \cdot 2^0 = 1$$

$$b = 2$$

$$b > 1$$

(satisfied)

$$c = \log_b a = \log_2 2^n = n$$

$$n^c = n^n$$

Case 2 satisfied because $f(n) = n^c$

$$T(n) = \Theta(n^n \log n)$$

5. $T(n) = 16 T(n/4) + n$

$$T(n) = a T(n/b) + f(n)$$

$$a = 16$$

$$(a \geq 1)$$

$$b = 4$$

$$(b > 1)$$

$$c = \log_b a = \log_4 16 = \log_4 4^2 = 2$$

$$n^c = n^2$$

Case 1 satisfied $f(n) < n^c$

$$T(n) = \Theta(n^2)$$

6. $T(n) = 2 T(n/2) + n \log n$

$$T(n) = a T(n/b) + f(n)$$

$$a = 2, b = 2$$

$$a \geq 1, b > 1$$

$$c = \log_b a = \log_2 2 = 1$$

$$n^c = n^1 = n$$

Case 3: $f(n) > n^c$

$$T(n) = \Theta(n \log n)$$

7. $T(n) = 2 T(n/2) + \frac{n}{\log n}$

$$T(n) = a T(n/b) + f(n)$$

$$a = 2, b = 2$$

$$(a \geq 1, b > 1)$$

$$c = \log_b a = \log_2 2 = 1$$

$$n^c = n^1 = n$$

Case 1: $f(n) < n^c$

$$T(n) = \Theta(n)$$

8. $T(n) = 2 T(n/4) + n^{0.51}$

$$T(n) = a T(n/b) + f(n)$$

$$a = 2$$

$$b = 4$$

$$a \geq 1, b > 1$$

$$c = \log_b a = \log_4 2 = 0.5$$

$$n^c = n^{0.5}$$

Case 2 = $f(n) = n^c$
 $T(n) = O(n^{0.5} \log n)$

9. $T(n) = 0.5T(n/2) + 1/n$
 $T(n) = aT(n/b) + f(n)$
 $a = 0.5$ $b = 2$ $(a > 1)$
(Master's Theorem can't be applied)
(Condition not satisfied)

10. $T(n) = 16T(n/4) + n!$
 $T(n) = aT(n/b) + f(n)$
 $a = 16$ $b = 4$
 $c = \log_b a = \log_4 16 = \log_4 4^2 = 2$
 $n^c = n^2$
 $n = 5$; $n^2 = 25$
 $n! = 5 \times 4 \times 3 \times 2 \times 1 = 120$

$f(n) > n^c$ case (3)
 $T(n) = O(n!)$

11. $T(n) = 4T(n/2) + \log n$
 $T(n) = aT(n/b) + f(n)$
 $a = 4$ $b = 2$
 $c = \log_b a = \log_2 4 = \log_2 2^2 = 2$
 $n^c = n^2$

$f(n) < n^c = \text{case 1}$
 $T(n) = O(n^2)$

12. $T(n) = \sqrt{n}T(n/2) + \log n$
 $a = \sqrt{n}$ $b = 2$ $a < 1, b > 1$
 $c = \log_2 \sqrt{n}$
 $n^c = n^{\log_2 \sqrt{n}}$

13. $T(n) = 3T(n/2) + n$
 $a = 3, b = 2$ $a > 1, b > 1$
 $c = \log_b a = \log_2 3 = 1.58$
 $n^{1.58} = n^c$
 $f(n) < n^c$ case 1
 $T(n) = O(n^{1.58})$

14. $T(n) = 4T(n/2) + cn$
 $a = 4$ $b = 2$
 $c = \log_b a = \log_2 4 = 2$

$n^c = n^2$ $f(n) = 2n$
Case 1 : $f(n) < n^c$
 $T(n) = O(n^2 \log n)$

(15) $T(n) = 3T(n/3) + \sqrt{n}$
 $a=3, b=3, a>1, b>1$

$$c = \log_a b = \log_3 3 = 1$$

$$n^c = n^1 = n$$

case 1 $f(n) < n$

$$T(n) = \Theta(n)$$

(16) $T(n) = 3T(n/4) + n \log n$
 $a=3, b=4, a>1, b>1$

$$c = \log_a b = \log_4 3 = 0.79$$

$$n^c = n^{0.79}$$

For $n=5, 15^{0.79} = 3.56$

$$5 \times \log 5 = 3.49$$

For $n=20 = (20)^{0.79} = 10.66$

$$= 20 \log 20 = 26.0204$$

$$f(n) > n^c = \text{case 3}$$

$$T(n) = \Theta(n \log n)$$

(17) $T(n) = 3T(n/2) + \frac{n}{2}$

$$a=3, b=2$$

$$c = \log_a b = \log_2 3 = 1.58$$

$$n^c = n^{1.58}$$

$f(n) < n^c$, case 1.

$$T(n) = \Theta(n^{1.58})$$

(18) $T(n) = 6T(n/3) + n^2 \log n$

$$a=6, b=3, (a>1, b>1) \checkmark$$

$$c = \log_a b = \log_3 6 = 1.63$$

$$n^c = n^{1.63}$$

$f(n) > n^c$ (case 3)

$$T(n) = \Theta(n^2 \log n)$$

(19) $T(n) = 4T(n/2) + \frac{n}{\log n}$

$$a=4, b=2, a>1, b>1$$

$$c = \log_a b = \log_2 4 = 2$$

$$n^c = n^2$$

For $n=10 = 10^2 = 100$

$$\frac{n}{\log n} = \frac{100}{\log 10} = 10$$

$f(n) < n^c$ (case 1)
 $T(n) = \Theta\left(\frac{n^2}{\log n}\right)$

- 20 $T(n) = 64T(n/8) - n^2 \log n$
Master Method cannot be applied, because
it cannot do negative work

21 $T(n) = 7T(n/3) + n^2$
 $a = 7, b = 3$
 $c = \log_b a = \log_3 7 = 1.77$
 $n^c = n^{1.77}$
 $f(n) > n^c$ case 3.
 $T(n) = \Theta(n^2)$

22 $T(n) = T(n/2) + n(2 - \cos n)$
 $a = 1, b = 2$ $a \geq 1, b > 1$

$c = \log_b a = \log_2 1 = 0$

$n^c = n^0 = 1$

$f(n) > n^c$ (case 3)

$T(n) = \Theta(n(2 - \cos n))$