toc: true title: 《从1到100深入学习Flink》—— Flink JobManager 有什么作用? date: 2019-02-25 tags:

- Flink
- 大数据
- 流式计算

前言

JobManager 是 flink 集群的中控节点,类似于 Apache Storm 的 Nimbus 以及 Apache Spark 的 Driver 的角色,它负责作业的调度、jar 包管理、checkpoint 的协调和发起等,为了后续章节的开展,本文将介绍 flink JobManager 中所部署的一些服务。

BolbServer

flink 用来管理二进制大文件的服务,flink JobManager 中启动的 BLOB Server 负责监听请求并派发线程去处理。更进一步,它将负责创建对应的目录结构去存储这些BLOBs 或者只是临时性地缓存。背后支持的文件系统:本底磁盘

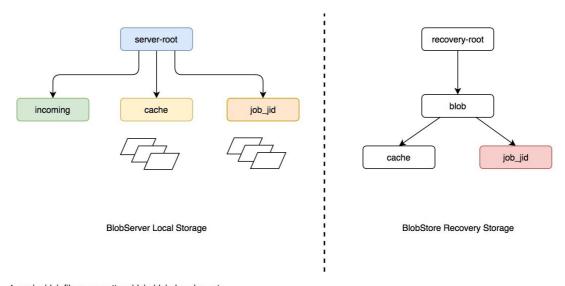
来看它的构造器:

- 第一步获取 RecoveryMode,一共两种 STANDALONE 和 ZOOKEEPER,后者 是有 JobManager leader 选举的高可用模式
- 获取文件系统存储的根目录,可配置,默认是从系统环境变量 System.getProperty("java.io.tmpdir") 中获取,其实就是本次磁盘存储
- 初始化 恢复存储 模块 BolbStore, STANDALONE 模式下为 VoidBlobStore, VoidBlobStore 是一个空实现;不会有任何持久化操作; ZOOKEEPER 模式下为 FileSystemBlobStore, FileSystemBlobStore 内部封装了磁盘文件的管理, 包括添加、删除、拷贝等, BlogStore 会备份 BlobServer 的本地存储,主要用于恢复模式下的作业磁盘状态恢复用
- 启动 ServerSocket
- 启动 BlobServer 服务线程

BlogServer 和 BlobStore

BlobStore 是 BlobServer 的组件之一,BolbStore 主要负责 BlobServer 本地存储的恢复【JobManager 重启】,这里只介绍 FileSystemBlobStore,FileSystemBlobStore 依据配置的不同支持两种文件系统存储:HDFS 和 本地文件系统

BlobServer 和 FileSystemBlobStore 的存储目录结构如下图所示:



1. cache blob file name pattern blob_blob_key_hex_str 2. job blob file name pattern blob_user_key_base64_str

下面以一次客户端连接请求的发起介绍两者的协同

来看 BolbServer 的核心 run 方法:

```
//BlobServer line230
public void run() {
   try {
      while (!this.shutdownRequested.get()) {
         BlobServerConnection conn = new
BlobServerConnection(serverSocket.accept(), this);
         try {
            synchronized (activeConnections) {
               while (activeConnections.size() >= maxConnections) {
                  activeConnections.wait(2000);
               activeConnections.add(conn);
            conn.start();
            conn = null;
         }
         finally {
            if (conn != null) {
               conn.close();
               synchronized (activeConnections) {
                  activeConnections.remove(conn);
            }
         }
     }
```

简要概括下逻辑:

- 当服务端收到一次存储的 request 时,会首先封装成对象 BlobServerConnection,并执行其 start()方法
- BlobServerConnection 本身也是一个 Thread, 封装了具体的存储逻辑
- 会接收 3 种客户端请求: PUT/GET/DELETE, 具体见:

```
//BlobServerConnection line111
switch (operation) {
  case PUT_OPERATION:
    put(inputStream, outputStream, buffer);
    break;
  case GET_OPERATION:
    get(inputStream, outputStream, buffer);
    break;
  case DELETE_OPERATION:
    delete(inputStream, outputStream, buffer);
    break;
  default:
    throw new IOException("Unknown operation " + operation);
}
```

