

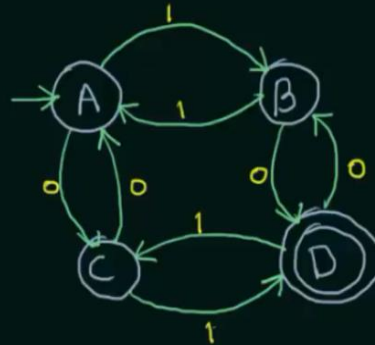
NFA - Non-deterministic Finite Automata

Deterministic Finite Automata



DETERMINISM

- >> In DFA, given the current state we know what the next state will be
- >> It has only one unique next state
- >> It has no choices or randomness
- >> It is simple and easy to design

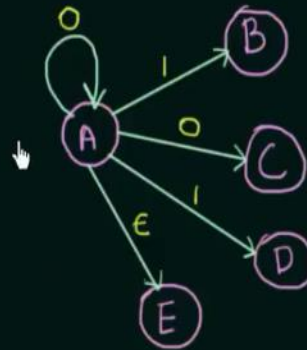


Non-deterministic Finite Automata

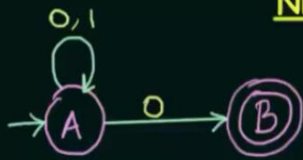


NON-DETERMINISM

- >> In NFA, given the current state there could be multiple next states
- >> The next state may be chosen at random
- >> All the next states may be chosen in parallel



NFA - Formal Definition



$L = \{ \text{Set of all strings that end with 0} \}$

$(Q, \Sigma, q_0, F, \delta)$

$Q =$ Set of all states

$\Sigma =$ inputs

$q_0 =$ start state / initial state

$F =$ Set of final states

$\delta = Q \times \Sigma \rightarrow \underline{\quad}$

- $\{A, B\}$

- $\{0, 1\}$

- A

- B

- ?

$A \times 0 \rightarrow A$

$A \times 0 \rightarrow B$

$A \times 1 \rightarrow A$

$B \times 0 \rightarrow \phi$

$B \times 1 \rightarrow \phi$

$A \xrightarrow{i} A, B, AB, \phi$

$(Q, \Sigma, q_0, F, \delta)$

$Q =$ Set of all states

$\Sigma =$ inputs

$q_0 =$ start state / initial state

$F =$ Set of final states

$\delta = Q \times \Sigma \rightarrow \underline{2^Q}$

- $\{A, B\}$

- $\{0, 1\}$

- A

- B

- ?

$A \times 0 \rightarrow A$

$A \times 0 \rightarrow B$

$A \times 1 \rightarrow A$

$B \times 0 \rightarrow \phi$

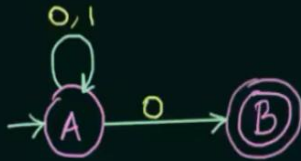
$B \times 1 \rightarrow \phi$

$A \xrightarrow{i} A, B, AB, \phi - 2^2 - 4$

3 states - A, B, C

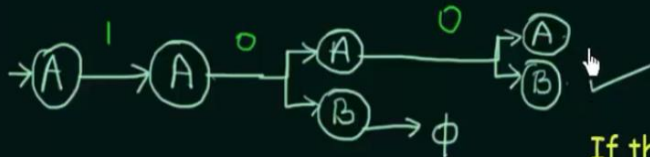
$A \xrightarrow{i} A, B, C, AB, AC, BC, ABC, \phi$
 $2^3 - 8$

NFA - Example-1

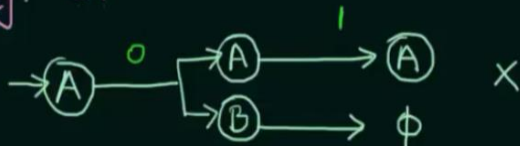


$L = \{ \text{Set of all strings that end with 0} \}$

Eg. 100



Eg. 01

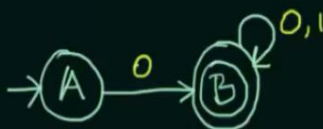


If there is any way to run the machine that ends in any set of states out of which atleast one state is a final state, then the NFA accepts

