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 NPTEL (<https://swayam.gov.in/explorer?ncCode=NPTEL>) » Design and analysis of algorithms (course)


## Course outline

How does an NPTEL online course work?

 Week 1 :  
Introduction

 Week 1 :  
Analysis of algorithms

Week 1 Quiz

 Week 2 :  
Searching and sorting

Week 2 Quiz

 Week 2  
Programming Assignment

Week 3 : Graphs

Week 3 Quiz

# Week 5 Quiz

The due date for submitting this assignment has passed.

Due on 2021-09-29, 23:59 IST.

Score: 10/10=100%

## Assignment submitted on 2021-09-29, 21:57 IST

All questions carry equal weightage. You may submit as many times as you like within the deadline. Your final submission will be graded.

1) Suppose we want to extend the union-find data structure to support the operation **2 points**  $\text{Reset}(c)$ , which takes as input the name of a component  $c$  and then breaks up  $c$  into singleton components, like  $\text{MakeUnionFind}()$ . For instance if  $c = 3$  and  $c$  currently consists of  $\{1, 3, 7\}$ , then  $\text{Reset}(c)$  will produce three components called 1, 3 and 7 consisting of  $\{1\}$ ,  $\{3\}$  and  $\{7\}$ , respectively.

Which of the following is correct about the cost of adding  $\text{Reset}(c)$  to the array and pointer implementations of union-find discussed in the lecture?

- ☐ Array representation:  $O(n)$ , Pointer representation:  $O(\text{size}(c))$
- ☐ Array representation:  $O(\text{size}(c))$ , Pointer representation:  $O(\text{size}(c))$
- ☒ Array representation:  $O(\text{size}(c))$ , Pointer representation:  $O(n)$
- ☐ Array representation:  $O(n)$ , Pointer representation:  $O(n)$

Yes, the answer is correct.

Score: 2

Feedback:

*In the array representation we have the list  $\text{Members}[c]$  which allows us to update the contents of  $c$  in time  $O(\text{size}(c))$ . In the pointer representation there is no easy way to identify all elements that belong to component  $c$  without scanning the entire set, so it takes time  $O(n)$*

Accepted Answers:

*Array representation:  $O(\text{size}(c))$ , Pointer representation:  $O(n)$*

### Week 3 Programming Assignment

### Week 4 : Weighted graphs

### Week 4 Quiz

### Week 4 Programming Assignment

### Week 5: Data Structures: Union-Find and Heaps

### Week 5 : Divide and Conquer

### Week 5 Quiz

#### ● Quiz: Week 5 Quiz (assessment? name=127)

### Week 6: Data Structures: Search Trees

### Week 6: Greedy Algorithms

### Week 6 Quiz

### Week 6 Programming Assignment

### Week 7: Dynamic Programming

### Week 7 Quiz

### Week 7 Programming Assignment

2) In the algorithm presented for the closest pair of points problem, we have assumed **2 points** that no two points have the same x or y coordinate. Which of the following steps would become more complicated to justify without this assumption.

- ☐ Arguing that every  $d/2$  side square in the d-band around the separator can have at most one point.
- ☐ Constructing  $S_Y$  from  $Q_Y$  and  $R_Y$  in time  $O(n)$  in the combine step.
- ☐ Constructing  $Q_X$  and  $R_X$  from  $P_X$  in time  $O(n)$  in the divide step.
- ☒ Constructing  $Q_Y$  and  $R_Y$  from  $P_Y$  in time  $O(n)$  in the divide step.

Yes, the answer is correct.

Score: 2

Feedback:

*All x-coordinates could be the same, in which case we may need  $O(n)$  time for each point in  $P_Y$  to assign it to  $Q_Y$  or  $R_Y$ .*

Accepted Answers:

*Constructing  $Q_Y$  and  $R_Y$  from  $P_Y$  in time  $O(n)$  in the divide step.*

3) Suppose we want to support the operations predecessor and successor in a heap. **2 points** Given a value  $v$  in the heap,  $\text{pred}(v)$  tells us the next smaller value currently in the heap and  $\text{succ}(v)$  tells us the next larger value currently in the heap.

- ☒ In both min heaps and max heaps, both operations take time  $O(n)$ .
- ☐ In both min heaps and max heaps, both operations take time  $O(\log n)$ .
- ☐ In a min heap,  $\text{pred}(v)$  takes time  $O(\log n)$  and  $\text{succ}(v)$  takes  $O(n)$  whereas in a max heap  $\text{pred}(v)$  takes time  $O(n)$  and  $\text{succ}(v)$  takes  $O(\log n)$ .
- ☐ In a min heap,  $\text{pred}(v)$  takes time  $O(n)$  and  $\text{succ}(v)$  takes  $O(\log n)$  whereas in a max heap  $\text{pred}(v)$  takes time  $O(\log n)$  and  $\text{succ}(v)$  takes  $O(n)$ .

Yes, the answer is correct.

Score: 2

Feedback:

*In either case, there is no obvious way to compute  $\text{pred}(v)$  and  $\text{succ}(v)$  without scanning all elements, so both cases take  $O(n)$ .*

Accepted Answers:

*In both min heaps and max heaps, both operations take time  $O(n)$ .*

4) Consider the min-heap [12, 27, 48, 30, 37, 79, 54, 43, 39] built by repeatedly inserting **2 points** values into an empty heap. Which of the following *could not* have been the last element inserted into this heap?

- ☐ 27
- ☐ 30
- ☒ 37
- ☐ 39

Yes, the answer is correct.

Score: 2

Feedback:

*The last position added was the one containing 39. The last element added must lie on the path from 39 to the root: {12, 27, 30, 39}.*

Accepted Answers:

**Week 8: Linear Programming and Network Flows****Week 8: Intractability****Week 8 Quiz****Text Transcripts****Books****Download Videos**

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5) Suppose we implement merge sort with a five-way split: divide the array into 5 equal parts, sort each part and do a 5 way merge. What would the worst-case complexity of this version be? **2 points**

- ☐  $O(n^2)$
- ☐  $O(n^2 \log_5 n)$
- ☒  $O(n \log_2 n)$
- ☐  $O(n (\log_2 n)^2)$

Yes, the answer is correct.

Score: 2

Feedback:

*The recurrence will yield  $O(n \log_5 n)$ , but  $\log_5 n = \log_2 n / \log_2 5 = O(\log_2 n)$ .*

Accepted Answers:

$O(n \log_2 n)$