这里重点介绍下 PUT 操作

- 获取本次存储操作是否带 JobID
- 在 BlobServer 的本地 incoming 文件夹中生成临时文件: temp-[auto increment integer]
- 读取将要存储的字节长度
- 读取该长度字节存储到临时文件 temp-[auto increment integer]
- 如果带 JobID,会将临时文件移动到 JobID 对应的存储目录,并将该存储文件 在 BlobStore 的对应 JobID恢复目录中备份,写 OK 消息到 Socket Client 端, 最终生成的路径和文件: job-id/blob_[base64 encode key]
- 如果不带 JobID,则依据传递的消息字节数组生成一个 key: BlogKey,并存储在 cache 文件夹下,同时在 BlobStore 的 cache 文件夹下做备份,将 OK 消息和 BlobKey 写回 Socket Client,最终生成的路径和文件: cache/blob_[unique hex string]

BlobServer 交互协议

与 BlobServer 通信的消息协议包括四段:操作类型【PUT/GET/DELETE】、存储类型【是否带 JobID】、内容长度、内容,如下图所示:

0 put 1 get 2 delete	0 conter 1 name	nt	
0/1/2	0/1	length of content	content
1 byte	1 byte	4 byte	

InstanceManager

flink 用来追踪当前存活的 TaskManager 的管理组件,实现比较简单,这里只简单 罗列下其功能:

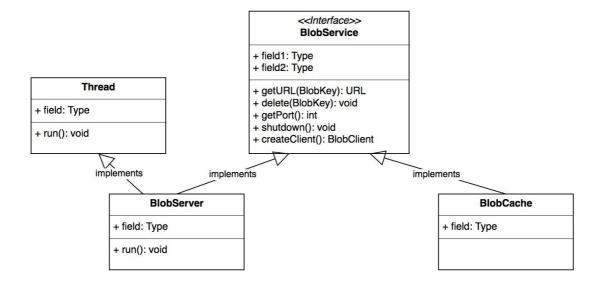
- book 下载 JobManager 中注册的所有 TaskManager
- 负责更新从 TaskManager 中上报的心跳及 metrics 信息
- 通知 InstanceListener TaskManager 的增加与死亡

BlobLibraryCacheManager

flink job 的 jar 包存储服务,使用上面的 BlobServer 完成,一个 JVM 里只会存在一个 BlobLibraryCacheManager,BlobLibraryCacheManager 负责管理 BlobService【这里为BlobServer】 中存储的 jars,并存储运行时 task 对 BlobService 中 jars 的引用计数,会清理不被使用任何 task 使用的 jars。

BlobCache 负责 jars 的下载,介绍 TaskManager 的时候会详细介绍

BlobLibraryCacheManager 与 BlobService 交互,而 BlobService 负责具体的文件管理,其具体实现有两个: BlobServer 和 BlobCache,具体见下图:



BlobServer 前面已经介绍过了,那么 BlobCache 的功能是什么呢?

来看 BlobCache 的构造器:

```
//BlobCache line60
public BlobCache(InetSocketAddress serverAddress, Configuration
configuration) {
   if (serverAddress == null || configuration == null) {
        throw new NullPointerException();
   }

   this.serverAddress = serverAddress;

   // configure and create the storage directory
   String storageDirectory =
configuration.getString(ConfigConstants.BLOB_STORAGE_DIRECTORY_KEY,
null);
   this.storageDir =
BlobUtils.initStorageDirectory(storageDirectory);
   LOG.info("Created BLOB cache storage directory " + storageDir);
```

这里传入的 serverAddress 其实是 BlobServer 的服务端口,在 TaskManager 中可以看到:

```
// start a blob service, if a blob server is specified TaskManager
line940
if (blobPort > 0) {
  val jmHost = jobManager.path.address.host.getOrElse("localhost")
  val address = new InetSocketAddress(jmHost, blobPort)

  log.info(s"Determined BLOB server address to be $address.
Starting BLOB cache.")

  try {
   val blobcache = new BlobCache(address, config.configuration)
   blobService = Option(blobcache)
   libraryCacheManager = Some(new
BlobLibraryCacheManager(blobcache, config.cleanupInterval))
  }
```

