

Complexity Analysis:

Divide:

We divide the process into two equal subarray. Thus if the original program takes $T(n)$ time then, the subdivided problem takes $T(n/2)$ amount of time.

Conquer:

In the conquer step we find the smallest distance between two subarray, this takes $O(1)$ time

Merge:

In the merge step we determine the smallest interval between two arrays. This needs sorting and a linear time searching. That means the complexity of this step is $O(n \ln n) + O(n)$ or $O(n \lg n)$

That means, the complexity of searching the smallest distance is $T(n) = 2T(n/2) + O(n \lg n)$

If we take the substitution approach then substituting the complexity $O(n \lg n)$

We can see that $T(n) = O(n \lg n)$

In the search of second smallest distance we used the algorithm three times, that means our time complexity becomes $O(n \lg n)$