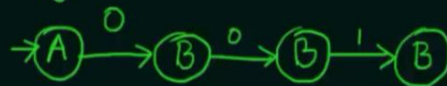
NFA - Example-2

$L = \{ \text{Set of all strings that start with 0} \}$

$= \{ 0, 00, 01, 000, \dots \}$



Eg. 001 ✓



Eg. 101 ✗



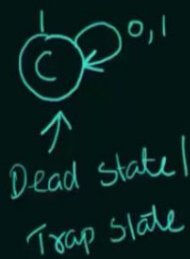
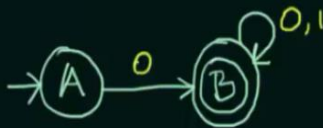
Dead configuration

[Pause]

NFA - Example-2

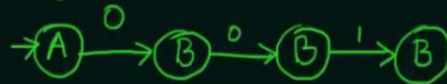
$L = \{ \text{Set of all strings that start with 0} \}$

$= \{ 0, 00, 01, 000, \dots \}$

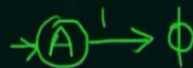


Dead state /
Trap state

Eg. 001 ✓



Eg. 101 X



Dead configuration

Eg. 101 X

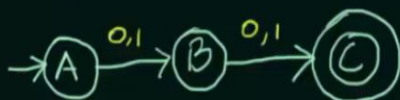


Dead configuration

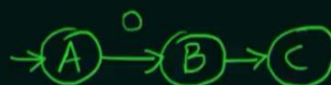
>> Construct a NFA that accepts sets of all strings over $\{0,1\}$ of length 2

$\Sigma = \{0,1\}$

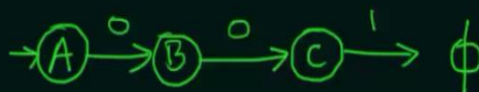
$L = \{00, 01, 10, 11\}$



Eg. 00 ✓

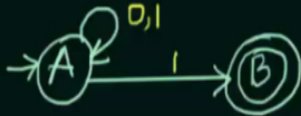


Eg. 001 X



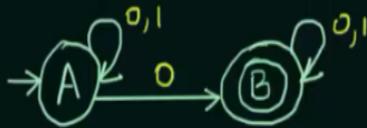
NFA - Example-3

Ex 1) $L_1 = \{ \text{Set of all strings that ends with '1'} \}$

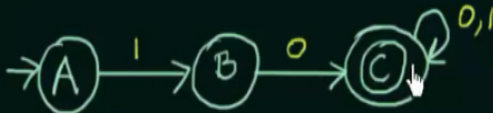


01, 001, 0001, 0^*1 , 1,
101, 1101,

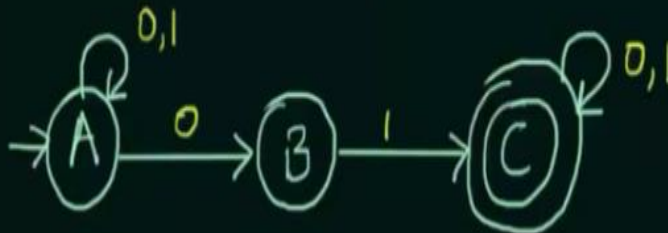
Ex 2) $L_2 = \{ \text{Set of all strings that contain '0'} \}$



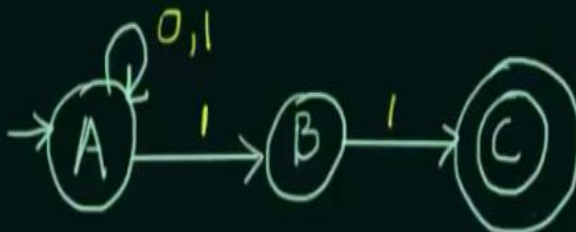
Ex 3) $L_3 = \{ \text{Set of all strings that starts with '10'} \}$