来看 BlobCache 的核心服务方法:

```
//BlobCache line97
```

```
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   if (requiredBlob == null) {
      throw new IllegalArgumentException("BLOB key cannot be
null.");
   }
   final File localJarFile =
BlobUtils.getStorageLocation(storageDir, requiredBlob);
   if (!localJarFile.exists()) {
      final byte[] buf = new byte[BlobServerProtocol.BUFFER_SIZE];
      // loop over retries
      int attempt = 0;
      while (true) {
         if (attempt == 0) {
            LOG.info("Downloading {} from {}", requiredBlob,
serverAddress);
         } else {
            LOG.info("Downloading {} from {} (retry {})",
requiredBlob, serverAddress, attempt);
         }
         try {
            BlobClient bc = null;
            InputStream is = null;
            OutputStream os = null;
            try {
               bc = new BlobClient(serverAddress);
               is = bc.get(requiredBlob);
               os = new FileOutputStream(localJarFile);
               while (true) {
                  final int read = is.read(buf);
                  if (read < 0) {
                     break;
                  }
                  os.write(buf, 0, read);
               // we do explicitly not use a finally block, because
we want the closing
              // in the regular case to throw exceptions and cause
the writing to fail.
               // But, the closing on exception should not throw
further exceptions and
```

```
// let us keep the root exception
os.close();
os = null;
is.close();
is = null;
bc.close();
bc = null;

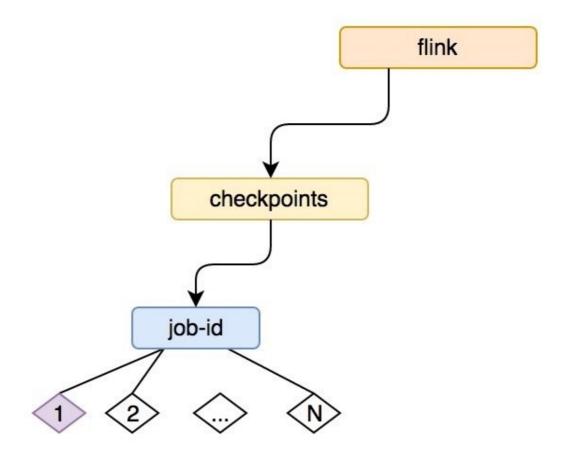
// success, we finished
break;
```

简要概括下其逻辑:

- 先从本地磁盘中获取,如果存在,直接返回
- 如果没有,生成 BlobClient 与 BlobServer 交互,并拉取文件到本地缓存,后 返回本地缓存的文件句柄

从这里我们可以看到 BlobCache 是 TaskManager 操作本地文件的工具,它负责从 JobManager 中的 BlobServer 同步所需的文件【jar包等】,而 BlobServer 和 BlobCache 的文件管理的入口,统一由对应 JVM 中的 BlobLibraryCacheManager 来控制【没有任务使用的 jar 定期清除等】。

task 拉取 jar包文件的过程如下:



ZooKeeperCompletedCheckpointStore

flink 做 checkpoint 【有关 checkpoint 会另起一节介绍】存储的组件,负责存储已完成的 Checkpoint ,实现了接口 CompletedCheckpointStore,

StandaloneCompletedCheckpointStore 和

ZooKeeperCompletedCheckpointStore 都实现了 CompletedCheckpointStore 接口,前者只在内存里存储 checkpoint,这里只介绍

ZooKeeperCompletedCheckpointStore 的实现。

ZooKeeperCompletedCheckpointStore 存储 checkpoint 的基本思路:

- 先在本地磁盘持久化指定数量的 checkpoint
- 将文件句柄更新到 ZK 的特定节点下
- 滑动更新 zk 的节点存储
- 在恢复的时候只取最近一次的更新值

先来看下 ZooKeeperCompletedCheckpointStore 用来和 ZK 存储交互的组件: ZooKeeperStateHandleStore,来看它的核心添加 state 的方法:

```
//ZooKeeperStateHandleStore line117
public StateHandle<T> add(
      String pathInZooKeeper,
      T state,
      CreateMode createMode) throws Exception {
   checkNotNull(pathInZooKeeper, "Path in ZooKeeper");
   checkNotNull(state, "State");
   StateHandle<T> stateHandle = storage.store(state);
   boolean success = false;
   try {
     // Serialize the state handle. This writes the state to the
backend.
      byte[] serializedStateHandle =
InstantiationUtil.serializeObject(stateHandle);
     // Write state handle (not the actual state) to ZooKeeper.
This is expected to be
     // smaller than the state itself. This level of indirection
makes sure that data in
     // ZooKeeper is small, because ZooKeeper is designed for data
in the KB range, but
     // the state can be larger.
      client.create().withMode(createMode).forPath(pathInZooKeeper,
serializedStateHandle);
      success = true;
     return stateHandle;
   }
   finally {
      if (!success) {
        // Cleanup the state handle if it was not written to
ZooKeeper.
         if (stateHandle != null) {
            stateHandle.discardState();
      }
   }
}
```

- 使用 StateStorageHelper 存储 state, ZK 模式下为 FileSystemStateStorageHelper, 方式为直接存储到本地磁盘
- 将 state 的句柄对象 StateHandle 序列化并持久化到 ZK 的节点

其在 zk 上的存储路径如下图所示:

现在来看 ZooKeeperCompletedCheckpointStore 的核心功能:添加 checkpoint 和从 checkpoint 做 recovery

添加 checkpoint

```
//ZooKeeperCompletedCheckpointStore line190
public void addCheckpoint(CompletedCheckpoint checkpoint) throws
Exception {
   checkNotNull(checkpoint, "Checkpoint");
   // First add the new one. If it fails, we don't want to loose
existing data.
   String path = String.format("/%s",
checkpoint.getCheckpointID());
   final StateHandle<CompletedCheckpoint> stateHandle =
checkpointsInZooKeeper.add(path, checkpoint);
   checkpointStateHandles.addLast(new Tuple2<>(stateHandle, path));
   // Everything worked, let's remove a previous checkpoint if
necessary.
   if (checkpointStateHandles.size() >
maxNumberOfCheckpointsToRetain) {
removeFromZooKeeperAndDiscardCheckpoint(checkpointStateHandles.remo
veFirst());
   }
   LOG.debug("Added {} to {}.", checkpoint, path);
}
```

简要概括其逻辑:

- 在本地磁盘存储该 checkpoint 的内容并返回句柄对象: StateHandle
- 以 checkpoint id 在 zk 上新建一个 node,并存储对应的序列化后的 StateHandle
- 检查存储的 checkpoint 个数是否超过限制,如果超过,删除本地磁盘及zk上最 旧的数据
- 如果添加失败,已有的 checkpoint 数据不会受影响,这里 flink 想最大化保留 作业的 checkpoint

从 checkpoint 中恢复

```
//ZooKeeperCompletedCheckpointStore line137
public void recover() throws Exception {
   LOG.info("Recovering checkpoints from ZooKeeper.");
   // Clear local handles in order to prevent duplicates on
   // recovery. The local handles should reflect the state
   // of ZooKeeper.
   checkpointStateHandles.clear();
   // Get all there is first
   List<Tuple2<StateHandle<CompletedCheckpoint>, String>>
initialCheckpoints;
   while (true) {
      try {
         initialCheckpoints =
checkpointsInZooKeeper.getAllSortedByName();
         break;
      }
      catch (ConcurrentModificationException e) {
         LOG.warn("Concurrent modification while reading from
ZooKeeper. Retrying.");
      }
   }
   int numberOfInitialCheckpoints = initialCheckpoints.size();
   LOG.info("Found {} checkpoints in ZooKeeper.",
numberOfInitialCheckpoints);
   if (numberOfInitialCheckpoints > 0) {
      // Take the last one. This is the latest checkpoints, because
path names are strictly
      // increasing (checkpoint ID).
      Tuple2<StateHandle<CompletedCheckpoint>, String> latest =
initial Checknoints
```

```
.get(numberOfInitialCheckpoints - 1);

CompletedCheckpoint latestCheckpoint =
latest.f0.getState(userClassLoader);

checkpointStateHandles.add(latest);

LOG.info("Initialized with {}. Removing all older
checkpoints.", latestCheckpoint);

for (int i = 0; i < numberOfInitialCheckpoints - 1; i++) {
    try {

removeFromZooKeeperAndDiscardCheckpoint(initialCheckpoints.get(i));
    }
    catch (Exception e) {
        LOG.error("Failed to discard checkpoint", e);
    }
    }
}</pre>
```

简要概括其逻辑:

- 清除内存中维护的句柄对象 StateHandle s
- 从 ZK 上拉取作业对应的所有的 checkpoint StateHandle 节点,并排序【从小 到大】
- 获取最新的一次快照并从本地磁盘恢复 checkpoint
- 删除其余所有的 checkpoint 信息【ZK 和本地磁盘】

ZooKeeperCompletedCheckpointStore 由

ZooKeeperCheckpointRecoveryFactory 负责实例化,一个 Job 会实例化一个 ZooKeeperCompletedCheckpointStore 负责快照。这里存储的只是个节点快照的 句柄,并不是真正的状态数据。

具体的启动流程见 JobManager

```
line1208 val completedCheckpoints =
checkpointRecoveryFactory.createCheckpointStore(jobId, userCodeLoader)
```

line1238 executionGraph.enableSnapshotCheckpointing



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