Ex 4) $L_4 = \{ \text{Set of all strings that contain '01'} \}$



Ex 5) $L_5 = \{ \text{Set of all strings that ends with '11'} \}$



Conversion of NFA to DFA

Every DFA is an NFA, but not vice versa

But there is an equivalent DFA for every NFA

DFA

$$\delta = \underline{Q \times \Sigma \rightarrow Q}$$

NFA

$$\delta = \underline{Q \times \Sigma \rightarrow 2^Q}$$

$$NFA \cong DFA$$

$L = \{ \text{Set of all strings over } (0,1) \text{ that starts with '0'} \}$

$$\Sigma = \{0,1\}$$

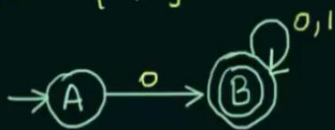


$$NFA \cong DFA$$

$L = \{ \text{Set of all strings over } (0,1) \text{ that starts with '0'} \}$

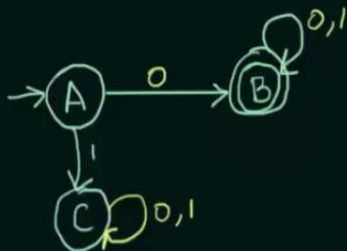
$$\Sigma = \{0,1\}$$

NFA



| | 0 | 1 |
|---|---|--------|
| A | B | ϕ |
| B | B | B |

DFA



| | 0 | 1 |
|---|---|---|
| A | B | C |
| B | B | B |
| C | C | C |

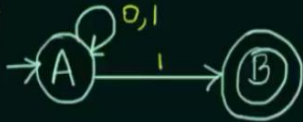
C - Dead state /
Trap state

Conversion of NFA to DFA - Examples (Part 1)

$L = \{ \text{Set of all strings over } (0,1) \text{ that ends with '1'} \}$

$\Sigma = 0, 1$

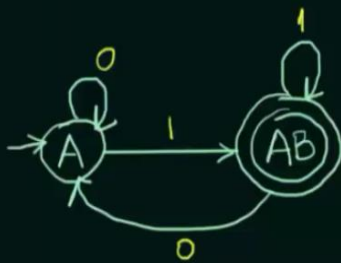
NFA



| | 0 | 1 |
|---|--------|--------|
| A | {A} | {A, B} |
| B | ϕ | ϕ |

Subset
construction
method

DFA



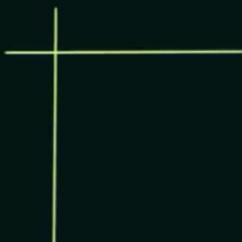
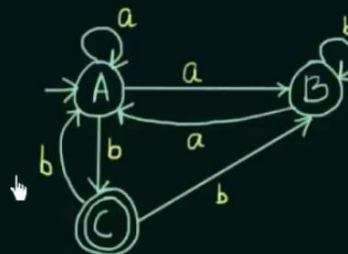
| | 0 | 1 |
|----|-----|------|
| A | {A} | {AB} |
| AB | {A} | {AB} |

AB - single
state

Conversion of NFA to DFA - Examples (Part-2)

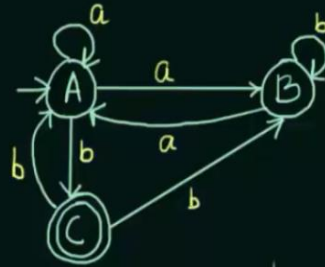
Find the equivalent DFA for the NFA given by $M = [\{A, B, C\}, (a, b), \delta, A, \{C\}]$ where δ is given by:

| | a | b |
|-----------------|------|------|
| $\rightarrow A$ | A, B | C |
| B | A | B |
| $\odot C$ | - | A, B |

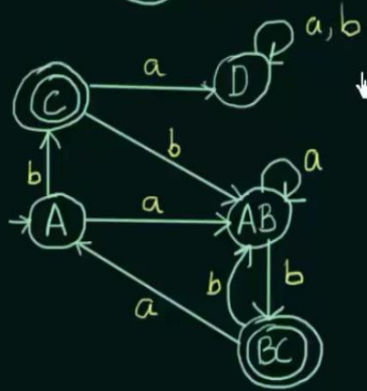


Find the equivalent DFA for the NFA given by $M = [\{A, B, C\}, \{a, b\}, \delta, A, \{C\}]$ where δ is given by:

| | a | b |
|-----------------|------|------|
| $\rightarrow A$ | A, B | C |
| B | A | B |
| $\odot C$ | - | A, B |



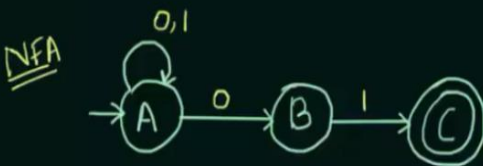
| | a | b |
|-----------------|----|----|
| $\rightarrow A$ | AB | C |
| AB | AB | BC |
| $\odot BC$ | A | AB |
| $\odot C$ | D | AB |
| D | D | D |



Conversion of NFA to DFA - Examples (Part-3)

Given below is the NFA for a language

$L = \{ \text{Set of all strings over } (0,1) \text{ that ends with '01'} \}$. Construct its equivalent DFA

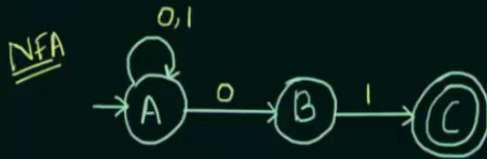


| | 0 | 1 |
|-----------------|--------|--------|
| $\rightarrow A$ | A, B | A |
| B | ϕ | C |
| $\odot C$ | ϕ | ϕ |

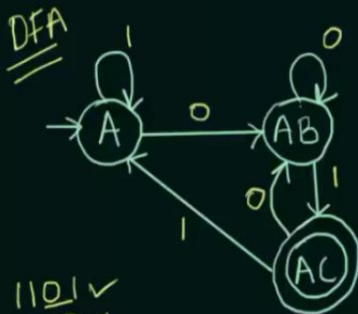
| | 0 | 1 |
|-----------------|----|----|
| $\rightarrow A$ | AB | A |
| AB | AB | AC |
| $\odot AC$ | AB | A |

Given below is the NFA for a language

$L = \{ \text{Set of all strings over } (0,1) \text{ that ends with '01'} \}$. Construct its equivalent DFA



| | 0 | 1 |
|-----|------|---|
| → A | A, B | A |
| B | φ | C |
| C | φ | φ |

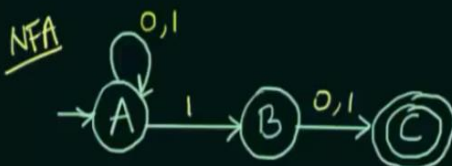


| | 0 | 1 |
|-----|----|----|
| → A | AB | A |
| AB | AB | AC |
| AC | AB | A |

[Pause]

Conversion of NFA to DFA - Examples (Part-4)

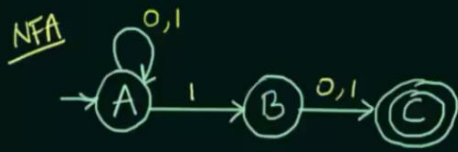
Design an NFA for a language that accepts all strings over $\{0,1\}$ in which the second last symbol is always '1'. Then convert it to its equivalent DFA.



| | 0 | 1 |
|-----|---|------|
| → A | A | A, B |
| B | C | C |
| C | φ | φ |

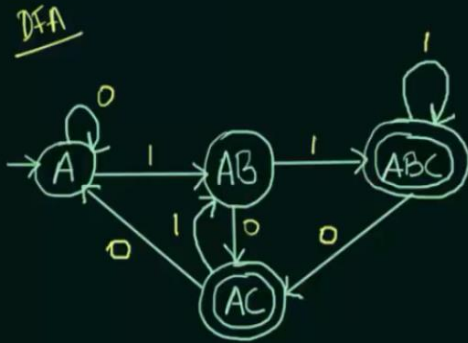
Eg. 1010
110
1101010

second last symbol is always 1. Then convert it to its equivalent DFA.



| | 0 | 1 |
|-----|---|------|
| → A | A | A, B |
| B | C | C |
| C | φ | φ |

Eg. 1010 ✓
 110 ✓
 1101010 ✓



| | 0 | 1 |
|-----|----|-----|
| → A | A | AB |
| AB | AC | ABC |
| AC | A | AB |
| ABC | AC | ABC |